Course 02263 Mandatory Assignment 1, 2016

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March 16, 2016

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1 Introduction

This document contains a solution for the problem of colouring the pieces of a puzzle such that no neighbours get same colors. The solution is proposed for unlimited number of colors.

2 Requirements Specification

This section is specifying the requirements for the colouring of a relation. ColouringBasics is defining functions which will be used in ColouringReq.

2.1 Auxiliary functions for requirements

The class *ColouringBasics* contains the methods *isRelation*, *isCorrectColouring* and the auxiliary functions we have used to implement them.

The function isRelation is used to check if a relation (i.e. (Piece >< Piece)-set) is well formed. We are checking that all couples of the relation are composed of two distinct pieces. If this is the case, the function will return true. We assume here that it is not a problem for us to have too many couples. For instance, having (P1, P2) and (P2, P1) in the same relation does not give more information and is not considered as an error.

The function *isCorrectColouring* is used to check if a coulouring (i.e. (Pieceset)-set) is well formed for a given relation. We are checking three different properties in the colouring:

- Unicity of each piece of the relation. Each piece can only have one colour, so each piece should be at most once in the colouring.
- Exhaustivity for each piece of the relation. Each piece of the relation should be coloured, so each piece should be at least once in the colouring.
- Neighborhood relations. Each piece of the relation should have a different colour from its neighbours.

Unicity and exhaustivity could be put together in a condition saying that each piece of the relation should be exactly once in the colouring. However, we are testing this properties separately in the code.

The following piece of code includes comments for each method:

```
 \begin{aligned} \mathbf{scheme} & \  \, \mathbf{ColouringBasics} = \\ \mathbf{class} & \\ \mathbf{type} & \\ & \  \, \mathbf{Piece} = \mathbf{Text}, \\ & \  \, \mathbf{Relation} = (\mathbf{Piece} \times \mathbf{Piece})\text{-}\mathbf{set}, \end{aligned}
```

```
Colour = Piece-set,
Colouring = Colour-set
```

value

```
/* The point is to check if a piece is in the colouring (i.e. (Piece-set)-set)
 * params
                 cn1 – piece we want to check
                 colouring - colouring which cotaining colours (Piece-set)
 * returns
                 true - if the piece is in the colour set
                  false - if the piece is not in the colour set
 */
isInColouring : Piece \times Colouring \rightarrow Bool
isInColouring(cn1, colouring) \equiv
  (\exists colour : Colour \bullet colour \in colouring \land
                 (\exists p1 : Piece \cdot p1 \in colour \land (p1 = cn1))),
/* Test if a couple in the relation has the same colouring
 * params
                 cn1 - first piece we want to test
                 cn2 - second piece we want to test
                 cols - colouring which contains the colours
 * returns
                 true - if Piece already in colour set
                  false - if Piece not in color set
isSameColour : Piece \times Piece \times Colouring \rightarrow Bool
isSameColour(cn1, cn2, cols) \equiv
    (\exists colour1 : Colour \cdot colour1 \in cols \land cn1 \in colour1
        \land (\existscolour2 : Colour • colour2 \in cols \land cn2 \in colour2 \land (colour1 = colour2))),
/* Defining that for all pieces (p1, p2) defining couples in a relation,
 * we have p1 different from p2.
 * The point is to check that the relation doesn't have any couple defined
 * with the two times the same piece.
 * For instance: {p1, p1}
 * params
                 r - relation we want to test
 * returns
```

```
true – if the relation is well-formated
                  false - if the relation is not well-formated
 */
is Relation \ : \ Relation \ \to \textbf{Bool}
isRelation(r) \equiv
   (\forall p1 : Piece \cdot p1 \in \{ \mathbf{let} \ (p1, \ p2) = e \ \mathbf{in} \ p1 \ \mathbf{end} \ | \ e : Piece \times Piece \cdot e \in r \} \Rightarrow
                  (\forall p2 : Piece \cdot p2 \in \{ let (p1, p2) = e in p2 end \mid e : Piece \times Piece \cdot e \in r \} \Rightarrow
                           ((p1, p2) \in r \Rightarrow (p1 \neq p2))
                  )
         ),
/* Defining two pieces p1 and p2 as parts of a couple which belongs to r,
 * then, we have the implication that if cn1 and cn2 are neighbour,
 * they should be equal to p1 and p2.
 * The point is to check that it exists (or not) two pieces (p1, p2)
 * from the relation which are neighbours.
 * Then, we check if these pieces are equals cn1 and cn2.
   params
                  cn1 - first piece we want to check
                  cn2 – second piece we want to check
                  r – relation where are all the neighborhoods relation are defined
 * returns
                  true - if pieces are neighbors according to the relation
                  false - if pieces are not neighbors according to the relation
 */
areNb : Piece \times Piece \times Relation \rightarrow Bool
areNb(cn1, cn2, r) \equiv
       (\exists (p1, p2) : Piece \times Piece \bullet (p1, p2) \in r \land
                  ((cn1 = p1 \land cn2 = p2) \lor (cn1 = p2 \land cn2 = p1)))
pre isRelation(r),
/* A colouring is correct if:
                  - all couples of neighbours have different colours
                  - all pieces of the relation are only once in the colouring
                    since one Piece can have only one colour
                  - all pieces of the relation are there in the colouring at least once
   params
                  cols – colouring used on the pieces
                  r – relation defining neighborhood relations
 * returns
                  true – if colouring is done correctly
```

```
false – if colouring is not done correctly
 */
isCorrectColouring : Colouring \times Relation \rightarrow \textbf{Bool}
isCorrectColouring(cols, r) \equiv
      neighborCondition(cols, r) \wedge unicityCondition(cols, r) \wedge exhaustivityCondition(cols, r)
pre isRelation(r),
/* Fonction used to check that all pieces
 * of the relation are in colouring */
exhaustivityCondition : Colouring \times Relation \rightarrow Bool
exhaustivityCondition(cols, r) \equiv
  (\forall p1 : Piece \cdot p1 \in \{let (p1, p2) = e in p1 end \mid e : Piece \times Piece \cdot e \in r\} \Rightarrow
                    (\forall p2 : Piece \cdot p2 \in \{let (p1, p2) = e in p2 end \mid e : Piece \times Piece \cdot e \in r\} \Rightarrow
                              (\exists col \; : \; Colour \; \bullet \; col \; \in \; cols \; \land \; p1 \in col) \; \land \\
                                        (\exists col : Colour \cdot col \in cols \land p2 \in col)
                    )
          ),
/* Function used to check that neighbour pieces have different colours */
neighborCondition : Colouring \times Relation \rightarrow Bool
neighborCondition(cols, r) \equiv
       (\forall p1 : Piece \cdot p1 \in \{ let (p1, p2) = e in p1 end \mid e : Piece \times Piece \cdot e \in r \} \Rightarrow
                    (\forall p2 : Piece \cdot p2 \in \{let (p1, p2) = e in p2 end \mid e : Piece \times Piece \cdot e \in r\} \Rightarrow
                              (areNb(p1, p2, r)) \Rightarrow \sim isSameColour(p1, p2, cols)
          ),
/* Functions used for the unicity of each Piece
 * from the relation in the colouring */
/* Unicity is proved by showing that for all pieces P in the relation,
 * it doesn't exist any couple of colours (c1, c2),
 * so that c1 and c2 are different and the piece P is in both of them.
unicityCondition : Colouring \times Relation \rightarrow Bool
unicityCondition(cols, r) \equiv
         (\forall p1 : Piece \cdot p1 \in \{let (p1, p2) = e in p1 end \mid e : Piece \times Piece \cdot e \in r\} \Rightarrow
                    (\forall p2 : Piece \cdot p2 \in \{let (p1, p2) = e in p2 end \mid e : Piece \times Piece \cdot e \in r\} \Rightarrow
                              (unicityConditionInside(cols, p1) ∧ unicityConditionInside(cols, p2))
          ),
/* Function used to prove that the piece cn1 is only in one colour of the colouring.
 */
```

```
\begin{split} & \text{unicityConditionInside}: \ Colouring \times \ Piece \to \textbf{Bool} \\ & \text{unicityConditionInside}(\ cols \ , \ cn1) \equiv \\ & \sim (\exists colour1: \ Colour \bullet \ colour1 \in \ cols \ \land \ cn1 \in colour1 \ \land \\ & (\exists colour2: \ Colour \bullet \ colour2 \in \ cols \ \land \ cn1 \in colour2 \land \ (colour1 \neq colour2))) \end{split}
```

end

2.2 Colouring Requirements

The class ColouringRes contains an implicit function using the isRelation and isCorrectColouring methods.

ColouringBasics

```
scheme ColouringReq = extend ColouringBasics with class value /* requirement spec */
col : Relation \stackrel{\sim}{\to} Colouring col(r) as cols post isCorrectColouring(cols, r) pre isRelation(r)
```

The purpose of this class is to create a correct colouring for a well formed relation, as explained in the description of the functions *isRelation* and *isCorrectColouring*.

This class has been type checked, but it has not been translated to RSL and no tests have been performed on it since the SML translation does not support implicit functions.

3 Explicit Specification

3.1 Explanation on the algorithm created

The following section provides an explicit declaration of a function returning a correct colouring for a well formed relation.

The function starts by performing a recursive treatment on the relation in the method treatmentOnRelation. In order to apply some specific treatment to one couple of the relation (i.e. a (Couple)-set), we extract the first element of the relation and then perform treatmentOnRelation recursively on the remaining tail. We have added an empty set verification to stop this recursion. This way, we can make recursive calls without knowing the number of couples inside the relation.

The treatment made on couple is done inside *treatmentOnCouple*. This function is performing a recursive call of *treatmentOnPiece* which is putting the current piece at the right place in the colouring considering the neighbourhood relationships.

This class is so composed of three types of functions:

- auxiliary functions used to perform small tasks or verifications
- treatment functions used to perform treatment in a recursive way on Relations, Couples of pieces and Pieces
- the function col the function starting the algorithm

In the code, comments are provided in the header of each function in order to explain its purpose:

3.2 ColouringEx.rsl

ColouringBasics

```
/* This class extends the ColouringBasics in order to
    * provide a method returning a correct colouring for a given relation.
    *
    * This class is composed of:
    * — auxiliary functions for the small tasks
    * — treatment functions which are used recursively on Pieces,
    * Couples of Pieces and Relations
    * — the main function col
    */

scheme ColouringEx =
extend ColouringBasics with
```

class

value

```
/* Auxiliary functions */
/* This functions checks if the piece can be added or not in the colour being tested.
 * We are checking if there is a neighbour of the piece in the colour
 * and if the piece is already in the colour.
   params
                  cn1 – piece we want to check
                  r – relation defining neighbourhood relations
                   colour - current colour we want to test
 * returns
                   true - if Piece can be added to the colour
                   false - if Piece cannot be added to the colour
 */
\operatorname{canAddPieceToColour}: \operatorname{Piece} \times \operatorname{Relation} \times \operatorname{Colour} \to \operatorname{\bf Bool}
canAddPieceToColour(cn1, r, colour) \equiv
    \sim ((\exists p1 : Piece \cdot p1 \in colour \land areNb(p1, cn1, r))) \land \sim (cn1 \in colour),
/* This function checks if the piece can be added to the existing colours of the colouring
 * params
                  cn1 – piece we want to check
                  r - relation defining neighbourhood relations
                   colouring - Colouring to test
 * returns
                   true – if the piece can be added in a color of the colouring
                   false – if the piece cannot be added to the colour of the colouring
 */
\operatorname{canAddPieceToExistingColours}: \operatorname{Piece} \times \operatorname{Relation} \times \operatorname{Colouring} \to \operatorname{\textbf{Bool}}
canAddPieceToExistingColours(cn1, r, colouring) \equiv
(\exists colour : Colour \cdot colour \in colouring \land canAddPieceToColour(cn1, r, colour))
         \land \sim \text{pieceInColouring(cn1, colouring)},
/* This function is used to add a colour to a new Colour (i.e. Piece-set)
 * params
                  cn1 – piece to add to the colour
                  r - relation defining neighbourhood relations
                   colour - current colour we want to add the piece
```

```
* returns
                 colour - colour in the parameters with potentially the new piece
 */
add Piece To Colour: Piece \times Relation \times Colour \times Colouring \rightarrow Colour
{\rm addPieceToColour}({\rm cn1},\,{\rm r,\,colour},\,{\rm colouring}) \equiv
     if canAddPieceToColour(cn1, r, colour) ∧ ~pieceInColouring(cn1, colouring)
        then colour \cup \{cn1\}
        else colour
        end,
pieceInColouring : Piece \times Colouring \rightarrow Bool
pieceInColouring(cn1, colouring) \equiv
        (\exists colour : Colour \bullet colour \in colouring \land cn1 \in colour),
/* Treatment functions */
/* This fonction shows the treatment made on a Piece
 * that is added to the Colouring (.i.e. {Piece-set}-set)
 *
   params
                 cn1 – piece to add to the colour
                 r – relation defining neighbourhood relations
                 colouring – current colouring where we want to add the piece
 * returns
                 colouring - colouring from the input with potentially a new Piece
 */
treatmentOnPiece: Piece \times Relation \times Colouring \rightarrow Colouring
treatmentOnPiece(cn1, r, colouring) \equiv
     if canAddPieceToExistingColours(cn1, r, colouring)
        then
                 let colour = hd { c \mid c : Colour \cdot c \in colouring \land canAddPieceToColour(cn1, r, c)},
                          colouring = colouring \setminus \{colour\},\
                          colouring = colouring \cup {addPieceToColour(cn1, r, colour, colouring)}
                 in
                          colouring
                 end
        else elseCondition(cn1, colouring)
        end.
/* This function extends the treatmentOnPiece function in the case of a Piece
 * that cannot be added to an existing Colour (i.e. Piece-set)
 * This Piece should so either not be added at all (if the Piece is already in the Colouring),
 * or should be added in a new Colour.
```

```
* params
                 cn1 – Piece to add to the colouring
                 colouring – Current colouring where we want to add the piece
 * returns
                 colouring - Colouring from the input with potentially a new Piece
 */
elseCondition : Piece \times Colouring \rightarrow Colouring
elseCondition(cn1, colouring) \equiv
   if \sim (\exists \text{colour} : \text{Colour} \bullet \text{colour} \in \text{colouring} \land \text{cn1} \in \text{colour}) \land \sim \text{pieceInColouring}(\text{cn1}, \text{colouring})
        then colouring \cup \{\{\text{cn1}\}\}\
        else colouring
        end,
/* This fonction shows the treatment made on a couple of Pieces
 * that should be added to the Colouring (.i.e. {Piece-set}-set)
   This fonction is recursive and calls treatmentOnPiece.
   params
                 couple – couple of Pieces to add to the colouring
                 r – relation defining the neighborhood relationships
                 colouring — Current colouring where the couple of Pieces should be added
 * returns
                 colouring - Colouring from the input with potentially a new couple of Piece
treatmentOnCouple: (Piece \times Piece) \times Relation \times Colouring \rightarrow Colouring
treatmentOnCouple(couple, r, colouring) \equiv
let (p1, p2) = couple
        in treatmentOnPiece(p1, r, treatmentOnPiece(p2, r, colouring))
/* This fonction shows the recursive treatment made on a relation of Pieces
 * that should be added to the Colouring (i.e. {Piece-set}-set)
 * First, we split the relation in a couple of Pieces and a tail using the hd method.
 * Then, we and apply the treatmentOnCouple on this couple.
 * Then, we apply recursively the treatmentOnRelation method on the remaining part of the relation.
 * This is done recursively until the relationSet is empty.
 * params
                 relationSet - relation used to define the set of pieces we are working on
                 r – relation defining the neighborhood relationships
                 colouring - Current colouring where the Pieces should be added
 * returns
```

```
colouring - Colouring from the input
 */
treatmentOnRelation: Relation \times Relation \times Colouring \rightarrow Colouring
treatmentOnRelation(relationSet, r, colouring) \equiv
  if relationSet \neq \{\}
         then let head = hd relationSet,
                           tail = relationSet \setminus \{head\}
                 in treatmentOnRelation(tail, r, treatmentOnCouple(head, r, colouring))
                 end
         else colouring
         end,
/* Main function */
tmpColour : Colouring = \{\{\}\},\
/* This function returns a correct colouring for the relation in input.
 * This function is calling: treatmentOnRelation and starting the use of recursive functions.
 * params
                 r - relation defining neighbourhood relations of piece we want to colour
 * returns
                 Colouring – colouring considering correctly the neighbourhood relations
 */
\operatorname{col} : \operatorname{Relation} \xrightarrow{\sim} \operatorname{Colouring}
col(r) \equiv
  treatmentOnRelation(r, r, tmpColour)
pre isRelation(r)
end
```

4 Testings for by Translation to SML

The following section provides tests for all the functions presented in the previous files.

4.1 Test Specification

4.1.1 Tests for ColouringBasics

The purpose of the following class is to perform unit tests on all functions of *ColouringBasics*. For all these functions, we tried to test the possible scenarios. The tests are expressed as booleans returning true if the results are the expected ones.

ColouringBasics

```
scheme testColouringBasics =
extend ColouringBasics with
        class
                value
                         P1: Piece = "Piece 1",
                         P2: Piece = "Piece 2",
                         P3: Piece = "Piece 3",
                         P4 : Piece = "Piece 4".
                         r1 : Relation = \{(P1,P2), (P1,P3), (P2,P4), (P3,P4)\},\
                         r2 : Relation = \{(P1,P1), (P1,P3), (P2,P4), (P3,P4)\},\
                         c1 : Colouring = \{\{P1, P2\}, \{P3\}\},\
                         c2 : Colouring = \{\{P1, P2\}, \{P3\}, \{P4\}\},\
                         c3 : Colouring = \{\{P1, P4\}, \{P2\}, \{P3\}\}\}
                test_case
                         [isARelation] isRelation(r1) = true,
                         [isNotRelation] isRelation(r2) = false,
                         [areNeighbours] areNb(P1, P2, r1) = \mathbf{true},
                         [areNotNeighbours] areNb(P1, P4, r1) = false,
                         [isInColouring] isInColouring(P1, c1) = \mathbf{true},
                         [isNotInColouring] isInColouring(P4, c1) = false,
                         [isSameColour] isSameColour(P1, P2, c1) = true,
                         [isNotSameColour] isSameColour(P1, P3, c1) = false,
                         [notCorrectedColoured] isCorrectColouring(c2, r1) = false,
                         [correctedColoured] is CorrectColouring(c3, r1) = true
```

end

4.1.2 Tests for ColouringEx - Auxiliary functions

The purpose of the following class is to perform unit tests on all the auxiliary functions that we have used in order to implement the colouring algorithm. For all these functions, we tried to test the possible scenarios.

Most of our tests are expressed as boolean. They return true if the result is the expected one, they return false if not. If they are not, they return a colouring. Comments are put next to the test name in order to explain the expected behaviours and results.

ColouringEx

```
/* This class extends the class ColouringEx in order to perform unit tests on
 * all auxiliary functions used for the colouring algorithm. */
scheme testColouringExAuxiliaryFunctions =
extend ColouringEx with
        class
               value
P1 : Piece = "Piece 1"
P2 : Piece = "Piece 2"
P3: Piece = "Piece 3",
P4 : Piece = "Piece 4",
P5: Piece = "Piece 5",
r : Relation = \{(P1, P2), (P2, P3), (P3, P4)\},\
r1 : Relation = \{(P1, P2), (P2, P3), (P3, P4), (P1, P5)\},\
c1 : Colour = \{P2\},\
c2 : Colour = \{P3, P4\},\
colouring : Colouring = \{\{P1, P3\}, \{P2\}\},
colouring : Colouring = \{\{P1,P3\}\}
test_case
/* canAddPieceToColour */
[canAddPieceToColour] /* False because P2 is in c1 and is neighbour with P1 */
        canAddPieceToColour(P1, r, c1) = false,
[canAddPieceToColour2] /* False because P2 is in c1 */
        canAddPieceToColour(P2, r, c1) = false,
[canAddPieceToColour3] /* Because no neighbour of P1 in c2 and P1 not in c2 */
        canAddPieceToColour(P1, r, c2) = true,
/* addPieceToColour */
[addPieceToColour] /* Should return {P2} because P1 already has neighbour in c1*/
```

```
addPieceToColour(P1, r, c1, colouring) = \{P2\},\
[addPieceToColour2] /* Should return {P2} because P2 already in c1 */
       addPieceToColour(P2, r, c1, colouring) = \{P2\},\
[addPieceToColour3] /* Should return {P2, P4} because P4 is added in c1 */
       addPieceToColour(P4, r, c1, colouring) = {P2, P4},
/* canAddPieceToOneColour */
[canAddPieceToExistingColours] /* True because P4 can be added to P2's colour */
       canAddPieceToExistingColours(P4, r, colouring) = true,
[canAddPieceToExistingColours2] /* False because P2 cannot be with P1 and P3 */
       canAddPieceToExistingColours(P2, r, colouring2) = false,
/* treatmentOnPiece */
[treatmentOnPiece] /* Should add P4 with P2 */
       treatmentOnPiece(P4, r, colouring) = \{\{P1, P3\}, \{P2, P4\}\},\
[treatmentOnPiece2] /* Should not do anything since P2 is already inside */
       treatmentOnPiece(P2, r, colouring) = colouring,
[treatmentOnPiece3] /* Should create a new colour for P2, colouring2 will be colouring */
       treatmentOnPiece(P2, r, colouring2) = colouring,
[treatmentOnPiece4] /* Should not change anything */
       treatmentOnPiece(P3, r, \{\{P1, P3\}, \{P2, P4\}\}) = \{\{P1, P3\}, \{P2, P4\}\},\
[treatmentOnPiece5] /* Should add P3 with P1 or with P5 */
       treatmentOnPiece(P3, r1, {{P2, P4}, {P1}, {P5}}),
/* treatmentOnCouple */
[treatmentOnCouple] /* should add P4 and P5 into the colouring */
       treatmentOnCouple((P5, P4), r1, colouring),
[treatmentOnCouple1] /* Should P2 and P1 in two different colours */
       treatmentOnCouple((P1, P2), r, {{}}),
[treatmentOnCouple2] /* Should add P3 with P1 */
       treatmentOnCouple((P2, P3), r, {{P1},{P2}}),
[treatmentOnCouple3] /* Should add P4 into the colour of P2 */
```

```
treatmentOnCouple((P3, P4), r, {{P1, P3},{P2}})
```

end

4.1.3 Tests for ColouringEx - Colouring Algorithm

The following class is used to test our colouring algorithm on:

- the relation of the assignment containing 9 pieces
- an other relation we created which contains 15 pieces

For each of these relations, we are performing two tests. The first one should return a colouring for the set. The second one should check if the colouring is correct -it should so return the boolean true.

ColouringEx

```
/* This class extends the class ColouringEx in order to
 * perform unit test on the colouring algorithm */
scheme testColouringEx =
extend ColouringEx with
       class
               value
P1 : Piece = "Piece 1"
P2 : Piece = "Piece 2"
P3: Piece = "Piece 3",
P4: Piece = "Piece 4",
P5: Piece = "Piece 5",
P6: Piece = "Piece 6",
P7: Piece = "Piece 7",
P8 : Piece = "Piece 8".
P9: Piece = "Piece 9".
P10 : Piece = "Piece 10"
P11 : Piece = "Piece 11".
P12 : Piece = "Piece 12",
P13 : Piece = "Piece 13",
P14 : Piece = "Piece 14",
P15 : Piece = "Piece 15",
/* Set presented in the assignment */
r : Relation = \{(P1,P2), (P1,P3), (P2,P4), (P2,P5), (P3,P4), \}
                                (P3,P7), (P4,P5), (P4,P6), (P4,P7), (P4,P8),
                                (P5,P6), (P6,P8), (P7,P8), (P7,P9), (P8,P9)
```

```
/* Other set used to test our algorithm */
r1 : Relation = \{(P1, P2), (P1, P3), (P2, P3), (P2, P4), \}
                                 (P3, P4),(P4, P5), (P4, P6), (P4, P9),
                                 (P4, P10), (P5, P6), (P7, P1), (P7, P8),
                                 (P7, P14), (P8, P3), (P8, P9), (P8, P12),
                                 (P8, P15), (P8, P14), (P9, P10), (P9, P12),
                                 (P10, P12), (P10, P11), (P11, P6), (P11, P13), (P13, P12)}
test\_case
[Assignment]
        col(r),
[isCorrectAssignment]
        isCorrectColouring(col(r), r) = true,
[Example]
        col(r1),
[isCorrectExample]
        isCorrectColouring(col(r1), r1) = true
end
```

4.2 Test Results

The results of excecuting the SML translation of the RSL test cases are:

4.2.1 Result for ColouringBasics

```
[isARelation] true
[isNotRelation] true
[areNeighbours] true
[areNotNeighbours] true
[isInColouring] true
[isNotInColouring] true
[isSameColour] true
[isNotSameColour] true
[notCorrectedColoured] true
[notCorrectedColoured2] true
[notCorrectedColoured3] true
[correctedColoured] true
```

4.2.2 Result for ColouringEx - Auxiliary functions

```
[canAddPieceToColour] true
[canAddPieceToColour2] true
[canAddPieceToColour3] true
[addPieceToColour] true
[addPieceToColour2] true
[addPieceToColour3] true
[canAddPieceToExistingColours] true
[canAddPieceToExistingColours2] true
[treatmentOnPiece] true
[treatmentOnPiece2] true
[treatmentOnPiece3] true
[treatmentOnPiece4] true
[treatmentOnPiece5] {{"Piece 1"},{"Piece 4","Piece 2"},{"Piece 5","Piece 3"}} [treatmentOnCouple] {{"Piece 3","Piece 1"},{"Piece 4","Piece 2","Piece 5"}}
[treatmentOnCouple1] {{"Piece 2"},{"Piece 1"}}
[treatmentOnCouple2] {{"Piece 2"},{"Piece 1","Piece 3"}}
[treatmentOnCouple3] {{"Piece 3","Piece 1"},{"Piece 2","Piece 4"}}
4.2.3 Result for ColouringEx - Colouring Algorithm
[Assignment] {{"Piece 8","Piece 1"},{"Piece 4"},
{"Piece 3", "Piece 9", "Piece 2", "Piece 6"}, {"Piece 7", "Piece 5"}}
[isCorrectAssignment] true
[Example] {{"Piece 11","Piece 12","Piece 2","Piece 14"},
{"Piece 4","Piece 8"},{"Piece 9","Piece 1","Piece 13","Piece 6","Piece 15"},
{"Piece 5","Piece 3","Piece 10","Piece 7"}}
[isCorrectExample] true
```

5 Conclusion

In this project, we started by specifying our requirements on relations and colourings using RSL. We used this two functions in an implicit function *Colour-Req* which could not be translated to SML. We implemented so an explicit function creating a correct colouring for a well formed relation.

To create this function, we had to create recursive methods and auxiliary ones to perform tasks. We also have been creating unit tests in order to test all of them and to detect problems while implementing these functions. The implementation of the *col* function in *ColouringEx* was beneficial for us since we were not used to functional programming and to the use of recursion in algorithm. We made then tests to check this *col* function. From this test, we can conclude that our results are matching the requirements we had explained previously.