possible binding site as a matrix. When all of the sequences have been added to the analysis, the one with the highest information content is the best guess for the pattern of the binding site. The method works well on typical prokaryotic binding sites, and it is robust enough to work even with some erroneous data included.^{20,21}

Summary

Matrices can provide realistic representations of protein/DNA specificity. In many cases simple mononucleotide-based matrices are adequate representations, but more complex matrices may be needed for other cases. Unlike simple consensus sequences, matrices allow for different penalties to be assessed for different changes to a binding site, a property that is essential for accurate description of a binding site pattern. When only a collection of binding site sequences is known, the best representation for the pattern is an information content formulation, based on both thermodynamic and statistical considerations. Quantitative data on relative binding affinities may be used to determine matrices that provide a best fit to the data. Matrix representations also provide an efficient method of aligning multiple sequences to identify binding site patterns that they have in common.

Acknowledgments

I thank Calvin Harley for sending me the *E. coli* promoter sequences on a floppy disk and George Hartzell for helping with the figures and Gerald Hertz for critical comments on the manuscript. Supported by National Institutes of Health Grant GM28755.

²¹ G. Z. Hertz, G. W. Hartzell III, and G. D. Stormo, Comput. Appl. Biosci., in press.

[14] Consensus Methods for DNA and Protein Sequence Alignment

By Michael S. Waterman and Robert Jones

Introduction

The increasing body of nucleic acid sequence data has created interest among many scientists in computational approaches to macromolecular sequence analysis. Several international databases have been created in order to store the data in a useful format, both for archival and analysis purposes.¹ Both DNA and protein sequences databases are maintained. The value of simply having easy access to all membrane protein sequences, for example, is not to be underestimated. The quantity of data has naturally led to the development of computer approaches to sequence analysis.² The purpose of this chapter is to present some of the tools that we have created in order to analyze multiple sequences in a rigorous, efficient, and systematic way.

Much computer analysis of molecular sequences is directed toward discovery of biologically significant patterns. These patterns include homologous genes, RNA secondary structure, tRNA or structural RNAs, palindromes in DNA sequences, regulatory patterns in promoter regions, and protein structural patterns. Once the patterns have been located they can often be tested by experiment, as in the case of promoter elements. Evolutionary relationships, however, cannot be directly tested, and increasing emphasis is being attached to the discovery and interpretation of sequence evolution.

Sequence alignment is a popular approach to pattern analysis.² Computer alignments are often based on an explicit optimization function, rewarding matches and penalizing mismatches, insertions, and deletions. Sequence alignment often gives useful information about evolutionary or functional relationships between sequences. Our approach is based on what we refer to as consensus analysis.²⁻⁴

Consensus sequence analysis is usually performed by visual inspection of the sequences and by experiment. Of course, a protein binding site can only be verified by experiment, and analysis by "eye" can be biased. Thus, it is useful to have computer methods that can find consensus patterns best fitting explicitly stated criteria. Some algorithms have been developed along these lines,²⁻⁴ and they are described here, along with some biological examples. Our earlier methods applied only to DNA; here we also describe recent extensions to protein sequences.

In 1970 Needleman and Wunsch⁵ published an approach sequence comparison (alignment) using a dynamic programming algorithm. Their algorithm find maximum similarity between two sequences, where matches score positive weight and mismatches, insertions, and deletions

¹ C. Burks, J. W. Fickett, W. B. Goad, M. Kanehisa, F. I. Lewitter, W. P. Rindone. C. D. Swindell, C.-S. Tung, and H. S. Bilofsky, *CABIOS* 1, 225 (1985).

² M. S. Waterman, ed., "Mathematical Methods for DNA Sequences." CRC Press, Boca Raton, Florida, 1988.

³ M. S. Waterman, D. Galas, and R. Arratia, Bull. Math. Biol. 46, 515 (1984).

⁴ D. J. Galas, M. Eggert, and M. S. Waterman, J. Mol. Biol. 186, 117 (1985).

⁵ S. B. Needleman and C. Wunsch, J. Mol. Biol. 48, 444 (1970).

score nonpositive weight. Mathematicians began to attempt to define a distance between sequences and so to construct a metric space. Sellers⁶ obtained these results for single insertions and deletions, and later workers extended the work to multiple insertions and deletions.⁷ While dynamic programming methods are very widespread in sequence analysis, there are severe restrictions in computation time with the extension of the dynamic programming methods to allow more than two sequences. A great many biological problems do involve more than two sequences. The consensus methods we have developed avoid the computational difficulties of dynamic programming by using a very different approach to sequence analysis.

The basis of the consensus method is an algorithm to find consensus words, with the degree of matching and alignment specified by the user of the program. We give the specifications of this method in the next section for DNA and protein sequences, along with examples. In this setting the consensus method finds patterns or words that are conserved in an unusual number of sequences. Then the basic method is extended, both for DNA and protein sequences, to an algorithm for sequence alignment. To illustrate the behavior of the algorithms, we have chosen two sequence sets, one DNA and the other protein. The DNA sequences are 19 promoters from the genome of vaccinia virus.8 The protein sequence set is 16 proteins related to the Escherichia coli ntrC gene product. We use these sequences to illustrate the use and power of the programs and the effect of varying certain parameters. We do not attempt to interpret the consensus pattern found in any biological context, but we invite anyone interested in these specific sequences and patterns to contact us for more detailed information.

Consensus Patterns

Now we give a general description of the consensus word algorithm. To begin, take a set of R sequences of length N

⁶ P. Sellers, SIAM J. Appl. Math. 26, 787 (1974).

⁷ M. S. Waterman, T. F. Smith, and W. A. Beyer, Adv. Math. 20, 367 (1976).

⁸ M. Mars and G. Beaud, J. Mol. Biol. 198, 619 (1987).

⁹ B. T. Nixon, C. W. Ronson, and F. M. Ausubel, *Proc. Natl. Acad. Sci. U.S.A.* 83, 7850 (1986).

These sequences can be taken to be initially aligned on some biologically or statistically determined feature. The true alignment is, of course, unknown except approximately. Now we give an algorithm for locating consensus words of a given size. By way of comparison, the usual methods of sequence analysis align on single letters, that is words of length 1.

Of course, a concept basic to our algorithm is that of consensus word. The definition has been given in earlier work²⁻⁴ and will be briefly reviewed here. First, take a fixed word size k and a word w of length k; there are 4^k such words in DNA and 20^k in proteins. Next, define the window width W which gives the width of sequence in which a consensus word can be found and thus defines the amount of shifting allowed in matching consensus words. The sequences starting at column j+1 with window width W appear as

To begin, we search the first sequence of the window for matches to our word w. An exact match to w is called a d=0 neighbor while a 1-letter mismatch from w is called a d=1 neighbor, and so on. For protein sequences, for example, it is desirable to distinguish the many types of d=1 mismatches by different weightings based on amino acid similarity. It is possible to include insertions and deletions in this list of neighbors. We may decide, e.g., to limit the amount of mismatch to d=0, 1, 2 and not find w in a portion of sequence unless it is within this neighborhood. Let $q_{w,d}$ equal the number of lines that the best occurrence of w is as a dth neighbor. Each of these occurrences receives weight λ_d . The score of word w in this window is

$$S_{j+1,j+W}(w) = \sum_{d} \lambda_{d} q_{w,d}$$

A best scoring word is word w^* satisfying

$$s_{j+1,j+W}(w^*) = \max_{w} s_{j+1,j+W}(w)$$

DNA Consensus Patterns

In the case of DNA all d = 1, d = 2, . . . mismatches are considered identical in weight. While more complex weighting schemes are easy to

incorporate, we have found that it is adequate to score a word by the fraction of letters matching the consensus word. Thus, for a *d*-letter mismatch to a *k*-letter consensus word, $\lambda_d = 1 - d/k$.

To perform the computations, each word in the window is read and its neighborhood calculated. The possible words in a neighborhood are found by a simple combinatoric scheme; since all d=2 mismatches receive the same weight it is only necessary to enumerate all k(k-1)/2 of these mismatches. When the portion of each sequence in the window has been examined, the best score each of the 4^k possible consensus words is retained.

Our consensus method for DNA sequences has been implemented in a program called RTIDE written in C and using the SunView window system. Figure 1 shows a typical screen display from this program. The aligned sequences are displayed at the bottom, the displayed consensus word and score of the word are shown above, along with parameters for the run. At top is a plot of the consensus word score against window position for the alignment. Peaks in this plot indicate regions of conservation that

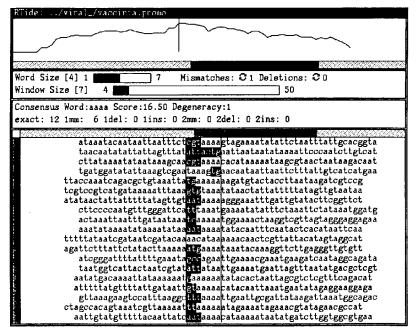


FIG. 1. Typical screen display of the program RTIDE. The sequence alignment is shown in the bottom window. At top is a plot of maximum score against window position in which peaks represent conserved regions.

may be of biological interest; these patterns can then be examined in detail and sequences realigned on the consensus words in any particular window. Through cycles of realignment and analysis it is possible to identify and refine conserved sequence patterns. The nature of the method involves variation of several parameter settings during an analysis session. It is possible to adjust window width, word size, and degree of matching required as well as alignment of the sequences. To facilitate this interaction all features of the program are controlled through the mouse.

The effect of varying the word size on the consensus score is shown in Fig. 2. No mismatches are allowed in these runs, and the window size is increased along with word size to keep the number of words per window constant. There are 16 words of size 2, and the score is the maximum scoring word over all sequences. Since there are so few words, all window positions have a high score and no features are clearly resolved. At word size 3 the graph shows a few features, with the highest score indicated by the dashed line. At word size 4 the central peak is resolved from the background. As word size increases further this peak becomes smaller and

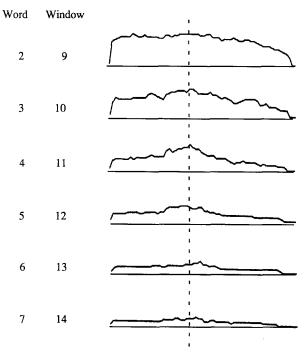


Fig. 2. Effect of varying word size from 2 to 7 for DNA sequence alignment. The window size is varied to maintain a fixed number of words per window. No mismatches are allowed.

merges into the background of nonconserved sequence. The word size at which a peak is most evident is a clear indication of the size of the feature that is conserved; in the present case that word size is 4 nucleotides.

In many cases the initial alignment in which sequences are supplied will not be optimal for a given conserved sequence feature. Varying the window size is a way to accommodate poorly aligned sequence sets. Figure 3 shows the effect of variation of window size from 4 to 20 for the DNA data set. Using a window equal to word size permits no misalignment and for our sequences gives a very low graph of scores. Widening the window brings more instances of a consensus word into a window. This is demonstrated in Fig. 3 by the appearance of a peak in the graphs, most clearly resolved at window sizes 10 and 12. Extending the window further may not bring any new instances of the consensus word, but it does increase the number of window positions that achieve a high score. This is shown by the plateau in the graphs for window sizes 16 and 20.

The effect of varying the neighborhood of words that can contribute to the score of a consensus word is shown in Fig. 4 for our DNA data set. We fix a window size of 12 and a word size of 6 and vary the number of mismatches permitted from 0 to 2, with no insertions or deletions. Requiring exact matches (0 mismatches) with the consensus word results in low scores with no distinct features. Permitting a single mismatch in general causes scores to increase, but the central conserved region emerges as a

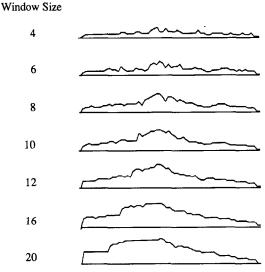


Fig. 3. Effect of varying window size from 4 to 20 for DNA sequence alignment. The word size is 4, and no mismatches are allowed.

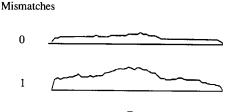


Fig. 4. Effect of varying the number of mismatches from 0 to 2 for DNA sequence alignment. The word size is 6, and the window size is 12.

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distinct feature. When the neighborhood is increased to 2 mismatches, however, the background scores are almost equivalent to that of the conserved feature.

A very useful feature of our program is that once a consensus word has been identified, the sequences can be realigned on that word and the new alignment reevaluated for additional conserved features. Figure 5 shows the refinement of a consensus word using this technique. The graph of the scores after cycle 1 shows a maximum score of 14.25 at the position marked by the solid line. The sequences were realigned on the words that

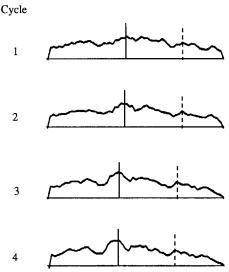


FIG. 5. Effect of realigning the sequences on a consensus word for DNA sequence alignment. The word size is 4, window size is 6, and 1 mismatch is allowed.

```
A C D E F G H I K L M N P Q R S T V W Y
Α
     8
C
     6
       8
     3
D
       3 8
     1
       2
E
          6 8
F
     4
       5 1 2 8
     7
       6 2 1
G
              4 8
       2 2 3 4
                1 8
Η
     1
            2
       5
                  2
Ι
     4
          1
              5
                4
       3
     2
         3 4
              3 2 6
                     3 8
K
       5 1 2 5 4
6 2 3 6 5
                    7
L
     4
              5 4
                  2
                       3
     5
Μ
                  3 6 4
                         6 8
     4
       5
         5
           4
              3 4
                  2
                     3
                       3
                         3
N
          3
            2
       5
              3 4
Ρ
     4
                  0
                     3
                       1
                         3
                           4
                             5 8
     3
       4
          4 5
              4
                3
                  3
                    4
Q
                      4
                         4
                           5
                             6 4 8
         4 5 2
     1
       2
               1
                  5
                     2
                      6
                         2
                           3 4 2 5 8
R
     5 4 4 3 2 5
                     2
                       2
                         2
                           3 6
S
                  1
                               4 5 3 8
     5
           3 4
                5
Т
       6 4
                  3 4
                           5
                                  5
                      4
                         4
                             6
                               4
     5 6 2 1
              4 5 1 6
                       2
                         6
                           5
                                  3
                                    1 3 5 8
V
                             4
                               4
W
     3 4 2 3
              6 3 5 4 4 4 5
                             4
                               2
                                  5
                                    3
                                      3 5 3 8
     3 4 2 3 6 3 5 4
                       4 4 5 4
                                2 5
Υ
     ACDEFGHIKLMNPQRSTVWY
```

Fig. 6. Similarity matrix used in weighting sequence mismatches, based on the representation of amino acid similarity of Taylor.¹⁰

contributed to that score and the program run again. In the second cycle the marked peak is more clearly resolved and its score has risen to 15.25. An additional change is that the minor peak marked by the dashed line has become more evident. At the end of the third cycle the main peak has become more distinct and the score is 16.75, resulting from the realignment bringing more related words into the conserved window. The final cycle of scoring does not increase the maximum score and indicates the end of the process.

Protein Consensus Patterns

In the case of protein sequences various mismatches are weighted according to a matrix (Fig. 6), which is derived from Taylor.¹⁰ The mismatches for each letter of a word are arranged according to weight, the nearest or smallest being first. Then in a systematic fashion we allocate mismatches until the limit or cutoff is reached. Then that letter is reduced to identity and the next letter is increased. The algorithm is similar to that of counting with the branch and bound feature we have described.

¹⁰ W. R. Taylor, J. Theor. Biol. 119, 205 (1986).

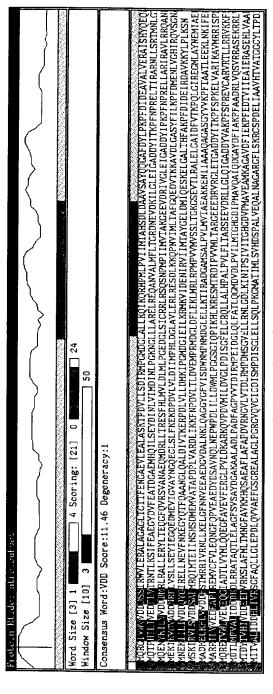


Fig. 7. Typical screen display from the program PRTIDE.

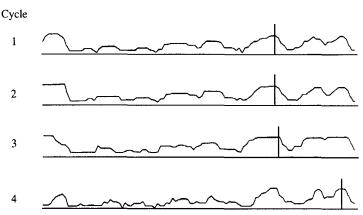


FIG. 8. Effect of realigning sequences on a consensus word over four cycles for the protein sequence data set. The word size is 3, window size is initially 10, and the neighborhood is set to 21, where 24 implies an exact match.

In the protein version of the consensus program, PRTIDE, the only substantial modification is in the definition of neighborhood. Whereas in the DNA version we specify the number of mismatches, insertions or deletions allowed, in the case of proteins we specify a similarity score that a word must attain when scored against a candidate consensus word in order to be included in the neighborhood. Figure 7 shows a typical display of this output.

In most practical aspects the DNA and protein programs are very similar. The larger alphabet of amino acids relative to nucleotides precludes use of words with more than 4 amino acids. Figure 8 shows an example of the protein version in which the sequences are realigned on a consensus word, with the result of resolving other conserved words. At top is shown the graph of scores for the protein data set in which the sequences are aligned at their left ends. The largest peak is found close to this end; aligning on the consensus word and rescoring creates the second plot in which three peaks near the right end of the plot have become more clearly resolved. Aligning on the largest of these and rescoring cause the formation of plateaus. This indicates too large a window size, and reducing the window size from 10 to 5 results in the final plot in which the three peaks are now well resolved.

Consensus Alignment

The idea of the algorithm builds on the previous section. We align on consensus words, attempting to maximize the sum of the scores of the

¹¹ M. S. Waterman and M. Eggert, Nucleic Acids Res. 14, 9095 (1986).

words. Before the practical algorithms are presented, a more general concept of alignment on words is presented.

We define a partial order on words. The words w_1 and w_2 satisfy $w_1 < w_2$ if the occurrences of w_1 in sequence i are to the left of the occurrences of w_2 in sequence i (and do not intersect) for i = 1 to R. It is not necessary for w_1 or w_2 to have occurrences in all sequences. Implicit in the definition is a window width W and neighborhood specification. The goal of an optimal alignment is to find words w_i which satisfy

$$\max \left\{ \sum_{i \ge 1} s(w_i) : w_1 < w_2 < \dots \right\}$$

(It is frequently desirable to require $s(w_i) \ge c$ for all i, where c is some cutoff value.) It is not possible to accomplish this goal in reasonable time, but it is possible to come quite close. We now define two practical algorithms.

Next, $w_1|w_2$ means that consensus words w_1 and w_2 can be found in nonoverlapping windows, each word satisfying as usual the window width and neighborhood constraints. The modified optimization problem is to satisfy

$$T = \max \left\{ \sum_{i \ge 1} s(w_i) : w_1 |w_2| \dots \right\}$$

There is a straightforward recursion to find T. Let T_i be the maximum sum for the sequences from base 1 to base i:

$$a_{1,1}$$
 \cdots $a_{1,i}$
 $a_{2,1}$ \cdots $a_{2,i}$
 \vdots
 \vdots
 \vdots
 \vdots
 \vdots
 \vdots

Then T_i satisfies

$$T_i = \max\{T_j + s_{j+1,i}: i - W + 1 \le j \le i - k\}$$

and $T_{-W} = T_{-W+1} = \cdots = T_0 = T_1 = \cdots = T_{k-1} = 0$. Also, $s_{x,y} = 0$ if y - x + 1 < k. This algorithm runs in time approximately proportional to NW^2RB where B is the neighborhood size. Here the factor WRB accounts for the consensus word algorithm with a window width W. (This is an overestimate since the actual windows vary from k to W in width.)

If much shifting is necessary to match the sequences, T is an underestimate and misses some of the relevant matching. To overcome this prob-

lem, the definition of T is modified to

$$S_i = \max\{S_j + \hat{s}_{j+1,i}: i - W + 1 \le j \le i - k\}$$

where $\hat{s}_{j+1,i}$ is the largest scoring consensus word in the window from j+1 to i such that all occurrences of the consensus word are to the right of the consensus words for S_j . This algorithm is not guaranteed to be equal to the global maximum, but it is much more useful than T. Alignment for each case, DNA and protein, proceeds as just described with the modifications given next.

DNA Alignment

We have written a program, RALIGN, to align multiple DNA or RNA sequences.¹¹ The sequences are supplied in some initial alignment, which usually consists of the sequences being left justified. As in all our consensus methods, the parameters for window size, word size, and neighborhood are set. We require a consensus word to have a score at least equal to one-half the number of sequences before it can be used for alignment. This requirement is to eliminate the "junk" regions in alignments that are common with programs that optimize a total score.

To illustrate RALIGN for nucleic acids we take a set of 15 tRNA sequences from Escherichia coli. These sequences are difficult to align as their relationship is largely determined by conserved helices (secondary and tertiary structure), not the primary sequence itself. Analyses that fold tRNAs by minimum free energy are not too successful, usually folding about 50% of the tRNAs into the correct cloverleaf shape. An analysis based on consensus helices is successful for many structural RNAs, and a study of tRNA by consensus folding appears in Ref. 2. The only universal primary sequence patterns in tRNAs are the CCA at the acceptor arm and the GTTC in the $T\psi C$ stem and loop. Our analyses (Fig. 9) find these invariant patterns along with other conserved words. Figure 9a has window size 7, word size 3, and up to 1 mismatch (total score equal to 179), while Fig. 9b has window size 8, word size 4, and up to 2 mismatches (total score equal to 127). The window is 7 in the first case and 8 in the second to allow shifts of 4 in each analysis. Notice that CCA is located at the 3' end of all the sequences; in the case of Fig. 9 CCA is generally found overlapping the pattern CACC. In Fig. 9a,b CCA is not found by the program in all the sequences. In several cases an earlier, equally strong pattern is chosen in accord with the algorithm. Of course in Fig. 9b the object is to optimize a 4-letter consensus word, and this weakens the contribution of the 3-letter pattern CCA. GTTC is always located in Fig. 9b between consensus words GGTT and CGAA. In 14 of 15 sequences in Fig. 9a GTTC is located

٦		

cca	cacca	,a	ccdcca	, a	acccca	tacca	sacca	acca	3a	3a	a c	2a	cacca	cacca	e c
	S	S	S	og oc	ő	ö	ö	ö	ŏ	အရှင်င	ca cc	ŭ	ŏ	ď	g Ö
.ctc	cgcatagct	acgegeet	gtccgtt	ccctgggggac		ctcagg									ctccgcccctg
.atc	atcc	ctccgga	gtccc	atc	atcccat	gtcca	atc	atccc	gtcca	atc	atc	atctgc	atccc	gtcca	gtct
.tcg	tcg	tcga	tcga	tcga	tcga	tcaa	tcga	tcga	tcga	tcga	tcga	tcga	tcg	tcga	tcga
.ggt	ggt	ggt	ggt	ggt	ggt	ggt	ggt	ggt	ggt	ggt	ggt	agt	ggt	ggt	ggagt
.tcg	tctgc	tagtcc	togog	agg	tcgtg	gtcggt	tcgca	tcaca	tcact	tccga	tcgga	tagga	tagga	tcggt	ttg
	caggagg tctgc	cgtc	caggggg	cggtaac	cagttg	gggtgag	tgg	tgggg	ccgtatg	cggcat	ടർ	agg	aggggg	tggggg	cgggtg
ctt	ctttgcacg	attgcaaatc	cctgtcacg	cttcacgg	attgtgattc	cctgataa	cttttaatcaat	ctcataatga	ctgttaat	attctgattc	ctgcgaaccgag	cttggtaagggtg	cttacaagg	cttgacatgg	ctccaaaac
. agc	Ę,	agcgg	acctg	accdcc	agccctgg	agegeace	agcagttga	agcacatca	aacggcgga	accgg	agtactcgg	agegeace	agcacctcc	agcaccac	aacaccggt
tag	gag	tatgt	teggttagaat	caggac	tag	tag	tag	tag	tag	taaggc	tag	tag	gag	tag	tag
tgg.	tgg	cggt	teggt	aggco	ttgg	tggt	tgg	tggt	tagg	cgg	tgg	tgg	tgg	tggt	tgg
cag			cag												
.tag	tagct	taa	tagtt	tegte	tagct	tagct	tagct	tagct	tagtt	taga	tagct	tagct	tagct	tagct	tagtt
.gct	ggg gcta tagct	gcgcgt	gcgg	scct	ggcta	gcttg	ggtcgt	gctacg	cctctg	ggta g	gcatccg	gctgata	ggtgat	gcgtccg	аддддсд
:	999	Б	gga	gtc	ggt	ag	p	р	ų	tgg			Б		agg

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																•
.cacc	cacca	ctcca	cdcca	cdcca	cacccca	tacca	cacca	cacca	gagcca	cagcca	cacca	cacca	cacca	cacca	ccctdcca	
.tccc	cccgcatagctc	tccggaacgcgc	tcccgtccgttc	tcccctggggga	tcccattagc	tccactcaggcc	tcctgcacgacc	tcccgtcgtagc	tccagtcagag	tcctcgtaccc	tcctcccggatg	tgcctatcag	cccgtcatcacc	tccactcggacg	tctcccgc	
cgaa	cgat	cgac	cgag	cgaa	cgaa	caag	cgaa	cgaa	cgag	cgaa	cgaa	cgaatc	cgat	cgag	cgag	·
.ggtt	ggtt	ggtt	ggtt	ggtt	ggtt	ggtt	ggtt	ggtt	ggtt	ggtt	ggtt	agtt	ggtt	ggtt	agtt	
ggtc	ggtctgc	agtcc	ggtcgcg	ggtaacagg	tgtcgtg	ggtcggt	ggtcgca	ggtcaca	tgtcact	attccga	ggtcgga	ggtcggc	ggtcggc	ggtcggt	ggtgttggg	1
.cacc	cgcctgctttgcacgcagga	tagoggattgcaaatcogtct	tacctgcctgtcacgcaggg	caccgccctttcacggc	ccctggattgtgattccagt	caccctgataagggtga	cagttgacttttaatcaatt	cacatcactcataatgatgg	cggcggactgttaatccgta	caccggattctgattccggc	tactcggctgcgaaccgagc	cacccttggtaagggtga	cacctcccttacaaggaggg	caccaccttgacatggtggg	caccggtctccaaaaccg	
agag	agag	atg	agaa	ccagga	agag	agagcg	agag	agag	agaa	aagg	agag	agagcg	agag	agag	agaa	
.tggtagag.	tggg	cggtta	cggtt	aggc	tggt agag	tggtt	tggt	tggtt	cggt	cggt	tggt	tggt	tggg	tggtt	tggt	
tcag	tcagc	aaag	tcagt	ctag	tcagt	_	_	_	_	Ŭ	_	-	tcagc	tcagt	tcaat	
tagc	tago	taac	tagt		atagc	tagc	tage	tagc	tagt	tatcg	tagc	tagc	tagc	tagc	tagt	
:	ggggcta	ggcgcgt	ggagcgg	gtcccct	ggtggct	aggettg tage	gggtcgt	ggctacg	tectetg	tgggg	gcatccg	gctgata	gggtgat tagc	gcgtccg	aggggcg	

Fig. 9. Multiple sequence alignment of 15 tRNA sequences by the program RALIGN. (a) The word size is 3 with up to 1 mismatch in a window of 7. (b) The word size is 4 with up to 2 mismatches in a window of 8.

between consensus words GGT and TCG. Note the failure in the last sequence where the alignment is

Consensus pattern: ggt . . . tcg Sequence 16: ggagt tcg

This can be accounted for by the "greedy" nature of our algorithm. To allow more chances the consensus matching, whenever there are ties in scoring the algorithm chooses the 5' or leftmost pattern. Therefore, gga is aligned rather than agt, the biologically correct alignment.

Protein Alignment

Figure 10 shows an example of the program PRALIGN applied to nine protein sequences. The sequences represent the amino-terminal 70 residues from a number of regulatory proteins related to the *E. coli ntrC* gene product. This region is fairly well conserved among the proteins and is believed to be responsible for interaction with proteins related to the *E. coli* gene product.

Figure 10 uses a word size of 3 and a window of 6. The neighborhood is limited to a similarity score of 18 or more. Since a perfect match receives a score of 8, an exact matching 3-letter word has score $24 = 3 \times 8$. As with Fig. 9, the top line of the alignment shows the conserved words identified at each position, and the instances of those words are shown in upper case in the sequence alignment below. With the neighborhood as specified, the conserved words are essentially those that contain conservative replacements from the consensus words. The majority of the amino acids in the sequences have been identified as part of conserved words and have been brought into alignment. The sequences are sufficiently dissimilar, however, that a number of short segments are not part of any conserved word.

Conclusion

The programs are available from Waterman. The programs RTIDE and PRTIDE are both written in C and run on a SUN using the SunView windows system. The programs RALIGN and PRALIGN are written in C and do not require the special graphics interface.

Several generalizations of these ideas are possible. One of the most obvious is to apply the methods to single rather than multiple sequences. Our programs for single sequences find the maximal nonoverlapping repeat pattern; as for the programs described in this chapter there are two programs, one for DNA and one for proteins. Elsewhere we have reported methods to find consensus palindromes in DNA, both for multiple and

mgr.kt]	.lvd.	dnc	irg.	lve	cln.	ded	fqv	qav	ena	ec1	lnk	vbd.	.111	dim.	.mpt
MQRgive	, wybd	DSS	IRW	VLEr	ALAge	gltct	Į.		ENGa	EVLeal	ASKt	PDV	LLS	DIR	MPG
MOTPHII	IVEde	ιlν	TRNt 1	ksi		FEAeg	YDVf	EAT	DGAemh	qi	MQTpHIL IVEdelv TRNtlksi FEAeg YDVf EAT DGAemhqi LSEydin LVIm DIN LPG		LVIm	DIN	LPG
MQEnyk i	. LWd	DDMr	LRA	LLEr	YLT	EQG	FQV	RSVana		EQMdr1	LTResfh		LM7	DLM	LPGedg
MEKiKVC	: VAD	DNRe1	vs]	LSEY		IEGged	MEV	IGVa	YNGq	ECLS	LFKekd	PDV	LVL	DII	MPH
MNE KI	TIVD	DQYg	IREI	LNE	VEN	KEG	YOTf	QAAng]		QALdi	VTKer	PDL	VLL	DMK	IPG
MSKiRVI	SVD	DSAl	MROi	MTE	IINS	sdm		EMVata	Ωı	DPLvaro	LIKkfn	PDV	LTLdve	a)	MPR
mad KEI	rkflvv	DDFst	mrr	IVRn]	llkelg	ıξ	NNV	EEAe	DGVda	LNKlqac	удуд	FVIsd	wnw	PNMdg	lel
MAR RII	, WE	DEAp	IREmv	cf	VLE	ONG	FQPve	aedy	DSAvnq		LNEpw	PDL	ILL	DWM	LPG
MOR EIV	/wLVEd	EOGlad	dt	LVY	MLO	OEG	FAVev	Ţ	ERGlpv	ldkarko	2.	PDV	MIL	DVG	LPD

KYLplksn EGMlayn ARV RTLlr ARI KAVmr EAV ALVERAIS ELL ARI RAVÌ ELTirARNllsr cng...lql...lrr....lkn...nip..vmm....ltv....hge....dei...glq..iga...dfa...spf..cpk...eic..akv.kgl EIRdavk QLGir EΓΛ EVC ATLeekl SYFilkpfdmenlvghir SPK SPR NPR KPF NPR KPFdid KPFdid KPF KPF KPF ELGaidf VVKpfta DYIP HEA DYIt DYVa DYIL DYLp LGAlt DLGa QGA£ IGAd IGAd IGAd TGAd SGY SEV TLRal GLO GLE GLE AYO GLE IAAaqaga DMIqeske EDVtkkav DRIV DRVr DRL1 HSDldaavs RSEev KGEev RGEee IMLtafgq MTA YGEl LTGrdnev VTA **WWWwssltgkg** LTA LIA LPVlmvTAEakkeni Š VLF HIM IΙΛ LMF MKVide NIR VII LREsd1kkgpnv DIP ALP LRSgsn PMP PMP IKOrhp MLP LREga NVA GAMsa FEL CROILS LHP LEKlmrlr IQPikhLKResmtr LKO LER LKR CRR LKTirad LLLare LAL ISI LDGlav MDG1df GSG MDG KNG ISG

Fig. 10. Multiple sequence alignment of nine protein sequences using the program PRALIGN. The word size is 3 in a window of 6 positions. The neighborhood is limited to words with a similarity score of at least 18, where a score of 24 represents an exact match.

single sequences.² In addition, these methods can also be applied to consensus secondary structure. Finally, although we have not done so, it is possible to include other ideas of consensus such as gap length to improve the alignments.

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[15] k-Tuple Frequency Analysis: From Intron/Exon Discrimination to T-Cell Epitope Mapping

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Introduction

To determine the function of a gene or a protein from mere inspection of its sequence is one of the ultimate goals of research in sequence analysis. In principle, the solution of the protein folding problem, among others, stands as a prerequisite to this accomplishment. Meanwhile, shortcuts have been found, and several methods, mostly based on the recognition of conserved features in genes or proteins with similar functions, can be used to give at least a partial answer to this general problem. For instance, a straightforward way to assign a function to a given sequence is to run exhaustive homology searches on available data banks. As a rule of thumb, an overall 30% amino acid identity (over several hundred residues) between two protein sequences is taken as strongly suggestive of a similar three-dimensional fold and, thus, function.^{1,2}

Alternatively, the simultaneous consideration of a set of divergent sequences of known similar function is used to establish characteristic consensus patterns which can be searched for in unassigned candidate sequences. These patterns can be defined on a small number of positions, be highly degenerate, and yet retain a good discriminative power. Such functional signatures usually correspond to key residues in the active (or binding) site of proteins or key nucleotides to be recognized within genes

¹ C. Chothia and A. M. Lesk, *EMBO J.* 5,823 (1986).

² R. F. Doolittle, in "Of URFs and ORFs." University Science Books, Mill Valley, CA, 1986.