

|

Build your own language

Building a extremely simple language

Supporting

- ▶ Numbers
- ▶ Addition
- ▶ Functions
- ▶ Typechecking

Examples

```
42
41 + 1
1 + 2 + 39
(\ x : Int . x + 1) 41
(\f : Int -> Int . f 41) (\x : Int . x + 1)
```

Examples

For a first version, we'll omit all types:

```
42
41 + 1
1 + 2 + 39
(\ x . x + 1) 41
(\f . f 41) (\x . x + 1)
1 + 41
```

ALGT

github.com/pietervdn/ALGT

All the files you need are in the **demo**-directory, download it entirely

ALGT

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Alpha

Beware of bugs and unclear error messages

II

Declaring the syntax

What is the syntax of the language?

Syntax = *what it looks like*

Backus-Naur-Formulation

What is BNF?

```
name ::= "literal1" | "literal2" | "literal3"
```

Backus-Naur-Formulation

What is BNF?

```
name ::= "literal1" | "literal2" | "literal3"
```

Possible files:

```
literal1
```

OR

```
literal2
```

OR

```
literal3
```

Syntactic form names are written with a **lowercase letter**

Backus-Naur-Formulation

What is BNF?

```
name ::= "literal1" "literal2"
```

Backus-Naur-Formulation

What is BNF?

```
name ::= "literal1" "literal2"
```

Possible files:

```
literal1 literal2
```

Whitespace is ignored by default! See the manual for other options

Backus-Naur-Formulation

```
name          ::= "literal1"  
otherName     ::= name
```

Backus-Naur-Formulation

```
name          ::= "literal1"  
otherName     ::= name
```

Possible files:

```
literal1
```

Backus-Naur-Formulation

```
name          ::= "literal0"  
otherName     ::= name name | "literal1" name | "literal2"
```

Backus-Naur-Formulation

```
name          ::= "literal0"  
otherName     ::= name name | "literal1" name | "literal2"
```

Possible files:

```
literal0 literal0
```

OR

```
literal1 literal0
```

OR

```
literal2
```


Backus-Naur-Formulation

Defining numbers:

```
int ::= "0" | "1" | "2" ...
```

Backus-Naur-Formulation

Defining variables:

```
var ::= "a" | "b" | "c" | ... | "someVariableName" | ...
```

Backus-Naur-Formulation

Too much work. . .

Special rules **Number** and **Identifier** have been provided as builtin

Builtins are written with an *uppercase*

How do we define a language?

In the file *Demo.language*

```
Demo
*****

Syntax
=====

int      ::= Number
```

Parsing!

```
./ALGT Demo.language demo.demo int -l
```

```
# "42" was parsed as:  
42      Number.0
```

Addition

Syntax

=====

```
int      ::= Number
expr     ::= int "+" int
```

Addition

```
41 + 1
```

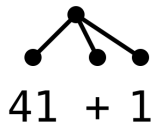
```
./ALGT Demo.language demo.demo expr -l
```

```
# "41 + 1" was parsed as:  
+ 41      Number.0  
| +       expr.0  
| 1       Number.0
```

Addition

Use extra flag `-ptsvg` Name to create an image of your parsetree:

expr.0



Addition

42

```
"demo.demo (line 0)" (line 1, column 3):  
unexpected end of input  
expecting "+"
```

Addition

```
Syntax  
=====
```

```
int      ::= Number  
expr     ::= int "+" int  
          | int
```

```
# "42" was parsed as:  
42      Number.0
```

Addition

What with $1 + 2 + 3$?

```
expr    ::= int "+" int  
        | int
```

Addition

Let's change expression to be recursive!

```
expr    ::= expr "+" expr  
         | int
```

Left recursion

```
int    ::= Number
expr   ::= expr "+" expr
        | int
```

```
Error:
  While checking file DemoDynB.language:
    While checking the syntax:
      Potential infinite left recursion detected in the syntax.
      Left cycles are:
        expr -> expr
```

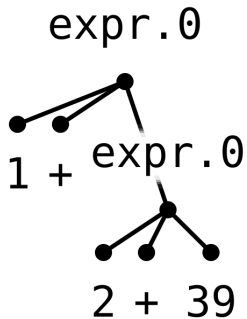
Left recursion

Simply use int for the first term:

```
int    ::= Number
expr   ::= int "+" expr
        | int
```

```
# "1 + 2 + 39" was parsed as:
+  1          Number.0
|  +          expr.0
|  +  2       Number.0
|  |  +       expr.0
|  |  39      Number.0
```

Left recursion



Adding Functions

$(\sqrt{x} \cdot x + 1)$

Adding Functions

```
( \ x . x + 1)
```

```
expr      ::= "(" " \ " var "." expr ")"  
            | ...
```

Adding Variables

```
var ::= Identifier
```

Expression

```
expr ::= "(" "\\" var "." expr ")" arg  
      | int  
      | var  
      | int "+" expr  
      | var "+" expr  
      | "(" "\\" var "." expr ")" arg "+" expr  
      | ...
```

Adding Terms

We'll want to introduce a syntactic form **term** , for variables, ints and functions:

```
term    ::= "(" "\\" var "." expr ")"  
        | "(" expr ")"  
        | int  
        | var
```

Adding Terms

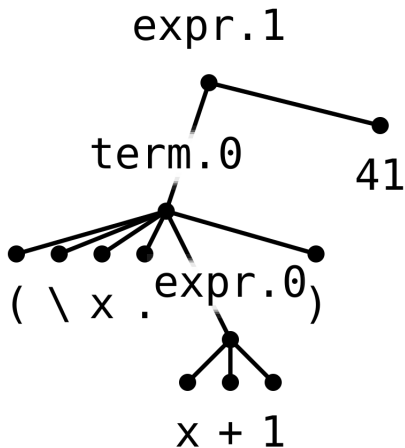
```
expr ::= term "+" expr | term expr | term
```

Parsing stuff

Still done with “./ALGT Demo.language demo.demo expr -l”

```
# "(\\ x . x + 1) 41" was parsed as:
+ + ( term.0
| | \ term.0
| | x Identifier.0
| | . term.0
| | + x Identifier.0
| | | + expr.0
| | | 1 Number.0
| | ) term.0
| 41 Number.0
```

Parsing stuff



III

Building the evaluator

Function or natural deduction?

ALGT supports two ways to perform computations:

Functions and natural deduction

For the evaluator, we'll use natural deduction

Declaring the relation

In a new section in the *.language*

```
Relations
```

```
=====
```

```
( $\rightarrow$ ) : expr (in), expr (out)      Pronounced as "Smallstep"
```

Declaring the relation

In a new section in the *.language*

```
Relations  
=====
```

```
( $\rightarrow$ ) : expr (in), expr (out)      Pronounced as "Smallstep"
```

This relation tells us *e0 becomes e1*

$1 + 1 \rightarrow 2$

$(\backslash x . x + 1) \ 41 \rightarrow 41 + 1$

Defining →

Relations are defined in another section:

```
Rules  
=====
```

Defining EvalPlus

```
----- [EvalPlus]  
... → ...
```

The conclusion goes beneath the line
The rulename goes on the right

Defining EvalPlus

```
----- [EvalPlus]  
i0 "+" i1 → ...
```

Defining EvalPlus

```
----- [EvalPlus]  
i0 "+" i1 → !plus(i0, i1)
```

Builtin functions do have an exclamation mark

Defining EvalPlus

```
i0:int i1:int  
----- [EvalPlus]  
i0 "+" i1 → !plus(i0, i1)
```


Running →

```
i0:int i1:int  
----- [EvalPlus]  
i0 "+" i1 → !plus(i0, i1)
```

./ALGT DemoDyn.language demodyn.demo expr -l -r →

```
# 41 + 1 applied to →  
# Proof weight: 3, proof depth: 2
```

```
41 : int      1 : int  
----- [EvalPlus]  
41 + 1 → 42
```

Function application

```
----- [EvalApp]  
function arg → ...
```

Function application

```
----- [EvalApp]  
("(" "\"\\\" x \".\" expr ")") arg → ...
```

Extra parentheses around function, to group the subterm!

Function application

```
----- [EvalApp]  
("(" " \" x \".\" expr ")") arg → !subs:expr(x, arg, expr)
```

Builtin function **!subs**: replace this, with that, everywhere in
For **!subs** is an explicit type needed

Function application

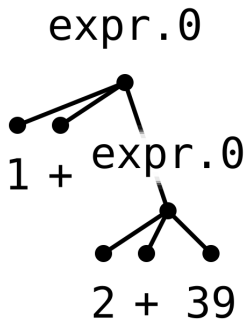
```
# ( \ x . x + 1 ) 41 applied to →  
# Proof weight: 1, proof depth: 1  
  
----- [EvalApp]  
( \ x . x + 1 ) 41 → 41 + 1
```

Contexts

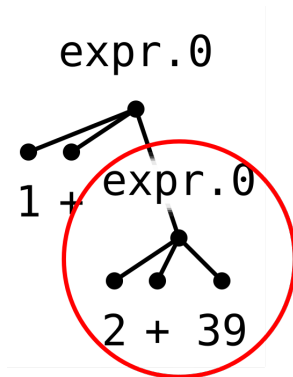
What with $1 + 2 + 39$?

$$\frac{\text{expr0} \rightarrow \text{expr1}}{\text{expr}[\text{expr0}] \rightarrow \text{expr}[\text{expr1}]} \quad [\text{EvalCtx}]$$

Contexts

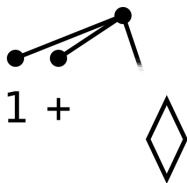


Contexts



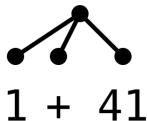
Contexts

expr.0



Contexts

expr.0



Contexts

```
# 1 + 2 + 39 applied to →  
# Proof weight: 4, proof depth: 3
```

```
2 : int      39 : int  
----- [EvalPlus]  
2 + 39 → 41  
----- [EvalCtx]  
1 + 2 + 39 → 1 + 41
```

Bigstep

(\rightarrow^*) : $\text{expr (in)}, \text{expr (out)}$ Pronounced as "Bigstep"

```
i:int
----- [BigStep Base]
i  $\rightarrow^*$  i
```

```
e0  $\rightarrow$  e1                    e1  $\rightarrow^*$  e2
----- [BigStep recursive]
e0  $\rightarrow^*$  e2
```

Bigstep

```
# 1 + 2 + 39 applied to  $\rightarrow^*$ 
# Proof weight: 11, proof depth: 4
```

2 : int	39 : int	1 : int	41 : int	42 : int
-----		-----		-----
2 + 39 \rightarrow 41		1 + 41 \rightarrow 42		42 \rightarrow^* 42
-----		-----		-----
1 + 2 + 39 \rightarrow 1 + 41		1 + 41 \rightarrow^* 42		

1 + 2 + 39 \rightarrow^* 42				

IV

Your turn!

Your turn

Now it's your turn to give these a try.

- ▶ github.com/pietervdn/ALGT
- ▶ Download the **demo**-directory
- ▶ Overview of commands and usefull stuff in **readme.md**

If there is still time, we'll also build a typechecker for the demo language. . .

V

The typechecker

VI

Introducing types

Introducing types

We'll need syntactic forms for types

Introducing types

We'll need syntactic forms for types

```
typeTerm      ::= "Int" | "(" type ")"  
type          ::= typeTerm "->" type | typeTerm
```

Introducing types

Explicit type tags on the input arguments:

```
term    ::= "(" "\\" var ":" type "." expr ")"  
        | "(" expr ")"  
        | int | var  
expr    ::= term "+" expr | term expr | term
```

Introducing types

Which also means we'll have to update **EvalApp**

Introducing types

Which also means we'll have to update **EvalApp**

```
----- [EvalApp]  
("(" " \"\\\" x ":" TArg "." expr ")") arg → !subs:expr(x, arg, expr)
```

VII

Functions

Domain and codomain

Function type	dom	cod
$\text{Int} \rightarrow \text{Int}$	Int	Int
$\text{Int} \rightarrow (\text{Int} \rightarrow \text{Int})$	Int	$\text{Int} \rightarrow \text{Int}$
$(\text{Int} \rightarrow \text{Int}) \rightarrow \text{Int}$	$\text{Int} \rightarrow \text{Int}$	Int

Domain

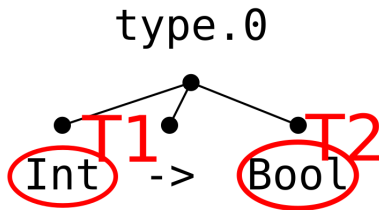
Functions

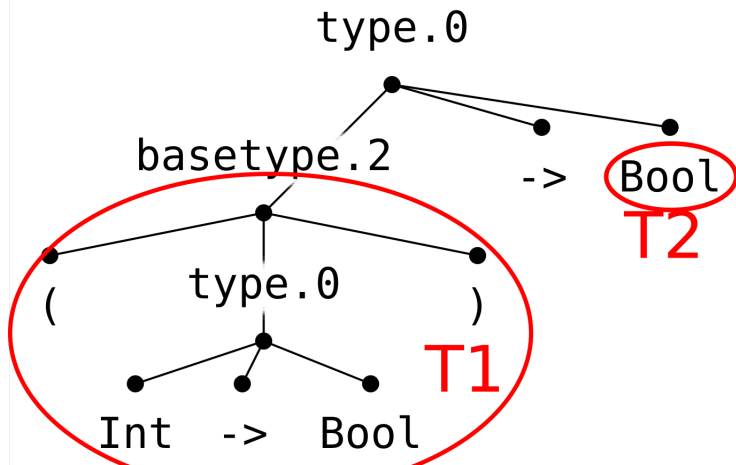
=====

`dom : type -> type`

`dom (T1 "->" T2) = T1`

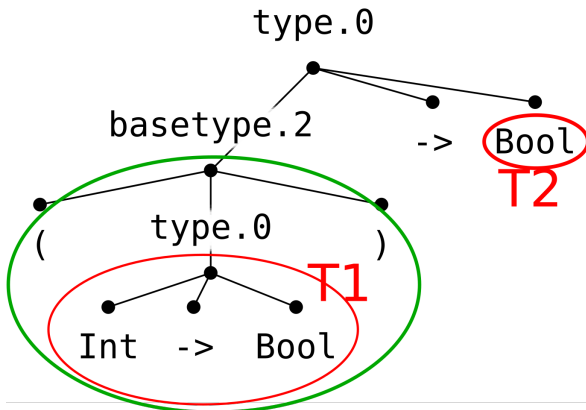
```
dom(T1 "->" T2) = T1
```





Domain

`dom (("(" T1 ")") "->" T2) = T1`



Domain

```
dom("(" T ")") = dom(T)
```

Domain

```
dom          : type -> type
dom("(" T ")") = dom(T)
dom("(" T1 ") "-> T2)
               = T1
dom(T1 "-> T2) = T1
```

Domain

Built-in totality check

```
While checking file Demo.language:  
  Warning:  
    While checking the totality of function "dom":  
      Following calls will fall through:  
        dom("Int")
```

Domain

Built-in totality check

```
While checking file Demo.language:  
  Warning:  
    While checking the totality of function "dom":  
      Following calls will fall through:  
        dom("Int")
```

```
dom("Int")      = !error("Undefined")
```


Codomain

```
cod          : type -> type
cod("(" T ")") = cod(T)
cod(T1 "->" (" T2 "))
              = T2
cod(T1 "->" T2) = T2
cod("Int")     = !error("Undefined")
```

VIII

The typing environment

Variables

```
(\x : Int . x + 1) 41
```

How to keep track of what type a variable (such as **x**) has?

Typing environment

```
typing      ::= var ":" type  
typings    ::= typing typings | "{}"
```

The typing environment will be denoted with Γ (U+393)

IX

Typing

The typing relation

$(\vdash) : \text{typings (in), expr (in), type (out)}$ Pronounced as "Is typed as"

(\vdash) is pronounced *entails*; U+22A2)

Typing constants

```
i:int  
----- [TConstant]  
Γ ⊢ i, "Int"
```

Typing constants

```
i:int  
----- [TConstant]  
⊢ i, "Int"
```

```
# 42 applied to ::  
# Proof weight: 3, proof depth: 3
```

```
42 : int  
----- [TConstant]  
{ } ⊢ 42, Int
```


Typing plus

----- [TPlus]

$\Gamma \vdash i0 \text{ "+" } i1, \text{ "Int"}$

Typing plus

```
Γ ⊢ i0, "Int"           Γ ⊢ i1, "Int"  
----- [TPlus]  
Γ ⊢ i0 "+" i1 , "Int"
```

Typing plus

```

Γ ⊢ i0, "Int"           Γ ⊢ i1, "Int"
----- [TPlus]
Γ ⊢ i0 "+" i1 , "Int"

```

41 + 1 applied to ::

Proof weight: 6, proof depth: 4

```

41 : int           1 : int
----- [TConstant] ----- [TConstant]
{} ⊢ 41, Int       {} ⊢ 1, Int
----- [TPlus]
{} ⊢ 41 + 1, Int

```

Typing variables

We lookup the variable in the **typingEnvironment**:

```
----- [Tx]  
⊢ [x ":" t] ⊢ x, t
```

Typing variables

We lookup the variable in the **typingEnvironment**:

```
----- [Tx]
Γ[x ":" t] ⊢ x, t
```

```
----- [ Tx ]
x : Int {} ⊢ x , Int
```

Typing functions

----- [TLambda]
 $\Gamma \vdash "(" \ " \backslash \backslash " \ x \ ":" \ TArg \ "." \ expr \ ")" \ , \ ???$

What type do we return?

Typing functions

$\Gamma \vdash \text{expr}, \text{TExpr}$

----- [TLambda]

$\Gamma \vdash "(" \text{"\\"} x ":" \text{TArg} "." \text{expr} ")" , ???$

Hmm, something is missing here...

Typing functions

$$(x ":" TArg) \Gamma \vdash expr, TExpr$$
$$\Gamma \vdash "(" \ " \backslash \backslash " \ x \ ":" \ TArg \ "." \ expr \ ")" \ , \ ???$$

----- [TLambda]

Typing environment syntax

typing ::= var ":" type

typings ::= typing typings | "{}"

Nearly done...

Typing functions

$$(x ":" TArg) \Gamma \vdash expr, TExpr$$

----- [TLambda]

$$\Gamma \vdash "(" \ " \backslash \backslash " \ x \ ":" \ TArg \ "." \ expr \ ")" \ , \ TArg \ "->" \ TExpr$$

There is a catch...

Typing functions

```
(x ":" TArg)  $\Gamma \vdash \text{expr}, \text{TExpr}$   
----- [TLambda]  
 $\Gamma \vdash "(" \text{"\\"} x ":" TArg "." \text{expr} ")" , TArg \rightarrow TExpr$ 
```

There is a catch...

$TArg = \text{Int} \rightarrow \text{Int}$

$TExpr = \text{Int}$

Typing functions

```
(x ":" TArg)  $\Gamma \vdash \text{expr}, \text{TExpr}$   
----- [TLambda]  
 $\Gamma \vdash "(" \text{"\\"} x ":" TArg "." \text{expr} ")" , TArg \rightarrow TExpr$ 
```

There is a catch...

$TArg = \text{Int} \rightarrow \text{Int}$

$TExpr = \text{Int}$

$TArg \rightarrow TExpr = \text{Int} \rightarrow \text{Int} \rightarrow \text{Int}$

Typing functions

$(x ":" TArg) \Gamma \vdash expr, TExpr$

----- [TLambda]

$\Gamma \vdash "(" \ " \backslash \backslash " \ x \ ":" \ TArg \ "." \ expr \ ")" \ , \ ("(" \ TArg \ ")") \ "->" \ TExpr$

Typing functions

```
(x ":" TArg)  $\Gamma \vdash \text{expr}, \text{TExpr}$ 
```

[TLambda]

```
 $\Gamma \vdash "(" \text{"\\"} x ":" TArg "." \text{expr} ")" , ( "(" TArg ")" ) \rightarrow \text{TExpr}$ 
```

```
# (\x : Int . x + 1) applied to ::
```

```
# Proof weight: 6, proof depth: 5
```

```

                                     1 : int
----- [Tx]           ----- [TConstant]
x : Int {}  $\vdash x, \text{Int}$       x : Int {}  $\vdash 1, \text{Int}$ 
----- [TPlus]
x : Int {}  $\vdash x + 1, \text{Int}$ 
----- [TLambda]
{}  $\vdash ( \backslash x : \text{Int} . x + 1 ), ( \text{Int} ) \rightarrow \text{Int}$ 
```

Typing application

 $\Gamma \vdash e1 \ e2, ???$

[Tapp]

Typing application

$$\Gamma \vdash e1, \text{Tfunc}$$
$$\Gamma \vdash e1 \ e2, ???$$

[Tapp]

Typing application

$$\Gamma \vdash e1, T_{\text{func}} \quad \Gamma \vdash e2, T_{\text{arg}}$$

[Tapp]

$$\Gamma \vdash e1 \ e2, ???$$

Typing application

$$\Gamma \vdash e1, T_{\text{func}} \quad \Gamma \vdash e2, T_{\text{arg}}$$
$$\Gamma \vdash e1 \ e2, \text{cod}(T_{\text{func}})$$

[Tapp]

Typing application

$$\frac{\Gamma \vdash e1, \text{Tfunc} \quad \Gamma \vdash e2, \text{Targ} \quad \text{Targ} = \text{dom}(\text{Tfunc}) : \text{type}}{\Gamma \vdash e1 \ e2, \text{cod}(\text{Tfunc})} \quad [\text{Tapp}]$$

Typing application

$$\frac{\Gamma \vdash e1, \text{Tfunc} \quad \Gamma \vdash e2, \text{Targ} \quad \text{Targ} = \text{dom}(\text{Tfunc}) : \text{type}}{\Gamma \vdash e1 \ e2, \text{cod}(\text{Tfunc})} \quad [\text{Tapp}]$$

```

                                1 : int
-----
x : Int {} ⊢ x, Int x : Int {} ⊢ 1, Int
-----
x : Int {} ⊢ x + 1, Int                                41 : int
-----
{} ⊢ ( \ x : Int . x + 1 ), ( Int ) -> Int {} ⊢ 41, Int Targ = Int = dom(Tfunc)
-----
{} ⊢ ( \ x : Int . x + 1 ) 41, Int

```

Typing: practically

Define relation $::$ to type in an empty environment

```
(::) : expr (in), type (out)
```

```
"{}" ⊢ e, T  
----- [Typing in empty]  
e :: T
```

`./ALGT Demo.language demo.demo expr -l -r ::`

One more thing

X

Syntax coloring

Syntax coloring

Extra section, just under the syntax

Syntax Style

=====

Number -> "constant"

type -> "type"

typeTerm -> "type"

Identifier -> "identifier"

Syntax coloring

–style **Terminal**

```
{( \ x : Int . x + 1 ) 41 }
```

–style **White**

```
{( \ x : Int . x + 1 ) 41 }
```