Chap 2: Single-Variable Optimization

2.1 * Definition and Optimality Conditions

— Higher Order Analysis

2.2 # Bracketing: f(x1)>f(x2) < f(x3) for x1< x1< x3

2.3 & Gradient Based Methods

- Newton, Secant, Cubic and 226

Quadratic Estimation

2.4 Region Elimination Methods

13-15

- Bisection and Interval Halving
- Fibonacci Search Method
- Golden Section Search Method

2.5 * TWO Contexts of Univariate Optimization

2.6 × Line Search

The complete problem itself

The line-search "sub-problem"

- Exact line search (Accurate)

Inexact line search (Inaccurate)

- Termination conditions

2.4 Region Elimination Methods

& Bisection and Interval Halving

Interval: [a,b]

$$\begin{cases}
\text{If } f(x_1) < f(x_m) & a x_1 x_m x_2 b \\
b \leftarrow x_m; \text{ break;} \end{cases}
\end{cases}$$
If $f(x_2) < f(x_m)$

$$\begin{cases}
a \leftarrow x_m; \text{ break;} \end{cases}$$

$$a \leftarrow x_1; b \leftarrow x_2; \end{cases}$$

Two new evaluations in every iteration.

Interval reduces to HALF in every iteration.

Other (better) possibilities?

Fibonacci Series

K	0	-	2	3	4	5	6	7	8	9
Fx	1	1	2	3	5	8	13	21	34	55
$F_j = F_{j \ge 1} + F_{j-1}$										
6 5-										
	6	5								
	4	3/2-	1~	1	9.01					

2.4 Region Elimination (Contd)

* Fibonacci Search Method $F_0 = 1$, $F_1 = 1$, $F_j = F_{j-2} + F_{j-1}$ FN = FN-2 + FN-1

Starting: A unimodal bracket [a,b] of length $L_1 = b - a$.

I teration 1: k=2

"Evaluate function of at $x_1 = b - L_2$ and $x_2 = a + L_2$ where $L_2 = \frac{F_{N-1}}{F_N} L_1$ / Of f(x1)>f(x2) Else [i.e. f(n2) > f(n1)]

Iteration 2: k=3

[with new [a,b], evaluate at 21 = b-13 and 21 = a+13 where $L_3 = \frac{F_{N-2}}{F_N} L_1 = \frac{F_{N-2}}{F_N} L_2$ Cleck: a+L3 = b-L1 + FN-2 L1 $= b - \frac{F_N - F_{N-2}}{C} L_1 = b - \frac{F_{N-1}}{F_N} L_1 = b - L_2$ I teration (k-1): Lk = FN-K+1 L1 Iteration (N-1): $L_N = \frac{F_1}{F_N}L_1 = \frac{L_1}{F_1} = E$ (tolerance)

2.4 Region Elimination (contd)

Fibonacci Search: Algorithmic Steps
Given: Function f
Unimodal bracket [a,b] of length 4
Tolerance E

From $L_1=b-a$, find $F_N > L_1/\epsilon & N$.

Evaluate f at $x_1=b-L_2$ & $x_2=a+L_2$ # Set k=2

of f(x1) > f(x2) {a → x1; x1 → x2; x2 → a + L x+1; Evaluate f(x2)}

Else { b+x2; x2+x1; x1+b-Lk+1; Evaluate f(x1) }

If k=N; STOP Update k+k+1

Golden Section Search

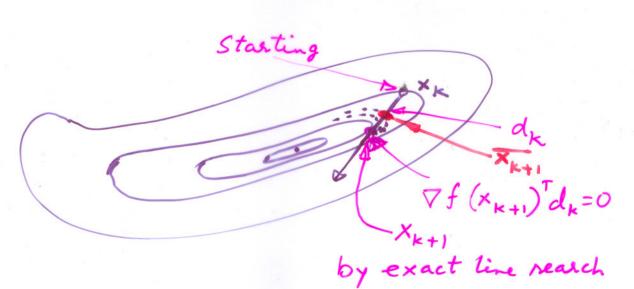
Explore: With N-00

 $\frac{F_{N-1}}{F_{N}} \approx \frac{F_{N-2}}{F_{N-1}} \approx \tau \Rightarrow \frac{F_{N-2}}{F_{N}} = \tau^{2}$ $\Rightarrow \tau^{2} + \tau = 1 \Rightarrow \tau \approx 0.618$ $\equiv \tau^{N-1}(b-a) \Rightarrow \tau^{N-1} = \frac{\epsilon}{b-a}$ $\Rightarrow N = \frac{\log(\frac{\epsilon}{b-a})}{\log \tau} + 1$ $N \text{ actually not } \propto 1$

ME 752 2.6 Line Search

For f(x), $x \in \mathbb{R}^n$,

after choosing d_k from x_k , $\phi(\alpha) = f(x) = f(x_k + \alpha d_k)$ Let ringle variable problem $\Rightarrow \phi(\alpha) \approx f(x_k) + \alpha \left[\nabla f(x_k) \right]^T d_k$ 1st order condition: $\phi'(\alpha) = 0$ $\Rightarrow \left[\nabla f(x_k + \alpha d_k) \right]^T d_k = 0$



or Pf(xx+1) dx = 0

From XK+1, a new line search will start.

- ? How much effort was spent in arriving at x_{k+1} exactly?
- ? How does it matter if the new line search starts at Xxx, instead?

2.6.3 Inexact Line Search (Try enough, but not too much)

Notion of sufficient decrease:

 $\phi(\alpha) \leq f(x_k) + \mu \alpha g_k^T d_k$, $\phi(\alpha) = f(x_k) + \mu \alpha g_k^T d_k$ $O \leq \mu \leq 1$ Aromigo's rule

Not sufficient alone cunvature condition $\phi(\alpha) = f(x_k) + \alpha g_k^T d_k$

Notion of "NOT much hope" further: $\phi'(\alpha_k) \gg \eta \phi'(0)$, $\eta < 1$ or, $[\nabla f(x_k + a_k d_k)]^T d_k > \eta g_k^T d_k$,

Curvature Condition

Caution: Inexact line search would NOT guarantee $d_k^T g_{k+1} = 0$, only exact line search will.

12