

# New bilevel programming approaches to the location of controversial facilities

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Martine Labb  <sup>1</sup>, Marina Leal<sup>2,3</sup>, Justo Puerto<sup>2,3</sup>.

<sup>1</sup> Computer Sciences Department, Universit   Libre de Bruxelles.

<sup>2</sup> Department of Statistics and Operational Research, Universidad de Sevilla.

<sup>3</sup> Institute of Mathematics of the University of Seville (IMUS).

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## 1 Introduction

## 2 Model Description

## 3 Model Formulations and Algorithm

- Examples:  $l_1$  and  $l_\infty$  norm

## 4 Computational Results

## 5 Extensions

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*“Location is an area of research aimed at determining the optimal locations of suppliers of services, facilities, structures or objects of a certain type, so that certain goals are achieved. The aims pursued depend on the location of the object may be, for example, maximizing the benefits from the operation of the facilities to be located, to minimize travel costs, maximization of the population covered by a service points or minimizing these facilities produced by unwanted or adverse effects.”*

Redloca, (Location and Related Problems Network in Spain).

# Controversy in location problems

## Controversy in location problems

Motivated by:

- Different decision-makers with distinct or opposite objectives.
- Conflict of interests between users and decision-makers.
- By the own nature of the facility or service.
- Etc.

## Controversy in location problems



Hammad, A. W., Rey, D., and Akbarnezhad, A. (2018). A Bi-level Mixed Integer Programming Model to Solve the Multi-Servicing Facility Location Problem, Minimising Negative Impacts Due to an Existing **Semi-Obnoxious** Facility. *Data and Decision Sciences in Action*, 381-395. Springer, Cham.



Heydari, R., and Melachrinoudis, E. (2012). *Location of a semi-obnoxious facility with elliptic maximin and network minisum objectives*. European journal of operational research, 223(2), 452-460.



Yapicioglu, H., Smith, A. E., and Dozier, G. (2007). *Solving the semi-desirable facility location problem using bi-objective particle swarm*. European Journal of Operational Research, 177(2), 733-749.



Melachrinoudis, E., and Xanthopoulos, Z. (2003). *Semi-obnoxious single facility location in Euclidean space*. Computers and Operations Research, 30(14), 2191-2209.



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## Controversy in location problems

### Semi-obnoxious facilities

useful but unwelcome facilities that produce environmental concerns. Facilities that population centres want them away, but there are some interests, or some interested people in locating them close the demand points.



Brimberg, J., and Juel, H. (1998) A bi-criteria model for locating a semi-desirable facility in the plane. European Journal of Operational Research 106 (1), 144–151.

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Akbari-Jafarabadi, M., Tavakkoli-Moghaddam, R., Mahmoodjanloo, M., and Rahimi, Y. (2017). A tri-level **r-interdiction** median model for a facility location problem under imminent attack. *Computers and Industrial Engineering*, 114, 151-165.



Aksen, D., and Aras, N. (2012). A bilevel fixed charge location model for **facilities under imminent attack**. *Computers and Operations Research*, 39(7), 1364-1381.



Church, R. L., and Scaparra, M. P. (2007). Protecting critical assets: the **r-interdiction** median problem with fortification. *Geographical Analysis*, 39(2), 129-146.



Scaparra, M. P., and Church, R. L. (2008). A bilevel mixed-integer program for **critical infrastructure** protection planning. *Computers and Operations Research*, 35(6), 1905-1923.



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These problems have been usually addressed via

- biobjective (multiobjective) approaches,
- difference of objectives,
- maximin or minmax objectives,
- bilevel optimization.

Bilevel Programming targets hierarchical optimization problems in which part of the constraints translate the fact that some of the variables constitute an optimal solution to another nested optimization problem.

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$$\begin{aligned} \max \quad & f(x, y) && \textit{leader objective function} \\ \text{st.} \quad & (x, y) \in X && \textit{leader constraints} \end{aligned}$$

$$\begin{aligned} y \in \arg \max_y g(x, y) \quad & \textit{follower objective function} \\ \text{st.} (x, y) \in Y \quad & \textit{follower constraints} \end{aligned}$$

First mathematical model by Bracken and McGill (1972).

## Model motivation



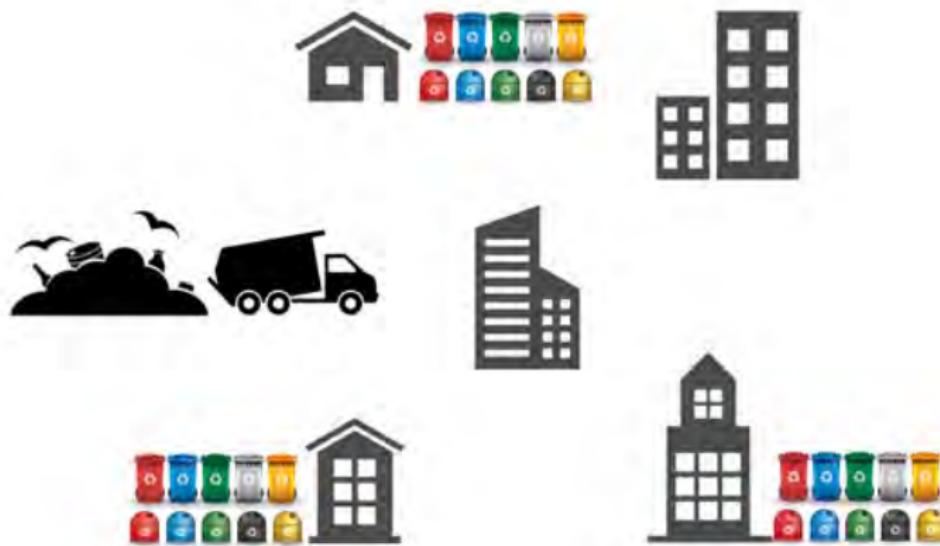
## Model motivation



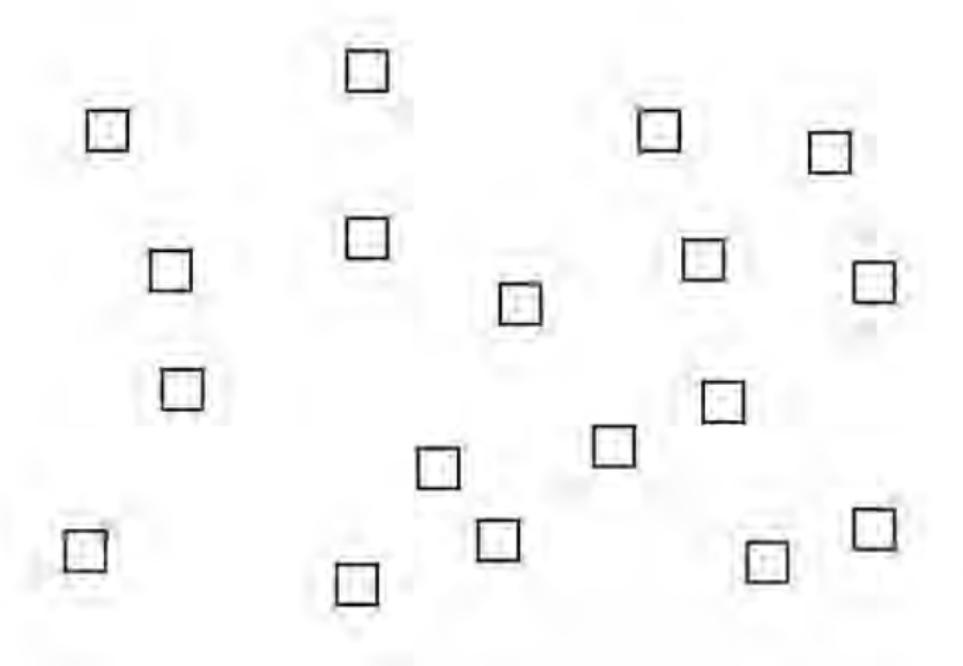
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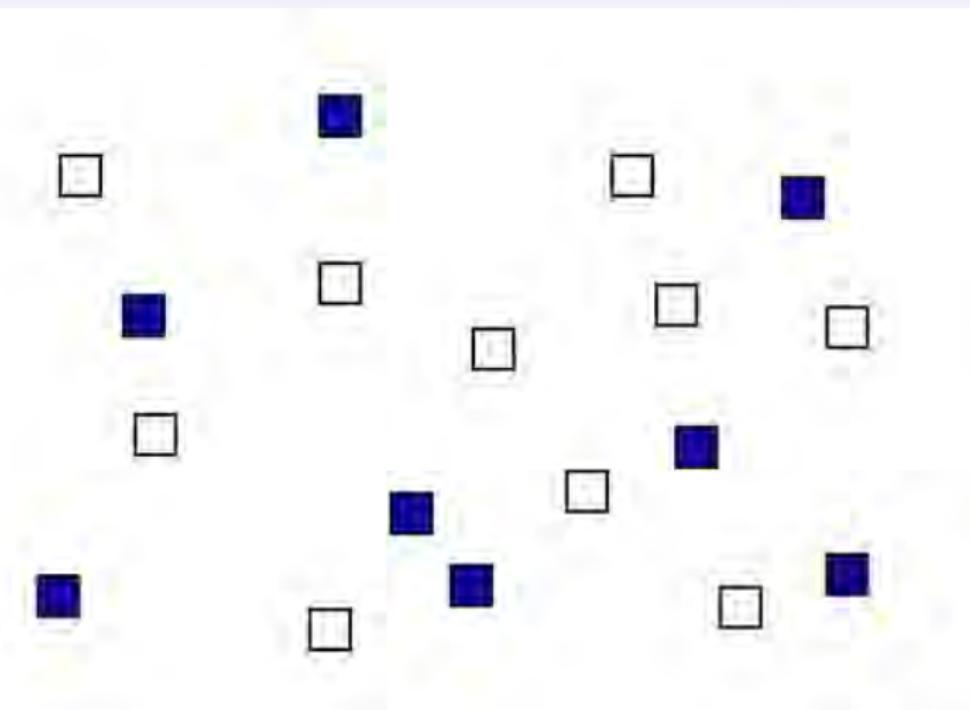
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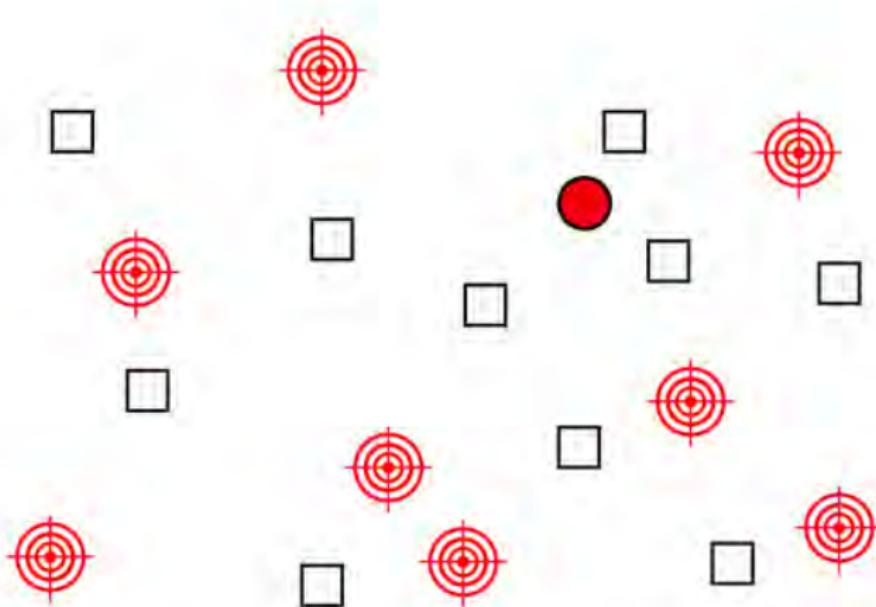
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## The model

Location entities

Hierarchy

Aim

## The model

### Location entities

- Primary facilities  
(critical infrastructures, goods to protect, demand points,...)

### Hierarchy

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### Hierarchy

- **Secondary facilities**

(attackers, thefts,  
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- **First**, from a discrete set of possible location

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### Hierarchy

- **First**, from a discrete set of possible location
- **Second**, in a continuous framework

### Aim

## The model

<u>Location entities</u>	<u>Hierarchy</u>	<u>Aim</u>
<ul style="list-style-type: none"><li>• <b>Primary facilities</b> (critical infrastructures, goods to protect, demand points,...)</li></ul>	<ul style="list-style-type: none"><li>• <b>First</b>, from a discrete set of possible location</li></ul>	<ul style="list-style-type: none"><li>• <b>Far</b> from secondary facilities</li></ul>
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# The Model

**Upper level:**

**Lower level:**

## The Model

### Upper level:

- $B$  the set of indexes of potential locations.
- $NB$  the set indexes of already fixed locations.
- $\forall j \in B \cup NB, f_j = (f_{j1}, \dots, f_{jn})$  the potential location, where  $n$  is the dimension of the space.
- Binary decision variables  $y_j = 1$  if  $f_j, j \in B$  is established, and  $y_j = 0$  otherwise.
- $c_j$  the cost of setting the primary point  $j$ , for all  $j \in B$ , and  $C$  the maximum budget.

### Lower level:

## The Model

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### Lower level:

- Decision variable  $x = (x_1, \dots, x_n) \in \mathbb{R}^n$  that determines the location of the secondary point

## The Model

$$\max \sum_{j \in B} w_j d_p(x, f_j) y_j + \sum_{j \in NB} w_j d_p(x, f_j) \quad (\text{BLM})$$

st.

$$\sum_{j \in B} c_j y_j \leq C,$$

$$y_j \in \{0, 1\} \quad \forall j \in B,$$

$$x \in \operatorname{argmin}_x \sum_{j \in B} w_j d_p(x, f_j) y_j + \sum_{j \in NB} w_j d_p(x, f_j),$$

where  $w_p \in \mathbb{R}$  and  $d_p(x, f_j)$  denotes any distance induced by a norm:

$$d_p(x, f_j) = \|x - f_j\|_p.$$

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### Theorem

*The BLM is NP-hard.*

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- $\text{ext}(P)$  is the set of extreme points, and
- $P^\circ$  is the polar set of  $P$ ,

$$P^\circ = \{x \in \mathbb{R}^n : \langle x, p \rangle \leq 1, \forall p \in P\}.$$

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## The model

### Norm induced by $P$

$$\|x\|_P = \min \sum_{u \in \text{ext}(P)} \mu_u,$$

$$st. x = \sum_{u \in \text{ext}(P)} \mu_u u,$$

$$\mu_u \geq 0, \quad u \in \text{ext}(P).$$

## The model

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### Norm induced by $P^\circ$

$$\|x\|_P = \max_{u \in \text{ext}(P^\circ)} \langle x, u \rangle.$$



Rockafellar, R. T. (1970). *Convex analysis*. Princeton university press.

## Formulation 1 (Norm induced by $P$ )

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$$\|x\|_p = \min r$$

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$$\max \sum_{j \in B} w_j y_j r_j + \sum_{j \in NB} w_j r_j$$

$$\text{s.t. } \sum_{j \in B} c_j y_j \leq C,$$

$$y_j \in \{0, 1\} \quad j \in B,$$

$$r \in \operatorname{argmin}_{x, r} \sum_{j \in B} w_j y_j r_j + \sum_{j \in NB} w_j r_j,$$

$$x_i \in \mathbb{R}, \quad i = 1, \dots, n,$$

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Discrete location leader problem  
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Dual:

$$\max \sum_{i=1}^n \sum_{j \in B \cup NB} \beta_{ji} f_{ji}$$

$$\text{s.t. } \alpha_j \leq w_j y_j, \quad j \in B,$$

$$\alpha_j \leq w_j, \quad j \in NB,$$

$$\sum_{j \in B \cup NB} \beta_{ji} = 0, \quad i = 1, \dots, n,$$

$$-\alpha_j - \sum_{i=1}^n u_i \beta_{ji} \leq 0, \quad j \in B \cup NB,$$

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Using strong duality

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$$\alpha_j \leq w_j y_j, \quad j \in B,$$

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$$\sum_{j \in B \cup NB} \beta_{ji} = 0, \quad i = 1, \dots, n,$$

$$-\alpha_j - \sum_{i=1}^n u_i \beta_{ji} \leq 0, \quad j \in B \cup NB,$$

$$u \in \text{ext}(P).$$

$$\max \sum_{j \in B} w_j y_j r_j + \sum_{j \in NB} w_j r_j$$

$$\text{s.t. } \sum_{j \in B} c_j y_j \leq C,$$

$$y_j \in \{0, 1\} \quad j \in B,$$

Linearising  $y_j r_j$ :  
MILP Formulation.

$$\sum_{j \in B} w_j y_j r_j + \sum_{j \in NB} w_j r_j = \sum_{i=1}^n \sum_{j \in B \cup NB} \beta_{ji} f_{ji}$$

$$x_i \in \mathbb{R}, \quad i = 1, \dots, n,$$

$$r_j = \sum_{u \in \text{ext}(P)} \mu_u^j, \quad j \in B \cup NB,$$

$$x_i = \sum_{u \in \text{ext}(P)} \mu_u^j u_i + f_{ji}, \quad j \in B \cup NB, \\ i = 1, \dots, n,$$

$$\mu_u^j \geq 0, \quad u \in \text{ext}(P), j \in B \cup NB, \\ r_j \geq 0, \quad j \in B \cup NB.$$

$$\alpha_j \leq w_j y_j, \quad j \in B,$$

$$\alpha_j \leq w_j, \quad j \in NB,$$

$$\sum_{j \in B \cup NB} \beta_{ji} = 0, \quad i = 1, \dots, n,$$

$$-\alpha_j - \sum_{i=1}^n u_i \beta_{ji} \leq 0, \quad j \in B \cup NB,$$

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## Formulation 2 (Norm induced by $P^0$ )

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$$\|x\|_p = \max_{u \in \text{ext}(P^0)} \langle x, u \rangle.$$

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$$\|x\|_p = \min r$$

$$\text{st. } \langle u, x \rangle \leq r, \forall u \in \text{ext}(P^0).$$

## Formulation 2 (Norm induced by $P^0$ )

$$\max \sum_{j \in B} w_j y_j r_j + \sum_{j \in NB} w_j r_j$$

$$\text{s.t. } \sum_{j \in B} c_j y_j \leq C,$$

$$y_j \in \{0, 1\} \quad j \in B,$$

$$\sum_{j \in B} w_j y_j r_j + \sum_{j \in NB} w_j r_j = \sum_{k \in \text{ext}(P_0)} \sum_{j \in B \cup NB} \left( \sum_{i=1}^n -f_{ji} u_{ki} \right) \alpha_{kj},$$

$$r_j \geq \sum_{i=1}^n u_{ki} (x_i - f_{ji}), \quad k \in \text{ext}(P^0), j \in B \cup NB,$$

$$r_j \geq 0, \quad j \in B \cup NB,$$

$$x_i \in \mathbb{R}, \quad i = 1, \dots, n,$$

$$\sum_{k \in \text{ext}(P^0)} \alpha_{kj} \leq w_j y_j, \quad j \in B,$$

$$\sum_{k \in \text{ext}(P^0)} \alpha_{kj} \leq w_j, \quad j \in NB,$$

$$\sum_{k \in \text{ext}(P^0)} \sum_{j \in B \cup NB} (-u_{ki}) \alpha_{kj} = 0, \quad 1 = 1, \dots, n,$$

$$\alpha_{kj} \geq 0 \quad k \in \text{ext}(P_0), j \in B \cup NB.$$

# Benders like algorithm

## Benders like algorithm

$$\max \sum_{j \in B} w_j y_j r_j + \sum_{j \in NB} w_j r_j$$

$$\text{s.t. } \sum_{j \in B} c_j y_j \leq C,$$

$$y_j \in \{0, 1\} \quad j \in B,$$

$$r \in \operatorname{argmin}_{x, r} \sum_{j \in B} w_j y_j r_j + \sum_{j \in NB} w_j r_j,$$

$$x_i \in \mathbb{R}, \quad i = 1, \dots, n,$$

$$r_j = \sum_{u \in \operatorname{ext}(P)} \mu_u^j, \quad j \in B \cup NB,$$

$$x_i = \sum_{u \in \operatorname{ext}(P)} \mu_u^j u_i + f_{ji}, \quad j \in B \cup NB, i = 1, \dots, n,$$

$$\mu_u^j \geq 0, \quad u \in \operatorname{ext}(P), j \in B \cup NB,$$

$$r_j \geq 0, \quad j \in B \cup NB.$$

Discrete location  
leader problem  
(budget constraint)

**Continuous location  
follower problem**

**Representation  
of the norm**

## Benders like algorithm

We denote by  $\mathbb{P}$  the set of extreme points of the follower problem (for a norm  $l_p$ ).

## Benders like algorithm

We denote by  $\mathbb{P}$  the set of extreme points of the follower problem (for a norm  $l_p$ ).

Solving such follower problem is equivalent to:

$$\max q$$

s.t.

$$q \leq \sum_{j \in B} \sum_{i=1}^n w_j y_j r_j^\tau + \sum_{j \in NB} \sum_{i=1}^n w_j r_j^\tau, \quad r^\tau \in \mathbb{P}.$$

## Benders like algorithm

### Model rewriting:

$$\max q$$

$$\sum_{j \in B} c_j y_j \leq C,$$

$$y_j \in \{0, 1\}, \quad j \in B,$$

$$q \leq \sum_{j \in B} \sum_{i=1}^n w_j y_j r_j^\tau + \sum_{j \in NB} \sum_{i=1}^n w_j r_j^\tau, \quad r^\tau \in \mathbb{P},$$

## Benders like algorithm

### Master Problem:

$$\max q$$

$$q \leq \sum_{j \in B} \sum_{i=1}^n w_j y_j r_j^\tau + \sum_{j \in NB} \sum_{i=1}^n w_j r_j^\tau, \quad \tau \in \mathcal{P},$$

$$\sum_{j \in B} c_j y_j \leq C,$$

$$y_j \in \{0, 1\} \quad \forall j \in B,$$

## Benders like algorithm

### Primal Problem:

$$\min q(y) = \sum_{j \in B} w_j y_j r_j + \sum_{j \in NB} w_j r_j,$$

$$r_j = \sum_{e \in \text{ext}(P)} \mu_e^j, \quad j \in B \cup NB,$$

$$x_i = \sum_{e \in \text{ext}(P)} \mu_e^j e_i + f_{ji}, \quad j \in B \cup NB, i = 1, \dots, n,$$

$$\mu_e^j \geq 0, \quad e \in \text{ext}(P), j \in B \cup NB,$$

$$r_j \geq 0, \quad j \in B \cup NB,$$

$$x_i \in \mathbb{R}, \quad i = 1, \dots, n.$$

## Benders like algorithm

ALGORITHM:

**Initialization** Choose a solution  $y^0$ , and solve the Primal Problem for the chosen  $y^0$ . Let  $z^0$  be an optimal solution for the Primal Problem. Take  $\mathcal{P} = \{0\}$  and go to iteration  $\nu = 1$ .

**Iteration**  $\nu = 1, 2, \dots$  Solve the Master Problem. Let  $y^*$  be an optimal solution of such problem and  $q^+$  the corresponding optimal value.

- If  $q^+ = q(y^*)$ . END.
- Otherwise, solve the Primal Problem for  $y = y^*$ . Let  $z^*$  be an optimal solution of such problem. Take  $z^\nu = z^*$ ,  $\mathcal{P} := \mathcal{P} \cup \{\nu\}$ , and go to iteration  $\nu := \nu + 1$ .

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5 Extensions

## Examples: $\ell_1$ and $\ell_\infty$ norm

### $\ell_1$ norm (rectangular norm)

$$\|x\|_1 = \sum_{i=1}^n |x_i|.$$

- $B_{l_1} = \{x \in \mathbb{R}^n : \|x\|_1 \leq 1\}$
- $\text{ext}(P) = \{e_1, \dots, e_n, -e_1, \dots, -e_n\}$

## Examples: $\ell_1$ and $\ell_\infty$ norm

### $\ell_1$ norm (rectangular norm)

$$\|x\|_1 = \sum_{i=1}^n |x_i|.$$

### $\ell_\infty$ norm (Chebyshev norm)

$$\|x\|_\infty = \max_{1 \leq i \leq n} |x_i|.$$

- $B_{\ell_1} = \{x \in \mathbb{R}^n : \|x\|_1 \leq 1\}$
- $ext(P) = \{e_1, \dots, e_n, -e_1, \dots, -e_n\}$
- $B_{\ell_\infty} = \{x \in \mathbb{R}^n : \|x\|_\infty \leq 1\}$
- $ext(P) = \{(1^{a_1}, \dots, 1^{a_n}) \in \mathbb{R}^n : a_i \in \{1, -1\}, i = 1, \dots, n\}$

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- $ext(P^0) = \{e_1, \dots, e_n, -e_1, \dots, -e_n\}$

## Examples: $\ell_1$ and $\ell_\infty$ norm

$$\max \sum_{j \in B} \sum_{i=1}^n w_j \hat{r}_{ij} + \sum_{j \in NB} \sum_{i=1}^n w_j r_{ji} \quad (\ell_1 \text{ norm Formulation})$$

$$\text{s.t. } \sum_{j \in B} c_j y_j \leq C,$$

$$y_j \in \{0, 1\} \forall j \in B,$$

$$\sum_{j \in B} \sum_{i=1}^n w_j \hat{r}_{ij} + \sum_{j \in NB} \sum_{i=1}^n w_j r_{ji} = \sum_{j \in B \cup NB} \sum_{i=1}^n -f_{ji} u_{ji} + \sum_{j \in B} \sum_{i=1}^n f_{ji} (w_j y_j - u_{ji}) + \sum_{j \in NB} \sum_{i=1}^n f_{ji} (w_j - u_{ji})$$

$$r_{ji} \geq x_i - f_{ji}, \quad j \in B \cup NB, i = 1, \dots, n,$$

$$r_{ji} \geq f_{ji} - x_i, \quad j \in B \cup NB, i = 1, \dots, n,$$

$$\hat{r}_{ji} \leq M y_j, \quad j \in B, i = 1, \dots, n,$$

$$\hat{r}_{ji} \leq r_{ji}, \quad j \in B, i = 1, \dots, n,$$

$$\hat{r}_{ji} \geq r_{ji} - (1 - y_j)M, \quad j \in B, i = 1, \dots, n,$$

$$\hat{r}_{ji} \geq 0, \quad j \in B, i = 1, \dots, n,$$

$$\sum_{j \in B} (-2u_{ji} + w_j y_j) + \sum_{j \in NB} (-2u_{ji} + w_j) = 0, \quad i = 1, \dots, n,$$

$$u_{ji} \geq 0, \quad j \in B \cup NB, i = 1, \dots, n,$$

$$w_j y_j - u_{ji} \geq 0, \quad j \in B, i = 1, \dots, n,$$

$$w_j - u_{ji} \geq 0, \quad j \in B, i = 1, \dots, n.$$

## Examples: $\ell_1$ and $\ell_\infty$ norm

The inner location problem can be decomposed into  $n$  independent linear programs, one per each coordinate.

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The inner location problem can be decomposed into  $n$  independent linear programs, one per each coordinate.

### MILP formulation:

$$\sum_{j \in B} w_j \hat{r}_{ji} + \sum_{j \in NB} w_j r_{ji} = \sum_{j \in B \cup NB} -f_{ji} \alpha_{ji} + \sum_{j \in B} f_{ji} (w_j y_j - \alpha_{ji}) + \sum_{j \in NB} f_{ji} (w_j - \alpha_{ji}), \quad i = 1, \dots, n.$$

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### Benders like algorithm:

$$\max \quad \sum_{i=1}^n q_i$$

s.t.

$$q_i \leq \sum_{j \in B} w_j y_j z_{ij}^\tau + \sum_{j \in NB} w_j z_{ij}^\tau \quad \forall \tau \in \mathcal{P}, \quad \forall i = 1, \dots, n,$$

$$\sum_{j \in B} c_j y_j \leq C,$$

$$y_j \in \{0, 1\} \quad \forall j \in B,$$

## Examples: $\ell_1$ and $\ell_\infty$ norm

$$\max \sum_{j \in B} w_j \hat{r}_j + \sum_{j \in NB} w_j r_j \quad (\ell_\infty \text{ norm Formulation})$$

$$\text{s.t. } \sum_{j \in B} c_j y_j \leq C,$$

$$y_j \in \{0, 1\} \forall j \in B,$$

$$\sum_{j \in B} w_j \hat{r}_j + \sum_{j \in NB} w_j r_j = \sum_{j \in B \cup NB} \sum_{i=1}^n f_{ji} (-u_{ji} + v_{ji}),$$

$$r_j \geq x_i - f_{ji}, \quad j \in B \cup NB, i = 1, \dots, n,$$

$$r_j \geq f_{ji} - x_i, \quad j \in B \cup NB, i = 1, \dots, n,$$

$$\sum_{j \in B \cup NB} (-u_{ji} + v_{ji}) = 0, \quad i = 1, \dots, n,$$

$$\sum_{i=1}^n (u_{ji} + v_{ji}) = w_j y_j, \quad j \in B,$$

$$\sum_{i=1}^n (u_{ji} + v_{ji}) = w_j, \quad j \in NB,$$

$$u_{ji} \geq 0, v_{ji} \geq 0, \quad j \in B \cup NB, i = 1, \dots, n,$$

$$\hat{r}_j \leq M y_j, \quad j \in B$$

$$\hat{r}_j \leq r_j, \quad j \in B,$$

$$r_j \geq r_j - (1 - y_j)M, \hat{r}_j \geq 0 \quad j \in B.$$

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## Computational Results

- $n$  (dimension): **2,3,10,20.**
- $\text{card}(B)$  (potential facilities): **1000, 2000, 5000, 10000.**
- $\text{card}(NB)$  (opened facilities):  $\{\frac{1}{4}, \frac{1}{3}, \frac{1}{2}\} \text{card}(B).$
- $w_j$  (proportional charge): **random in [0, 1].**
- $c_j$  (opening cost): **random in [0, 1].**
- $C$  (budget):  $\{\frac{1}{4}, \frac{1}{3}\} \text{card}(B).$
- $f_{ij}$  (coordinates): **random in [-1000, 1000].**

## Computational Results

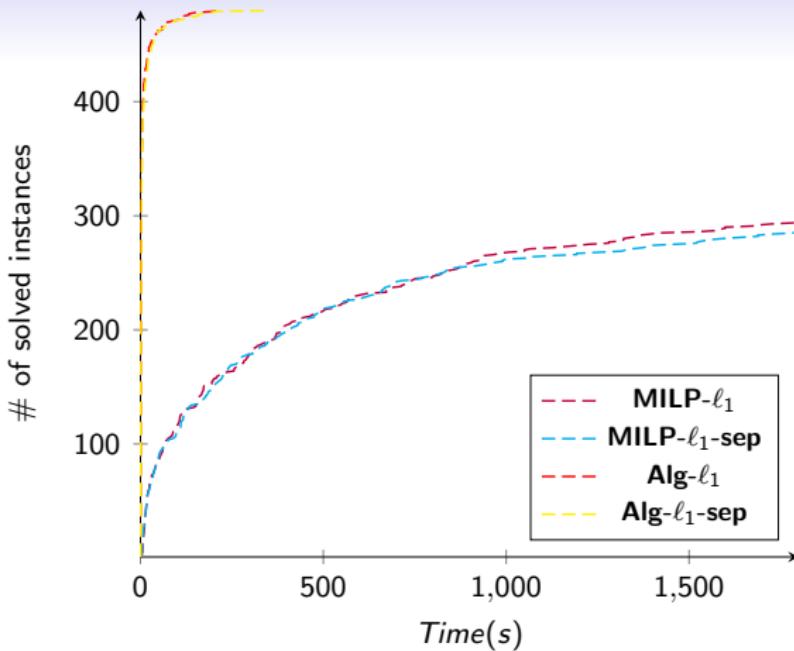
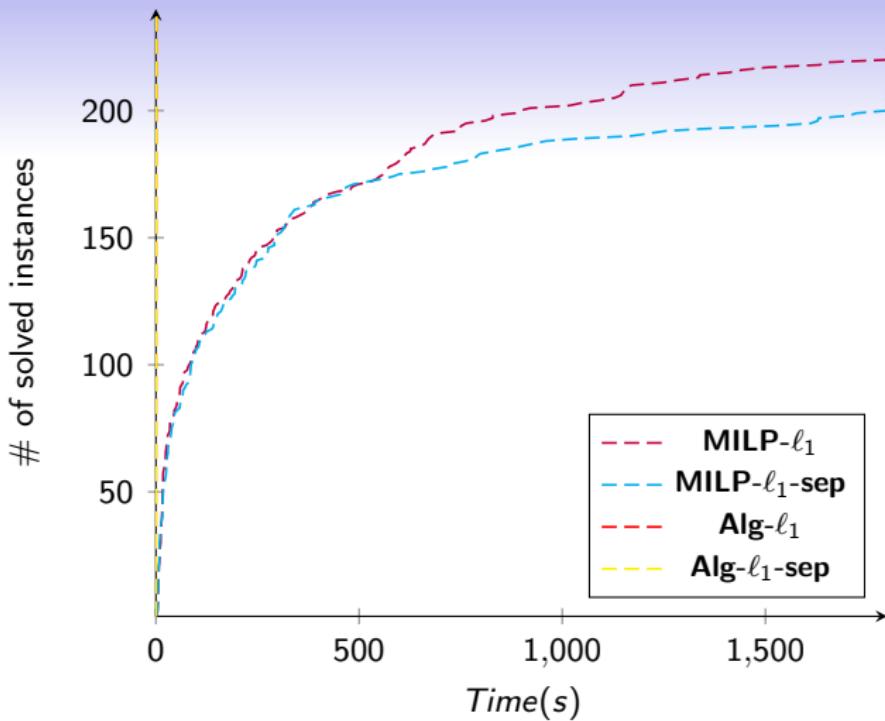


Figure: Performance profile graph of #solved instances for the different proposed models for the  $\ell_1$  norm.



**Figure:** Performance profile graph of #solved instances for the different proposed models for *small* instances ( $n = 2, 3$ ) for the  $\ell_1$  norm.

## Computational Results

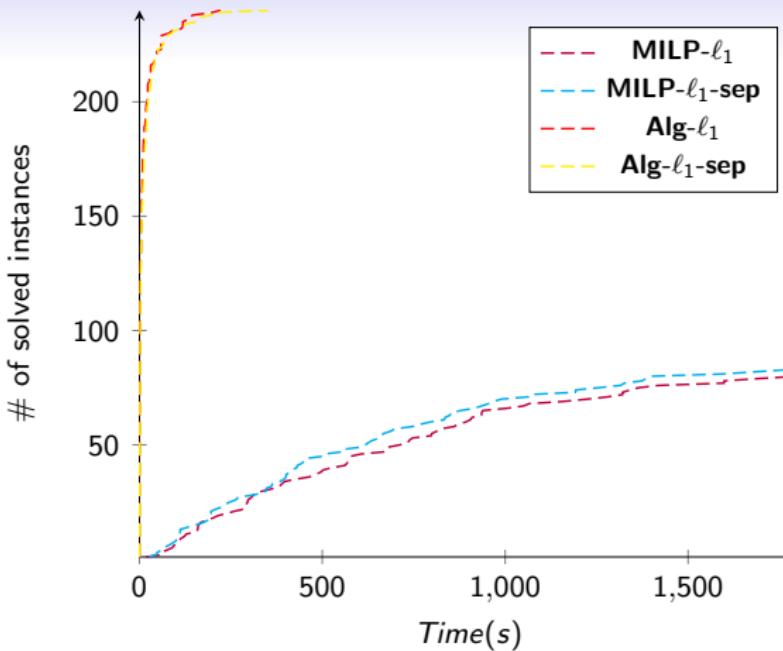


Figure: Performance profile graph of #solved instances for the different proposed models for *big* instances ( $n = 10, 20$ ) for the  $\ell_1$  norm.

## Computational Results

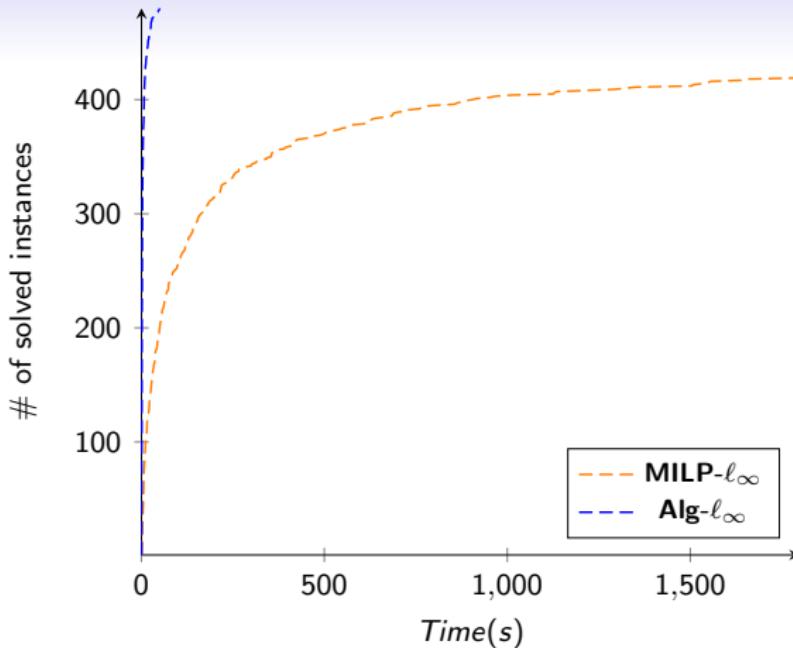


Figure: Performance profile graph of #solved instances for the different proposed models for the  $\ell_\infty$  norm.

INSTANCES				MILP- $\ell_1$		MILP- $\ell_1$ -sep		Alg- $\ell_1$		Alg- $\ell_1$ -sep		MILP- $\ell_\infty$		Alg- $\ell_\infty$	
n	B	NB	C'	#OPT	CPU	#OPT	CPU	#OPT	CPU	#OPT	CPU	#OPT	CPU	#OPT	CPU
2	1000	250	3	5	10,82	5	10,36	5	0,41	5	0,50	5	2,13	5	0,56
2	1000	250	4	5	4,80	5	5,89	5	0,62	5	0,52	5	2,22	5	1,03
2	1000	333	3	5	13,39	5	10,07	5	0,35	5	0,41	5	2,19	5	0,65
2	1000	333	4	5	7,57	5	4,98	5	0,55	5	0,62	5	2,47	5	1,05
2	1000	500	3	5	11,70	5	12,76	5	0,38	5	0,43	5	8,53	5	0,81
2	1000	500	4	5	7,86	5	6,35	5	0,54	5	0,52	5	3,02	5	0,95
2	2000	500	3	5	36,59	5	37,27	5	0,75	5	1,44	5	6,34	5	0,79
2	2000	500	4	5	13,36	5	11,66	5	0,65	5	0,73	5	9,15	5	1,36
2	2000	667	3	5	217,01	5	152,79	5	0,68	5	0,85	5	7,20	5	0,92
2	2000	667	4	5	20,58	5	14,54	5	0,45	5	0,60	5	6,14	5	1,17
2	2000	1000	3	5	56,48	5	78,63	5	0,45	5	0,61	5	15,29	5	0,81
2	2000	1000	4	5	29,85	5	20,60	5	0,52	5	0,65	5	12,74	5	1,56
2	5000	1250	3	5	204,62	5	508,34	5	1,19	5	1,58	5	48,51	5	1,62
2	5000	1250	4	5	124,68	5	97,10	5	1,49	5	2,36	5	52,24	5	1,70
2	5000	1667	3	5	381,53	5	268,18	5	1,03	5	1,32	5	45,06	5	1,47
2	5000	1667	4	5	135,94	5	198,89	5	2,08	5	1,66	5	37,01	5	1,96
2	5000	2500	3	5	416,01	5	369,43	5	0,80	5	1,12	5	139,15	5	2,34
2	5000	2500	4	5	107,17	5	405,78	5	1,23	5	1,89	5	38,32	5	1,92
2	10000	2500	3	4	708,80	3	955,21	5	1,68	5	2,57	5	137,97	5	3,29
2	10000	2500	4	5	325,03	5	355,51	5	2,44	5	3,75	5	81,66	5	3,72
2	10000	3333	3	5	390,82	5	296,18	5	1,94	5	2,78	5	121,14	5	3,42
2	10000	3333	4	5	329,09	5	439,71	5	3,09	5	3,46	5	159,98	5	3,16
2	10000	5000	3	5	506,03	4	671,22	5	2,16	5	2,09	4	530,49	5	3,38
2	10000	5000	4	5	516,73	5	416,86	5	2,83	5	3,27	5	141,33	5	3,77

Table: Numerical results for  $n = 2$  under the  $\ell_1$  and  $\ell_\infty$  norm.

INSTANCES			MILP- $\ell_1$		MILP- $\ell_1$ -sep		Alg- $\ell_1$		Alg- $\ell_1$ -sep		MILP- $\ell_\infty$		Alg- $\ell_\infty$		
n	B	NB	C'	#OPT	CPU	#OPT	CPU	#OPT	CPU	#OPT	CPU	#OPT	CPU	#OPT	CPU
3	1000	250	3	5	9,67	5	19,03	5	0,21	5	0,27	5	5,95	5	0,70
3	1000	250	4	5	58,74	5	43,20	5	0,63	5	0,71	5	3,09	5	1,18
3	1000	333	3	5	21,51	5	50,24	5	0,36	5	0,37	5	6,99	5	0,60
3	1000	333	4	5	29,05	5	32,52	5	0,65	5	0,90	5	4,68	5	1,21
3	1000	500	3	5	26,14	5	28,70	5	0,34	5	0,49	5	10,17	5	0,57
3	1000	500	4	5	33,46	5	47,28	5	0,38	5	0,52	5	6,86	5	0,73
3	2000	500	3	5	120,41	5	144,41	5	0,38	5	0,53	5	63,77	5	1,19
3	2000	500	4	5	389,05	5	301,49	5	0,76	5	1,22	5	21,33	5	1,41
3	2000	667	3	5	316,33	5	232,75	5	1,04	5	0,92	5	31,18	5	1,04
3	2000	667	4	5	129,52	5	224,55	5	0,47	5	0,56	5	21,81	5	1,48
3	2000	1000	3	5	113,90	5	261,64	5	0,65	5	0,53	5	81,42	5	1,38
3	2000	1000	4	5	179,31	5	249,42	5	0,62	5	0,71	5	11,05	5	0,88
3	5000	1250	3	4	1014,05	3	1310,24	5	1,95	5	1,85	5	131,33	5	1,94
3	5000	1250	4	3	1352,73	2	1442,08	5	2,43	5	2,24	5	106,38	5	2,12
3	5000	1667	3	4	763,66	4	669,57	5	1,14	5	1,28	5	180,78	5	1,94
3	5000	1667	4	3	1102,02	4	1047,36	5	1,56	5	1,91	5	213,59	5	2,08
3	5000	2500	3	5	636,19	5	689,33	5	0,76	5	1,14	5	391,97	5	2,07
3	5000	2500	4	5	592,53	2	1223,05	5	1,42	5	1,58	5	77,89	5	1,90
3	10000	2500	3	1	1778,05	1	1800,00	5	2,42	5	3,38	5	899,86	5	3,09
3	10000	2500	4	1	1727,79	0	—	5	5,59	5	5,16	5	198,48	5	4,18
3	10000	3333	3	3	1127,45	1	1732,91	5	2,67	5	3,91	5	832,12	5	3,17
3	10000	3333	4	3	1253,69	2	1321,56	5	4,79	5	6,15	4	797,61	5	3,91
3	10000	5000	3	3	1379,12	1	1771,16	5	2,86	5	4,41	4	470,05	5	4,01
3	10000	5000	4	0	—	0	—	5	4,18	5	5,79	5	325,39	5	4,87

Table: Numerical results for  $n = 3$  under the  $\ell_1$  and  $\ell_\infty$  norm.

INSTANCES			MILP- $\ell_1$		MILP- $\ell_1$ -sep		Alg- $\ell_1$		Alg- $\ell_1$ -sep		MILP- $\ell_\infty$		Alg- $\ell_\infty$		
n	B	NB	C	#OPT	CPU	#OPT	CPU	#OPT	CPU	#OPT	CPU	#OPT	CPU	#OPT	CPU
10	1000	250	3	5	175,69	5	262,45	5	0,52	5	0,84	5	23,53	5	0,73
10	1000	250	4	5	286,71	5	199,67	5	0,58	5	1,02	5	10,04	5	0,60
10	1000	333	3	5	170,27	5	168,75	5	0,29	5	0,57	5	50,58	5	0,68
10	1000	333	4	5	380,23	5	308,91	5	0,79	5	0,98	5	16,31	5	0,66
10	1000	500	3	5	365,20	5	412,48	5	0,56	5	0,64	5	40,66	5	0,80
10	1000	500	4	4	838,00	5	608,46	5	0,74	5	1,17	5	26,14	5	1,16
10	2000	500	3	3	1538,42	3	1310,11	5	0,81	5	1,90	5	104,57	5	1,31
10	2000	500	4	3	1197,88	4	1020,28	5	1,04	5	1,44	5	64,34	5	1,62
10	2000	667	3	2	1406,81	2	1365,28	5	1,23	5	1,20	5	123,13	5	1,22
10	2000	667	4	3	1092,84	4	928,63	5	1,01	5	1,24	5	58,13	5	1,16
10	2000	1000	3	4	1155,24	4	1125,38	5	0,72	5	0,83	5	183,55	5	1,59
10	2000	1000	4	3	980,54	3	1215,99	5	1,11	5	1,84	5	91,77	5	1,66
10	5000	1250	3	1	1673,56	1	1796,25	5	2,66	5	3,00	3	1265,41	5	4,80
10	5000	1250	4	0	—	0	—	5	9,46	5	14,29	5	421,92	5	5,76
10	5000	1667	3	0	—	0	—	5	5,08	5	10,21	4	840,57	5	3,86
10	5000	1667	4	2	1430,78	2	1690,69	5	2,77	5	3,19	4	627,48	5	5,12
10	5000	2500	3	1	1568,60	0	—	5	3,91	5	5,34	4	754,36	5	5,44
10	5000	2500	4	1	1727,35	1	1649,57	5	6,04	5	8,31	3	1125,20	5	5,97
10	10000	2500	3	0	—	0	—	5	9,37	5	17,15	2	1211,56	5	10,24
10	10000	2500	4	0	—	0	—	5	16,19	5	26,46	4	955,69	5	8,48
10	10000	3333	3	0	—	0	—	5	17,88	5	20,49	1	1649,16	5	13,44
10	10000	3333	4	0	—	0	—	5	19,62	5	18,73	3	1152,58	5	10,75
10	10000	5000	3	1	1800,00	0	—	5	11,24	5	23,00	2	1800,00	5	10,91
10	10000	5000	4	0	—	0	—	5	28,00	5	16,26	2	1314,08	5	14,30

Table: Numerical results for  $n = 10$  under the  $\ell_1$  and  $\ell_\infty$  norm.

INSTANCES				MILP- $\ell_1$		MILP- $\ell_1$ -sep		Alg- $\ell_1$		Alg- $\ell_1$ -sep		MILP- $\ell_\infty$		Alg- $\ell_\infty$	
n	B	NB	C	#OPT	CPU	#OPT	CPU	#OPT	CPU	#OPT	CPU	#OPT	CPU	#OPT	CPU
20	1000	250	3	4	797,87	5	477,74	5	0,70	5	1,38	5	41,66	5	1,25
20	1000	250	4	5	852,01	5	548,28	5	0,64	5	1,31	5	20,22	5	0,88
20	1000	333	3	4	608,99	5	296,04	5	0,54	5	1,03	5	29,01	5	0,99
20	1000	333	4	4	1026,53	4	773,41	5	1,01	5	1,93	5	39,01	5	1,55
20	1000	500	3	4	881,29	5	412,55	5	0,51	5	1,13	5	55,68	5	1,42
20	1000	500	4	2	1499,86	3	1478,83	5	0,93	5	1,55	5	36,02	5	1,48
20	2000	500	3	1	1629,43	2	1800,00	5	1,93	5	2,21	5	212,07	5	2,08
20	2000	500	4	1	1623,38	2	1634,03	5	3,91	5	2,82	5	217,86	5	2,42
20	2000	667	3	2	1612,20	3	1801,02	5	1,78	5	2,05	5	223,80	5	2,40
20	2000	667	4	3	1421,62	2	1416,62	5	1,34	5	2,44	5	134,67	5	2,64
20	2000	1000	3	0	—	2	1800,00	5	2,45	5	1,86	5	181,95	5	3,02
20	2000	1000	4	1	1800,00	1	1800,00	5	5,09	5	3,08	5	345,70	5	3,44
20	5000	1250	3	0	—	0	—	5	7,20	5	9,37	3	1188,15	5	8,67
20	5000	1250	4	0	—	0	—	5	14,87	5	17,02	5	868,00	5	9,05
20	5000	1667	3	0	—	0	—	5	7,92	5	30,75	2	1513,57	5	8,88
20	5000	1667	4	0	—	0	—	5	7,28	5	18,85	5	1034,49	5	9,46
20	5000	2500	3	0	—	0	—	5	16,23	5	13,53	2	2447,37	5	15,67
20	5000	2500	4	0	—	0	—	5	20,15	5	10,45	2	1585,53	5	13,13
20	10000	2500	3	0	—	0	—	5	52,42	5	65,91	2	1745,86	5	23,08
20	10000	2500	4	0	—	0	—	5	96,02	5	63,30	1	1800,00	5	27,02
20	10000	3333	3	0	—	0	—	5	28,93	5	25,17	2	1750,37	5	22,32
20	10000	3333	4	0	—	0	—	5	125,00	5	93,66	0	—	5	29,79
20	10000	5000	3	0	—	0	—	5	40,39	5	85,61	2	1800,00	5	38,66
20	10000	5000	4	0	—	0	—	5	88,29	5	167,54	1	1800	5	27,76

Table: Numerical results for  $n = 20$  under the  $\ell_1$  and  $\ell_\infty$  norm.

# Contents

1 Introduction

2 Model Description

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- Examples:  $l_1$  and  $l_\infty$  norm

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## K-followers

$$\begin{aligned} & \max \sum_{k \in K} \left( \sum_{j \in B} w_j^k d(x^k, f_j) y_j + \sum_{j \in NB} w_j^k d(x^k, f_j) \right) \\ \text{s.t. } & \sum_{j \in B} c_j y_j \leq C, \\ & y_j \in \{0, 1\}, \quad j \in B, \\ & x^k \in \arg \min_{x^k} \sum_{j \in B} w_j^k d(x^k, f_j) y_j + \sum_{j \in NB} w_j^k d(x^k, f_j) \quad \forall k = 1, \dots, K. \end{aligned}$$

## K-followers

$$\max \sum_{k \in K} q^k$$

$$\text{s.t. } q^k \leq \sum_{j \in B} \sum_{i=1}^n w_j^k y_j r_j^{k\tau} + \sum_{j \in NB} \sum_{i=1}^n w_j^k r_j^{k\tau} \quad \forall \tau \in \mathcal{P}, \forall k \in K,$$

$$\sum_{j \in B} c_j y_j \leq C,$$

$$y_j \in \{0, 1\} \quad \forall j \in B.$$

**$\ell_\tau$ -norms with  $\tau \in \mathbb{Q}, \tau \geq 1$** 

$$\begin{aligned} & \max \sum_{j \in B} w_j \|x - f_j\|_\tau y_j + \sum_{j \in NB} w_j \|x - f_j\|_\tau \\ \text{st. } & \sum_{j \in B} c_j y_j \leq C, \\ & y_j \in \{0, 1\} \quad j \in B, \\ & x \in \arg \min_x \sum_{j \in B} w_j \|x - f_j\|_\tau y_j + \sum_{j \in NB} w_j \|x - f_j\|_\tau. \end{aligned}$$

## $\ell_\tau$ -norms with $\tau \in \mathbb{Q}$ , $\tau \geq 1$

$$\begin{aligned} & \max \sum_{j \in B} w_j \|x - f_j\|_\tau y_j + \sum_{j \in NB} w_j \|x - f_j\|_\tau \\ \text{st. } & \sum_{j \in B} c_j y_j \leq C, \\ & y_j \in \{0, 1\} \quad j \in B, \\ & x \in \arg \min_x \sum_{j \in B} w_j \|x - f_j\|_\tau y_j + \sum_{j \in NB} w_j \|x - f_j\|_\tau. \end{aligned}$$



Blanco, V., Puerto, J., Ali, SEHB. (2014). Revisiting several problems and algorithms in continuous location with  $\ell_\tau$  norms. Computational Optimization and Applications. 58(3), 563-595.

## $\ell_\tau$ -norms with $\tau \in \mathbb{Q}, \tau \geq 1$

$$\max \sum_{j \in B} w_j ||x - f_j||_\tau y_j + \sum_{j \in NB} w_j ||x - f_j||_\tau$$

$$\text{s.t. } \sum_{j \in B} c_j y_j \leq C,$$

$$y_j \in \{0, 1\} \quad j \in B,$$

$$\sum_{j \in B} w_j ||x - f_j||_\tau y_j + \sum_{j \in NB} w_j ||x - f_j||_\tau = \sum_{j \in B \cup NB} \sum_{i=1}^n V_{ji} f_{ji},$$

$$x^+ - x^- - Z_j = f_j, \quad \forall j \in B \cup NB,$$

$$||Z_j||_\tau \leq r_j, \quad \forall j \in B \cup NB,$$

$$x^+, x^- \in \mathbb{R}_+^n, Z_j \in \mathbb{R}^n, r \in \mathbb{R}^n,$$

$$\sum_{j \in B \cup NB} V_{ji} + \lambda_i^1 = 0, \quad \forall i = 1, \dots, n,$$

$$- \sum_{j \in B \cup NB} V_{ji} + \lambda_i^2 = 0, \quad \forall i = 1, \dots, n,$$

$$- V_{ji} + \mu_{ji} = 0, \quad \forall j \in B \cup NB, i = 1, \dots, n,$$

$$||\mu_j||_\rho \leq \gamma_j, \quad \forall j \in B \cup NB,$$

$$\gamma_j = w_j y_j, \text{ if } i \in B,$$

$$\gamma_j = w_j, \text{ if } i \in NB,$$

$$\lambda^1, \lambda^2 \in \mathbb{R}_+^n, \mu_j \in \mathbb{R}^n, \forall j \in B \cup NB, \gamma \in \mathbb{R}^n.$$

# New bilevel programming approaches to the location of controversial facilities

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Martine Labb  <sup>1</sup>, Marina Leal<sup>2,3</sup>, Justo Puerto<sup>2,3</sup>.

<sup>1</sup> Computer Sciences Department, Universit   Libre de Bruxelles.

<sup>2</sup> Department of Statistics and Operational Research, Universidad de Sevilla.

<sup>3</sup> Institute of Mathematics of the University of Seville (IMUS).

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