

VI International Workshop on Locational Analysis and Related Problems

New products supply chains: the effect of short lifecycles on the supply chain network design.



REDLOCA 2015

Work in-progress by

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The elevator speech

- Node-demand is stochastic
- Facility location
 - Transportation_Cost
- Capacity decisions
 - Newsboy-type problem
 - Cost of Undercapacity: $C_u = (\text{Unit_Revenue} - \text{Unit_Cost}) - \text{Margin_Outsourcing}$
 - Cost Overcapacity: $C_o = \text{Capacity_Cost}$
 - Critical Ratio: $CR = C_u / (C_u + C_o)$
- Difficulty
 - Unit_Cost(Capacity_Cost, Transportation_Cost)

Agenda

- Motivation
- The problem
- The cost of capacity
- The formulation
- The single facility case
- The multi-facility case
- Bounds on the objective function

Overview

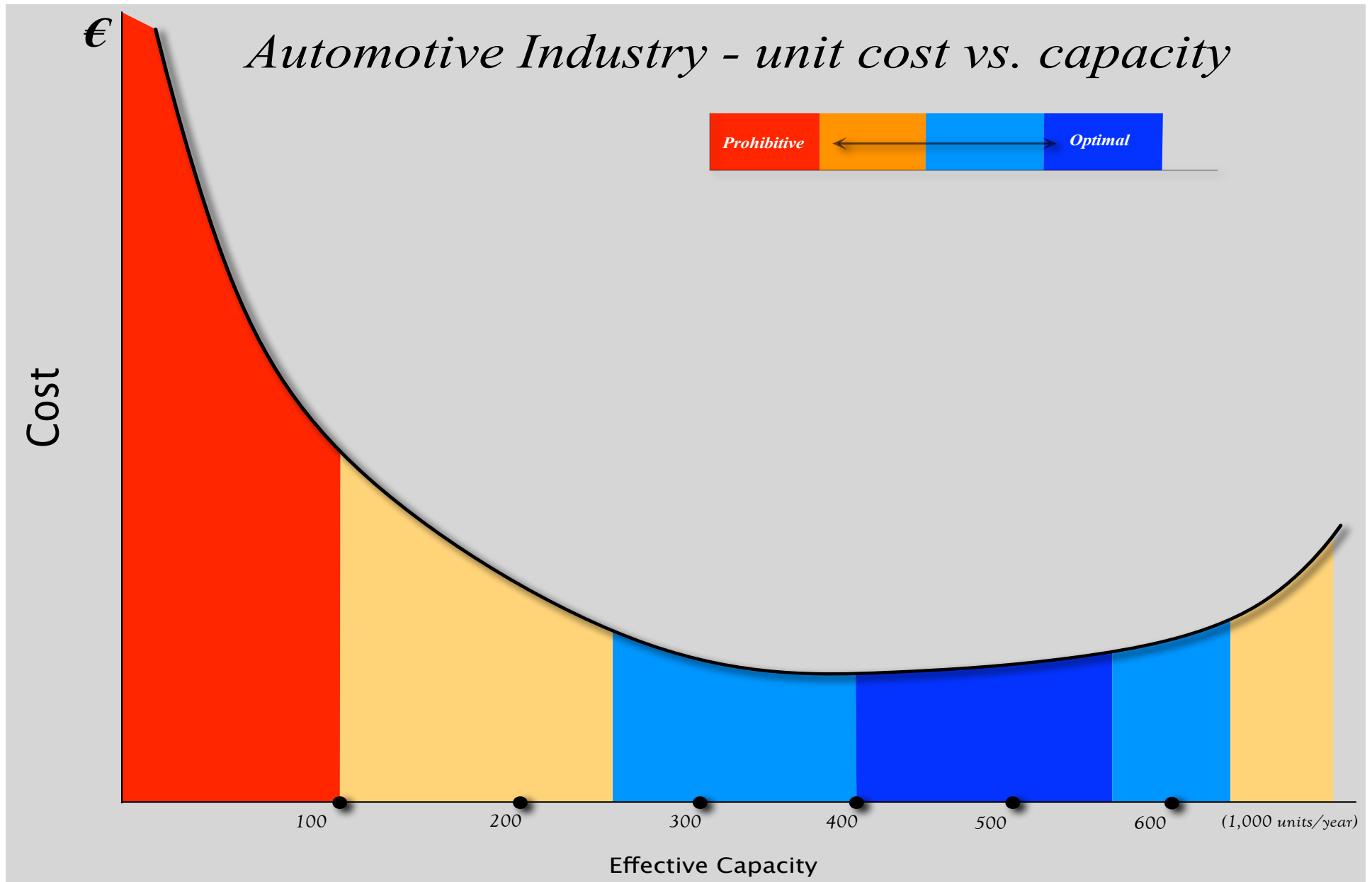
- H. Frank (1966): Optimum Location on a Graph with Probabilistic Demands
- ***Facility Location as an strategic level problem: capacity***
- Consider a problem where capacity has to be built up or reserved/bought from third-party
- Planning period is well ahead of product-introduction period
- Demand is unknown and described by scenarios and their probabilities
- Result: Facility Location w/ Capacity Decision (newsvendor type)

The Problem (α)

- Consider a vector of demand scenarios: $\partial=(D_1, D_2,..., D_y)$ and a network $G(N,A)$. $|N|=n$. D_s is the demand vector for all node-demand in scenario s .
- Each scenario ∂_i , happens with probability $f(\partial_i)$
- Facilities are labeled 1, 2, ..., are located according to a location set $S=\{S_1, S_2,..., S_p\}$, where $p \leq n$. Vector C is the capacity vector corresponding to each facility.
- For every pair node-facility (i, j) , there is a corresponding unit margin for serving node i from facility at j .
- α denotes an allocation matrix where element a_{ij}^d defines the quantity sent from facility at j to node at i when demand is d .

The Problem (*b*)

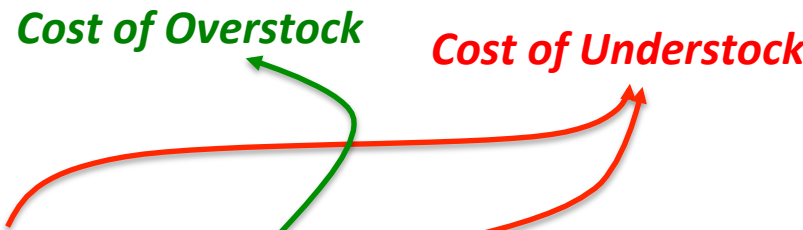
- A feasible allocation a is such that $\sum_i a_{ij}^d \leq C_j; \forall j \in \{1, \dots, p\}, d \in \partial$ and $\sum_j a_{ij}^d \leq d_i; \forall i \in N, d \in \partial$
- Cost parameters
 - m_0 is unit margin when outsourcing, m_{ij} is the margin from serving node i from facility at j .
- Quantity outsourced $L(C) = \max\{\sum_i D_i - \sum_j C_j, 0\}$
 - Cost of outsourcing is $m_0 L(C)$
- Capacity cost: $\varphi(C)$



The Problem (c)

- The deterministic cost function

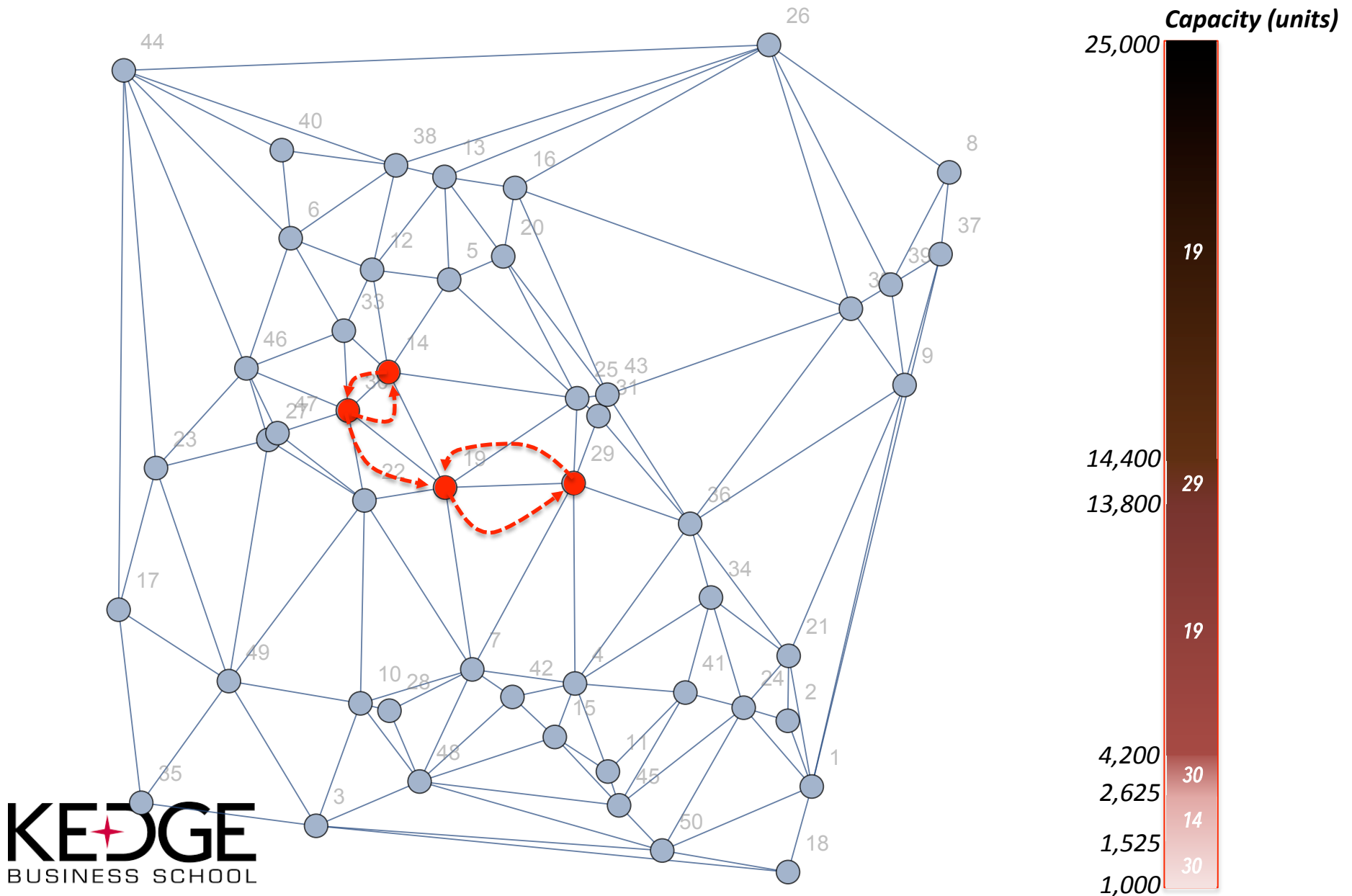
$$Z(S, C, a \mid d) = \underbrace{\left(\sum_i \sum_j m_{ij} a_{ij} \right)}_{\text{revenue}} - \underbrace{\psi(C)}_{\text{capacity cost}} + \underbrace{m_0 L(C)}_{\text{outsourcing margin}}$$

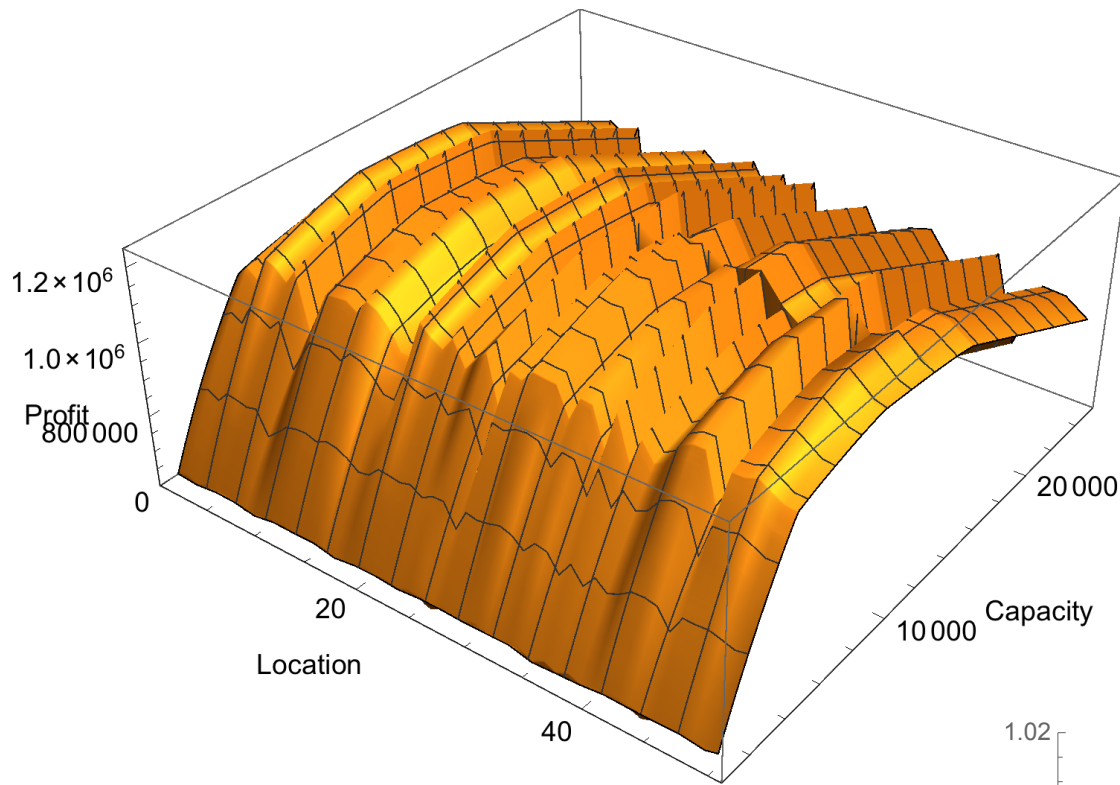


- The problem

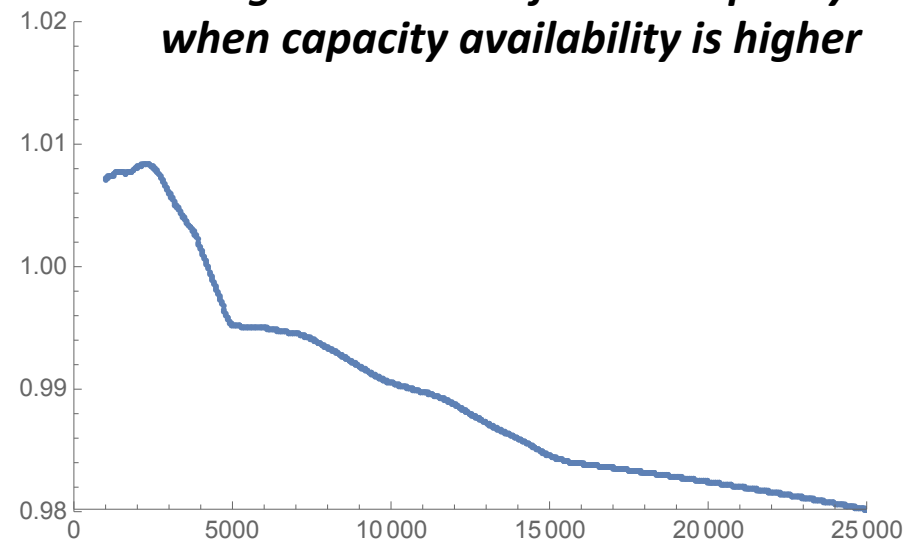
$$\max_{S, C, a} V(S, C, a) = E_d[Z(S, C, a \mid d)]$$

One facility case: sensitivity analysis on capacity

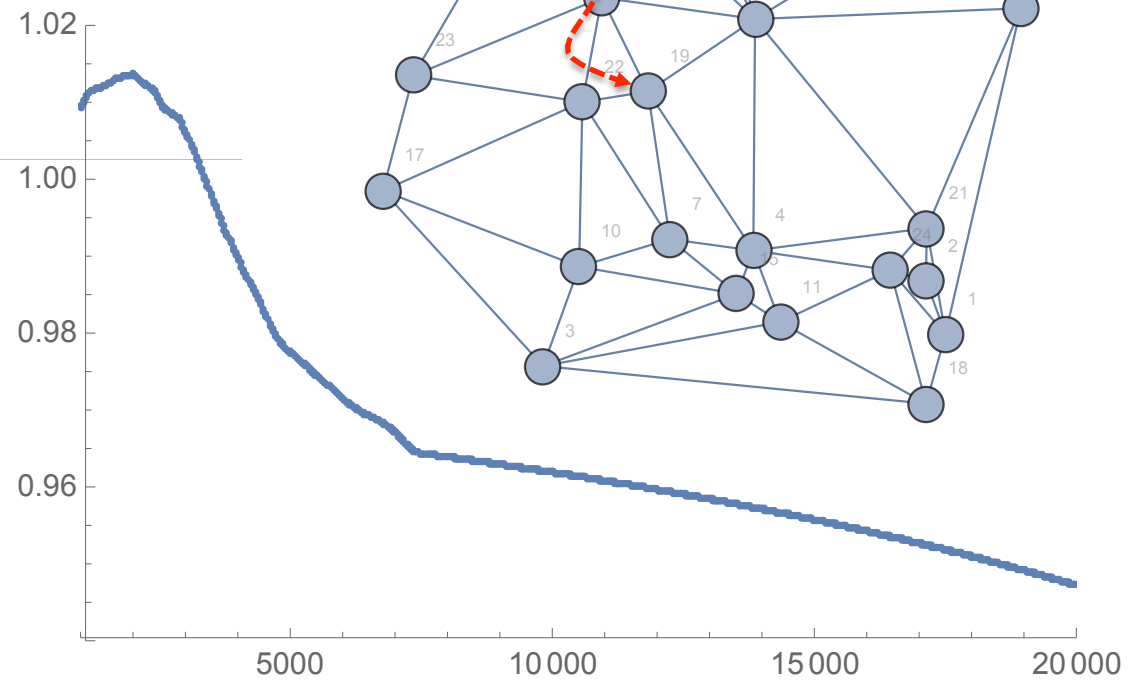
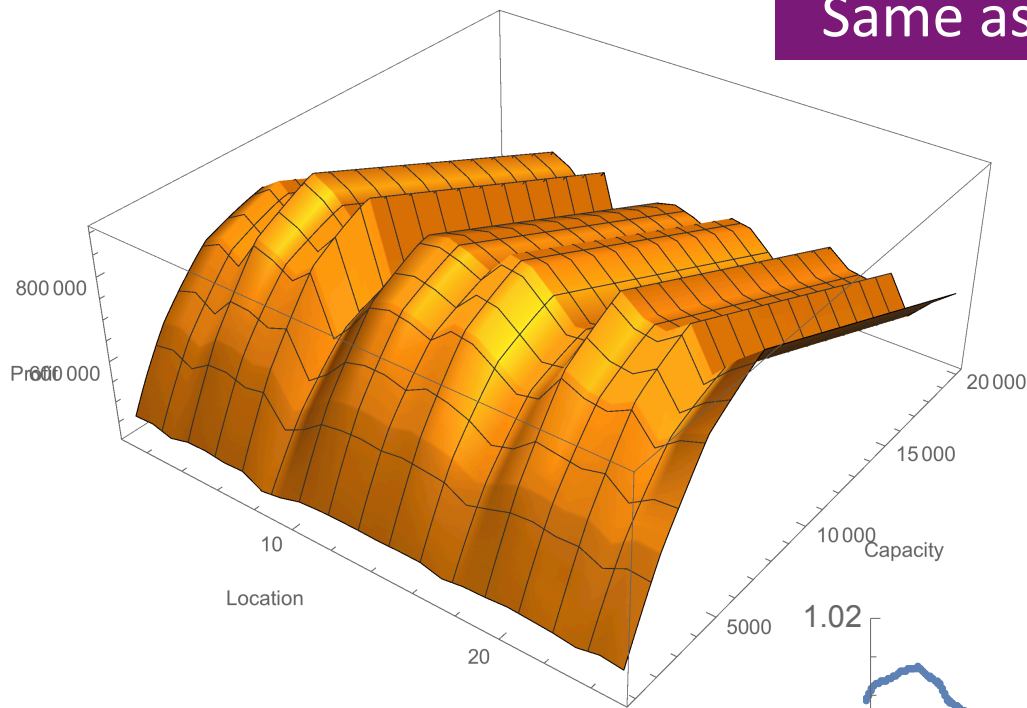




***Using the location for low capacity
when capacity availability is higher***

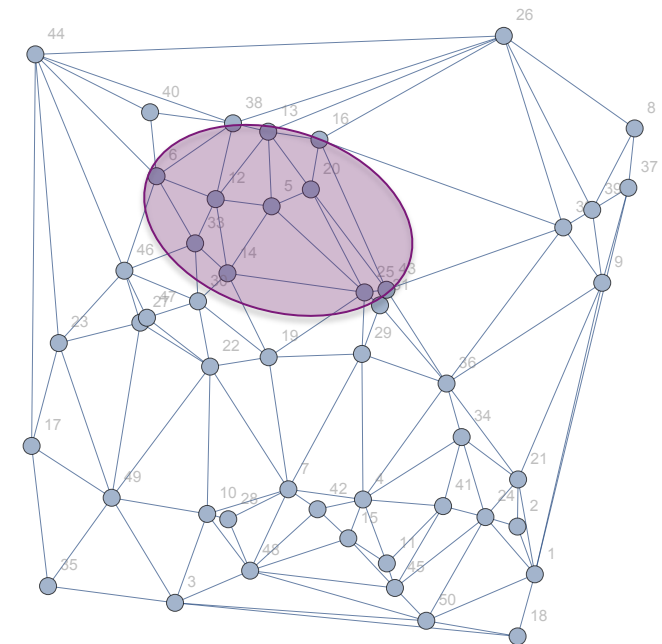


Same as previous with a smaller network



1-Facility Location Solution is Robust

- When independently generating node weights randomly from distributions with same mean value when building scenarios...
 - If the network is large enough then transportation cost dominates and induce solutions among the “usual suspects”; i.e. in the vicinity of each other.



Observations suggest...

- Solutions are robust for changes in critical ratios
- For a fixed set of facility location profit is concave in capacity vector
- Objective is(?) submodular on location set F and $C^*(F)$
- *Proposed Greedy Heuristic*
 - a) For a fix location set F
 - b) $|F|=f$, find node $k_f = \arg_k \text{ maximizes } Z(C,a | F \cup k)$
 - c) Stop when profit is not improved; otherwise go to (a) above.

Still hard to compute optimality?

Bounds to the objective function value

- Consider dynamic setting where total capacity is chosen a priori...
- ... but distribution of the capacity is done after scenario is realized
- The bound is easy to find for a fixed set S .
- **Recall: Objective function seems to inherit Submodularity property from other location problems (to be proven).**

THANKS!