

# KRP – Assignment 2

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★ This assignment, due on 21st April at 23:59, contributes to 10% of the final marks for this course. Please be advised that only Questions 1 – 8 are mandatory. Nevertheless, students can earn up to one bonus mark by completing Question 9. This bonus mark can potentially augment a student's overall marks but is subject to a maximum total of 100 for the course. By providing bonus marks, we aim to incentivize students to excel in their studies and reward those with a remarkable grasp of the course materials.

## Question 1. Bisimulation Invariance

In the lecture, we defined bisimulation for  $\mathcal{ALC}$  and showed bisimulation invariance of  $\mathcal{ALC}$  (Theorem 3.2).

- Define a notion of “ $\mathcal{ALCN}$ -bisimulation” that is appropriate for  $\mathcal{ALCN}$  in the sense that bisimilar elements satisfy the same  $\mathcal{ALCN}$ -concepts.
- Use the definition to show that  $\mathcal{ALCQ}$  is more expressive than  $\mathcal{ALCN}$ .

## Question 2. Bisimulation over Filtration

Let  $C$  be an  $\mathcal{ALC}$ -concept,  $\mathcal{T}$  an  $\mathcal{ALC}$ -TBox,  $\mathcal{I}$  an interpretation and  $\mathcal{J}$  its filtration w.r.t.  $\text{sub}(C) \cup \text{sub}(\mathcal{T})$  (see Definition 3.14 for the definition of filtration). Show the truth or falsity of the following statement:

- the relation  $\rho = \{(d, [d]) \mid d \in \Delta^{\mathcal{I}}\}$  is a bisimulation between  $\mathcal{I}$  and  $\mathcal{J}$ .

Hint: If the above relation  $\rho$  were a bisimulation, why do we have to explicitly prove Lemma 3.15 in the lecture? Wouldn't Lemma 3.15 then be a consequence of Theorem 3.2?

## Question 3. Bisimulation within the Same Interpretation (2 marks)

We define “bisimulations on  $\mathcal{I}$ ” as bisimulations between an interpretation  $\mathcal{I}$  and itself. Let  $d, e \in \Delta^{\mathcal{I}}$  be two elements. We write  $d \approx_{\mathcal{I}} e$  if they are bisimilar, i.e., if there is a bisimulation  $\rho$  on  $\mathcal{I}$  such that  $d \rho e$ .

- Show that  $\approx_{\mathcal{I}}$  is a bisimulation on  $\mathcal{I}$ .

Consider the interpretation  $\mathcal{J}$  defined like the filtration, but with  $\approx_{\mathcal{I}}$  instead of  $\simeq$ .

- Show that  $\rho = \{(d, [d]_{\approx_{\mathcal{I}}}) \mid d \in \Delta^{\mathcal{I}}\}$  is a bisimulation between  $\mathcal{I}$  and  $\mathcal{J}$ .
- Show that, if  $\mathcal{I}$  is a model of an  $\mathcal{ALC}$ -concept  $C$  w.r.t. an  $\mathcal{ALC}$ -TBox  $\mathcal{T}$ , then so is  $\mathcal{J}$ .
- Why can't we use the previous result to show the finite model property for  $\mathcal{ALC}$ ?

## Question 4. Closure under Disjoint Union

Recall Theorem 3.8 from the lecture, which says that the disjoint union of a family of models of an  $\mathcal{ALC}$ -TBox  $\mathcal{T}$  is again a model of  $\mathcal{T}$ . Note that the disjoint union is only defined for concept and role names.

- Extend the notion of disjoint union to individual names such that the following holds: for any family  $(\mathcal{I}_\nu)_{\nu \in \Omega}$  of models of an  $\mathcal{ALC}$ -knowledge base  $\mathcal{K}$ , the disjoint union  $\biguplus_{\nu \in \Omega} \mathcal{I}_\nu$  is also a model of  $\mathcal{K}$ .

**Question 5. Closure under Disjoint Union**

Let  $\mathcal{K} = \{\mathcal{T}, \mathcal{A}\}$  be a consistent  $\mathcal{ALC}$ -KB. We write  $C \sqsubseteq_{\mathcal{K}} D$  if  $C^{\mathcal{I}} \subseteq D^{\mathcal{I}}$  holds for every model  $\mathcal{I}$  of  $\mathcal{K}$ .

- Prove that for all  $\mathcal{ALC}$ -concepts  $C$  and  $D$  we have  $C \sqsubseteq_{\mathcal{K}} D$  iff  $C \sqsubseteq_{\mathcal{T}} D$ .

Hint: Use the modified definition of disjoint union from the previous question.

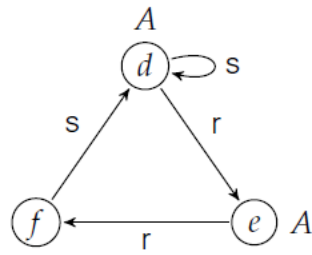
**Question 6. Finite Model Property (2 marks)**

Let  $C$  be an  $\mathcal{ALC}$ -concept that is satisfiable w.r.t. an  $\mathcal{ALC}$ -TBox  $\mathcal{T}$ . Show truth or falsity of the following statement:

- for all  $m \geq 1$  there is a finite model  $\mathcal{I}_m$  of  $\mathcal{T}$  such that  $|C^{\mathcal{I}_m}| \geq m$ .
- Does it hold if the condition “ $|C^{\mathcal{I}_m}| \geq m$ ” is replaced by “ $|C^{\mathcal{I}_m}| = m$ ”?

**Question 7. Unravelling**

Draw the unraveling of the following interpretation  $\mathcal{I}$  at  $d$  up to depth 5, i.e., restricted to  $d$ -paths of length at most 5 (see Definition 3.21):



**Question 8. Tree Model Property**

- Show the truth or falsity of the following statement: if  $\mathcal{K}$  is an  $\mathcal{ALC}$ -KB and  $C$  an  $\mathcal{ALC}$ -concept such that  $C$  is satisfiable w.r.t.  $\mathcal{K}$ , then  $C$  has a tree model w.r.t.  $\mathcal{K}$ .

**Question 9 (with 1 bonus mark). Bisimulation Invariance**

Interpretations of  $\mathcal{ALC}$  can be represented as graphs, with edges labeled by roles and nodes labeled by sets of concept names. More precisely, in such a graph:

each node corresponds to an element in the domain of the interpretation and it is labeled with all the concept names to which this element belongs in the interpretation;

an edge with label  $r$  between two nodes says that the corresponding two elements of the interpretation are related by the role  $r$ .

- Show that: the description logic  $\mathcal{S}$  (i.e.,  $\mathcal{ALC}$  with transitive roles) is more expressive than  $\mathcal{ALC}$ .