

## 1. INTRODUCTION

In [4, p. 80], Erdős wrote:

*I just want to state without proof a special result in this direction, namely*

$$\frac{\log \log n}{\log n}n + c_9 \frac{n}{(\log n)^2} \leq g_3(n) \leq \frac{\log \log n}{\log n}n + c_{10} \frac{n}{(\log n)^2}. \quad (1.1)$$

*It is not clear whether (1.1) can be sharpened.*

This leads to the following problem, which also appears as Problem #796 on Bloom's Erdős Problems website [1].

**Problem 1.1.** *Let  $k \geq 2$  and let  $g_k(n)$  be the largest possible size of  $A \subseteq \{1, \dots, n\}$  such that every  $m$  has  $< k$  solutions to  $m = a_1 a_2$  with  $a_1 < a_2 \in A$ .*

*Is it true that*

$$g_3(n) = \frac{\log \log n}{\log n}n + (c + o(1)) \frac{n}{(\log n)^2}$$

*for some constant  $c$ ?*

In this note, we show that (1.1) is false, thereby providing a counterexample to Problem 1.1. We also give a corrected version of the problem.

## 2. A LOWER BOUND

Fix  $n \geq 2$ . Then  $g_3(n)$  denote the largest cardinality of a set  $A \subseteq \{1, 2, \dots, n\}$  with the property that for every integer  $m$  the equation

$$m = a_1 a_2, \quad a_1 < a_2, \quad a_1, a_2 \in A,$$

has at most 2 solutions.

Let  $M$  be the Meissel–Mertens constant. Then we have the following lower bound.

**Proposition 2.1.** *For all sufficiently large  $n$ ,*

$$g_3(n) \geq \frac{n \log \log n}{\log n} + (M + 1 + o(1)) \frac{n}{\log n}.$$

*Proof.* Let

$$P := \{p \leq n : p \text{ is prime}\}, \quad S := \{pq \leq n : p, q \text{ primes and } q > \sqrt{n}\},$$

and define

$$A := P \cup S \subseteq \{1, 2, \dots, n\}.$$

**1.  $A$  has at most two multiplicative representations.** Call a prime *large* if it exceeds  $\sqrt{n}$ .

**Claim 1.** *Every element  $a \in A$  has at most one large prime divisor (counted without multiplicity).*

*Proof of Claim 1.* If  $a \in P$  then  $a$  is prime, so the claim is trivial. If  $a \in S$ , then  $a = pq$  with  $q > \sqrt{n}$  prime and  $p \leq n/q < \sqrt{n}$ , so  $q$  is the unique large prime divisor of  $a$ .  $\square$

**Claim 2.** *For every integer  $m$ , the number of solutions of  $m = a_1 a_2$  with  $a_1 < a_2$  and  $a_1, a_2 \in A$  is at most 2.*

*Proof of Claim 2.* Fix  $m$  and suppose  $m = a_1 a_2$  with  $a_1 < a_2$  and  $a_1, a_2 \in A$ . By Claim 1, each  $a_i$  contributes at most one large prime divisor, so  $m$  has at most two large prime divisors (counted with multiplicity).

*Case 0:  $m$  has no large prime divisor.* Then neither  $a_1$  nor  $a_2$  can lie in  $S$ , hence  $a_1, a_2 \in P$  and  $a_1, a_2 \leq \sqrt{n}$ . Thus  $m$  is the product of two primes and the representation is unique up to order; with the constraint  $a_1 < a_2$ , there is at most one solution.

*Case 1:  $m$  has exactly one large prime divisor  $q > \sqrt{n}$ .* Then exactly one of  $a_1, a_2$  is divisible by  $q$ . The factor divisible by  $q$  is either  $q \in P$  or  $pq \in S$  with  $p \leq \sqrt{n}$  prime. The other factor has no large prime divisor, hence must be a prime  $\leq \sqrt{n}$  (an element of  $P$ ). Consequently,  $m$  has the form

$$m = qpr$$

where  $p, r \leq \sqrt{n}$  are primes (not necessarily distinct), and any solution corresponds to deciding whether  $q$  is paired with  $p$  or with  $r$ . This yields at most 2 solutions (with fewer if  $p = r$ ).

*Case 2:  $m$  has exactly two large prime divisors (counted with multiplicity).* Let the large prime divisors of  $m$  (with multiplicity) be  $q_1, q_2 > \sqrt{n}$ . By Claim 1, each  $a \in A$  contains at most one large prime divisor, hence in any representation  $m = a_1 a_2$  with  $a_1 < a_2$  and  $a_1, a_2 \in A$ , the two large primes  $q_1, q_2$  must be split between  $a_1$  and  $a_2$ , i.e. each of  $a_1, a_2$  is divisible by exactly one of  $q_1, q_2$ . Moreover, by the definition of  $A = P \cup S$ , every element of  $A$  is either a prime  $\leq n$ , or a product  $pq$  with  $q > \sqrt{n}$  prime and  $p$  prime (necessarily  $p < \sqrt{n}$ ). In particular, if  $a \in A$  is divisible by a large prime  $q > \sqrt{n}$ , then

$$a = q \cdot u \quad \text{with } u \in \{1\} \cup \{\text{primes } \leq \sqrt{n}\}.$$

Therefore any solution  $m = a_1 a_2$  forces

$$a_1 = q_1 u_1, \quad a_2 = q_2 u_2$$

(up to swapping  $q_1, q_2$ ), where  $u_1, u_2 \in \{1\} \cup \{\text{primes } \leq \sqrt{n}\}$ .

If  $q_1 \neq q_2$ , then the only possible ambiguity comes from swapping which small factor ( $u_1$  or  $u_2$ ) is paired with  $q_1$  or  $q_2$ . Since each  $u_i$  is either 1 or a single prime, there are at most two such pairings:

$$(q_1 u_1)(q_2 u_2) \quad \text{or} \quad (q_1 u_2)(q_2 u_1),$$

and after imposing  $a_1 < a_2$  this gives at most 2 solutions (and fewer if  $u_1 = u_2$ ).

If  $q_1 = q_2 = q$ , then every admissible factor must contain exactly one copy of  $q$ , hence  $a_1 = q u_1, a_2 = q u_2$  with  $u_1, u_2$  as above. Swapping yields the same unordered pair, so with the constraint  $a_1 < a_2$  there is at most one solution.

Thus in all subcases, the number of solutions is at most 2.

In all cases, the number of solutions is at most 2. □

## 2. Size of $A$ . Clearly

$$|A| = |P| + |S| = \pi(n) + |S|.$$

Thus by Proposition A.2, we know that

$$|A| = \pi(n) + |S| = \frac{n \log \log n}{\log n} + (M+1) \frac{n}{\log n} + o\left(\frac{n}{\log n}\right).$$

Claim 2 shows that  $A$  is admissible, hence  $g_3(n) \geq |A|$ . This completes the proof. □

### 3. CONCLUDING REMARKS

In [3, p. 261], Erdős stated that his [3, Theorem 3] could be sharpened to

$$g_3(n) \leq \frac{\log \log n}{\log n} n + O\left(\frac{n}{(\log n)^{1+c}}\right),$$

where  $c > 0$  is a suitable positive constant. In view of Proposition 2.1, we now conjecture that the second term in the upper bound is also of order  $n/\log n$ :

**Conjecture 3.1.** *There exists an absolute constant  $c > 0$  such that, for all sufficiently large  $n$ ,*

$$g_3(n) \leq \frac{n \log \log n}{\log n} + c \frac{n}{\log n}.$$

More strongly, we also conjecture the following modified version of Problem 1.1:

**Conjecture 3.2.** *Is it true that*

$$g_3(n) = \frac{\log \log n}{\log n} n + (c + o(1)) \frac{n}{\log n}$$

*for some constant  $c$ ?*

### REFERENCES

- [1] T. F. Bloom, Erdős Problem #796, <https://www.erdosproblems.com/796>, accessed 2025-12-30.
- [2] Crisan, Dragos, and Radek Erban. "On the counting function of semiprimes." arXiv preprint arXiv:2006.16491 (2020).
- [3] P. Erdős, On the multiplicative representation of integers. Israel Journal of Mathematics (1964), 251–261.
- [4] P. Erdős, Some applications of graph theory to number theory. The Many Facets of Graph Theory (Proc. Conf., Western Mich. Univ., Kalamazoo, Mich., 1968) (1969), 77–82.

### APPENDIX A. PRELIMINARIES

Let  $n \geq 3$ . Define

$$S := \{pq \leq n : p, q \text{ are primes and } q > \sqrt{n}\}.$$

Let  $\pi(x)$  be the prime-counting function, and let  $\pi_2(x)$  denote the *semiprime counting function* defined by

$$\pi_2(x) := \#\{(p, q) : p \leq q, p, q \text{ prime}, pq \leq x\}.$$

**Lemma A.1.** *For every  $n \geq 3$ ,*

$$|S| = \pi_2(n) - \#\{(p, q) : p \leq q \leq \sqrt{n}, p, q \text{ prime}\} = \pi_2(n) - \frac{\pi(\sqrt{n})(\pi(\sqrt{n}) + 1)}{2}.$$

*Proof.* Every pair  $(p, q)$  of primes with  $p \leq q$  and  $pq \leq n$  falls into exactly one of the two disjoint classes:

$$\mathcal{A} := \{(p, q) : p \leq q \leq \sqrt{n}\}, \quad \mathcal{B} := \{(p, q) : p \leq q, pq \leq n, q > \sqrt{n}\}.$$

Hence  $\pi_2(n) = |\mathcal{A}| + |\mathcal{B}|$ .

If  $(p, q) \in \mathcal{A}$  then automatically  $pq \leq (\sqrt{n})^2 = n$ , so  $\mathcal{A}$  is simply the set of unordered pairs of primes  $\leq \sqrt{n}$  with repetition allowed. Therefore

$$|\mathcal{A}| = \binom{\pi(\sqrt{n})}{2} + \pi(\sqrt{n}) = \frac{\pi(\sqrt{n})(\pi(\sqrt{n}) + 1)}{2}.$$

If  $(p, q) \in \mathcal{B}$  then  $q > \sqrt{n}$  and  $pq \leq n$ ; this is equivalent to the condition that the integer  $pq$  lies in  $S$ . Moreover, because  $q > \sqrt{n}$  implies  $p < n/q < \sqrt{n} < q$ , each element of  $S$  corresponds to a unique pair  $(p, q) \in \mathcal{B}$  with  $p \leq q$ . Thus  $|\mathcal{B}| = |S|$ , and the claimed identity follows.  $\square$

Clearly, we have the asymptotic formula (see, for example, [2, Theorem 2.3])

$$\pi_2(x) = \frac{x \log \log x}{\log x} + M \frac{x}{\log x} + o\left(\frac{x}{\log x}\right). \quad (\text{A.1})$$

We also recall the prime number theorem  $\pi(x) \sim x/\log x$ .

**Proposition A.2.** *We have*

$$|S| = \frac{n \log \log n}{\log n} + M \frac{n}{\log n} + o\left(\frac{n}{\log n}\right).$$

*Proof.* By Lemma A.1,

$$|S| = \pi_2(n) - \frac{\pi(\sqrt{n})(\pi(\sqrt{n}) + 1)}{2}.$$

By the prime number theorem,  $\pi(\sqrt{n}) = O(\sqrt{n}/\log n)$ , hence

$$\pi(\sqrt{n})^2 = O\left(\frac{n}{(\log n)^2}\right) = o\left(\frac{n}{\log n}\right), \quad \pi(\sqrt{n}) = o\left(\frac{n}{\log n}\right).$$

Therefore

$$\frac{\pi(\sqrt{n})(\pi(\sqrt{n}) + 1)}{2} = o\left(\frac{n}{\log n}\right).$$

Combining this with (A.1) at  $x = n$  yields

$$|S| = \frac{n \log \log n}{\log n} + M \frac{n}{\log n} + o\left(\frac{n}{\log n}\right).$$

□