

Merge 2 Sorted Array

Symmetry Algorithm:

- # When we have 2 sorted Array

$$\text{Arr1} = [1 \ 3 \ 4 \ 7 \ 10 \ 12]$$

$$\text{Arr2} = [2 \ 3 \ 6 \ 15]$$

Our final array will look like

$$[1 \ 2 \ 3 \ 3 \ 4 \ 6 \ 7 \ 10 \ 12 \ 15]$$

$n=10$

Now if we have to create a symmetry basically 2 halves then

$$[1 \ 2 \ 3 \ 3 \ 4] \quad [6 \ 7 \ 10 \ 12 \ 15]$$

$$1 \ 2 \ 3 \ 3 \ 4 \mid 6 \ 7 \ 10 \ 12 \ 15$$

we will have 5 elements
on left

5 element on
right

Note: If len of our imaginary final array is odd,
we can assume our left array to have one
more value than right half.

Algorithm:

Step 1: first step would be to identify how many
elements need to be on the left half of
array.

Step 2:

Since we have 2 array. Check how many can we pick from arr1 & how many

Can we pick from arr2 to form left side of the array.

Eg: Consider we pick 4 ele from arr1 & 1 elem from arr2

| | | l_1 | | r_1 | |
|-----------|--|-------|-------|-------|----|
| Arr1-left | | 1 | 3 | 4 | 7 |
| Arr2-left | | 2 | l_2 | 10 | 12 |

$Arr1-right$

| | | l_2 | | r_2 | |
|------------|--|-------|---|-------|--|
| Arr2-right | | 3 | 6 | 15 | |

In order to get a proper sorted final array our $l_1 < r_2$ and $l_2 < r_1$

In our Eg: $l_1 > r_2$, that means we should reduce l_2 count, because 7 should really be on right side.

So Ideal condition are.

perfect : $l_1 < r_2$

$l_2 < r_1$

If $l_1 > r_2$: [$right = mid - 1$]

reduce l_1 count $\rightarrow l_1 \downarrow$

Increase l_2 count $\rightarrow l_2 \uparrow$

If $l_2 > r_1$ [$left = mid + 1$]

reduce l_2 count $\rightarrow l_2 \downarrow$

Increase l_1 count $\rightarrow l_1 \uparrow$