

a)

0	-50
1	-17
2	-12
3	-7
4	
5	-35
6	-46
7	-45
8	-26
9	-55

b)

0	-50
1	-7
2	-2
3	-55
4	-26
5	-35
6	-46
7	-17
8	
9	-45

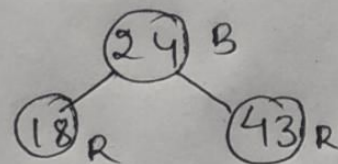
2) insert 24

(24) B.

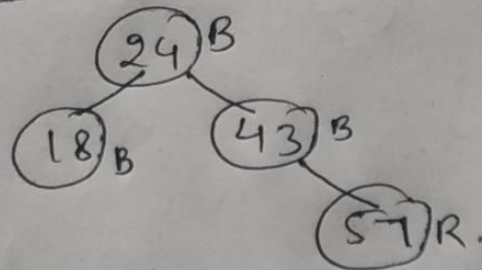
insert 18



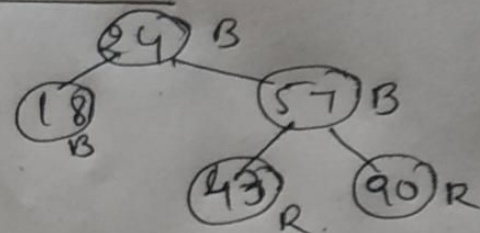
insert 43



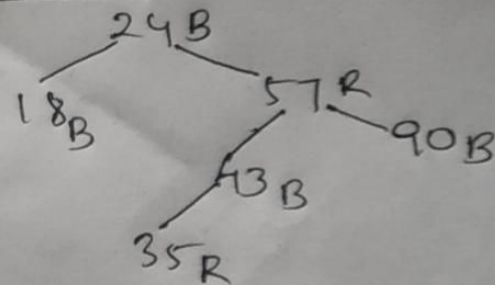
insert 57



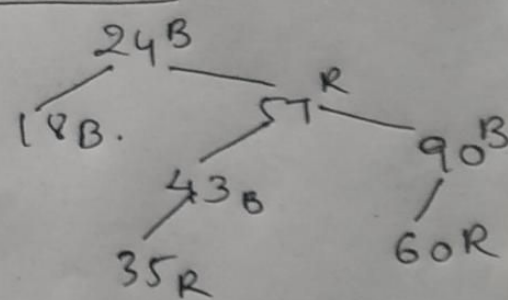
insert 90

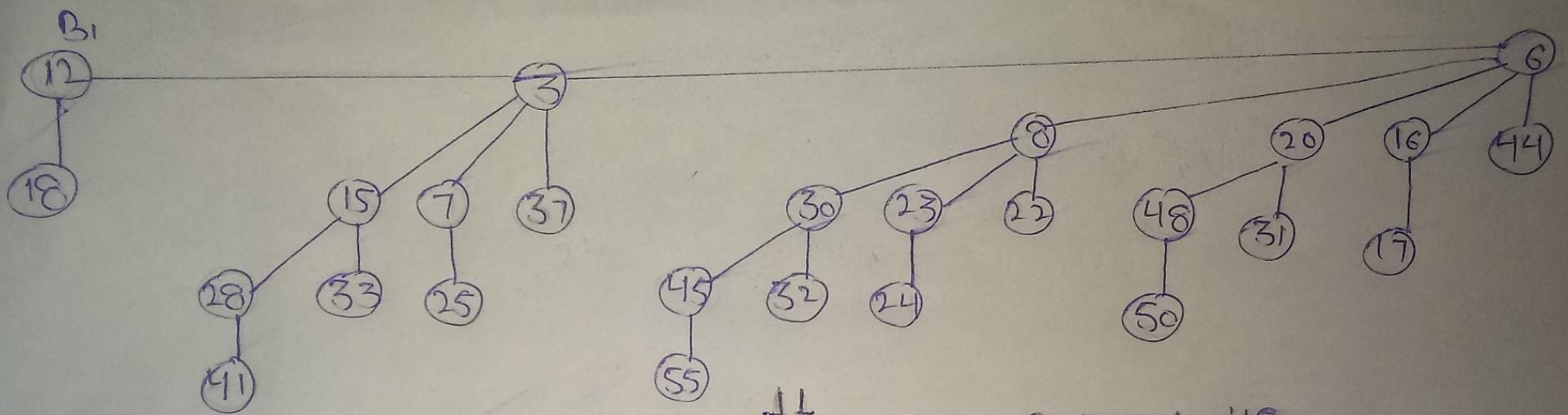


insert 35

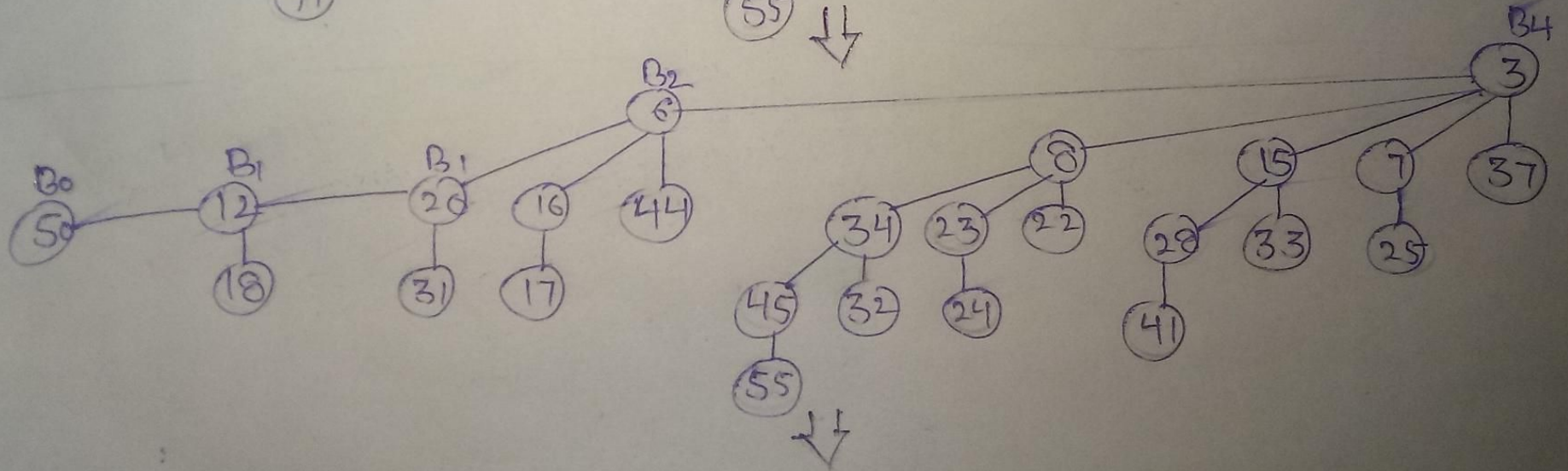
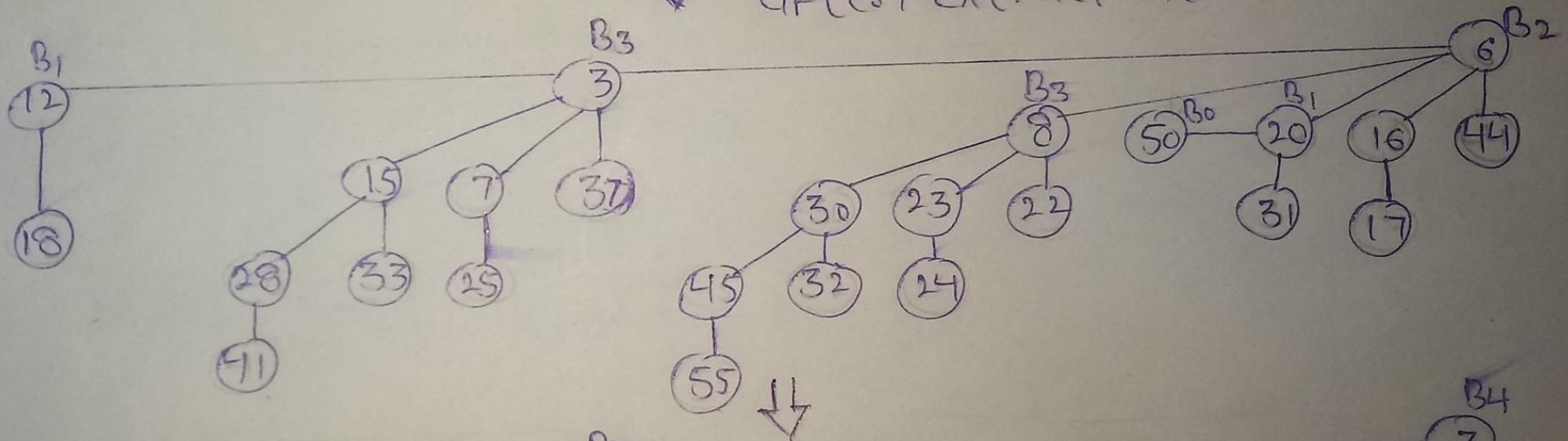


insert 60





after extract 48





Answer :-

There are mainly two major parts in Huffman coding

- 1) Build a Huffman Tree from input characters.
- 2) Traverse the Huffman Tree & assign codes to characters.

Steps to build Huffman Tree

→ Input is an array of unique characters along with their frequency of occurrences and output is Huffman Tree.

- 1) Create a leaf node for each unique character and build a min heap of all leaf nodes (Min Heap is used as a priority queue. The value of frequency field is used to compare two nodes in min heap. Initially the least frequency character is at root.

- 2) Extract two nodes with the minimum frequency from the min heap.
- 3) Create a new internal node with a frequency equal to the sum of the two nodes frequencies. Make the first extracted node as its ~~right~~ <sup>left</sup> child and the other extracted node as its right child. Add this node to the min heap.
- 4) Repeat steps #2 & #3 until the heap contains only one node. The remaining node is the root node & the tree is complete.

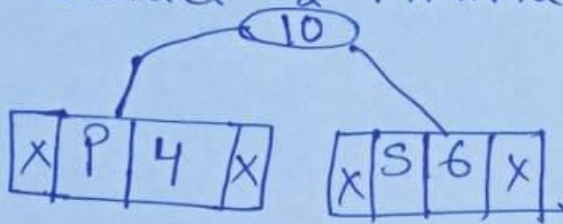
Character	Frequency
P	4
S	6
M	8
Y	9
K	12
E	14
A	16



P S H Y K E A

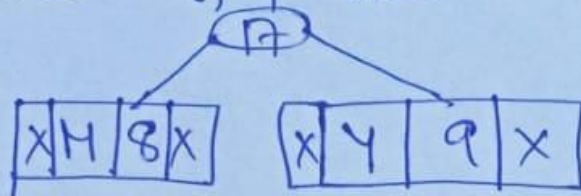
4 6 8 9 12 14 16

→ Extract 2 minimum nodes



(10), 8, 9, 12, 14, 16

→ now 8, 9 are 2 minimum nodes

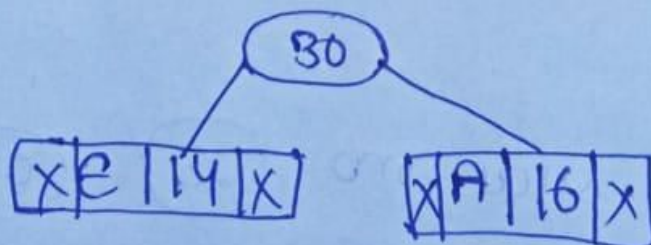


(10), (17), 12, 14, 16

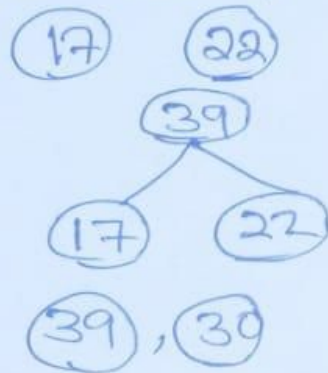
→ now (10), 12 are minimum nodes



→ ~~(10)~~ now (17), (22), 14, 16 are minimum

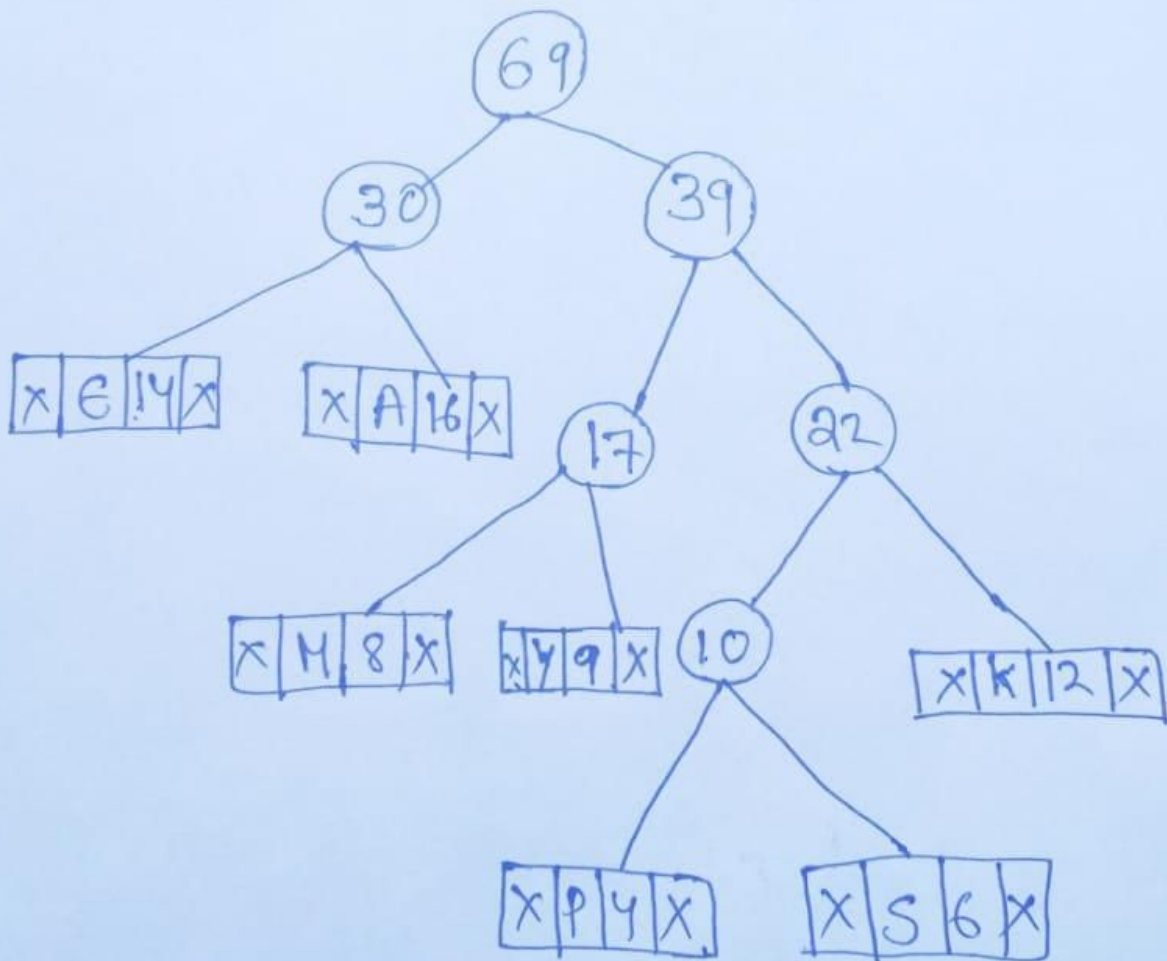


→ now 17, 22, 30 are minimum nodes



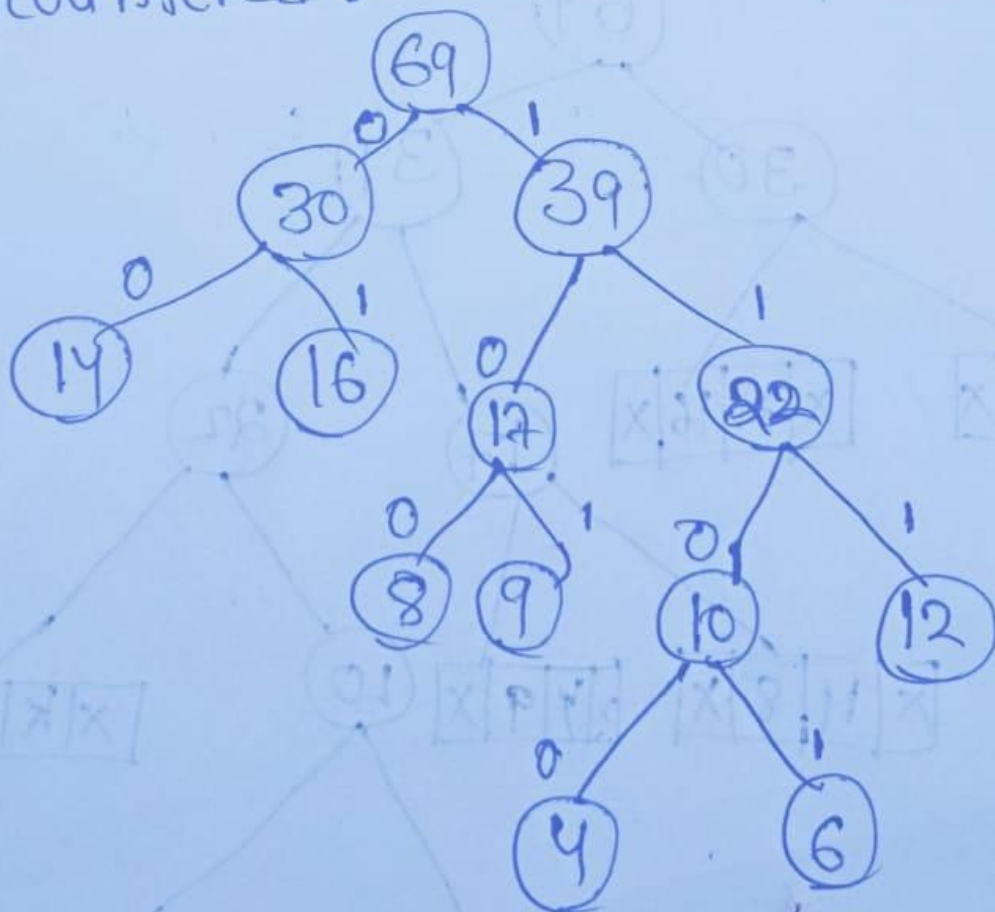
→ now add 30, 39

→ The final Huffman Tree is



## Steps to print codes from Huffman Tree

Traverse the tree formed starting from the root. Maintain an auxiliary array, while moving to the left child write 0 to the array, while moving to the right child write 1 to the array. Print the array when a leaf node is encountered.



**Pre-Order Traversal:** In this traversal,

1. Visit Root Node
2. Visit Left Subtree and
3. Visit Right Subtree

**Post-Order Traversal:** In this traversal,

1. Visit Left Subtree
2. Visit Right Subtree and
3. Visit Root Node

**In-Order Traversal:** In this traversal,

1. Visit Left Subtree
2. Visit Root Node and
3. Visit Right Subtree

**Given Problem is:**

**In-Order Traversal: 6 4 1 3 0 5 9 7 8 2**

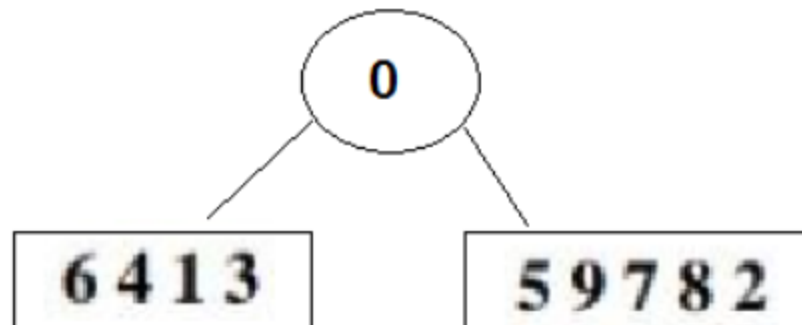
**Pre-Order Traversal: 0 1 4 6 3 2 5 7 9 8**



In Pre-order Traversal : Traversal order is Root, Left, Right

So, first number 0 is Root Node

From the list given In-Order traversal 0 is selected as Root and the list is divided like this according to its traversals order that is Left, Root and Right.

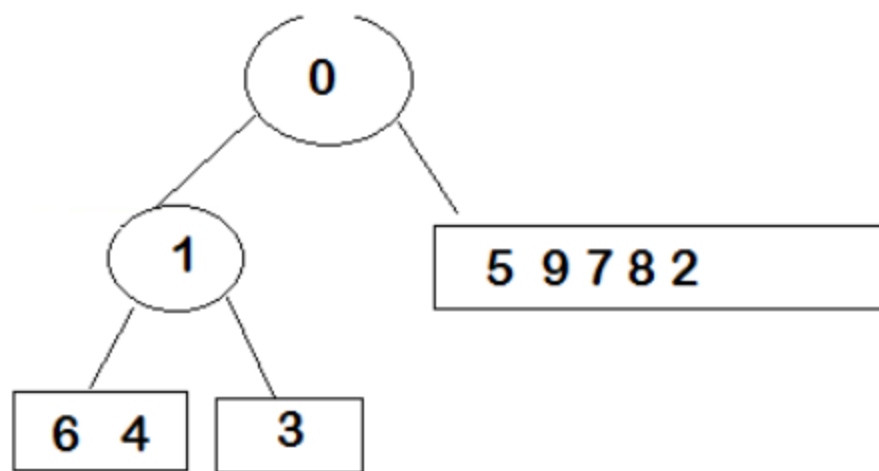


Now Considering List 1: 6 4 1 3

Now, finding root node among these numbers, we have check pre-order traversal list : 0 1 4 6 3 2 5 7 9 8

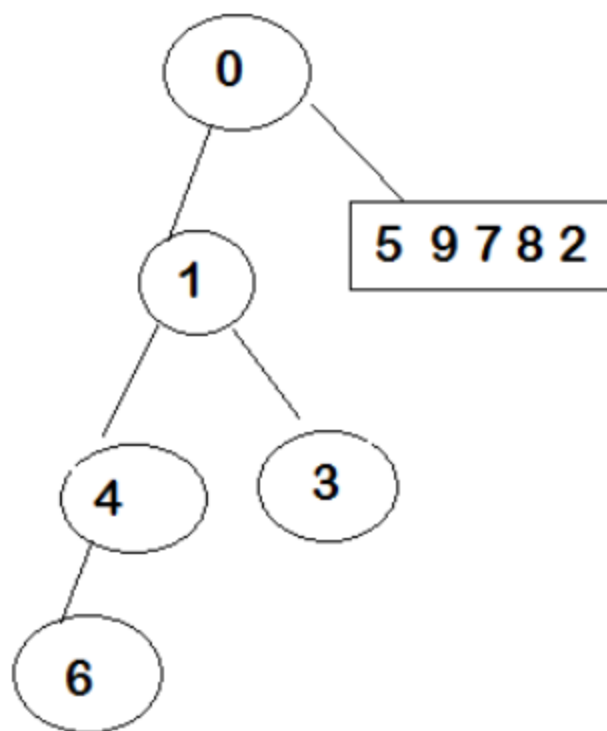
Among 6, 4, 1, 3 number 1 is visited first. So, number 1 becomes root.

So, list 6 4 1 3 is divided in this way: 6,4 becomes left sublist and 3 becomes right sublist.(1 is root Node).  
The tree becomes like this:



In the same considering 6, 4 list finding root element. For this check preorder traversal, Node which is visited first (among 6, 4) is treated as rooted. Here 4 is visited first in preorder traversal. 4 is root node. Left list is 6.

Now, tree is



Perform same procedure for right sublist also,

Right sublist is 5, 9, 7, 8, 2. To find root check pre-order traversal, 2 is visited first among 5, 9, 7, 8, 2.

So, number 2 is root node. Left list is 5, 9, 7, 8.

Again, to find root from Left list 5, 9, 7, 8 check pre-order traversal, 5 is visited, which becomes root node.

So, number 5 is root node. Right List is 9, 7, 8.

Again, to find root from Right list 9, 7, 8 check pre-order traversal 7 becomes root node, 9 left node and 8 right node.

The complete steps are given below:

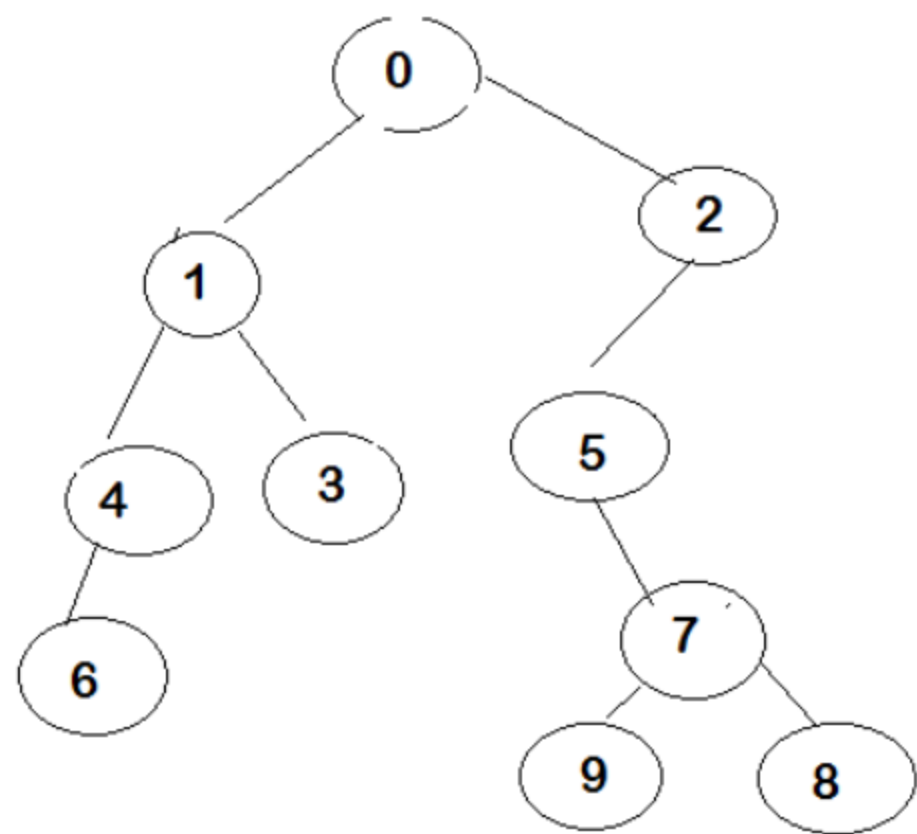


Resultant binary tree formed is

**In-Order Traversal: 6 4 1 3 0 5 9 7 8 2**

**Pre-Order Traversal: 0 1 4 6 3 2 5 7 9 8**





**Post order traversal is 6 4 3 1 9 8 7 5 2**

## 96) AVL Tree :-

50, 40, 25, 10, 15, 5, 45, 35, 20, 30

rules:-

- every tree has a root node (at the top)
- the root node has zero, one or two child nodes.
- and also child nodes has zero, one or two child nodes and so on
- The each node has up to two children.
- its left descendants are less than the current node, which is less than the right descendants.

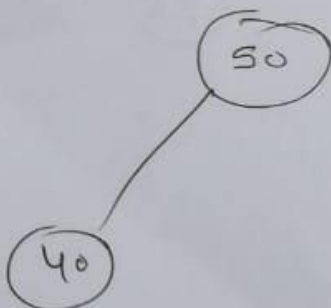
$\angle < \text{root} > = \angle$

Step 1:-

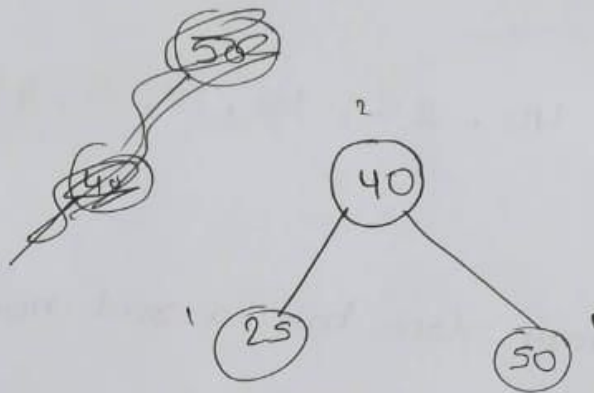
Insert 50:-

(50)

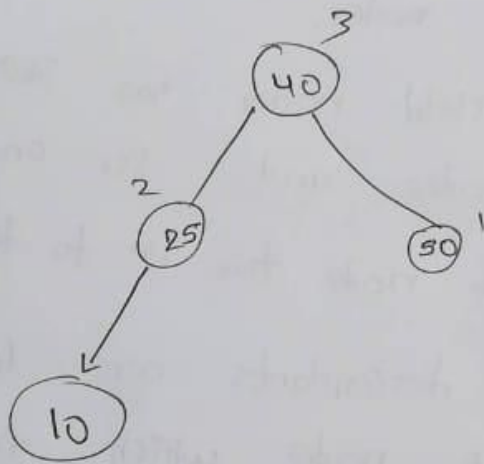
Insert 40:-



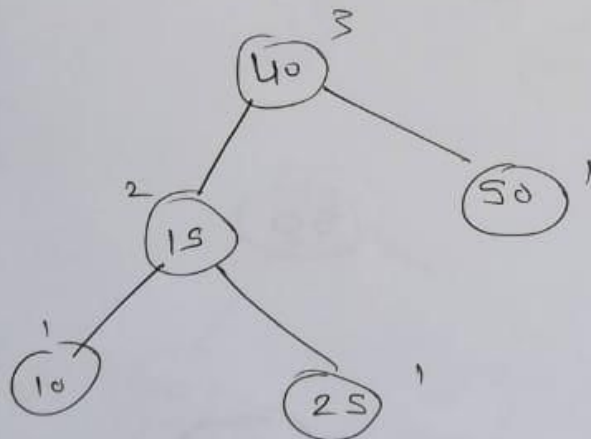
Insert 25:-



Insert 10:-

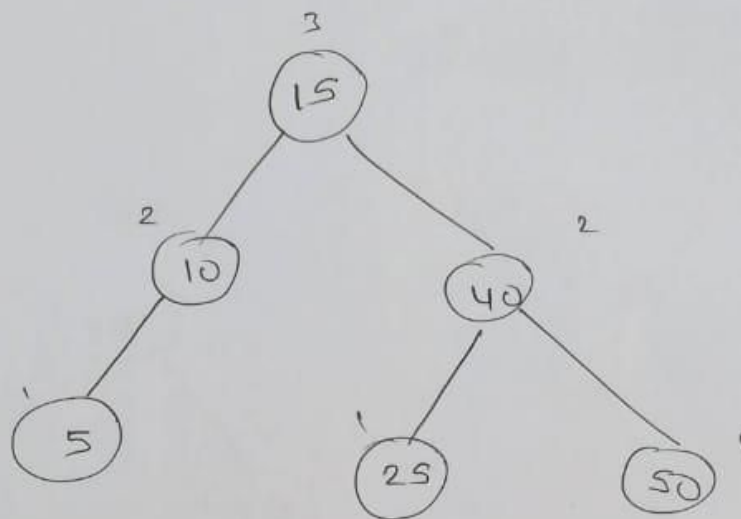


Insert 15:-

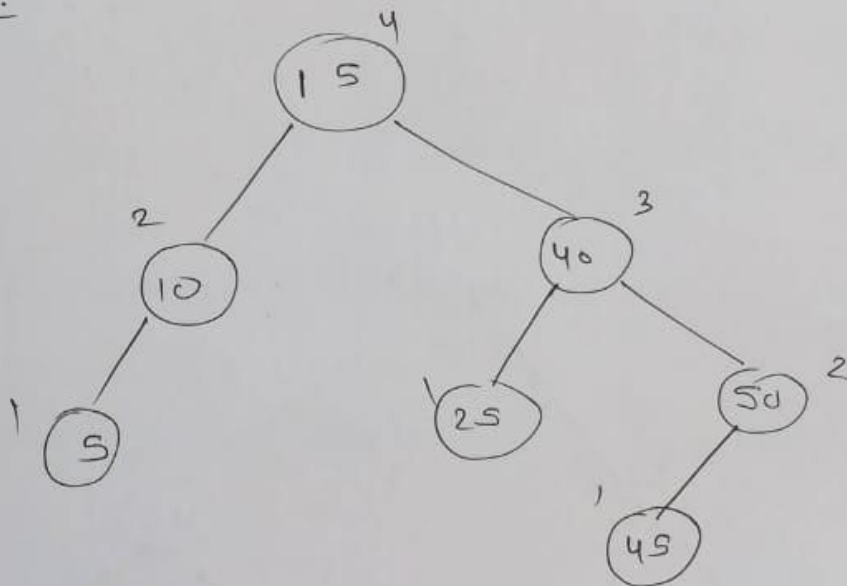




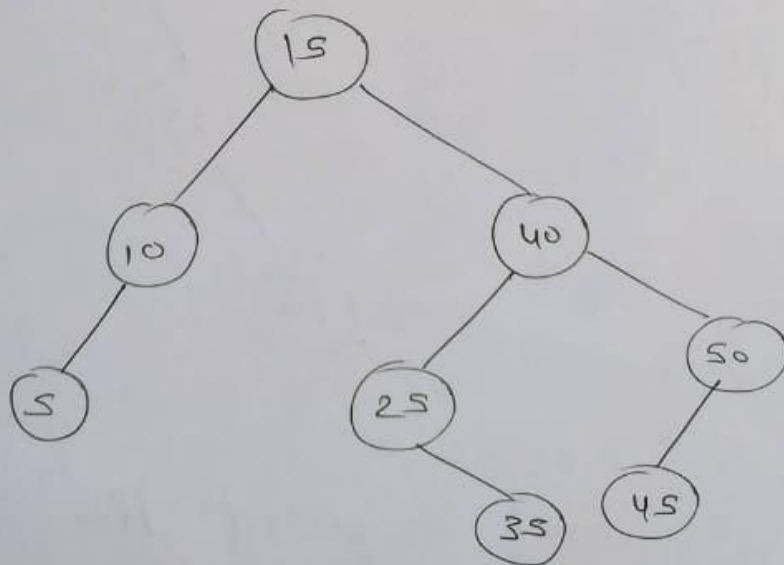
Insert 5:-



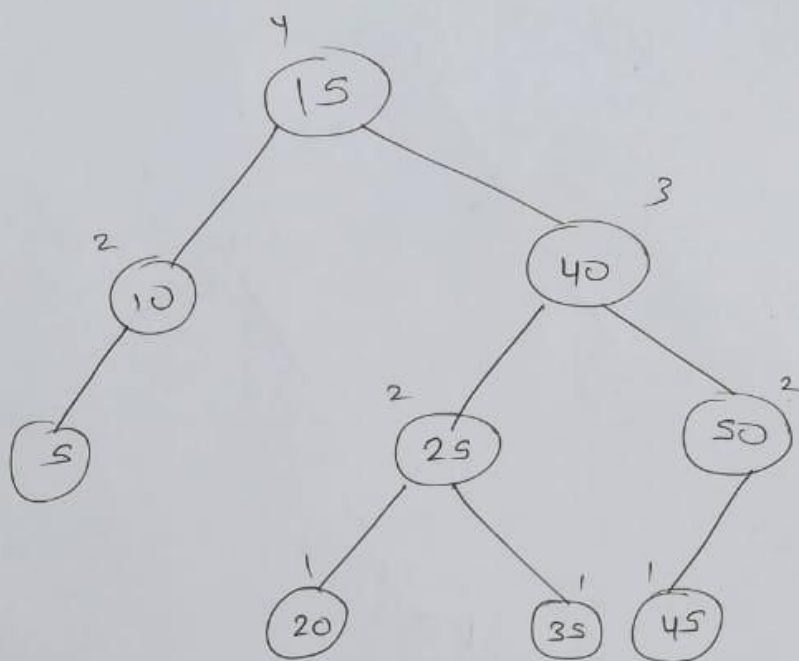
Insert 45:-



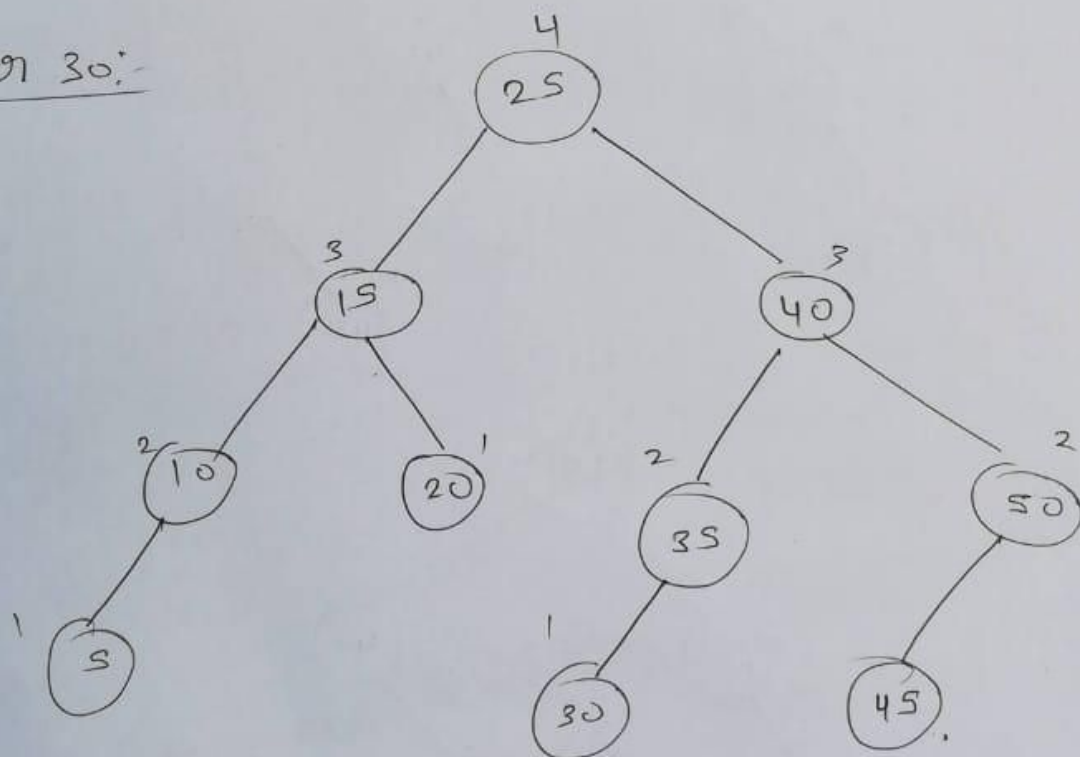
Insert 35:-



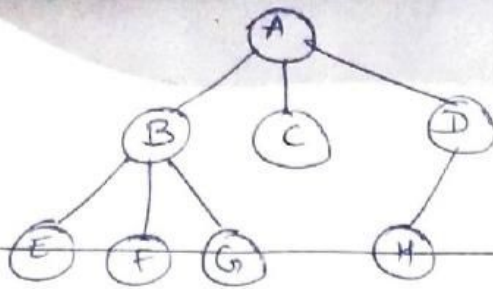
Insert 20:-



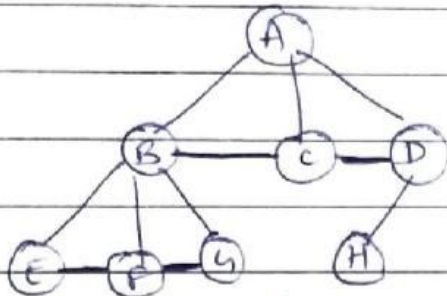
Insert 30:-



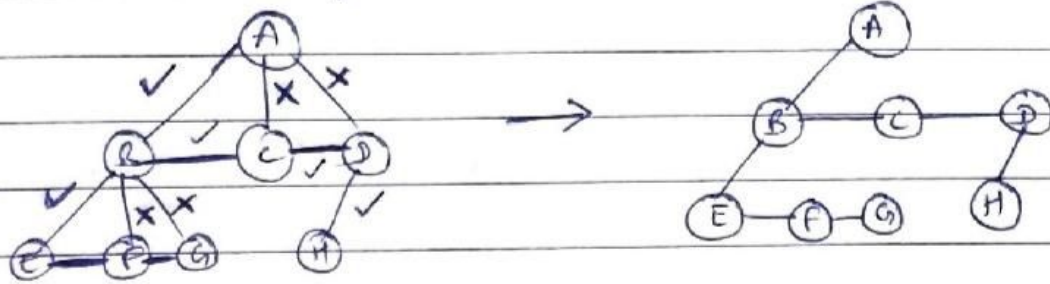
final required AVL Tree.



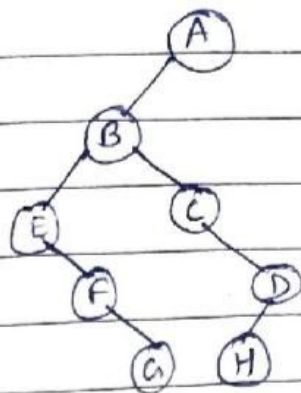
Step 1:- Insert edge connecting all siblings of the tree.



Step 2:- Delete all parent edges from parent to its children except the one to its leftmost child.



Step 3:- Rotate the resulting tree by  $45^\circ$  to differentiate b/w its left & right subtree.





Inserting 25

---

0025

Inserting 15:

0015    0025

Inserting 5:



Inserting 30:



Inserting 20:

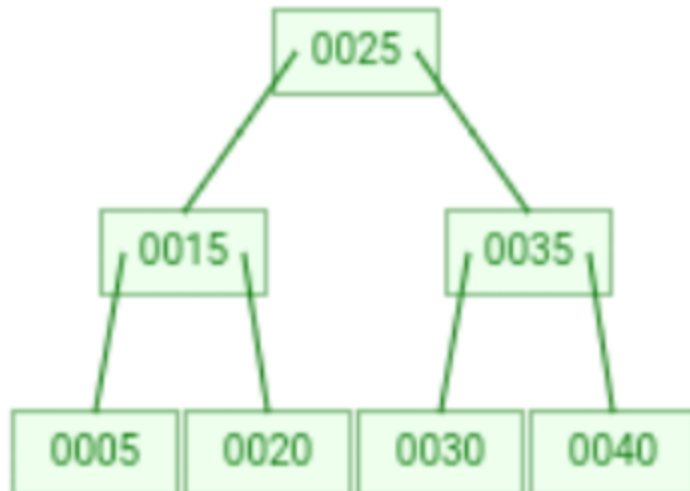


Inserting 35:





Inserting 40:



Deleting 30:

