

Programming Languages Using PLCC

What a programming course that uses PLCC might look like,
and what's possible

Common Approaches to Programming Languages

- Tour of languages:
 - all breadth, no depth
 - Language *du jour* changes (but the fundamentals don't)
- Tour of language constructs (with example languages):
 - A little more depth, more focus on language fundamentals
- Build compiler/interpreter (assembly and/or a simulated machine):
 - Requires more details than is necessary, which means less time for the fundamentals
- Build language implementations using industrial strength tools
lex, yacc, javacc, antlr
 - Generated code is complicated, so tools must be used without understanding

PLCC Approach

- Design and implement languages using PLCC
 - Tokens are defined using regular expressions
 - Syntactic rules defined using BNF
 - Generated parser is LL(1)
 - Comprehensible Left-Recursive parser with a 1-token look ahead.
 - Token matching is equally explainable: longest, first match.
 - Java is used for defining semantics
- For each construct
 - Syntax defined in PLCC's grammar language (see above)
 - Semantics is defined in BNF
 - Semantics defined in Java

Students gain a deep understanding of languages without having to dive into assembly or work with tools that hide too much details.

Learning Goals

- Know some of the history of languages (what's been tried).
- Learn and evaluate new language constructs more easily.
- Become familiar with some of the common tools for language design and implementation.
- Become better developers through a deeper understanding of programming languages.

Learning Objectives

- Differentiate between lexical, syntactic, and semantic analysis.
- Use regex to define tokens.
- Use BNF to define small language constructs.
- Use/implement operations that evaluate across a parse tree
- Draw environment diagrams that accurately illustrate the environment that a given expression will evaluate within.
- Accurately evaluate expressions containing nested static scopes, closures, recursion, selection, and polymorphism, all within the context of different call-semantics.
- Identify and explain the risks of side effecting code and how to mitigate them.

Context

- Semester: 15 week course (+1 for finals)
- 2-3 sessions per week, at least 150 minutes per week
- 3 exams, each focused on material since the last exam
- 6-12 homeworks, depending on size
- Upper-level in CS

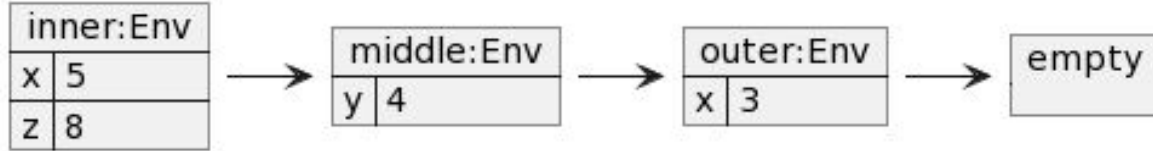
Schedule

- Weeks 1-2: Foundations
 - Lexical Analysis
 - Syntactic Analysis
 - Semantic Analysis
 - Environments
- Weeks 3-6: Build a functional language
 - Literals, symbols, let-expressions
 - Static scoping
 - If expression
 - Functions
 - Higher-ordered
 - Closures
 - Static scoping
 - Letrec
- Weeks 7-10: Call semantics
 - Side-effects
 - Pass by value
 - Pass by reference
 - Pass by name
 - Pass by need
- Week 10-11: Static Type System
- Week 11-13: Object-Oriented Language
- Week 14-15: Topics
 - Properties
 - Concurrency
 - Logic-Based Language
 - etc.

Weeks 1-2: Lexical, Syntactic, and Semantic Analysis

- Lexical analysis
 - Tokens
 - Regex
 - Longest, first match
- Syntactic analysis
 - BNF
 - Repetition Rule
 - LL(1) Parser
- Semantic analysis
 - Parse tree
 - Java snippets to walk tree

Weeks 1-2: Environments



Env.apply() semantics:

- `inner.apply("x");` // 5
- `inner.apply("y");` // 4
- `inner.apply("z");` // 8
- `middle.apply("x");` // 3
- `middle.apply("y");` // 4
- `middle.apply("z");` // Exception

Weeks 3-6: Define functional language: V0-V1

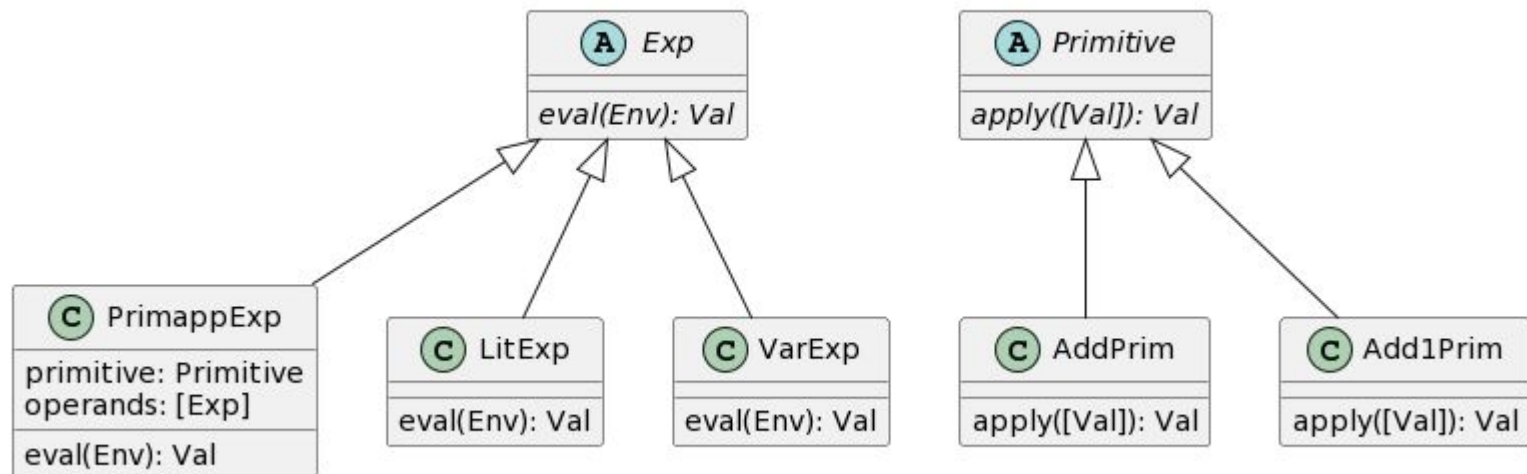
`<prog> ::= <exp>`

`<exp>:VarExp ::= <VAR>`

`<exp>:LitExp ::= <INT>`

`<exp>:PrimappExp ::= <prim> LPAREN <operands> RPAREN`

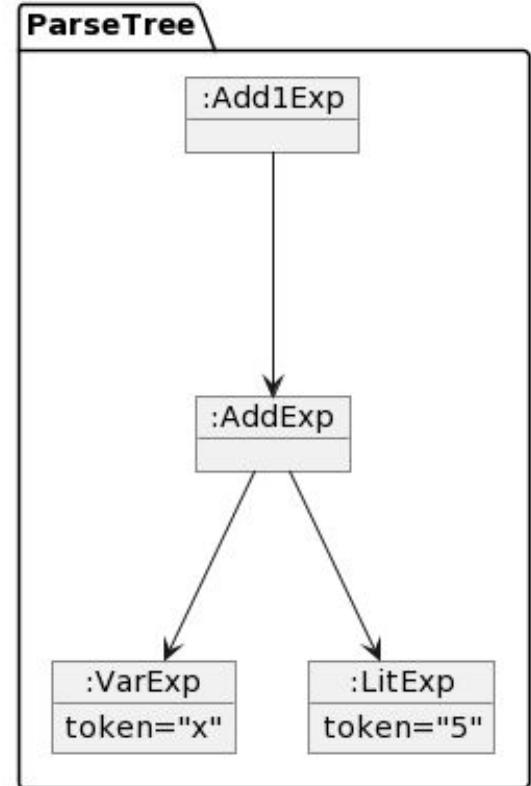
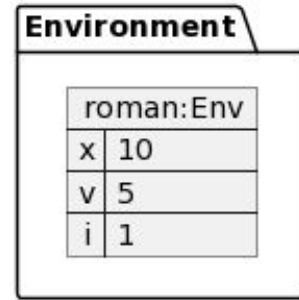
Classes



Weeks 3-6: Define functional language: V0-V1

`add1(+(x,5)) % => 15`

Note that, during semantic analysis, code in the parse tree is executed that searches the environment (not in parse tree) for the value of **x**.

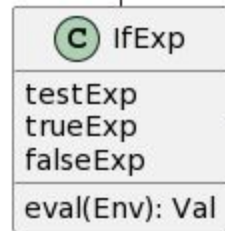
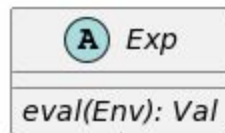


Weeks 3-6: Define functional language: V2

`<exp>:IfExp ::=`

`IF <exp>testExp THEN <exp>trueExp ELSE <exp>>falseExp`

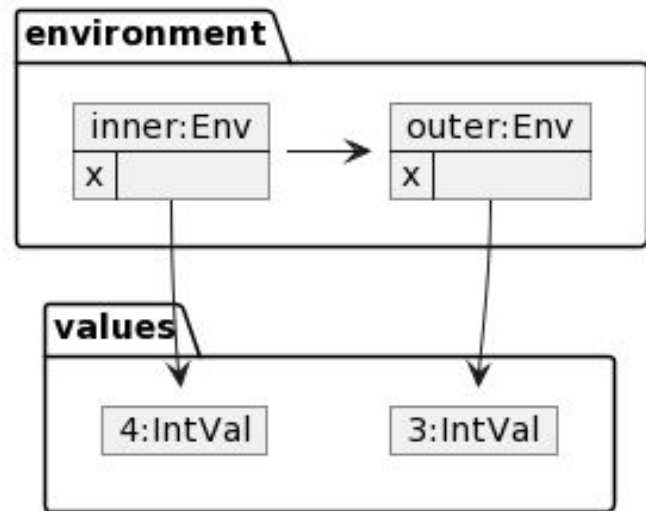
```
if 1 then 3 else 4    % => 3
if 0 then 3 else 4    % => 4
```



```
public Val eval(Env env) {
    Val testVal = testExp.eval(env);
    if (testVal.isTrue())
        return trueExp.eval(env);
    else
        return falseExp.eval(env);
}
```

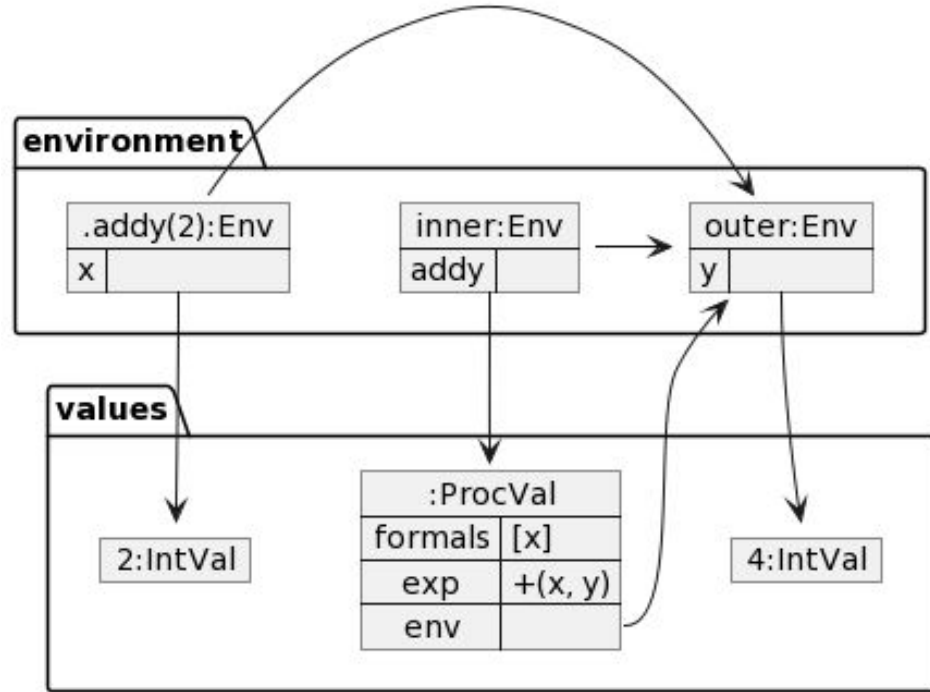
Weeks 3-6: Define functional language: V3

```
let      % outer
  x = 3
in
  let    % inner
    x = add1(x)
  in
    +(x, x) % => ?
```



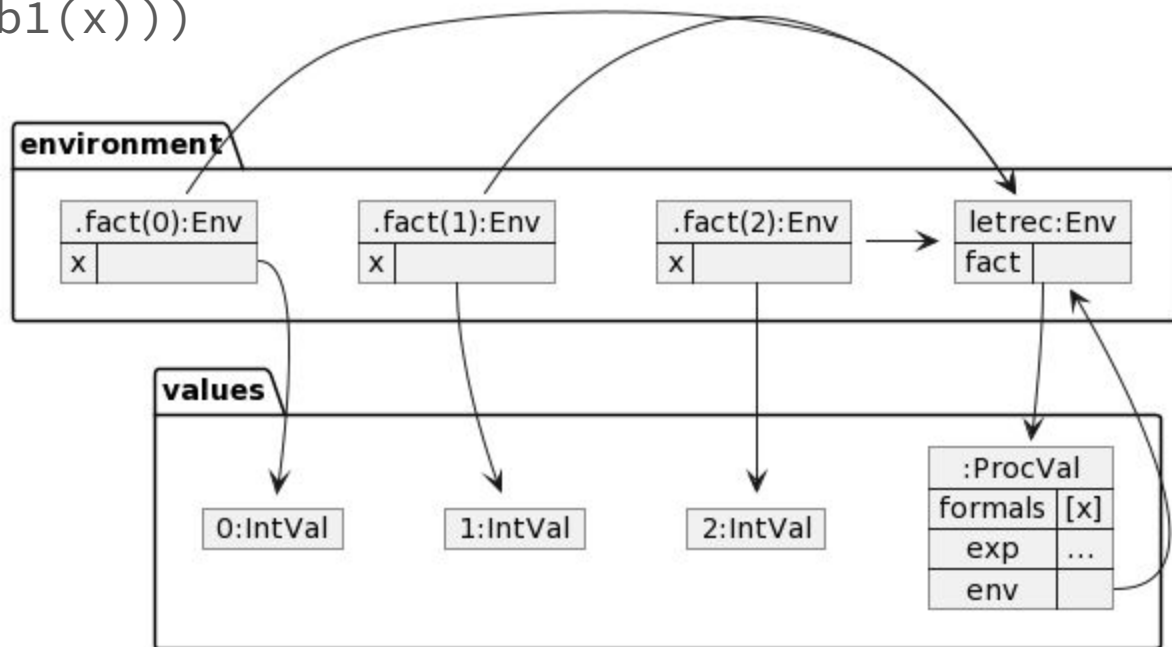
Weeks 3-6: Define functional language: V4

```
let
  y = 4
in
  let
    addy = proc(x) +(x, y)
  in
    . addy(2)      % => 6
```



Weeks 3-6: Define functional language: V5-V6

```
letrec
  fact = proc(x)
    if zero?(x) then 1
    else *(x, .fact(sub1(x)))
in
  .fact(2)
```



V0-V6: Summary

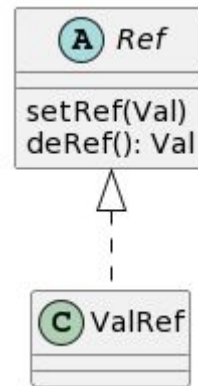
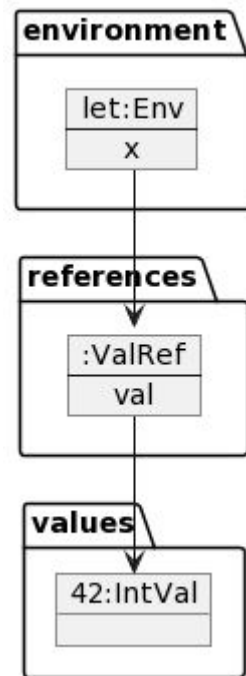
- Functional, expression language
- Closures
- Static-scoping
- Supports recursion
- Selection through if expressions
- Integer primitives
- Boolean values defined as 0 is False, non-0 is True
- Two types: integers and procs

Students practice understanding language concepts in terms of regex, BNF, semantics in Java, parse trees, and environment diagrams.

Weeks 7-10: Call Semantics: SET

```
let
  x = 42
in {
  set x = +(x, 2) ;
  x
}
```

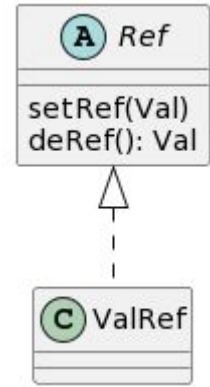
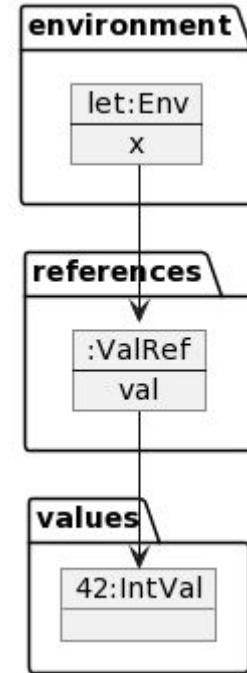
Symbols are now bound to **mutable** Refs instead of Vals.
Enables side effects.



Weeks 7-10: Call Semantics: SET

```
let
  x = 42
in {
  set x = +(x, 2) ;
  x
}
```

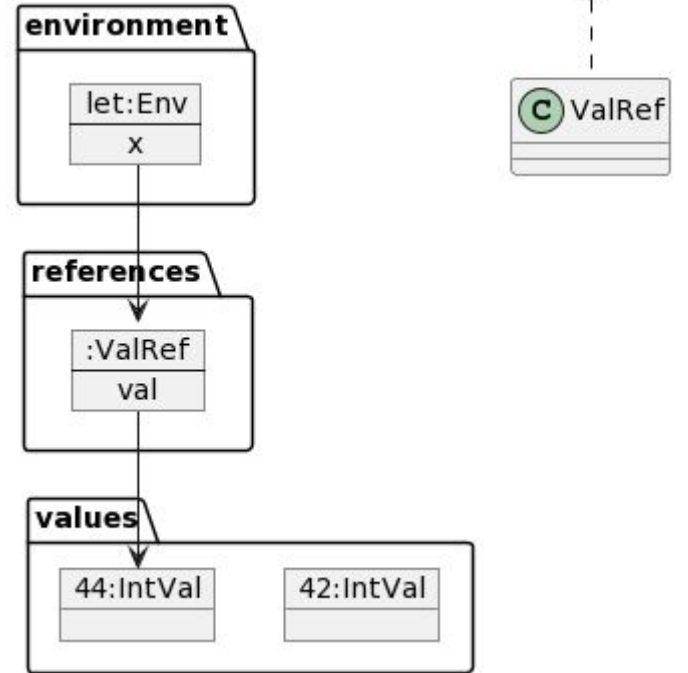
$xRef.deRef()$



Weeks 7-10: Call Semantics: SET

```
let
  x = 42
in {
  set x = +(x, 2) ;
  x
}
```

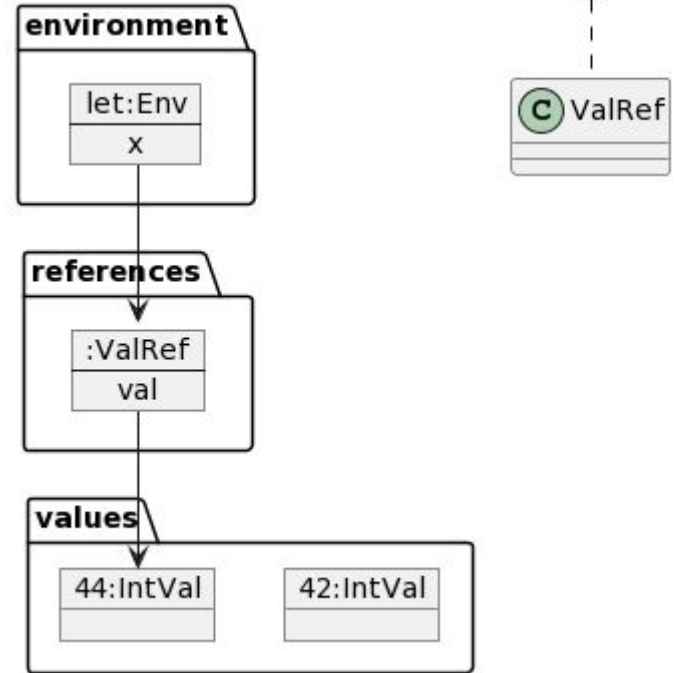
xRef.setRef(val);



Weeks 7-10: Call Semantics: SET

```
let
  x = 42
in {
  set x = +(x, 2) ;
  x
}
```

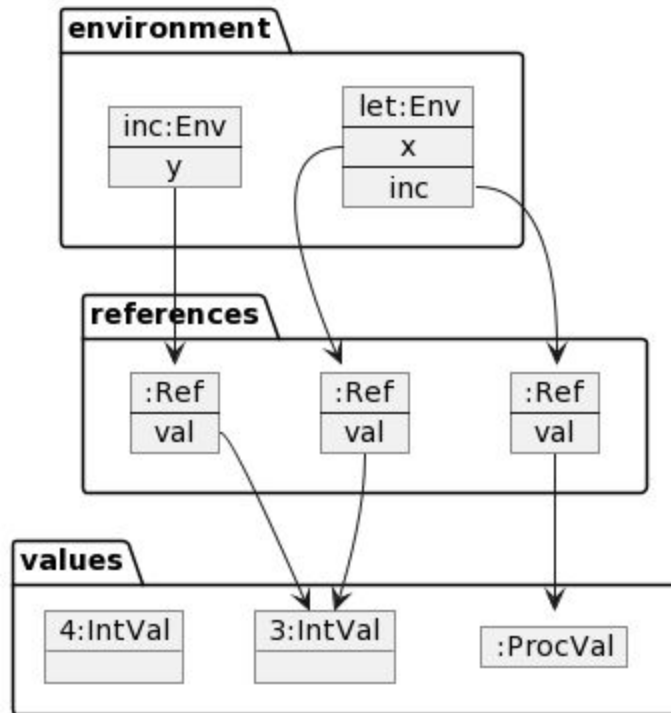
xRef.deRef()



Weeks 7-10: SET - Pass-by-Value

```
let
  x = 3
  inc = proc(y) set y = add1(y)
in {
  . inc(x) ;
  x % => ??
}
```

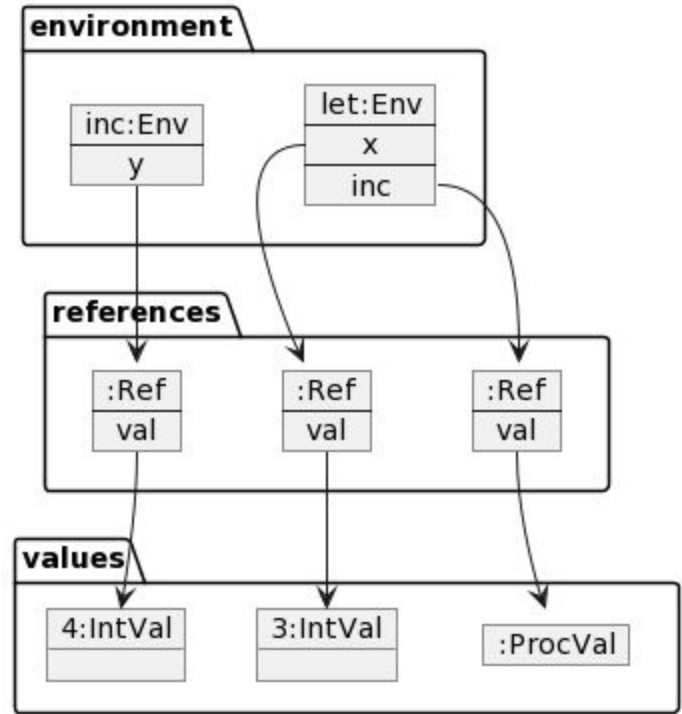
Just after the call, **before** inc's body evals.



Weeks 7-10: SET - Pass-by-Value

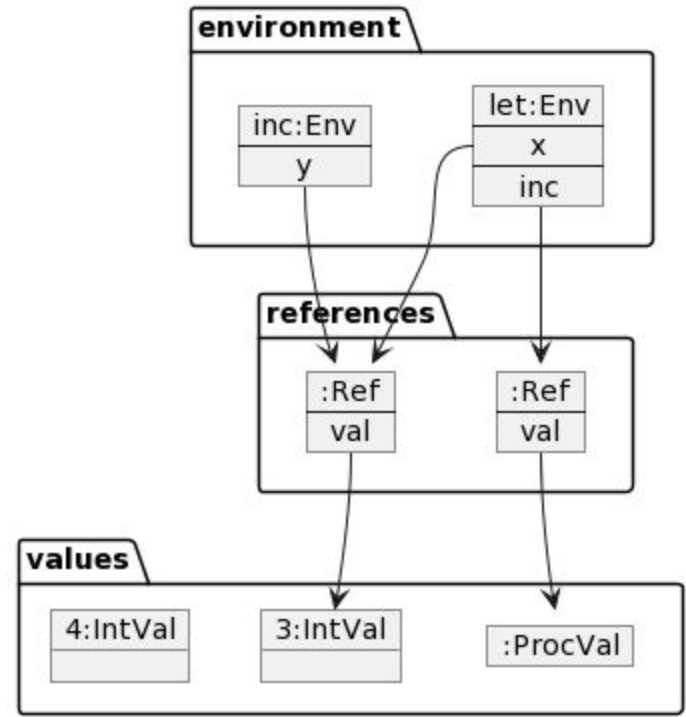
```
let
  x = 3
  inc = proc(y) set y = add1(y)
in {
  . inc(x) ;
  x % => 3
}
```

After .inc(x)



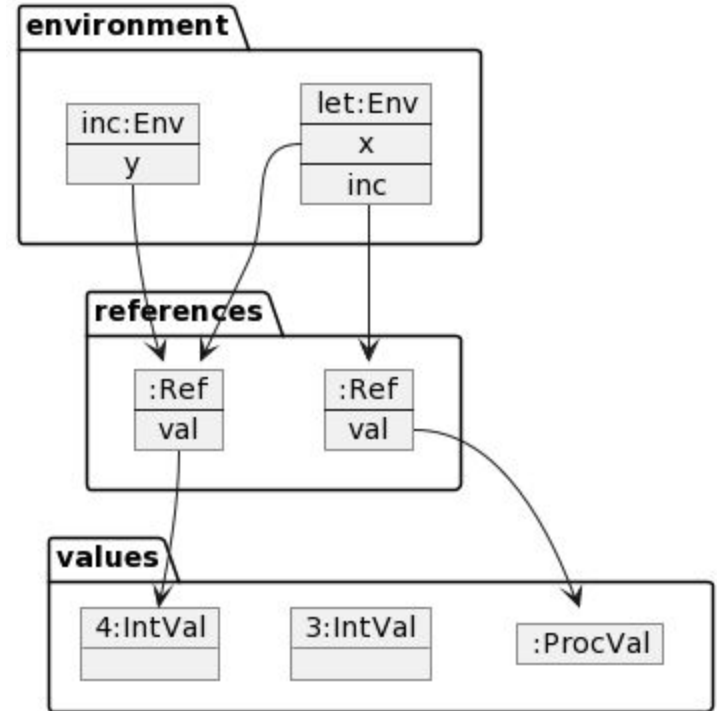
Weeks 7-10: REF - Pass-by-Reference

```
let
  x = 3
  inc = proc(y) set y = add1(y)
in {
  . inc(x) ;
  x % => ??
}
```



Weeks 7-10: REF - Pass-by-Reference

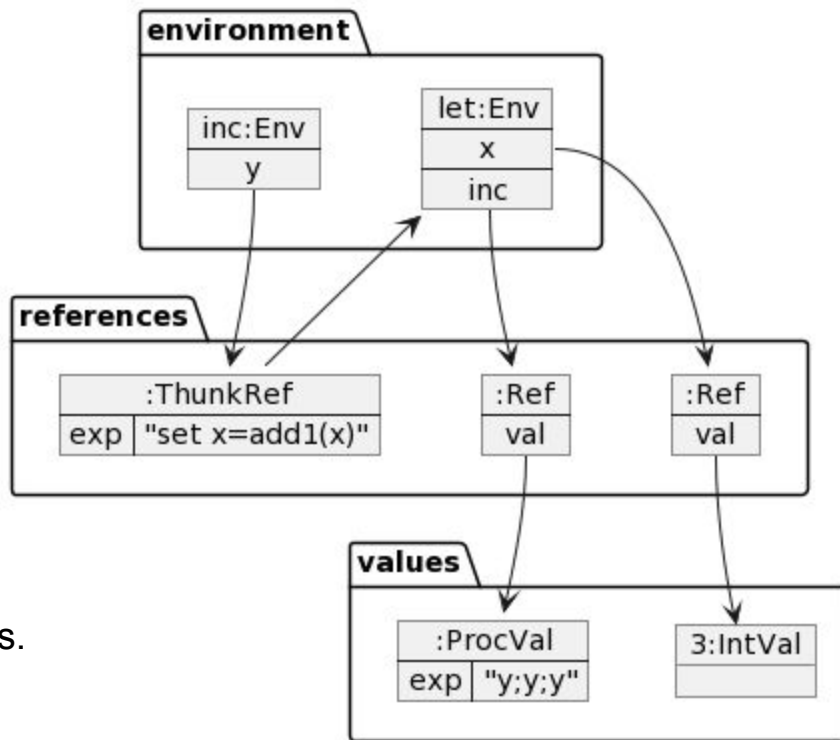
```
let
  x = 3
  inc = proc(y) set y = add1(y)
in {
  . inc(x) ;
  x % => 4
}
```



Weeks 7-10: NAME - Pass-by-Name (thunk)

```
let
  x = 3
  inc = proc(y) { y; y; y }
in {
  . inc(set x = add1(x)) ;
  x % => ??
}
```

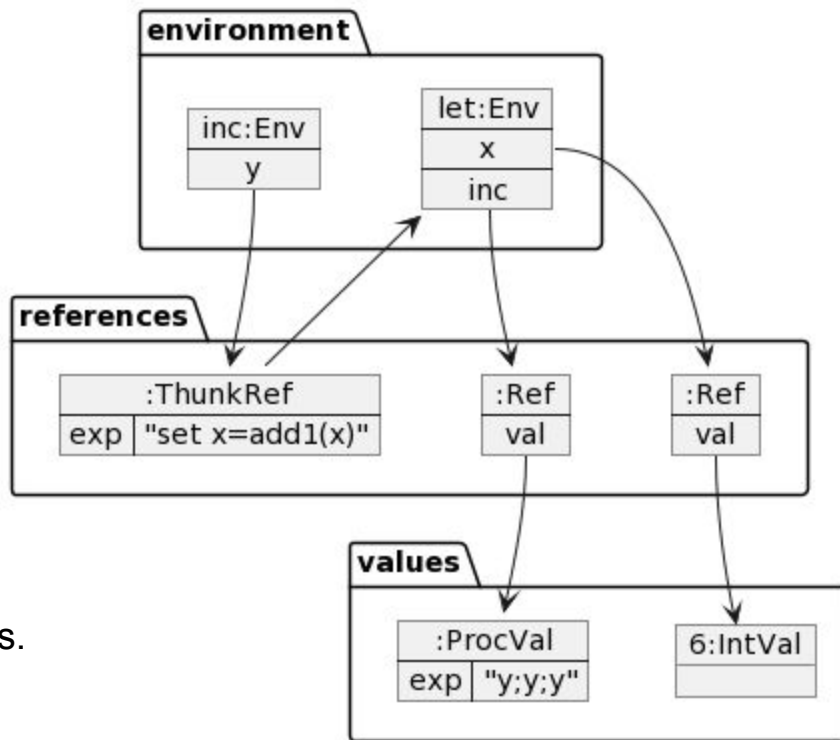
A thunk is similar to a proc without formals.
But it's a Ref.
thunk.deRef() evals its exp in calling env.



Weeks 7-10: NAME - Pass-by-Name (thunk)

```
let
  x = 3
  inc = proc(y) { y; y; y }
in {
  . inc(set x = add1(x)) ;
  x % => 6
}
```

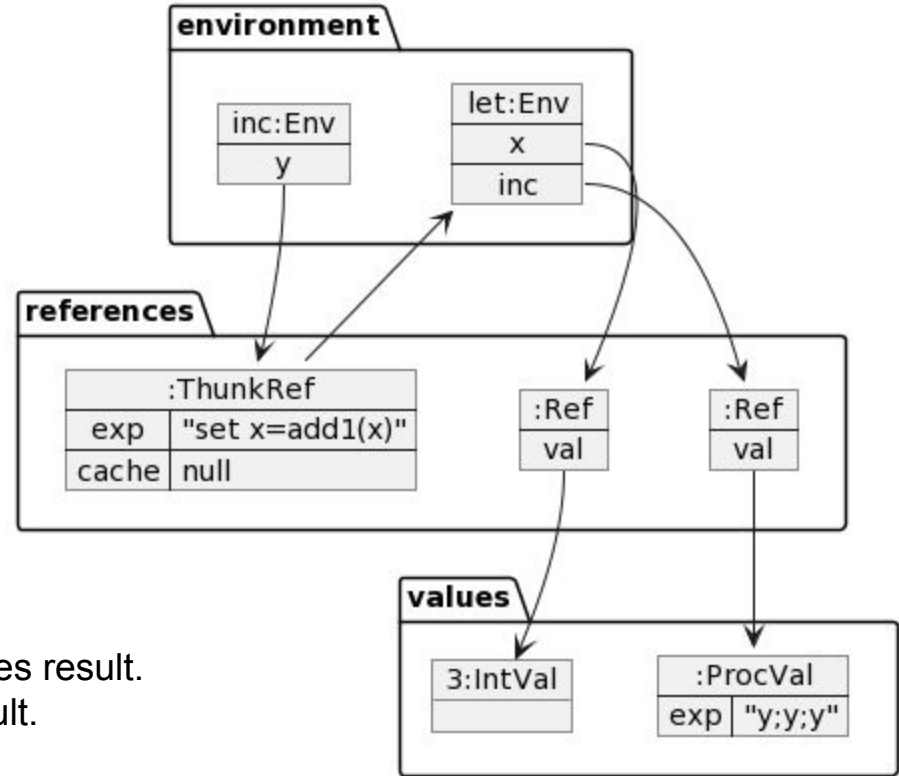
A thunk is similar to a proc without formals.
But it's a Ref.
thunk.deRef() evals its exp in calling env.



Weeks 7-10: NEED - Pass-by-Need

```
let
  x = 3
  inc = proc(y) { y; y; y }
in {
  . inc(set x = add1(x)) ;
  x % => ??
}
```

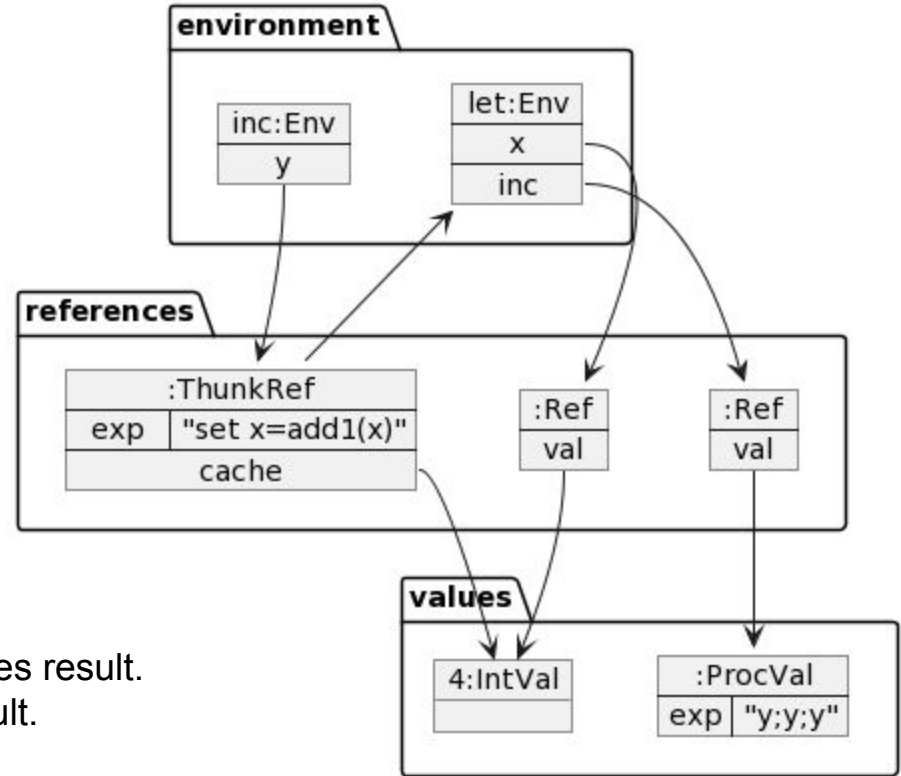
On first deRef(), think evaluates exp and caches result.
Subsequent deRef(), think returns cached result.
Lazy evaluation.



Weeks 7-10: NEED - Pass-by-Need

```
let
  x = 3
  inc = proc(y) { y; y; y }
in {
  . inc(set x = add1(x)) ;
  x % => 4
}
```

On first deRef(), think evaluates exp and caches result.
Subsequent deRef(), think returns cached result.
Lazy evaluation.



Weeks 10-11: TYPE0 and TYPE1: Static Type Checking

letrec

```
fact = proc(x: int): int  
  if zero?(x) then 1  
  else *(x, .fact(sub1(x)))
```

in

```
.fact(5)
```

let

```
twice = proc(f: [int => int], x: int): int {  
  f(f(x))  
}
```

```
add5 = proc(x: int): int +(5, x)
```

in

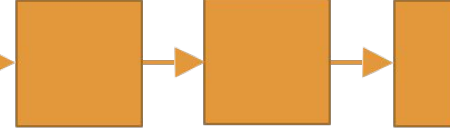
```
twice(add5, 10) % => 20
```

Type Checking Design

Interpreting an
Expression

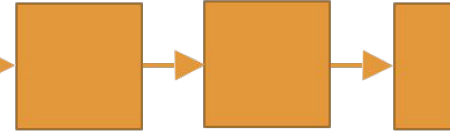
Halts if (1) fails.

1. Evaluate the
expression's type



type bindings

2. Compute the
expression's value



value bindings

4/30/21

Weeks 11-14: OBJ (shown) and PROP (not shown)

```
define shape = class
  method area = proc() -1
end
```

```
define rectangle = class extends shape
  field lenn % length
  field widd % width
```

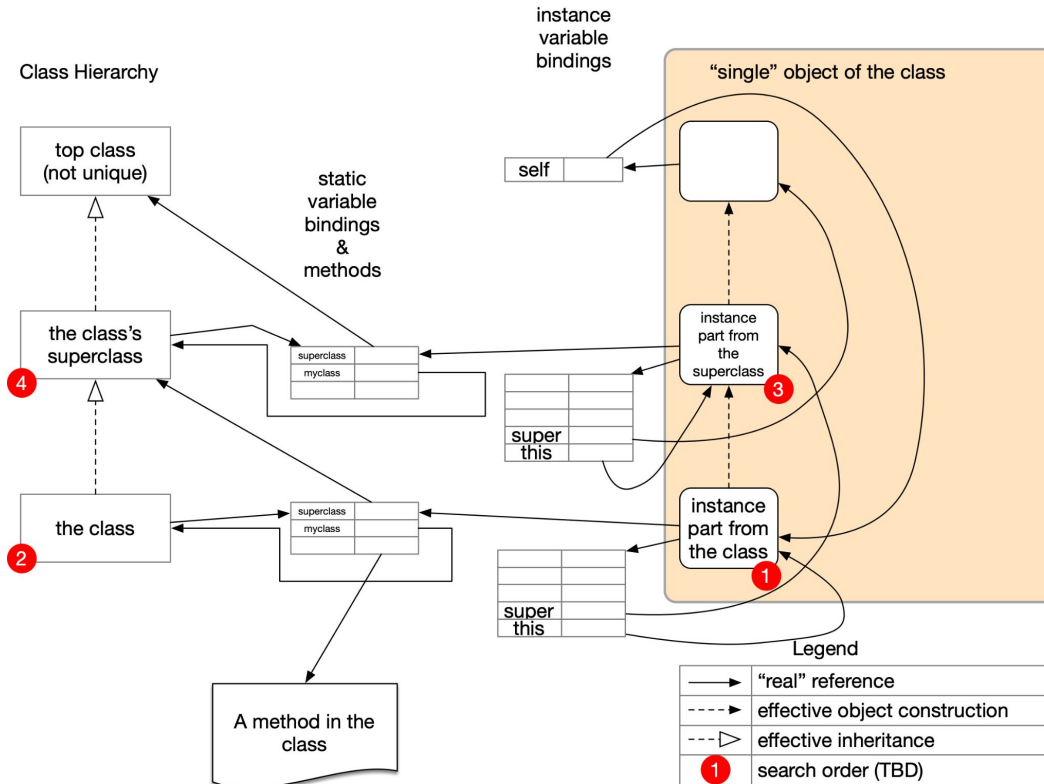
```
  method init = proc(lenn,widd) {set <self>lenn=lenn ; set <self>widd=widd ; self}
  method area = proc() *(lenn,widd)
end
```

```
define s = new shape
define r = .<new rectangle>init(4,5)
.<r>area() % => 20
.<s>area() % => -1
```

What is an Object?
An Environment!!!!

Weeks 11-14: OBJ: Making objects

"Environments, Environments, Environments"



ABC: Logic-Based-Language

PLCC implementation of ABCDatalog <<http://abcdatalog.seas.harvard.edu/>>
a subset of Prolog

```
bear(fuzzy).
```

```
bear(wuzzy).
```

```
bear(X)? % yields "bear(fuzzy)" and "bear(wuzzy)"
```

Course Materials

<https://github.com/ourPLCC/course/>
<https://plcc.python.net/>

Let's look at what's there...