

5> We know 'a' is loaded already in R0.

So, we use the following table:

Line No.	Operation	Assembly Code	Cost
101	$x = a + x$	ADD x, R0	2
Now R0 contains $x = a + x$.			
102	$y = b * 2$	Mov b, R1	2
		MUL 2, R1	2
Now, y in R1 ($y = b * 2$)			
103	$a = x - y$	SUB R0, R1 SUB R1, R0 MOV R0, a	1 2
So, finally, we have <u>a in R0</u> .			
			<u>9</u>

Total Cost = 9 ops

17

Grammar

$$S' \rightarrow S$$

$$S \rightarrow CB$$

$$S \rightarrow BC$$

$$C \rightarrow Cad$$

$$C \rightarrow d$$

$$B \rightarrow BaC$$

$$B \rightarrow a$$

LR(1) Table

goto

State

closure

0

$\{[S' \rightarrow S, \$], [S \rightarrow CB, \$]$
 $[S \rightarrow \cdot BC, \$], [C \rightarrow \cdot Cad, a]$
 $[C \rightarrow \cdot d, a], [B \rightarrow \cdot BaC, d/a]$
 $[B \rightarrow \cdot a, d/a]\}$

goto(0, s)

1

$\{[S' \rightarrow S, \$]\}$

goto(0, C)

2

$\{[S \rightarrow C \cdot B, \$], [C \rightarrow C \cdot ad, a]$
 $[B \rightarrow \cdot BaC, \$/a]$
 $[B \rightarrow \cdot a, \$/a]\}$

goto(0, B)

3

$\{[S \rightarrow B \cdot C, \$], [B \rightarrow B \cdot aC, d/a]$
 $[C \rightarrow \cdot Cad, \$/a]$
 $[C \rightarrow \cdot d, \$/a]\}$

goto(0, d)

4

$\{[C \rightarrow d \cdot, a]\}$

goto(0, a)

5

$\{[B \rightarrow a \cdot, d/a]\}$

- | | | |
|-------------|----|--|
| goto(2, b) | 6 | $\{ [C \rightarrow cb, \$], [B \rightarrow B.ac, \$ a] \}$ |
| goto(2, a) | 7 | $\{ [C \rightarrow ca.d, a], [B \rightarrow B.a, \$ a] \}$ |
| goto(3, c) | 8 | $\{ [C \rightarrow bc, \$], [C \rightarrow cad, \$ a] \}$ |
| goto(3, d) | 10 | $\{ [C \rightarrow d, \$ a] \}$ |
| goto(6, a) | 11 | $\{ [B \rightarrow B.ac, \$ a], [C \rightarrow cad, \$ a], [C \rightarrow d, \$ a] \}$ |
| goto(7, b) | 12 | $\{ [C \rightarrow cad, a] \}$ |
| goto(8, a) | 13 | $\{ C \rightarrow cad, \$ a \}$ |
| goto(9, c) | 14 | $\{ [B \rightarrow B.ac, d a], [C \rightarrow cad, d a] \}$ |
| goto(9, d) | 15 | $\{ [C \rightarrow d, d a] \}$ |
| goto(11, c) | 16 | $\{ [B \rightarrow B.ac, \$ a], [C \rightarrow cad, \$ a] \}$ |
| goto(11, d) | 10 | |
| goto(13, d) | 17 | $\{ [C \rightarrow cad, \$ a] \}$ |
| goto(14, a) | 18 | $\{ [C \rightarrow cad, d a] \}$ |
| goto(16, a) | 13 | |
| goto(18, d) | 19 | $\{ [C \rightarrow cad, d a] \}$ |

Now, first

$$S' \rightarrow \{b, a\}$$

$$C \rightarrow \{d\}$$

$$S \rightarrow \{b, a\}$$

$$B \rightarrow \{a\}$$

LR Table

State

Action

goto

	a	d	\$	S'	S	C	B
0	S5	S4		1	2	3	
1							
2			acc				
3	S7						
4	S9	S10					6
5	r4					8	
6	r6	r6					
7	S11		r1				
8	r6	S12	r6				
9	S13		r2				
10		S15	\$			14	
11	r4		r4				
12		S10				16	
13	r3						
14		S17					
15	S8/r5	r5					
16	r4	r4					
17	S13/r5		r5				
18	r3		r3				
19		S9					
	r3	r3					

2. > 3 address code for the given C-code
 can be written as:

$i := 0$ ----- leader

L1 : if $i < n$ goto S1 ----- leader
 goto L2 ----- leader
 ----- leader

S1 : $t1 := a[i]$
 $t2 := a[n]$
 $t3 := t2[t1]$
 $t3 := t3 / \text{env}$
 $t3 := t3 \% 10$
 $t3 := t3 * 4$
 $t1 := \text{count}$
 $t2 := t1[t3]$
 $t2 := t2 + 1$
 $t1[t3] := t2$
 $i := i + 1$
 goto L1

L2 : $i = 1$ ----- leader
 L3 : if $i < 10$ goto S2 ----- leader
 goto L4 ----- leader
 ----- leader

S2 : $t1 := i - 1$
 $t1 := t1 * 4$
 $t2 = \text{count}$
 $t3 := t2[t1]$
 $t4 = t4$
 $t4 = t4 + t3$
 $t2[t1] = t4$
 $i := i + 1$

goto L3

L4: $i := n$
 $i := i - 1$

leader

L5: if $i > 0$ goto S3
 goto L6

leader

leader

S3: $t1 := 4 \times i$

leader

$t2 := \text{arr}$

$t3 := t2[t1]$

$t3 := t3 / \text{emp}$

$t3 := t3 \cdot 10$

$t1 := \text{count}$

$t2 := t1[t3]$

$t2 := t2 - 1$

$t2 := t2 \cdot 4$

$t1 := i \cdot 4$

$t3 := \text{arr}$

$t4 := t3[t1]$

$t1 := \text{output}$

$t1[t2] := t4$

$t1 := 4 \cdot 1$

$t2 := \text{arr}$

$t3 := t2[t1]$

$t3 := t3 / \text{emp}$

$t3 := t3 \cdot 10$

$t1 := \text{count}$

$t2 := t1[t3]$

$t2 := t2 - 1$

$t1[t2] := t2$

$i := i - 1$

goto L5

L6: $i := 0$ ———→ leader
 L7: if $i < n$ goto S4 ——— leader
 goto L8 ——— leader

S4: $t1 := i \times 4$ ——— leader
 $t2 := \text{output}$
 $t3 := t2[t1]$
 $t2 := \text{arr}$
 $t2[t1] := t3$
 $i := i + 1$
 goto L7

L8:

Now, we have

16 Blocks

$$37 \quad x := b[i] + c^d * t - 2$$

TAC can be written as:-

$$t_1 = i$$

$$t_2 = t_1 * 4$$

$$t_3 = b[t_2]$$

$$t_4 = d * t$$

$$t_5 = t_3 + c$$

$$t_6 = t_4 - 2$$

$$t_7 = t_5 \text{ XOR } t_6$$

$$x = t_7$$

Triples Representation:-

#	Op	Arg 1	Arg 2
(0)	Assign	i	
(1)	*	4	(0)
(2)	[]	b	(1)
(3)	*	d	t
(4)	+	c	(2)
(5)	-	(3)	2
(6)	XOR	(4)	(5)
(7)	assign	x	(6)

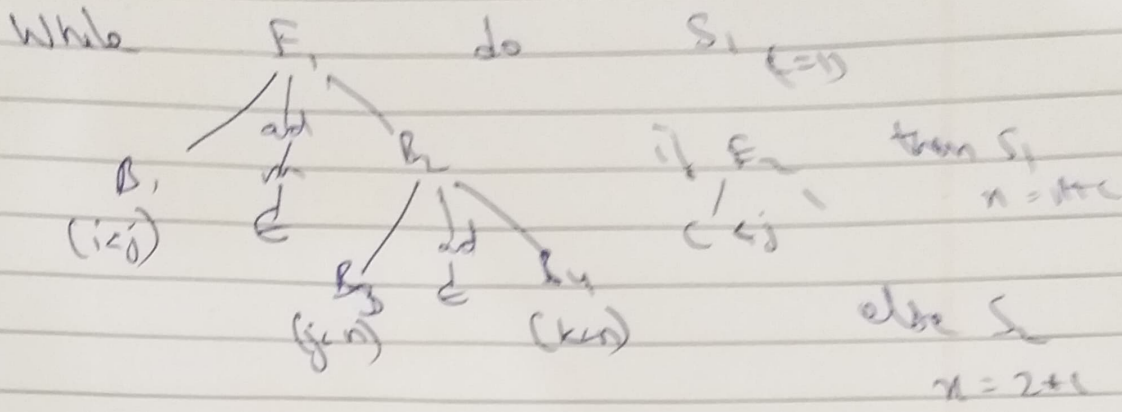
4.7

TAC without backpadding

Filling goto

1:	if $i < j$ goto	3
2:	goto	21
3:	if $j < n$ goto	7
4:	goto	5
5:	if $k < n$ goto	7
6:	goto	21
7:	$C = 1$	
8:	goto	9
9:	if $C < j$ goto	10 11
10:	goto	13
11:	$x = x + C$	
12:	goto	15
13:	$x = 2 * C$	
14:	goto	15
15:	$C = C + 1$	
16:	goto	17
17:	$j = j + 1$	
18:	goto	19
19:	$K = K * 2$	
20:	goto	1
21:	$a = x + n$	

Abstracted Tree:



S_1 $a = n + n$

S_2 $c = c + 1$

S_3 $f = j + 1$

S_4 $k = k + 1$

Using these, we can fill the gap