

format readable format without knowing P.T > D Eigher the plaintext of the runage without having access to the key. prove technical définition can be, Cryptanalysis is used to break cryptographic security systems and gain access to the contents of the encrypted menages, even if oryptographie key is inknown. There are various cryptanalytic attacks

(i) Ciphertext only attack

(ii) Ciphertext only attack

(iii) Knows only ciphertext.

Cryptanalysis is used to break cryptographic security systems and gain access to the contents of the encrypted menages, even if oryptographie key is inknown. There are various cryptanalytic attacks (i) Ciphertext only attack.
altacker knows only ciphertext. (ii) Known plaintext only attack attacker knows some combination of Pi, Ci attacker knows some combination of Pi, Ci cend based on these, he try to decrypt the nessages. sender ciphertext decrypt pleceiver AHacker

chosen plaintext attack

model of oryptanelysis which arrunes that the
attacker can choose random plaintexts to be
enerypted and obtain the covers ponding D LE DE 62 82 23 92 92 92 22 27 12 02 61 81 21 91 51 FL E1 21 11 01 6 8 2 9 5 W E 2 1 SSJIMIW cipher texts. ST YAAUNAL The goal of attacker is to gain further Enfo. which reduces the security of the encryption whene. S-cipher, [] In the moust case, this attack can expose the secret key secret information after calculating the secret key attacker. (iv) Chasen Ciphertext attach

Can analyze any

attacher decryptions - plaintexts Pi.

What is Cryptology in Network Security, this attack can expose the secret information after calculating the secret key. (iv) Chosen Ciphertext attack

can analyze any
attacher decryptions - plaintexts Pi. His goal is to acquire a secret key or to get as many info. about the attacked ciphen Receiver system as possible. to make the victim Attacker -The attacker has capability decrypt any ciphertext and send him back the chosen ciphertext o result analy zung 7:27 / 10:02

ElGamal Cryptography: Asymmetric Key. " Key Generation: 1) Select Large Prime no. (P) => P=11 11) Select decryption Key/Private Key(D)=3 iii) Select second paut of encryption key iii) Decryption: 08 public Key (E1) = 2

IV) Third part of the encryption key or public Key (E2). E2 = E1 mod P. = 8 v) Public Key = (E1, E2, P), Private Key = D

(125) mod 11s →(2,8,11), L=3 (125 x x) mod 11 = 1

iii Encryption:

1) Select Random Integer (R) \$4 ii) C1 = E1 mod P, C1= 24 mod 11=5 iii) C2 = (PT x E2R) mod P=(7x84) iv) C.T= ((1,(2)) => 28672 mod 11

PT= [C2x((1))-1] mod P

P= 11, D=3, E1=2 P.T=7

E2 = (2) mod 11 = 8 mod 11 = (8)

C. T= (5,6) {P.T= (6x(53)) mod 11 1 = 1 (53) mod 11 + 3 Decryption Encryption.

(6x3) mod 11 = 18 mod 11