

# Introduction to the Relational Model

7 November 2018

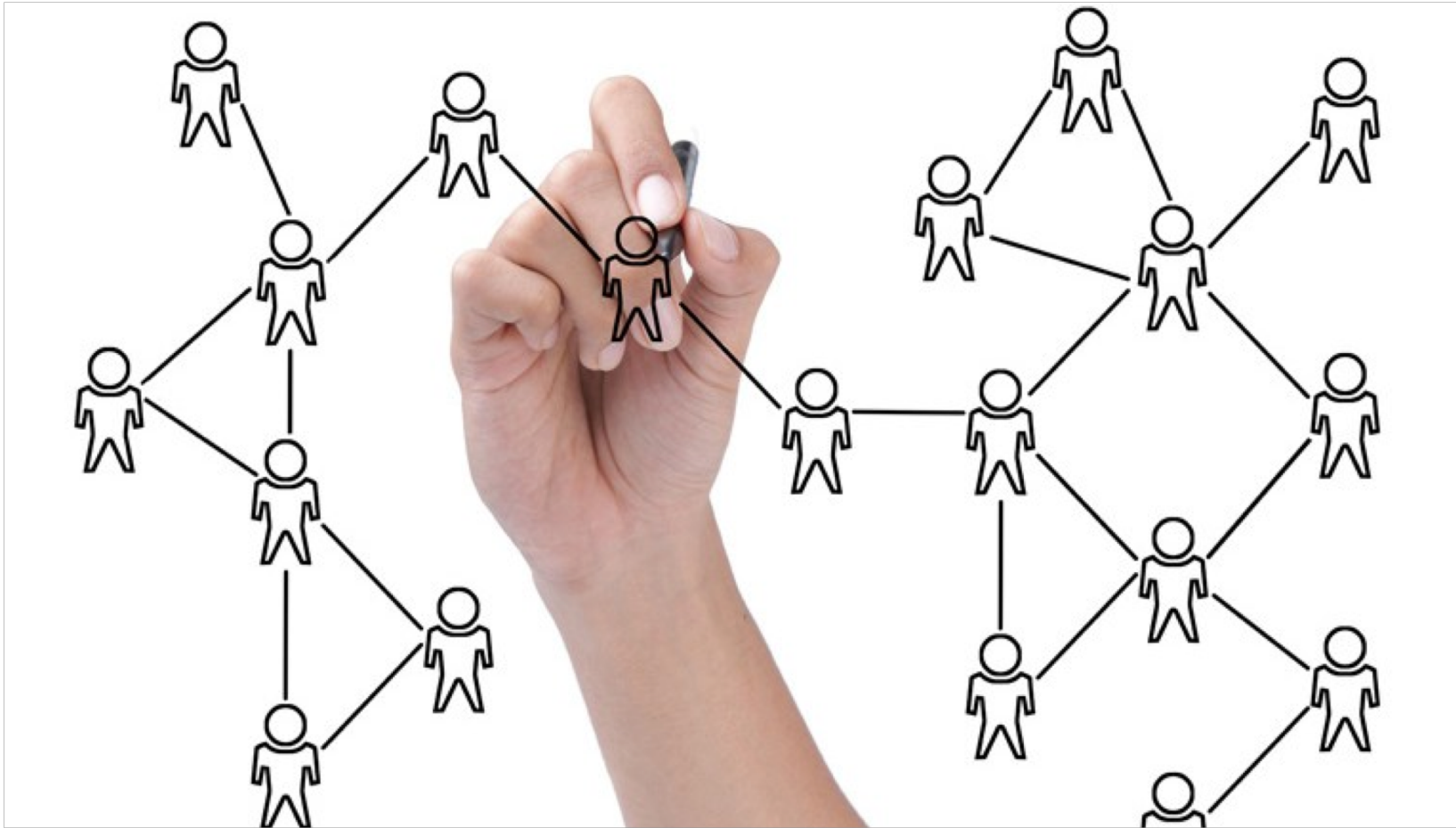
# Entities

- Something that can exist *independently* and that can be identified *uniquely*.
- A real world object such as a car or an employee.
- Can be thought of as nouns that come up during the problem description.

# Attributes

- Properties of entities are called attributes.
- Attributes represent a subgroup of information of the object represented by the entity.
- Attributes define the individual instances and help to differentiate between each instance by describing their characteristic.

# Relations



# Relational Model

<i>SID</i>	<i>Name</i>	<i>Major</i>	<i>GPA</i>
546007	Susan	CS	3.80
546100	Bob	CoE	3.65
546500	Bill	CS	3.70

***Students***

<i>CID</i>	<i>Name</i>
CS 1555	DB
CS 1530	SW
CS 1550	OS

***Courses***

<i>SID</i>	<i>CID</i>	<i>Grade</i>
546007	CS 1550	A
546007	CS 1530	B+
546100	CS 1550	B

***Enrollment***

- It is the most popular implementation model
  - Simplest, most uniform data structures, and is the most formal of all data model
- Both entity types and relationship types are represented by ***relations***, i.e., **tables**

# Relation Schema

**STUDENT**

<i>SID</i>	<i>LName</i>	<i>Name</i>	<i>Class</i>	<i>Major</i>
123	Smith	John	3	CS
395	Aiken	Mary	4	CS

Schema

## ◆ *What is the meaning?*

- A **relation schema**  $R$  specifies
  - The name of the relation
  - the attribute names  $A_i$  of  $R$
  - the domain  $D_i$  (data type + format) for each attribute  $A_i$
- data type is a set of **atomic data** values:
  - no attribute is a set-valued (**1st Normal Form**, 1-NF)
  - no attribute is composite
- format specifies the representation of a data value

# Example Table Schema

Schema of STUDENT(SID, Name, Major, GPA)

```
CREATE TABLE STUDENT  
(  SID    INTEGER,  
   Name  CHAR(20),  
   Major  CHAR(4),  
   GPA    DEC(3, 2)  
);
```

# SQLite Data Types

Each value stored in an SQLite database (or manipulated by the database engine) has one of the following storage classes:

**NULL.** The value is a NULL value.

**INTEGER.** The value is a signed integer, stored in 1, 2, 3, 4, 6, or 8 bytes depending on the magnitude of the value.

**REAL.** The value is a floating point value, stored as an 8-byte IEEE floating point number.

**TEXT.** The value is a text string, stored using the database encoding (UTF-8, UTF-16BE or UTF-16LE).

**BLOB.** The value is a blob of data, stored exactly as it was input.



# Date and Time in SQLite

SQLite does not have a storage class set aside for storing dates and/or times. Instead, the built-in date and time functions of SQLite are capable of storing dates and times as TEXT, REAL, or INTEGER values:

**TEXT** as ISO8601 strings ("YYYY-MM-DD HH:MM:SS.SSS").

**REAL** as Julian day numbers, the number of days since noon in Greenwich on November 24, 4714 B.C. according to the proleptic Gregorian calendar.

**INTEGER** as Unix Time, the number of seconds since 1970-01-01 00:00:00 UTC.

Applications can chose to store dates and times in any of these formats and freely convert between formats using the built-in functions.

# NULL in SQLite

SQLite **NULL** is the term used to represent a missing value. A NULL value in a table is a value in a field that appears to be blank.

A field with a NULL value is a field with no value. It is very important to understand that a NULL value is different than a zero value or a field that contains spaces.

We'll come back to this in our demos.

# Relational Database Schema

- ❑ A *database schema* is a set of relation schemas and a set of ***integrity constraints***

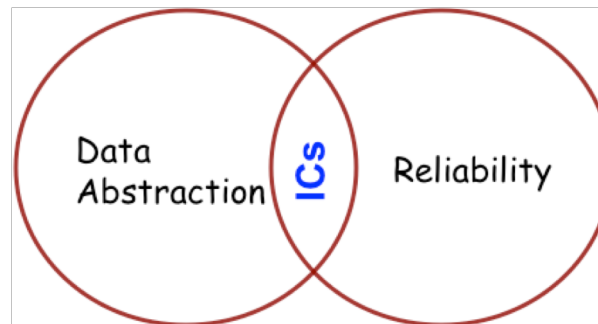


- ❑ ***Structural*** Integrity Constraints

- **key** constraints: uniqueness of keys
- **entity integrity** constraint:  
no primary key value can be **NULL**
- **referential integrity** constraint

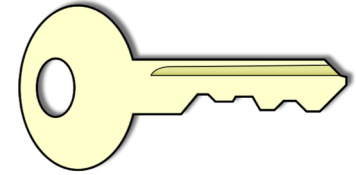
# Integrity Constraints (ICs)

- **IC**: condition that must be true for *any* instance of the database (e.g., domain constraints)
  - A **legal** instance of a relation is one that satisfies all specified ICs
  - ICs are specified when schema is **defined**
  - ICs are enforced when tables are **modified**



# Primary Key Constraint

- ❑ A set of fields is a **key** for a relation if :
  - No two distinct tuples can have same values in all key fields



- ❑ If there is more than one key for a relation:
  - Each is called a candidate key
  - One candidate key is designated as the **primary key**
  - Other candidate key(s) are designated as **alternative** or **unique key(s)**



# Example of Keys

<i>SID</i>	<i>Name</i>	<i>Login</i>	<i>Age</i>	<i>GPA</i>
546007	Jones	jones@cs	18	3.4
546100	Smith	smith@ee	18	3.2
546500	Smith	smith@math	19	3.8

- **Candidate Keys:** *SID*, and *Login*
- **Primary Key:** *SID*
- **Unique Key:** *Login*

# Example Table Schema in SQL

Schema of STUDENT(SID, Login, Name, Major, GPA)

```
CREATE TABLE STUDENT
(  SID      INTEGER NOT NULL,
   Login    CHAR(15),
   Name     CHAR(20),
   Major    CHAR(4),
   GPA      DEC(3,2),

   CONSTRAINT STUDENT_PK
       PRIMARY KEY (SID),
   CONSTRAINT STUDENT_UN
       UNIQUE (Login)           -- UNIQUE can take NULL values

);
```

# Identifying the Key

- What is the key in relation  
GRADUATE=(SID, Degree, Major, Year) ?

<i>SID</i>	<i>Degree</i>	<i>Major</i>	<i>Year</i>
123	BS	CS	1992
123	MS	CS	1993
064	BA	History	1991
445	PhD	CS	1999



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# Foreign Keys



- ❑ **Foreign key** (FK) in relation  $R_2$  is a set of attributes of  $R_2$  that forms a primary key (PK) of another relation  $R_1$ .
  - Attributes in FK and PK have the **same domain**

## STUDENT

<u>SID</u>	Degree	Major	Year
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PK

## COURSE

<u>CID</u>	Name
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PK

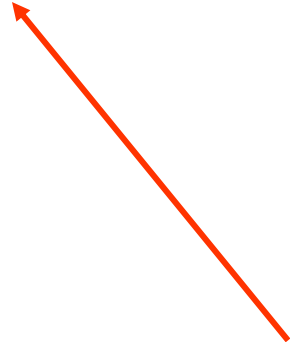
## Enrolled

<u>SID</u>	<u>CID</u>	Grade
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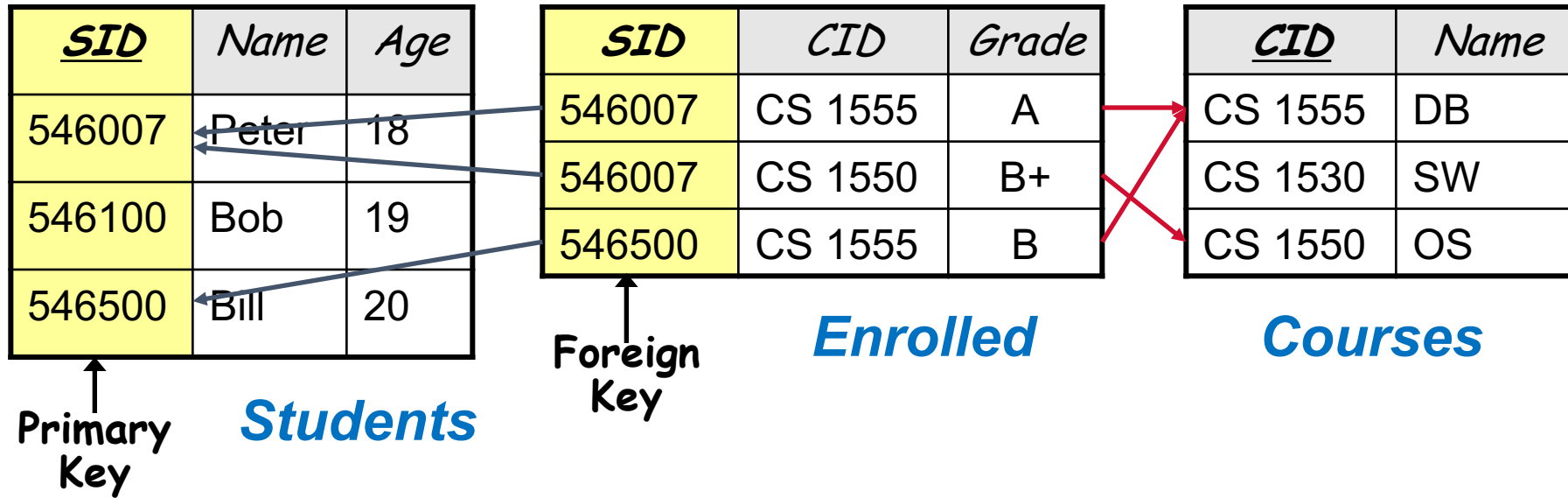
PK

FK

FK



# Foreign Key & Primary Key



- **Foreign key:** Set of fields in one relation that is used to “refer” to a tuple in another relation
  - Must correspond to primary key of the referred relation
  - E.g. *SID* is a foreign key referring to *Students*

# Foreign Key Constraints

- If foreign key constraints are enforced, **referential integrity** is achieved
  - E.g.: Only students can enroll in a class
    - Only students listed in the “Students” relation should be allowed to enroll for courses
- Like a “*logical pointer*”
  - There shouldn’ t be dangling references
    - Either valid PK or NULL

# Any Attribute can be a Foreign Key

*Faculty*

<u><i>FID</i></u>	<i>Name</i>	<i>Area</i>
007	Panos	DB
100	Daniel	OS
500	Adriana	AI

↑  
Primary  
Key

*Courses*

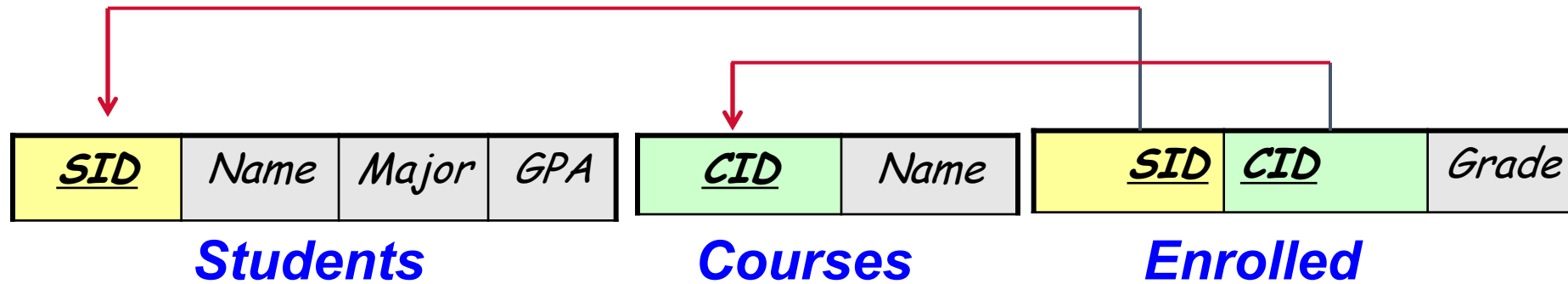
<u><i>CID</i></u>	<i>Name</i>	<i>Instructor</i>
CS 1555	DB	007
CS 1530	SW	NULL
CS 1550	OS	100

↑  
Primary  
Key

↑  
Foreign  
Key

- **Foreign key:** Set of fields in one relation that is used to “refer” to a tuple in another relation
  - Must correspond to primary key of the referred relation
  - If not part of a key, it could be NULL

# Foreign Keys in SQL



```
❑ CREATE TABLE Enrolled (  
    SID CHAR(20), CID CHAR(20), Grade CHAR(2),  
    CONSTRAINT Enrolled_PK PRIMARY KEY (SID, CID),  
    CONSTRAINT Enrolled_FK_sid  
        FOREIGN KEY (SID) REFERENCES Students (SID),  
    CONSTRAINT Enrolled_FK_cid  
        FOREIGN KEY (CID) REFERENCES Courses  
);
```

Now let's move on to using the relational model with SQLite.

Questions?