



SAT solvers

Computational Models of Argumentation

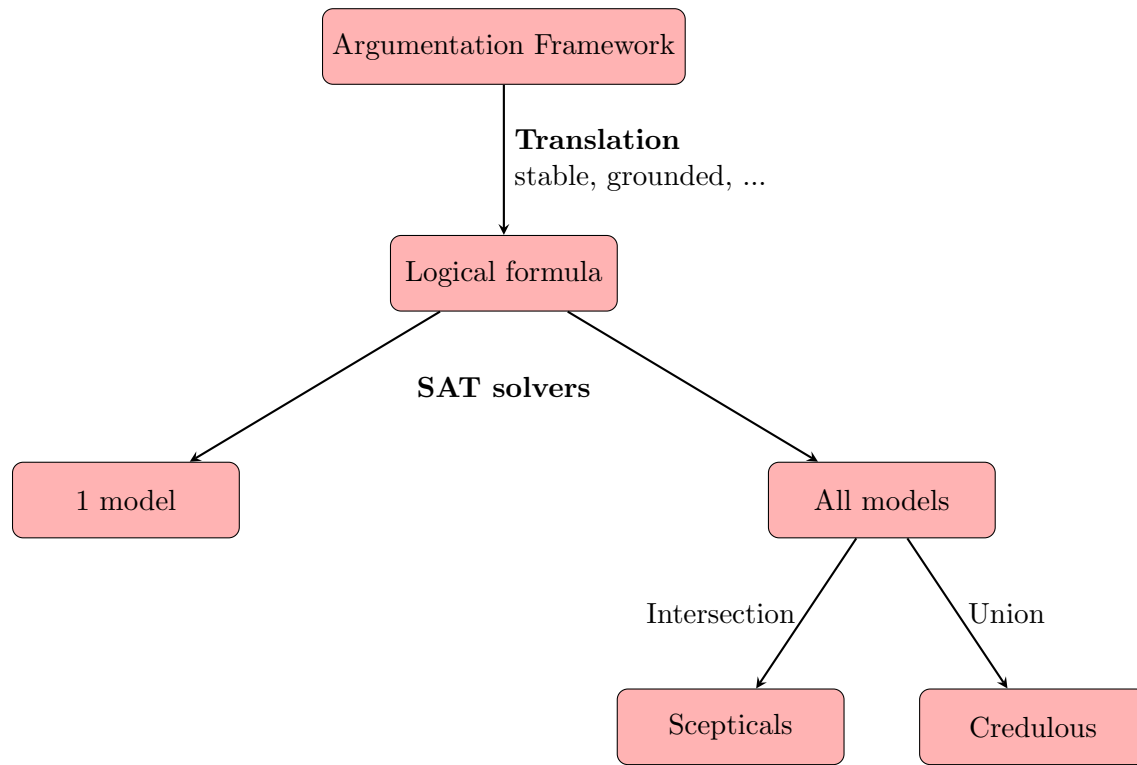
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Contents

1	Flowchart	3
2	Extensions	4
3	Schema	5

1 Flowchart



2 Extensions

Let $F := (A, R)$ be an Argumentation Framework and $S \subset A$, then S is

- **conflict-free** if for all (a, b) in S^2 , (a, b) is not in R .

$$\phi_{cf}(F) = \bigwedge_{(a,b) \in R} \neg(a \wedge b)$$

- **admissible** if S is conflict-free and defends all its arguments.

$$\phi_{ad}(F) = \phi_{cf}(F) \wedge \bigwedge_{a \in A} \left(a \implies \bigwedge_{(b,a) \in R} \bigvee_{(c,b) \in R} c \right)$$

- **complete** if S is admissible and contains all the arguments that it defends.

$$\phi_{co}(F) = \phi_{cf}(F) \wedge \bigwedge_{a \in A} \left(a \iff \bigwedge_{(b,a) \in R} \bigvee_{(c,b) \in R} c \right)$$

- **grounded** if S is complete and that there is no complete extension $S' \neq S$ such that $S' \subset S$.
Apply unit propagation to ϕ_{co} .
- **preferred** if S is complete and that there is no complete extension $S' \neq S$ such that $S \subset S'$.
There is no encoding to SAT. Instead, use ϕ_{co} and if the solution is not preferred, forbid it in the logical formula and try again.
- **stable** if S is conflict-free and attacks all the arguments it does not contain.

$$\phi_{st}(F) = \bigwedge_{a \in A} \left(a \iff \bigwedge_{(b,a) \in R} \neg b \right)$$

3 Schema

In the schema below, the arrow is an implication.

