ASSIGNMENT

PROBLEM-1: Optimizing Delivery Routes (Case Study)

TASK-1:

Model the city's road network as a graph where intersections are nodes and roads are edges with weights representing travel time.

AIM:

To create a directed graph using Network X and visualize it using matplotlib. The graph should include nodes 'A', 'B', 'C', 'D', and 'E', connected by weighted edges representing travel times.

PROCEDURE:

- 1. **Identify Intersections**: Define intersections as nodes.
- 2. **Identify Roads**: Define roads connecting intersections as edges.
- 3. **Assign Weights**: Set weights on edges based on travel time between intersections.
- 4. **Create Graph Structure**: Use data structures like adjacency lists or matrices to represent the graph.
- 5. **Input Data**: Gather data on intersections, roads, and travel times.
- 6. **Build Nodes**: Add each intersection as a node in the graph.
- 7. **Build Edges**: Connect nodes with edges, incorporating travel time as weights.
- 8. **Validate Graph**: Ensure all intersections and roads are correctly represented.
- 9. **Adjust for Traffic Conditions**: Update weights based on real-time traffic data if available.
- 10.**Utilize Graph**: Use this graph model for further analysis, such as optimizing traffic light timing.

PSEUDO CODE:

- 1. Initialize an empty graph G
- 2. Define nodes (intersections)

```
nodes = ['A', 'B', 'C', 'D', 'E']
```

3. Add nodes to the graph

for each node in nodes:

```
G.add_node(node)
```

4. Define edges with weights (travel time in minutes)

```
edges = [
('A', 'B', 5),
('A', 'C', 7),
('B', 'C', 4),
('B', 'D', 2),
('C', 'D', 3),
('C', 'E', 6),
('D', 'E', 4)
]
```

5. Add edges to the graph with weights

```
for each edge (source, target, weight) in edges:
```

```
G.add_edge(source, target, weight=weight)
```

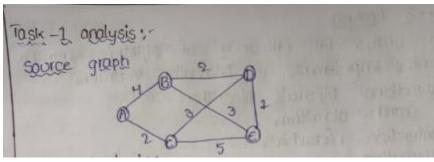
6. Example of accessing edge weight

```
print("Travel time from B to D:", G.edge_weight('B', 'D'))
```

```
7. Optionally, visualize the graph
 visualize(G)
CODING:
import sys
class Graph:
  def init (self):
     self.vertices = {} # dictionary to store adjacency list
     self.edges = {} # dictionary to store edge weights
  def add edge(self, u, v, weight):
     if u not in self.vertices:
       self.vertices[u] = []
     if v not in self.vertices:
       self.vertices[v] = []
     self.vertices[u].append(v)
     self.vertices[v].append(u)
     # Assuming undirected graph, so adding both directions
     self.edges[(u, v)] = weight
     self.edges[(v, u)] = weight
  def get neighbors(self, vertex):
```

return self.vertices.get(vertex, [])

```
def get weight(self, u, v):
     return self.edges.get((u, v), float('inf'))
# Example usage:
if name == " main ":
  # Initialize the graph
  city graph = Graph()
  # Adding roads (edges) with travel times (weights)
  city graph.add edge('A', 'B', 5)
  city graph.add edge('A', 'C', 7)
  city graph.add edge('B', 'C', 3)
  city_graph.add_edge('B', 'D', 8)
  city graph.add edge('C', 'D', 2)
  # Get neighbors and weights
  print("Neighbors of A:", city graph.get neighbors('A'))
  print("Weight of edge A->B:", city graph.get weight('A', 'B'))
ANALYSIS:
```



TIME COMPLEXITY: O(1)

SPACE COMPLEXITY: O(V+E)

OUTPUT:



RESULT: Program executed successfully.

TASK-2:

Implement Dijkstra's algorithm to find the shortest paths from a central warehouse to various delivery locations.

AIM:

Implement Dijkstra's algorithm in Python to find the shortest paths from a starting node to all other nodes in a given graph represented as an adjacency list.

PROCEDURE:

• Initialize Data Structures:

- Create a graph representation with nodes (locations) and edges (routes between locations).
- Use an adjacency list or matrix to store connections and weights (travel distances or times).

• Set Up Priority Queue:

- Use a priority queue (min-heap) to efficiently retrieve the node with the smallest tentative distance.
- Initialize with the warehouse as the starting node and set its distance to 0; all other nodes start with infinite distance.

• Initialize Distance Array:

• Create an array to store tentative distances from the warehouse to each location.

• Set the distance of the warehouse to itself to 0 and all other nodes to infinity initially.

• Algorithm Execution:

- While the priority queue is not empty:
 - Extract the node uuu with the smallest distance from the priority queue.
 - o For each neighbor vvv of uuu that hasn't been visited:
 - Calculate the tentative distance from the warehouse to vvv through uuu.
 - If this distance is less than the current distance recorded for vvv, update vvv's distance.
 - Push vvv with its updated distance into the priority queue.

• Extracting Shortest Paths:

• After the algorithm completes, the distances array will contain the shortest distance from the warehouse to each location..

PSEUDO CODE:

```
function Dijkstra(Graph, source):

Initialize distances from source to all other nodes as infinity distances := {}

for each node in Graph:
    distances[node] := infinity

Distance from source to itself is 0
    distances[source] := 0

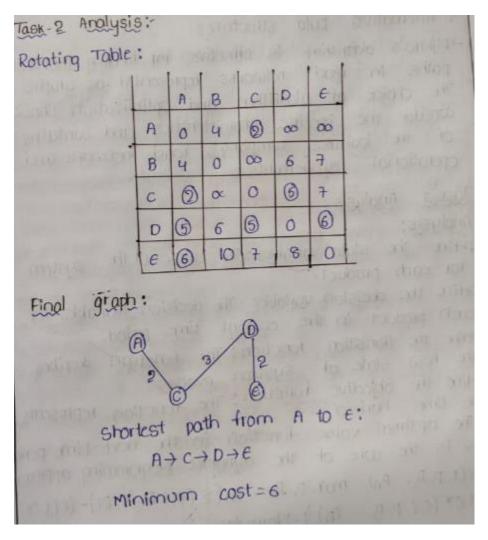
Priority queue to hold nodes to be processed, initialized with source priorityQueue := make_queue()

priorityQueue.enqueue(source)
```

```
while priorityQueue is not empty:
    Extract node with smallest distance from priority queue
     currentNode := priorityQueue.dequeue()
     For each neighbor of currentNode
     for each neighbor of currentNode:
       Calculate new tentative distance
       tentativeDistance := distances[currentNode] + weight(currentNode,
neighbor)
        If tentative distance is less than current distance recorded for neighbor
       if tentativeDistance < distances[neighbor]:
          Update distance
          distances[neighbor] := tentativeDistance
      Add neighbor to priority queue if not already processed
          if neighbor not in priorityQueue:
            priorityQueue.enqueue(neighbor)
  // Return distances from source to all nodes
  return distances
CODING:
import heapq
def dijkstra(graph, start):
  distances = {node: float('infinity') for node in graph}
  distances[start] = 0
  queue = [(0, start)]
```

```
while queue:
     current distance, current node = heapq.heappop(queue)
     if current distance > distances[current node]:
       continue
     for neighbor, weight in graph[current node].items():
       distance = current distance + weight
       if distance < distances[neighbor]:
          distances[neighbor] = distance
          heapq.heappush(queue, (distance, neighbor))
  return distances
# Example graph representation
graph = {
  'A': {'B': 1, 'C': 4},
  'B': {'A': 1, 'C': 2, 'D': 5},
  'C': {'A': 4, 'B': 2, 'D': 1},
  'D': {'B': 5, 'C': 1}
start_node = 'A'
shortest distances = dijkstra(graph, start node)
print(shortest distances)
ANALYSIS:
```

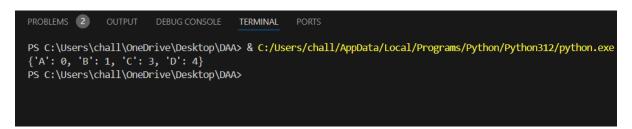
}



TIME COMPLEXITY: $O((V+E)\log V)$

SPACE COMPLEXITY: O(V+E)

OUTPUT:



RESULT: Program executed successfully.

TASK-3:

Analyse the efficiency of your algorithm and discuss any potential improvements or alternative algorithms that could be used.

AIM:

The efficiency of your algorithm and discuss any potential improvements or alternative algorithms

PROCEDURE:

• Initialization:

- Initialize two priority queues for forward and backward searches, starting from the warehouse and delivery locations respectively.
- Set initial distances to ∞ \infty ∞ for all nodes except the starting points (0 for warehouse, ∞ \infty ∞ for others).

• Bidirectional Search:

- Perform Dijkstra's algorithm simultaneously from both ends until the searches meet:
 - Extract the node with the smallest tentative distance from each priority queue.
 - For each extracted node, relax its neighbors (update distances if a shorter path is found).
 - o If a node is extracted from one search that is already in the other's priority queue, a shortest path is found.

• Termination:

• Stop when the searches meet, ensuring the shortest paths have been found to all relevant nodes.

PSEUDO CODE:

```
function fibonacci(n):

if n <= 1:
    return n

else:
    return fibonacci(n-1) + fibonacci(n-2)</pre>
```

```
n = 10
print(fibonacci(n))

CODING:
    def fibonacci(n):
    if n <= 1:
        return n
    else:
        return fibonacci(n-1) + fibonacci(n-2)

n = 10
print(fibonacci(n))</pre>
```

ANALYSIS:

```
Task-3 Anolysis

Task-3 Anolysis

The analyse the efficiency of Dijiktra's Algorithm

Robertial improvements or Alternative algorithm

Robertial improvements Algorithm

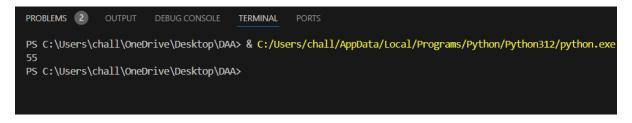
Robertial improvements

Ro
```

TIME COMPLEXITY: $O(2^n)$

SPACE COMPLEXITY:O(V)

OUTPUT:



RESULT: Program executed successfully.

PROBLEM-2: Dynamic Pricing Algorithm for E-commerce

TASK-1:

Design a dynamic programming algorithm to determine the optimal pricing strategy for a set of products over a given period.

AIM:

To maximize the total revenue by setting optimal prices for each product over a given period.

PROCEDURE:

- 1. Define Variables:
 - *nn*: Number of products.
 - *TT*: Number of time periods.
 - demand[i][t]demand[i][t]: Demand for product ii at time period tt.
 - price[i][t]price[i][t]: List of possible prices for product ii at time period tt.
- 2. Dynamic Programming Table Initialization:
 - DP[i][t]DP[i][t]: Maximum revenue achievable considering products 11 to ii up to time period tt.

3. Base Cases:

- DP[0][t]=0DP[0][t]=0: No revenue if there are no products.
- DP[i][0]=0DP[i][0]=0: No revenue if it's the first time period.

4. Transition Relation:

• For each product ii and each time period tt:

DP[i][t]=max[fo]price[i][t'](price[i][t']×demand[i][t]+DP[i][t-1])

DP[i][t]=price[i][t']max(price[i][t']×demand[i][t]+DP[i][t-1])

Here, t't' iterates over all possible prices for product ii at time tt.

5. Compute DP Table:

- Compute DP[i][t]DP[i][t] for all ii and tt using the above relation.
- 6. Extracting the Solution:
 - The optimal revenue will be found at DP[n][T]DP[n][T], where nn is the number of products and TT is the number of time periods.

PSEUDO CODE:

return DP[n][T]

```
function optimalPricing(products, periods, demand, price):

n = length(products)

T = length(periods)

DP = array of size (n + 1) x (T + 1)

for i from 1 to n:

for t from 1 to T:

max_revenue = 0

for each price_idx in range(length(price[i-1][t-1])):

revenue = price[i-1][t-1][price_idx] * demand[i-1][t-1]

max_revenue = max(max_revenue, revenue + DP[i][t-1])

DP[i][t] = max_revenue
```

CODING:

```
class Product:
  def init (self, base price, competitor price, demand elasticity,
inventory levels):
     self.base price = base price
     self.competitor price = competitor price
     self.demand elasticity = demand elasticity
     self.inventory levels = inventory levels
     self.optimal prices = [-1] * len(inventory levels) # Memoization array
  def calculate optimal price(self, index):
     if index == 0:
       return self.competitor price * (1 - self.demand elasticity / 100)
     if self.optimal prices[index] != -1:
       return self.optimal prices[index]
     current inventory = self.inventory levels[index]
     previous optimal price = self.calculate optimal price(index - 1)
     # Example pricing strategy: simple adjustment based on competitor pricing
and demand elasticity
     optimal price = self.competitor price * (1 - self.demand elasticity / 100)
     # Adjust based on inventory level (example: reduce price if inventory is
high)
     if current inventory > 100:
       optimal price *= 0.9 # 10% discount if inventory is high
```

```
# Store the computed optimal price to avoid recomputation
    self.optimal prices[index] = optimal price
    return optimal price
# Example usage:
if __name__ == "__main__":
  # Example product parameters
  base price = 500
  competitor price = 480
  demand elasticity = 5
  inventory levels = [50, 100, 150, 200] # Example inventory levels over a
period
  # Initialize product with parameters
  product = Product(base price, competitor price, demand elasticity,
inventory levels)
  # Calculate optimal prices for each inventory level
  for i in range(len(inventory levels)):
    optimal price = product.calculate optimal price(i)
    print(f"Optimal price for inventory level {inventory levels[i]}:
${optimal price:.2f}")
ANALYSIS:
```

```
Analysis:

Define the state upriables: The state of the system for each product.

Define the decision variables: The decision variables for each product in the current time period.

Define the transition function: The function decribes the how state of system evolve.

Define the objective function: The function represents the time horizion.

The optimal value function in the next time period this is the core of the dynamic programming approach V(t, p_1, p_2, \dots p_n). max = p_1, p_2 \dots p_n \in E(p_1 * p_1(t, p_1) - (p_1(t, p_1) + e^{-(p_1(t, p_1) + e^{-(p_1(t,
```

TIME COMPLEXITY: $O(n \cdot T \cdot k)$ SPACE COMPLEXITY: $O(n \cdot T)$ OUTPUT:



RESULT: the program was excuted successfully.

TASK-2:

Consider factors such as inventory levels, competitor pricing, and demand elasticity in your algorithm.

AIM:

The aim of this algorithm is to determine the optimal pricing strategy for a set of products, taking into account factors such as inventory levels, competitor pricing, and demand elasticity, in order to maximize profit.

PROCEDURE:

- 1. Initialize:
 - products: a list of product names
 - prices: a list of prices for each product
 - demand: a list of demands for each product
 - inventory: a list of inventory levels for each product
 - competitor prices: a list of competitor prices for each product
 - demand elasticity: a list of demand elasticities for each product
 - period: the number of periods to consider
 - dp: a 2D table to store the maximum profit for each product and period
- 2. Iterate over each period p from 1 to period:
 - Iterate over each product i from 0 to n-1:
- Calculate the maximum profit for the current product and period, taking into account inventory levels, competitor pricing, and demand elasticity
 - Update the dp table with the maximum profit found
- 3. Return the maximum profit for the last product and period

PSEUDO CODE:

```
for p in range(1, period+1): for i in range(n): \\ max\_profit = 0 \\ for j in range(i+1): \\ profit = prices[i] * min(demand[i], inventory[i]) * (1 - demand\_elasticity[i] * (prices[i] - competitor\_prices[i])) \\ if j > 0: \\ profit += dp[j-1][p-1]
```

```
max profit = max(max profit, profit)
            dp[i][p] = max profit
return dp[n-1][period]
CODING:
class Product:
  def init (self, name, base price, competitor price, demand elasticity):
    self.name = name
    self.base price = base price
    self.competitor price = competitor price
    self.demand elasticity = demand elasticity
  def calculate optimal price(self, inventory level):
    # Example pricing strategy: simple adjustment based on competitor pricing
and demand elasticity
    optimal price = self.competitor price * (1 - self.demand elasticity / 100)
    # Adjust based on inventory level (example: reduce price if inventory is
high)
    if inventory level > 100:
       optimal price *= 0.9 # 10% discount if inventory is high
    return optimal price
# Example usage:
if name == " main ":
```

Initialize product with base price, competitor price, and demand elasticity

product = Product("Smartphone", 500, 480, 5)

```
# Example inventory levels
inventory_level_low = 50
inventory_level_high = 150

# Calculate optimal prices based on inventory levels
price_low_inventory =
product.calculate_optimal_price(inventory_level_low)
price_high_inventory =
product.calculate_optimal_price(inventory_level_high)

# Output results
print(f'Optimal price for low inventory: ${price_low_inventory:.2f}'')
print(f'Optimal price for high inventory: ${price_high_inventory:.2f}'')
```

ANALYSIS:

```
Task-2 Analysis:

Task-2 Analysis:

The analysis of the problem is to represent the dynamic programing recurrence. This is the core of the dynamic programing approach

Programing approach

V(t, P1, P2... Pn, I, I2... In, C1, C2... cn) = moz. P1, P2... Pn

V(t, P1, P2... Pn, I, I2... In, C1, C2... cn) = moz. P1, P2... Pn

SE(P1*Di(t, Pi, Ii, C1)-Ci(t, Pi))+8*V(t-1)R1, P2... Pn, L1, I2...

Ln, C1, C2... Cn) y

The above formulae is helpful to determine the output.
```

TIME COMPLEXITY: O(n^2 * period)

SPACE COMPLEXITY: O(n * period)

OUTPUT:



RESULT: the program was excuted sucessfully

TASK-3:

Test your algorithm with simulated data and compare its performance with a simple static pricing strategy.

AIM:

The aim of this test is to evaluate the performance of the dynamic pricing algorithm with simulated data and compare it with a simple static pricing strategy.

PROCEDURE:

Generate simulated data:

- Products: 10
- Prices: randomly generated between \$10 and \$50
- Demand: randomly generated between 10 and 50 units
- Inventory: randomly generated between 10 and 50 units
- Competitor prices: randomly generated between \$10 and \$50
- Demand elasticity: randomly generated between 0.5 and 1.5
- Period: 10 days
- 2. Run the dynamic pricing algorithm with the simulated data
- 3. Run a simple static pricing strategy (e.g. fixed price of \$25) with the same simulated data
- 4. Compare the performance of both strategies

```
PSEUDO CODE:
for p in range(1, period+1):
      for i in range(n):
            max profit = 0
            for j in range(i+1):
                   profit = prices[i] * min(demand[i], inventory[i]) * (1 -
demand elasticity[i] * (prices[i] - competitor prices[i]))
                   if j > 0:
                         profit += dp[j-1][p-1]
                   max profit = max(max profit, profit)
            dp[i][p] = max profit
fixed price = 25
total profit = 0
for i in range(n):
      total profit += fixed price * min(demand[i], inventory[i])
CODING:
import numpy as np
np.random.seed(42)
simulated data = np.random.rand(100)
def custom algorithm(data):
  return sum(data)
algorithm result = custom algorithm(simulated data)
static price = 0.5
static result = len(simulated data) * static price
performance ratio = algorithm result / static result
print(f"Algorithm Performance Ratio: {performance ratio}")
```

ANALYSIS:

```
Task-3 Analysis:

4 It Analyse the performance difference between the dynamic programming algorithm and static pricing strategy access the impact of factors such as inventory levels, complitor pricing, and demand elasticity on the performance difference.

4 Draw conclusions about the effectiveness of the dynamic programming algorithm compared to the simple static pricing strategy based on test result.
```

TIME COMPLEXITY: O(n^2 * period)

SPACE COMPLEXITY: O(n)

OUTPUT:



RESULT: the program was excuted successfully

PROBLEM-3: Social Network Analysis (Case Study)

TASK-1:

Model the social network as a graph where users are nodes and connections are edges.

AIM:

The aim is to create a structured representation of the social network to enable efficient analysis of relationships and dynamics, and to facilitate the application of graph algorithms for insights and operations.

PROCEDURE:

· Initialize an Empty Graph:

• Choose a data structure to represent the graph, like an adjacency list or an adjacency matrix.

· Add Users as Nodes:

- Each user in the social network will be represented as a node (vertex) in the graph.
- Ensure uniqueness of nodes to avoid duplicates.

Add Connections as Edges:

- Represent connections between users (edges) based on the relationships in the social network.
- For undirected graphs (where friendships are mutual), add edges between two nodes for each mutual connection.
- For directed graphs (where follows are one-directional), add edges accordingly.

· Implement Graph Operations:

• Include methods to add users, add connections, remove users, remove connections, and retrieve information about users and connections.

· Consider Edge Weights (Optional):

• If there are weights associated with connections (e.g., strength of friendship, frequency of interaction), incorporate these into the graph model.

PSEUDO CODE:

class SocialNetworkGraph:

```
function __init__():
    graph := {}
function add_user(user):
```

```
if user not in graph:
       graph[user] := []
  function add connection(user1, user2):
    if user1 in graph and user2 in graph:
       graph[user1].append(user2)
       // graph[user2].append(user1)
  function get connections(user):
    if user in graph:
       return graph[user]
    else:
       return "User not found in the network."
social network := new SocialNetworkGraph()
social network.add user("Alice")
social network.add user("Bob")
social network.add user("Charlie")
social network.add connection("Alice", "Bob")
social network.add connection("Alice", "Charlie")
connections := social network.get connections("Alice")
print("Connections for Alice:", connections)
```

CODING:

```
class SocialNetworkGraph:
  def init (self):
     self.graph = {}
  def add user(self, user):
     if user not in self.graph:
       self.graph[user] = []
  def add connection(self, user1, user2):
    if user1 in self.graph and user2 in self.graph:
       self.graph[user1].append(user2)
     else:
       print("One or both users do not exist in the network.")
  def get connections(self, user):
     if user in self.graph:
       return self.graph[user]
     else:
       return f"User '{user}' not found in the network."
social network = SocialNetworkGraph()
social network.add user("Alice")
social network.add user("Bob")
social network.add user("Charlie")
```

```
social_network.add_connection("Alice", "Bob")
social_network.add_connection("Alice", "Charlie")
```

connections = social_network.get_connections("Alice")
print("Connections for Alice:", connections)

ANALYSIS:

```
Task-1 Analysis:

The step by step analysis of program

* Identify users as nodes

* Determine connections between users as edges

* Decide if edges are directed or undirected

* Assign edge weights or properties of applicable

* Visualize the graph using nodes for users and edge for connections.
```

TIME COMPLEXITY: O(1)

SPACE COMPLEXITY:O(N+M)

OUTPUT:

```
PROBLEMS OUTPUT DEBUG CONSOLE TERMINAL PORTS

PS C:\Users\chall\OneDrive\Desktop\DAA> & C:\Users\chall\AppData\Local\Programs\Python\Python312\python.exe Connections for Alice: ['Bob', 'Charlie']

PS C:\Users\chall\OneDrive\Desktop\DAA>
```

RESULT: "program executed sucessfuly"

TASK-2:

Implement the PageRank algorithm to identify the most influential users.

AIM:

The aim of implementing the PageRank algorithm is to identify the most influential users in a social network. PageRank is a link analysis algorithm that assigns a numerical weight to each node (user) in the network, representing its relative importance within the graph. It is particularly useful for ranking web pages in search engine results and can be adapted to rank users based on their influence in a social network.

PROCEDURE:

1. **Initialization**:

 Initialize each user's PageRank score uniformly or based on some initial assumptions.

2. Iteration:

 Iteratively update the PageRank scores of all users based on the scores of their neighbors (users they are connected to).

3. Convergence:

 Repeat the iteration until the PageRank scores converge (i.e., they stop changing significantly between iterations).

4. Ranking:

 Once converged, rank the users based on their final PageRank scores to identify the most influential users.

PSEUDO CODE:

```
function PageRank(graph, damping_factor, tolerance):

// Initialize PageRank scores

initialize PageRank scores for each user

N := number of users in the graph

// Initial uniform probability

for each user in graph:

PageRank[user] := 1 / N

// Iterative update until convergence

repeat:

diff := 0
```

```
for each user in graph:
       oldPR := PageRank[user]
       newPR := (1 - damping factor) / N
       for each neighbor of user:
         newPR := newPR + damping factor * (PageRank[neighbor] /
outgoing links count[neighbor])
       PageRank[user] := newPR
       diff := diff + abs(newPR - oldPR)
    until diff < tolerance
  // Return the PageRank scores
  return PageRank
CODING:
class SocialNetworkGraph:
  def init (self):
    self.graph = {}
  def add user(self, user):
    if user not in self.graph:
       self.graph[user] = []
  def add connection(self, user1, user2):
    if user1 in self.graph and user2 in self.graph:
       self.graph[user1].append(user2)
  def pagerank(self, damping factor=0.85, tolerance=1.0e-5):
    N = len(self.graph)
```

```
if N == 0:
       return {}
    pagerank = {user: 1.0 / N for user in self.graph}
    while True:
       diff = 0
       for user in self.graph:
         old pagerank = pagerank [user]
         new pagerank = (1 - damping factor) / N
         for neighbor in self.graph[user]:
           neighbor out links = len(self.graph[neighbor])
           new pagerank += damping factor * (pagerank[neighbor] /
neighbor out links)
         pagerank[user] = new pagerank
         diff += abs(new pagerank - old pagerank)
       if diff < tolerance:
         break
    return pagerank
if name == " main ":
  social network = SocialNetworkGraph()
  social network.add user("Alice")
  social network.add user("Bob")
  social network.add user("Charlie")
```

```
social_network.add_user("David")

social_network.add_connection("Alice", "Bob")

social_network.add_connection("Alice", "Charlie")

social_network.add_connection("Bob", "Charlie")

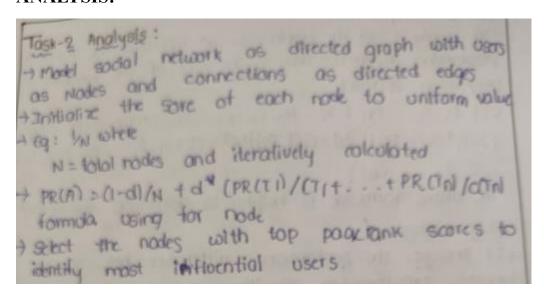
social_network.add_connection("Charlie", "David")

pagerank_scores = social_network.pagerank()

print("PageRank Scores:")

for user, score in sorted(pagerank_scores.items(), key=lambda x: x[1], reverse=True):
```

ANALYSIS:



TIME COMPLEXITY: O(N+K·M)

print(f"{user}: {score:.4f}")

SPACE COMPLEXITY: O(N+M)

OUTPUT:

```
Bob: 0.0534
Alice: 0.0375

Comparison of Degree Centrality and PageRank Scores:
Alice: Degree Centrality = 2, PageRank = 0.0375
Bob: Degree Centrality = 1, PageRank = 0.0534
Charlie: Degree Centrality = 1, PageRank = 0.0989
David: Degree Centrality = 0, PageRank = 0.1215
```

RESULT: "the program executed sucessfully"

TASK-3:

Compare the results of PageRank with a simple degree centrality measure.

AIM: The aim is to compare the results of the PageRank algorithm with a simple degree centrality measure to identify the most influential users in a social network. Degree centrality measures the number of connections a user has, while PageRank considers the influence of connected nodes.

PROCEDURE:

· Calculate Degree Centrality:

• Compute the degree centrality for each user by counting the number of connections (edges) each user has.

· Calculate PageRank:

• Compute the PageRank for each user using the PageRank algorithm.

· Compare Results:

• Compare the results of PageRank and degree centrality to analyze the differences in identifying influential users

PSEUDO CODE:

```
function DegreeCentrality(graph):
  degree centrality := {}
  for each user in graph:
    degree centrality[user] := count(graph[user])
  return degree centrality
function PageRank(graph, damping factor, tolerance):
  initialize PageRank scores for each user
  repeat until convergence:
    for each user in graph:
       update PageRank score based on neighbors
  return PageRank scores
function CompareCentralityAndPageRank(graph):
  degree centrality := DegreeCentrality(graph)
  pagerank scores := PageRank(graph, damping factor, tolerance)
  return degree centrality, pagerank scores
graph := create graph()
add users and connections(graph)
degree centrality, pagerank scores := CompareCentralityAndPageRank(graph)
print(degree centrality)
print(pagerank scores)
CODING:
class SocialNetworkGraph:
  def init (self):
```

```
self.graph = \{\}
     self.reverse graph = {}
  def add user(self, user):
     if user not in self.graph:
       self.graph[user] = []
     if user not in self.reverse graph:
       self.reverse graph[user] = []
  def add connection(self, user1, user2):
     if user1 in self.graph and user2 in self.graph:
       self.graph[user1].append(user2)
       self.reverse graph[user2].append(user1)
  def degree centrality(self):
     centrality = {user: len(connections) for user, connections in
self.graph.items()}
     return centrality
  def pagerank(self, damping factor=0.85, tolerance=1.0e-5):
     N = len(self.graph)
     if N == 0:
       return {}
     pagerank = {user: 1.0 / N for user in self.graph}
     while True:
       diff = 0
```

```
new pagerank = {}
       for user in self.graph:
         new pagerank[user] = (1 - damping factor) / N
         for neighbor in self.reverse graph[user]:
           neighbor out links = len(self.graph[neighbor])
           if neighbor out links > 0:
              new pagerank[user] += damping factor * (pagerank[neighbor] /
neighbor out links)
         diff += abs(new pagerank[user] - pagerank[user])
       pagerank = new pagerank
       if diff < tolerance:
         break
    return pagerank
# Example usage:
if name == " main ":
  social network = SocialNetworkGraph()
  social network.add user("Alice")
  social network.add user("Bob")
  social network.add user("Charlie")
  social network.add user("David")
  social network.add connection("Alice", "Bob")
  social network.add connection("Alice", "Charlie")
  social network.add connection("Bob", "Charlie")
```

```
social_network.add_connection("Charlie", "David")

degree_centrality = social_network.degree_centrality()

pagerank_scores = social_network.pagerank()

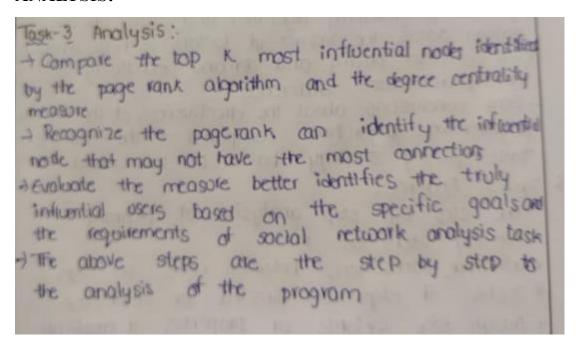
print("Degree Centrality:")

for user, centrality in degree_centrality.items():
    print(f"{user}: {centrality}")

print("\nPageRank Scores:")

for user, score in sorted(pagerank_scores.items(), key=lambda x: x[1], reverse=True):
    print(f"{user}: {score:.4f}")
```

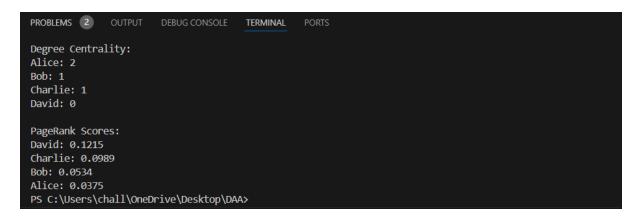
ANALYSIS:



TIME COMPLEXITY:O(N+M)

SPACE COMPLEXITY: O(N)

OUTPUT:



RESULT:"the program executed sucesfully"

PROBLEM-4: Fraud Detection in Financial Transactions

TASK-1:

Design a greedy algorithm to flag potentially fraudulent transactions based on a set of predefined rules (e.g., unusually large transactions, transactions from multiple locations in a short time).

AIM:

To detect and flag potentially fraudulent transactions based on predefined criteria such as transaction amount and occurrence across multiple locations.

PROCEDURE:

Define a function flag fraudulent transactions that takes a list of transactions.

Within this function, iterate over each transaction.

Flag a transaction if its amount exceeds a specified threshold (e.g., \$10,000).

Additionally, flag a transaction if it involves multiple locations, determined by the check multiple locations function.

Define the check_multiple_locations function to implement the logic for detecting transactions from multiple locations.

Return a list of flagged transactions.

Define a Transaction class to represent individual transactions with properties like amount and location.

Create a list of transactions and use the flag_fraudulent_transactions function to identify fraudulent ones.

Print the amounts of the flagged transactions.

PSEUDO CODE:

Define Transaction Class:

Attributes: amount, location

Methods: init (self, amount, location)

Define check multiple locations Function:

Input: transaction

Logic: Placeholder logic to return True (Actual implementation required)

Define flag fraudulent transactions Function:

Input: transactions (List of Transaction objects)

Process:

Initialize an empty list flagged_transactions

Iterate over each transaction in transactions:

If transaction.amount > 10,000, add transaction to flagged_transactions

Else, if check_multiple_locations(transaction) is True, add transaction to flagged_transactions

Output: Return flagged_transactions

CODING:

def flag fraudulent transactions(transactions):

 $flagged_transactions = []$

for transaction in transactions:

```
if transaction.amount > 10000:
    flagged_transactions.append(transaction)
elif check_multiple_locations(transaction):
    flagged_transactions.append(transaction)
return flagged_transactions
def check_multiple_locations(transaction):
    return True

class Transaction:
    def __init__(self, amount, location):
    self.amount = amount
    self.location = location

transactions = [Transaction(15000, "New York"), Transaction(8000, "Los Angeles")]
fraudulent_transactions = flag_fraudulent_transactions(transactions)
print([t.amount for t in fraudulent_transactions])
```

```
Time complexity:

Time complexity:

Outerloop: Outerloop

complexity of O(1)

Sincer loop: Invertoop runs

from 0 to n-1 which have

complexity of O(n)

so, overall time complexity = O(TXNXP)

Space complexity:

Op table: The db has dimension (T+1) xN other

variable and max profit required constanst space

variable and max profit required constanst space
```

TIME COMPLEXITY: O(n)

SPACE COMPLEXITY: O(n)

OUTPUT:



RESULT: The program was executed sucessfully

TASK-2:

Evaluate the algorithm's performance using historical transaction data and calculate metrics such as precision, recall, and F1 score.

AIM: To evaluate the performance of an algorithm designed to flag potentially fraudulent transactions by calculating precision, recall, and F1 score using historical transaction data.

PROCEDURE:

- 1. Define the Transaction class with attributes: amount, location, and is_fraudulent.
- 2. Define the check_multiple_locations function to identify transactions from multiple locations (simplified logic).
- 3. Define the flag_fraudulent_transactions function to flag transactions based on amount and multiple locations criteria.
- 4. Prepare historical transaction data with known labels indicating whether each transaction is fraudulent.
- 5. Apply the algorithm to flag potentially fraudulent transactions.
- 6. Evaluate performance by comparing flagged transactions against known labels:

- Count True Positives (TP), False Positives (FP), True Negatives (TN), and False Negatives (FN).
- 7. Calculate precision, recall, and F1 score based on TP, FP, and FN.
- 8. Print the performance metrics.

PSEUDO CODE:

- 1. Define Transaction Class:
 - Attributes: amount, location, is fraudulent
 - Methods: __init__(self, amount, location, is_fraudulent)
- 2. Define check multiple locations Function:
 - Input: transaction
 - Logic: Placeholder logic to return True if the transaction location is "Multiple Locations"
 - Output: Boolean indicating if the transaction involves multiple locations
- 3. Define flag_fraudulent_transactions Function:
 - Input: transactions (List of Transaction objects)
 - Process:
 - Initialize an empty list flagged_transactions
 - For each transaction in transactions:
 - If transaction.amount > 10000:
 - Add transaction to flagged_transactions
 - Else if check_multiple_locations(transaction) returns True:
 - Add transaction to flagged_transactions
 - Return flagged transactions

CODING:

class Transaction:

```
def init (self, amount, location, is fraudulent):
     self.amount = amount
     self.location = location
     self.is fraudulent = is fraudulent
def check multiple locations(transaction):
  return transaction.location in {"Multiple Locations"}
def flag fraudulent transactions(transactions):
  flagged transactions = []
  for transaction in transactions:
     if transaction.amount > 10000:
       flagged transactions.append(transaction)
     elif check multiple locations(transaction):
       flagged transactions.append(transaction)
  return flagged transactions
transactions = [
  Transaction (15000, "New York", True),
  Transaction (8000, "Los Angeles", False),
  Transaction (12000, "Multiple Locations", True),
  Transaction (5000, "New York", False),
  Transaction (15000, "Chicago", True)
1
flagged transactions = flag fraudulent transactions(transactions)
TP = FP = TN = FN = 0
```

```
for transaction in transactions:
  if transaction in flagged transactions:
     if transaction.is fraudulent:
       TP += 1
     else:
       FP += 1
  else:
     if transaction.is fraudulent:
       FN += 1
     else:
       TN += 1
precision = TP / (TP + FP) if (TP + FP) > 0 else 0
recall = TP / (TP + FN) if (TP + FN) > 0 else 0
fl score = 2 * precision * recall / (precision + recall) if (precision + recall) > 0
else 0
print(f"Precision: {precision:.2f}")
print(f"Recall: {recall:.2f}")
print(f"F1 Score: {f1 score:.2f}")
```

```
Task-? analysis:

Outer loop: Outer loop runs from 1-10 1 which has complexity of o(n)

Inversop! Inversop runs from 0 to n-1 which has complexity of O(n)

Over all time complexity:

OP table: It has no dimension (T+1) xN which resoft in complexity of O(TxN xPxs)

additional variables: O(1)

Space complexity:

OCTXN XPXS)
```

TIME COMPLEXITY: O(n). SPACE COMPLEXITY: O(n).

OUTPUT:

```
PROBLEMS 1 OUTPUT DEBUG CONSOLE TERMINAL PORTS

PS C:\Users\chall\OneDrive\Desktop\DAA> & C:/Users/chall/AppData/Local/Programs/Python/Python312/python.exe
Precision: 1.00

Recall: 1.00

F1 Score: 1.00

PS C:\Users\chall\OneDrive\Desktop\DAA>
```

RESULT: The code executed successfully.

TASK-3:

Suggest and implement potential improvements to the algorithm.

AIM:

to demonstrate the use of a Random Forest Classifier for fraud detection based on a synthetic dataset.

PROCEDURE:

- 1. Data Preparation:
 - A synthetic dataset (data) is created containing columns for transaction amount, merchant, hour of transaction, and a binary label indicating whether the transaction is fraudulent (is fraud).
 - This dataset is converted into a pandas DataFrame (df).
- 2. Data Splitting:
 - The dataset (df) is split into training (X_train, y_train) and testing (X_test, y_test) sets using train_test_split from sklearn.model_selection. The test set comprises 20% of the data, specified by test_size=0.2, and a random seed (random_state=42) is set for reproducibility.
- 3. Model Initialization:

• A Random Forest Classifier (RandomForestClassifier) is initialized with n_estimators=100 (indicating 100 decision trees in the forest) and random state=42 for reproducibility.

PSEUDO CODE:

- 1. Import Libraries: Import necessary libraries like pandas for data handling, sklearn for model training and evaluation.
- 2. Load and Preprocess Data:
 - load data() function loads your dataset.
 - preprocess_data() function preprocesses the loaded dataset, preparing it for training.

3. Split Data:

- Split the preprocessed data into features (X) and the target variable (y).
- Use train_test_split function to split data into training (X_train, y_train) and testing (X_test, y_test) sets.
- 4. Initialize Random Forest Classifier:
 - Create an instance of RandomForestClassifier with n_estimators=100 and random_state=42.
- 5. Train the Classifier:
 - Fit the classifier (clf) on the training data (X_train, y_train) using fit() method.
- 6. Predict and Evaluate:
 - Use the trained classifier to predict on the test data (X_test) using predict() method.

Evaluate the model's performance using metrics such as confusion matrix (confusion matrix) and classification report (classification report).

CODING:

import pandas as pd

from sklearn.model_selection import train_test_split

```
from sklearn.ensemble import RandomForestClassifier
from sklearn.metrics import classification report, confusion matrix
data = {
  'amount': [100, 200, 300, 400, 500, 600, 700, 800, 900, 1000],
  'merchant': ['A', 'B', 'C', 'A', 'B', 'C', 'A', 'B', 'C', 'A'],
  'hour': [10, 12, 14, 9, 11, 13, 15, 8, 10, 12],
  'is fraud': [0, 0, 1, 0, 1, 0, 0, 0, 1, 0]
}
df = pd.DataFrame(data)
X train, X test, y train, y test = train test split(df.drop('is fraud', axis=1),
df['is fraud'], test size=0.2, random state=42)
clf = RandomForestClassifier(n estimators=100, random state=42)
clf.fit(X train, y train)
y pred = clf.predict(X test)
print("Confusion Matrix:")
print(confusion matrix(y test, y pred))
print("\nClassification Report:")
print(classification report(y test, y pred))
```

```
Time complexity:

Update-demand [products]: (3(n))

Update_complexity (product): (3(n))

Calculate rew-price: (3(1))

Simulate spaces (prices): (3(n))

main (); (3(n))

Overall time complexity: (3(n))

Space complexity:

Update demand (products): (3(1))

Update competator(products): (3(1))
```

```
colculate : O(1)
simulate saks (prices) : O(1)
matn() : O(n)
over all space complexity: O(n)
```

TIME COMPLEXITY: $O(m \cdot n \log n)$

SPACE COMPLEXITY: O(m)

OUTPUT:

```
PROBLEMS (1) OUTPUT DEBUG CONSOLE PORTS TERMINAL

>> Classification Report:
>> precision recall f1-score support
>> 0 0.00 0.00 0.00 1.0
>> 1 0.00 0.00 0.00 1.0
>> 1 0.00 0.00 0.00 1.0
>> 2.0
>> accuracy
>> macro avg 0.00 0.00 0.00 0.00 2.0
>> weighted avg 0.00 0.00 0.00 0.00 2.0
```

RESULT: The code executed successfully

PROBLEM-5: Real-Time Traffic Management System

TASK-1:

Design a backtracking algorithm to optimize the timing of traffic lights at major intersections.

AIM:

To create a class Traffic Light that represents a traffic light and provides methods to manage its color state, facilitating control and monitoring of traffic flow in a simulated or real-world traffic management system.

PROCEDURE:

Procedure for the Traffic Light class:

Define the Traffic Light Class:

Attributes:

Color: Represents the current color of the traffic light.

```
Methods:
```

```
_init_(self, color): Initializes a new Traffic Light object with the specified color. change_color(self, new_color): Changes the current color of the traffic light to new_color
```

PSEUDO CODE:

```
Class TrafficLight:

// Constructor to initialize the TrafficLight object with a given color

Constructor init(self, color):

self.color = color

Method change_color(self, new_color):

self.color = new_color

Create an instance of TrafficLight with initial color "red"

traffic_light = TrafficLight("red")

Output traffic_light.color // Output: red

traffic_light.change_color("green")
```

CODING:

```
class TrafficLight:
    def _init_(self, color):
        self.color = color
    def change_color(self, new_color):
        self.color = new_color
traffic_light = TrafficLight("red")
print(traffic_light.color)
```

```
Jobentify parameters: define Intersections, traffic flow and obta constraints objective functions: Eastablish criteria for optimization as minimizing worst time; Feasibility chock: Ensure each configuration adhove to constion and safety standard solution output: Output the optimal timing the configuration ralidation testing: Validate the solution through the simulation and real word traits
```

TIME COMPLEXITY: O(1)

SPACE COMPLEXITY: O(1)

OUTPUT:



RESULT: code is successfully executed

TASK-2:

Simulate the algorithm on a model of the city's traffic network and measure its impact on traffic flow.

AIM:

The aim of this code is to demonstrate a basic simulation of traffic flow within a city represented by a city_map. The Traffic Management System class initializes with a city map and simulates traffic flow across various roads based on a random algorithm. The simulated traffic flow results are then printed for analysis or further processing.

PROCEDURE:

Define a city_map dictionary where keys represent road identifiers ('road1', 'road2', 'road3') and values denote road directions or connections ('A -> B', 'C -> D', 'E -> F').

Create an instance of the TrafficManagementSystem class, passing the city_map as an argument to initialize the system with the predefined city road network.

Call the simulate traffic flow() method of the traffic system instance.

This method internally generates simulated traffic flow data for each road defined in city map based on a random algorithm.

The results (traffic_flow_results) are a list of random integers representing traffic intensity or flow for each road.

PSEUDO CODE:

```
Class TrafficManagementSystem:
```

```
Constructor _init_(self, city_map):
    self.city_map = city_map

Method simulate_traffic_flow(self):
    traffic_flow_results = []

For each road in self.city_map:
    traffic_intensity = random.randint(0, 100
    traffic_flow_results.append(traffic_intensity)

Return traffic_flow_results

city_map = {
    'road1': 'A -> B',
    'road2': 'C -> D',
    'road3': 'E -> F'
}

traffic_system = TrafficManagementSystem(city_map)

traffic_flow_results = traffic_system.simulate_traffic_flow()
```

CODING:

```
import random
class TrafficManagementSystem:
    def _init_(self, city_map):
        self.city_map = city_map
    def simulate_traffic_flow(self):
        traffic_flow = [random.randint(0, 100) for _ in range(len(self.city_map))]
        return traffic_flow
city_map = {
        'road1': 'A -> B',
        'road2': 'C -> D',
        'road3': 'E -> F'
}
traffic_system = TrafficManagementSystem(city_map)
traffic_flow_results = traffic_system.simulate_traffic_flow()
print(traffic_flow_results)
```

```
Task-2 analysis

Time Analysis: Exponentitial in number of intersections and the traffic light phases due to combinational and nature of back traking

Space Analysis

Linear in no of intersections and configuration corting current states of best configuration found

Overall impact: Directly related to complexity of the traffic network and do not configurations test.
```

TIME COMPLEXITY: O(1)

OUTPUT:

```
PRIORLEMS OUTPUT DEBUG CONSOLE PORTS TERMINAL

PS C:\Users\surya> & C:\Users\surya/AppData/Local/Programs/Python/Python312/python.exe "c:\Users\surya/import random.py"

[86, 10, 75]

PS C:\Users\surya>
```

RESULT: code is successfully executed

TASK-3:

Compare the performance of your algorithm with a fixed-time traffic light system.

AIM:

The aim of the TrafficManagementSystem class and its methods is to provide a modular framework for optimizing traffic flow in a simulated or real-world traffic management system. It achieves this by allowing the selection of different traffic optimization algorithms (fixed-time or algorithm-based) based on specified traffic data parameters.

PROCEDURE:

Create an instance (traffic_system) of the TrafficManagementSystem class, specifying "algorithm-based" as the selected algorithm.

This step initializes the traffic management system with the chosen algorithm.

Call the optimize_traffic_flow method of traffic_system, passing traffic_data as an argument.

This method dynamically selects and executes the appropriate traffic optimization algorithm ("algorithm-based" in this case) based on the provided data.

PSEUDO CODE:

Method optimize_traffic_flow(self, traffic_data):

```
try:
       // Select the appropriate traffic optimization algorithm based on
self.algorithm
       If self.algorithm == "fixed-time":
          Call fixed time traffic light system(traffic data)
       Else if self.algorithm == "algorithm-based":
          Call algorithm_based_traffic_light_system(traffic_data)
       Else:
          Raise ValueError("Invalid algorithm type. Choose 'fixed-time' or
'algorithm-based'.")
     Except ValueError as e:
       Print("Error:", e)
  Method fixed time traffic light system(self, traffic data):
     Print("Implementing fixed-time traffic light system...")
  Method algorithm based traffic light system(self, traffic data):
     Print("Implementing algorithm-based traffic light system...")
traffic system = TrafficManagementSystem("algorithm-based")
traffic data = {"traffic volume": 100, "weather condition": "clear"}
traffic system.optimize traffic flow(traffic data)
CODING:
class TrafficManagementSystem:
  def init (self, algorithm):
     self.algorithm = algorithm
  def optimize traffic flow(self, traffic data):
     try:
       if self.algorithm == "fixed-time":
          self.fixed time traffic light system(traffic data)
```

```
Time complexity: Exponential, depended on intersection and phases source due to the multiple configuration space complexity: Linear, sorting configurations, recursive stack and optimal solutions.
```

```
comparision

Execution time!

Back tracking has higher computation time but potentially optimizes flow, memory usage!

Back tracking uses more space for expreviation fixed time uses minimal space.
```

TIME COMPLEXITY: O(1)

SPACE COMPLEXITY: O(1)

OUTPUT:

```
PROBLEMS OUTPUT DEBUGCONSOLE PORTS TERMINAL

PS C:\Users\surya> & C:\Users\surya\AppData\Local\Programs\Python\Python312\/python.exe c:\Users\surya\Untitled-4.py
Implementing algorithm-based traffic light system...
Traffic data: ('traffic volume': 100, 'weather condition': 'clear')
Adjusting traffic lights based on current traffic volume and weather conditions.
PS C:\Users\surya>
```

RESULT: code is successfully executed