

key	BFS	DFS
Definition	Stands for Breadth First Search	Stands for depth first search.
Data Struct	It uses Queue to find the Shortest path	It uses stack to find the shortest path.
Source	It is better when target is closer to source.	It is better when target is far from source.
Suitable for decision tree	It considers all neighbours. So it is not suitable for decision tree used in puzzle games.	It is more suitable as with one decision, we need to traverse further to augment the decision.
Speed	It is slower than DFS	It is faster than BFS
Time Comp with	$O(V+E)$ where V is vertices & E is edges	$O(V+E)$ where V is vertices & E is edges

Q2 Stack is used to implement DFS, because in it we first traverse the whole branch of the tree & later on visit the adjacent branch, since this is similar to LIFO, therefore stack is used.

Queue is used to implement BFS, it is because queue is used as a FIFO. Instead because BFS is to test the immediate children first & after all immediate children are tested, to then return to those children & check their children & so forth.

Q3 Sparse graph - graph where no. of edges is much less than the possible number of edges.

Dense graph - where number of edges is much ~~more~~ close to maximal number of edges.

If graph is dense it should be represented by adjacency matrix.

If graph is sparse it should be represented by adjacency list.

Q4 BFS

In undirected graph, do a BFS traversal on given graph, for each visited vertex v , if there is an adjacent ' u ' such that ' v ' is already visited & ' u ' is not parent of ' v ', then there is cycle in a graph.

DFS

Run DFS from a node and mark this node as visited, now for any other vertex if its neighbour is already visited & that neighbour is not the parent of that current node then there exists a cycle in the graph.

Q5 Disjoint set data structure

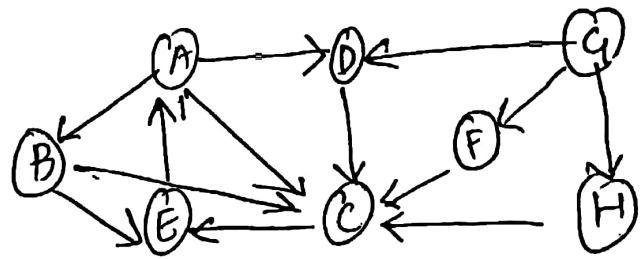
The disjoint set can be defined as the subsets where there is no common element b/w two sets.

Operations are

- i) Union
- ii) make new set
- iii) find

Q6 BFS

$A \rightarrow B \rightarrow C \rightarrow D \rightarrow E$,
 $G \rightarrow \cancel{D} \rightarrow H \rightarrow F \rightarrow \cancel{E}$



DFS

$A \rightarrow D \rightarrow C \rightarrow B$, $G \rightarrow F \rightarrow H$

Q7 connected components = 4
vertices = 10

Q8 Topological sort $\rightarrow 0-1-2-3-4-5$
DFS $\rightarrow 5 \rightarrow 2 \rightarrow 3 \rightarrow 1 \rightarrow 0$
4 can't be reached

Q9) Yes, heap data structure can be used to create priority queue.

- Dijkstra's to find shortest path
- Prim's Algo
- Huffman Algo

Q10) min heap \rightarrow root element is the smallest
max heap \rightarrow root element is the largest