

Blockchain System

1 Definitions

Definition 1 (Set). Let \mathbf{Elt} be the set of (concrete) elements. Let \emptyset be an empty set and $e \in \mathbf{Elt}$. A set of elements is expressed as the following syntax:

$s ::= \emptyset \mid e \mid s :: s$

Definition 2 (Account). An account is expressed as a tuple $\langle \mathbf{pak}, \mathbf{puk} \rangle$, where \mathbf{pak} is its private key and \mathbf{puk} is its public key.

Definition 3 (Manager). A manager manages an account on the blockchain. It is expressed as a tuple $\langle \mathbf{puk}, \mathbf{bal}, \mathbf{cou} \rangle$, where \mathbf{puk} is the public key of an account, \mathbf{bal} is its balance and \mathbf{cou} is the counter for its operations that could be T or F .

Let \mathbf{K} be a set of managers. The following functions interact with \mathbf{K} .

```
let rec checkBal K puk n m =
  match K with
  | 0 -> true
  | < puk', bal, cou > :: K' ->
    if (puk = puk') and (n + m) <= bal then true
    else checkBal (K', puk, n, m)

let rec checkExi K puk =
  match K with
  | 0 -> false
  | < puk', bal, cou > :: K' ->
    if (puk = puk') then true
    else checkExi (K', puk)

let rec checkCou K puk =
  match K with
  | 0 -> false
  | < puk', bal, cou > :: K' ->
    if (puk = puk') and (cou = T) then true
    else checkCou (K', puk)

let rec updateCou K puk =
  match K with
  | 0 -> 0
  | < puk', bal, cou > :: K' ->
    if (puk = puk') then < puk', bal, F > :: K'
    else < puk', bal, cou > :: updateCou (K', puk)
```

```

let rec updateSuc K puk puk' m n =
  match K with
  | 0 -> 0
  | < puk'', bal, cou > :: K' ->
    if (puk = puk'') then < puk'', bal - (n + m), T >
      :: updateSuc (K', puk, puk', n, m)
    else if (puk' = puk'') then < puk'', bal + n, cou > :: K'
    else < puk'', bal, cou >
      :: updateSuc (K', puk, puk', n, m)

```

Definition 4 (Operation). *An operation is expressed as the following syntax:*
 $\mathbf{op} ::= \text{transfer } n \text{ from } \mathbf{puk} \text{ to } \mathbf{puk}' \text{ fee } m$

Let \mathbf{C} be the set of accounts and \mathbf{O} be a set of operations.

Definition 5 (State of a node). *The state of a node is expressed as a pair [\mathbf{C}, \mathbf{O}].*

When an operation is injected in a node, it enters in a pending pool (and called a pending move).

Definition 6 (Pending operation). *A pending operation is expressed as a pair $\langle \mathbf{op}, \mathbf{t} \rangle$, where \mathbf{op} is an operation and \mathbf{t} is the time when it is injected.*

After sometime, a pending operation could be included in the blockchain as a accepted operation.

Definition 7 (Accepted operation). *An accepted operation is expressed as a tuple $\langle \mathbf{op}, \mathbf{t} \rangle$, where \mathbf{h} is the hash of the block head that includes the operation.*

Let \mathbf{P} be a set of pending operations, \mathbf{A} be a set of accepted operations, and \mathbf{t} is the current time of the blockchain.

Definition 8 (Blockchain). *The state of a blockchain is expressed as a tuple [$\mathbf{P}, \mathbf{A}, \mathbf{K}, \mathbf{t}$].*

Definition 9 (Blockchain system). *A blockchain system $\mathbf{S} \triangleq \langle \mathbf{M}, \mathbf{B} \rangle$ consists of*

1. $\mathbf{M} \equiv [C, O]$ is the state of a node, and
2. $\mathbf{B} \equiv [P, A, K, t]$ is the state of a blockchain such as $\forall c \in C \implies \exists k \in K, k.puk = c.puk$.

2 Rules

Rule 1 [injected]:

$$\frac{\text{checkBan}(K, puk, m, n) \wedge \text{checkCou}(K, puk) \wedge \text{checkExi}(K, puk')}{\langle [< pak, puk > :: C, O], [P, A, K, t] \rangle \rightarrow \langle [C, O], [(< \text{transfer } n \text{ from } puk \text{ to } puk' \text{ fee } m, t >) :: P, A, \text{updateCou}(K, puk), t] \rangle} \quad (1)$$

Rule 2 [included]:

$$\frac{[< \text{transfer } n \text{ from } puk \text{ to } puk' \text{ fee } m, t > :: P, A, K, t'] \rightarrow [P, < \text{transfer } n \text{ from } puk \text{ to } puk' \text{ fee } m, t' > :: A, \text{updateSuc}(K, puk, puk', n, m), t' + 1]}{(2)}$$

Rule 3 [timeout]:

$$\frac{t' - t \geq 60}{[< \text{transfer } n \text{ from } puk \text{ to } puk' \text{ fee } m, t > :: P, A, K, t'] \rightarrow [P, A, \text{updateFai}(K, puk), t']} \quad (3)$$