

# Introduction

- Random bit streams used in key generation and encryption...
- Two strategies:
  1. Compute bits deterministically using an algorithm- Pseudo-Random Number Generators (PRNGs) or Deterministic Random Bit Generators (DRBGs).
  2. Produce bits non-deterministically using some physical source that produces some sort of random output- True Random Number Generators (TRNGs) or Non-deterministic Random Bit Generators (NRBGs).

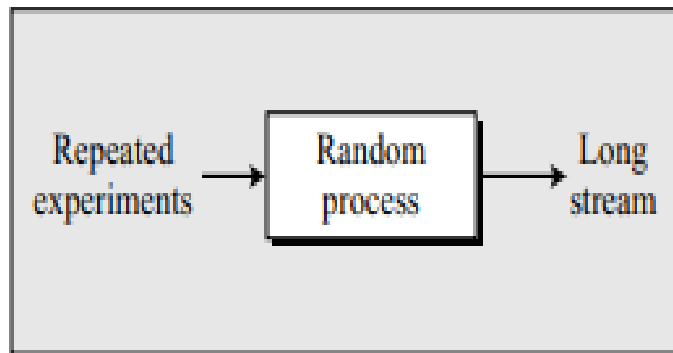
# The use of Random Numbers

- A number of network security algorithms and protocols based on cryptography make use of random binary numbers:
  1. Key distribution and reciprocal (mutual) authentication schemes
  2. Session key generation
  3. Generation of keys for the RSA public-key encryption algorithm
  4. Generation of a bit stream for symmetric stream encryption
- Two requirements for a sequence of random numbers: randomness and unpredictability

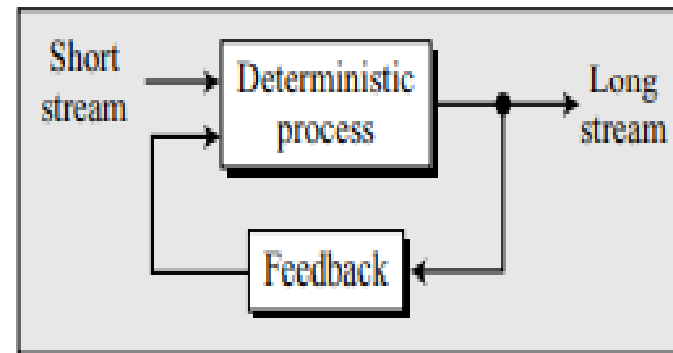
# The Two Approaches

- There are two approaches to generating a long stream of random bits:
  - Using a natural random process, such as flipping a coin many times and interpreting heads and tails as 0-bits and 1-bits,
  - Or using a deterministic process with feedback.
- The first approach is called a **True Random Number Generator (TRNG)**;
- The second is called a **Pseudo-Random Number Generator (PRNG)**.

# The Two Approaches



a. TRGN



b. PRNG

# PRNG

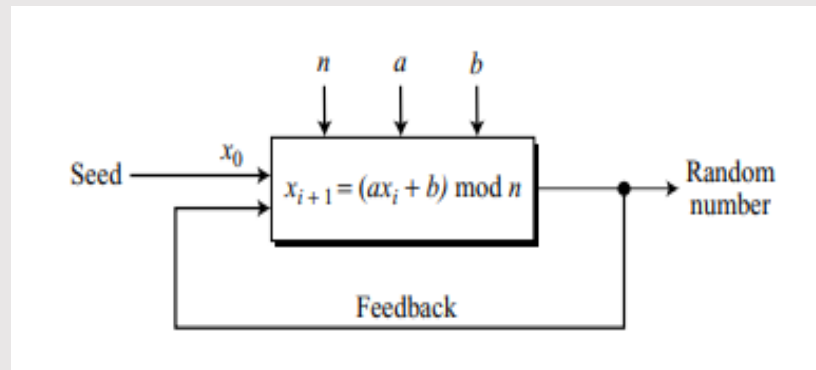
- A reasonably random stream of bits can be achieved using a deterministic process with a short random stream as the input (seed).
- The generated number is not truly random because the process that creates it is deterministic.
- PRNGs can be divided into **two broad categories**:
  - **Congruential generators** and
  - **Generators using cryptographic ciphers**

# Congruential Generators: Linear Congruential Generator

- This method recursively creates a sequence of pseudorandom numbers using a linear congruence equation of the form:

$$x_{i+1} = (ax_i + b) \bmod n,$$

where  $x_0$ , called the seed, is a number between 0 and  $n - 1$ .



# Congruential Generators:

## Linear Congruential Generator

- The sequence is periodic, where the period depends on how carefully the coefficients,  $a$  and  $b$ , are selected. The ideal is to make the period as large as the modulus  $n$ .
- **Example**
  - Assume that  $a = 4$ ,  $b = 5$ ,  $n = 17$ , and  $x_0 = 7$ .
  - Find out the sequence and the period

# Congruential Generators:

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- **Example**
  - Assume that  $a = 4$ ,  $b = 5$ ,  $n = 17$ , and  $x_0 = 7$ .
  - The sequence is 16, 1, 9, 7, 16, 1, 9, 7, ..., which is definitely a poor pseudorandom sequence; the period is only 4.



# Congruential Generators:

## Linear Congruential Generator

- **Criteria for an acceptable PRNG:**

1. The period must be equal to  $n$  (the modulus). This means that, before the integers in the sequence are repeated, all integers between 0 and  $n - 1$  must be generated.
2. The sequence in each period must be random.
3. The generating process must be efficient. Most computers today are efficient when arithmetic is done using 32-bit words.

# Congruential Generators:

## Linear Congruential Generator

- **Recommendation:**

For selecting the coefficients of the congruence equation and the value of the modulus:

1. For the modulus,  $n$ , choose the largest prime number close to the size of a word in the computer being used. The recommendation is to use the thirty-first Mersenne prime as the modulus:  $n = M_{31} = 2^{31} - 1$ .
2. To create a period as long as the modulus, the value of the first coefficient,  $a$ , should be a primitive root of the prime modulus. Although the integer 7 is a primitive root of  $M_{31}$ , it is recommended to use  $7^k$ , where  $k$  is an integer coprime with  $(M_{31} - 1)$ . Some recommended values for  $k$  are 5 and 13. This means that  $(a = 7^5)$  or  $(a = 7^{13})$ .
3. For the second recommendation to be effective, the value of the second coefficient,  $b$ , should be zero.

# Congruential Generators:

## Linear Congruential Generator

- Linear Congruential Generator:

$$x_{i+1} = (ax_i + b) \bmod n,$$

where  $n = 2^{31} - 1$  and  $(a = 7^5)$  or  $(a = 7^{13})$ .

- **Security**
  - Shows reasonable randomness if the previous recommendations are followed.
  - The sequence is useful in some applications where only randomness is required (such as simulation);
  - It is useless in cryptography where both randomness and secrecy are desired.

# Congruential Generators:

## Linear Congruential Generator

- Because  $n$  is public, the sequence can be attacked by Eve using one of the two strategies:
  - If Eve knows the value of the seed ( $x_0$ ) and the coefficient  $a$ , she can easily regenerate the whole sequence.
  - If Eve does not know the value of  $x_0$  and  $a$ , she can intercept the first two integers and use the following two equations to find  $x_0$  and  $a$ :

$$x_1 = ax_0 \bmod n$$

$$x_2 = ax_1 \bmod n$$

# Other Congruential Generators

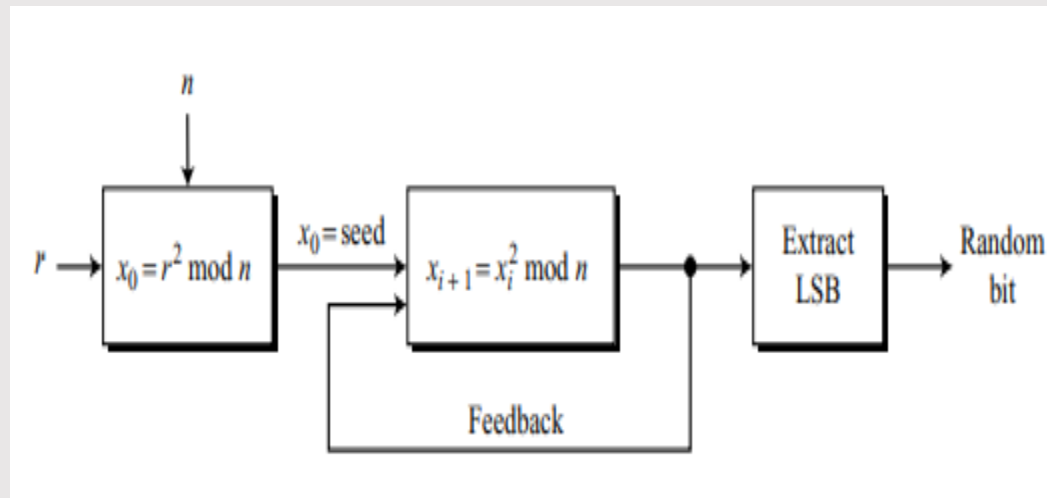
- **Quadratic Residue Generator**

- To make the pseudorandom sequence less predictable.
- $x_{i+1} = x_i^2 \bmod n$ , where  $x_0$ , called the seed, is a number between 0 and  $n - 1$ .

- **Blum Blum Shub(BBS) Generator**

- After the names of its three inventors.
- BBS uses quadratic residue congruence, but it is a pseudorandom bit generator instead of a pseudorandom number generator;
- It generates a sequence of bits (0 or 1).

# Congruential Generators: Blum Blum Shub(BBS) Generator



# Congruential Generators:

## Blum Blum Shub(BBS) Generator

- **Steps:**

1. Find two large primes numbers  $p$  and  $q$  in the form  $4k + 3$ , where  $k$  is an integer (both  $p$  and  $q$  are congruent to 3 modulo 4).
2. Select the modulus  $n = p \times q$ .
3. Choose a random integer  $r$  which is coprime to  $n$ .
4. Calculate the seed as  $x_0 = r^2 \bmod n$ .
5. Generate the sequence as  $x_{i+1} = x_i^2 \bmod n$ .
6. Extract the least significant bit of the generated random integer as the random bit

# Congruential Generators:

## Blum Blum Shub(BBS) Generator

- **Security**

- If  $p$  and  $q$  are known, the  $i^{\text{th}}$  bit in the sequence can be found as the least significant bit of
- $x_i = x_0^{2^i \bmod [(p-1)(q-1)]} \bmod n$
- If Eve knows the value of  $p$  and  $q$ , she can find the value of the  $i^{\text{th}}$  bit by trying possible values of  $x_0$  (the value of  $n$  is usually public).
- This means that the complexity of this generator is the same as the factorization of  $n$ .
- If  $n$  is large enough, the sequence is secure (unpredictable).
- It has been proved that with a very large  $n$ , Eve cannot guess the value of the next bit in the sequence even if she knows the values of all previous bits.
- The probability of each bit being 0 or 1 is very close to 50 %.

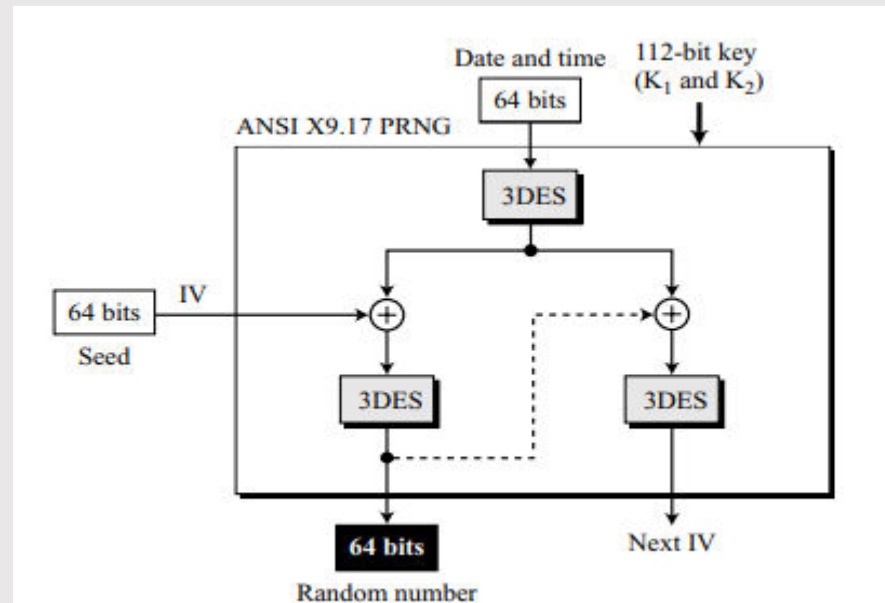


# Cryptosystem-Based Generators

- A cryptosystem such as an encryption cipher or a hash function can also be use to generate a random stream of bits.
- **ANSI X9.17 PRNG**
- A cryptographically strong pseudorandom number generator.
- The generator uses three 3DES with two keys (encryption-decryption-encryption).

# Cryptosystem-Based Generators: ANSI X9.17 PRNG

- **Fig: The Cipher-Block Chaining (CBC) mode**
  - Note that the first pseudorandom number uses a 64-bit seed as the initial vector (IV);
  - the rest of the pseudorandom numbers use the seed shown as the next IV.
  - The same 112-bit secret key ( $K_1$  and  $K_2$  in 3DES), are used for all three 3DES ciphers.



# Cryptosystem-Based Generators:

## ANSI X9.17 PRNG

- Uses two stages of the block chaining.
- The plaintext for each stage comes from the output of the first 3DES, which uses the 64-bit date and time as the plaintext.
- The ciphertext created from the second 3DES is the random number.
- The ciphertext created from the third 3DES is the next IV for the next random number.

# Cryptosystem-Based Generators: ANSI X9.17 PRNG

- The strength of X9.17 can be due to the following facts:
  1. The key is 112 ( $2 \times 56$ ) bits.
  2. The date-and-time input of 64 provides a good timestamp preventing replay attack.
  3. The system provides an excellent confusion-diffusion effect with six encryptions and three decryptions.

# Cryptosystem-Based Generators:

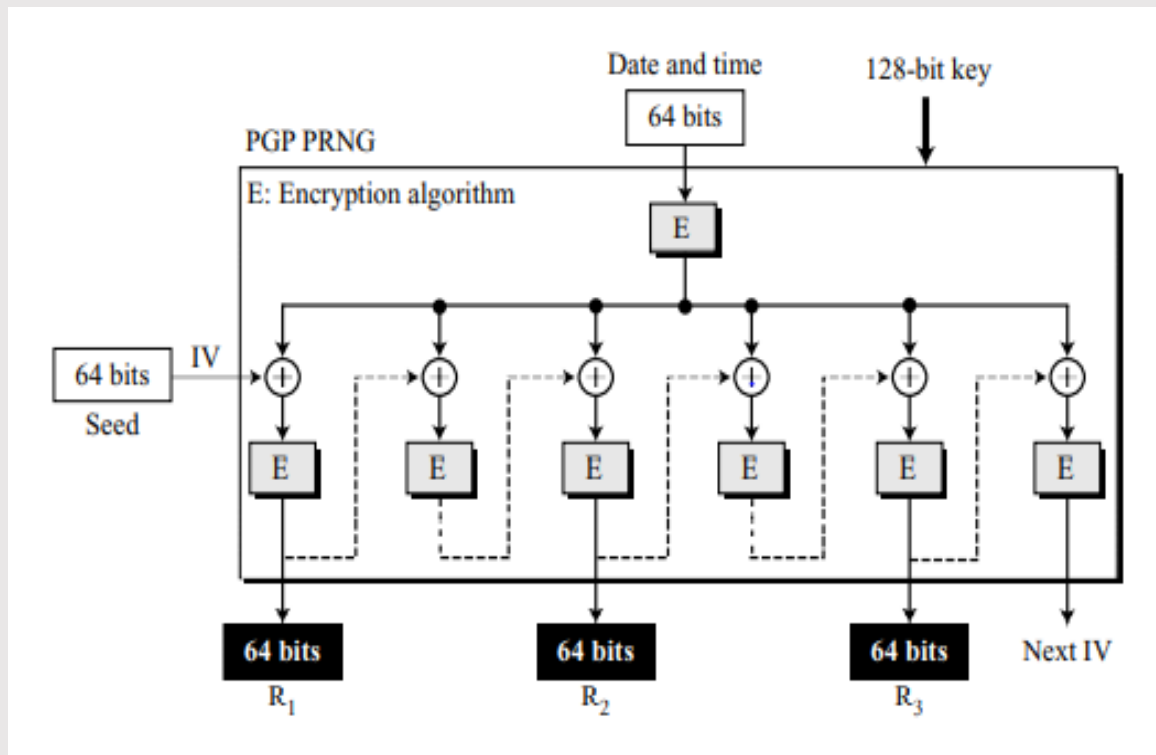
## PGP PRNG

- **PGP PRNG**

- Uses the same idea as X9.17 with several changes:

1. PGP PRNG uses seven stages instead of two.
2. The cipher is either IDEA or CAST-128.
3. The key is normally 128 bits. PGP PRNG creates three 64-bit random numbers: the first is used as the IV secret (for communication using PGP, not for PRNG), the second and the third are concatenated to create a 128-bit secret key (for communication using PGP).

# Cryptosystem-Based Generators: PGP PRNG



- The strength of PGP PRNG is in its key size and in the fact that the original IV (seed) and the 128-bit secret key can be generated from a 24-byte true random variable.