

# The generalized discrete ( $r|p$ )-centroid problem

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## Introduction

We consider a discrete competitive location model which is called ( $r|p$ )-centroid problem, leader-follower problem or Stackelberg problem in locations. The model represents a situation where two players, the leader and the follower make decisions sequentially in order to reach certain objectives. Leader and follower want to determine the locations for  $p$  and  $r$  facilities respectively. The objective of the follower, who makes the decision once the leader has selected its locations, is to maximize the demand captured by its facilities. The objective of the leader is to minimize the maximum demand that the follower could capture, as demand is assumed to be essential, this objective is equivalent to maximize the demand captured by its facilities. We study a generalized model in which the customer's choice rule is defined using a non-increasing capture function, the demand captured by the players is given by the value of this function for the difference between the distance from the demand point to the follower and the distance from the demand point to the leader. Given the set of  $p$  locations for the leader,  $X_p$ , the solution for the follower is an ( $r|X_p$ )-medianoid. The solution for the leader is an ( $r|p$ )-centroid.

We present a linear programming formulation for the generalized ( $r|p$ )-centroid problem and an exact solution procedure. The model is illustrated with an example.

## Statement of the problem

Let  $V$  be a set of points and  $C, L$  be subsets of  $V$  with cardinality  $|C|=n$  and  $|L|=m$ .  $C$  is the set of demand points or clients, and  $L$  is the set of potential locations for facilities. Let  $d(c,x)$  be the distance between point  $c \in C$  and point  $x \in L$ . For  $c \in C$  and  $X$  subset of  $L$ ,  $d(c,X)$  denotes the distance between  $c$  and  $X$ . Every point  $c \in C$  has a weight  $w(c)$  which represents the demand at point  $c$ . Denote  $C=\{c_k; k \in [1..n]\}$  and  $L=\{l_i; i \in [1..m]\}$ ,  $w_k=w(c_k)$ ,  $d_{ki}=d(c_k, l_i)$ , and, for  $X$  subset of  $L$ ,  $d_k(X)=d(c_k, X)$ . Consider a market of essential goods, which means that the sum of demands served by the firms operating in the market is equal to the total existing demand.

Assume that two competing firms, A and B, operate in the market with  $p$  and  $r$  facilities located at  $X_p$  and  $Y_r$ , respectively. The demand at point  $c_k$  captured by the firms depends on the difference  $d_k=d_k(Y_r)-d_k(X_p)$ . The market share for firms A and B are given by

$$W_A = W_A(X_p, Y_r) = \sum_{k=1}^n w_k (1 - f_k(d_k))$$

$$W_B = W_B(X_p, Y_r) = \sum_{k=1}^n w_k f_k(d_k)$$

where  $f_k(d)$  is a non-negative and non-increasing function such that  $0 \leq f_k(d) \leq 1$  for  $d \geq 0$ . The total demand is  $W_T=W_A+W_B$ .

Initially, no firm is operating in the market, firm A, the leader, wants to enter the market with  $p$  facilities taking into account that firm B, the follower, will enter later the market installing  $r$  facilities in the locations where the market share of B is maximum. Firm A wants to determine the  $p$  locations that minimize the demand captured by the competing firm B. As goods are essential, to minimize the demand captured by the competitor is equivalent to maximize the own market share.

If the leader has  $p$  facilities opened at  $X_p$ , the problem of the follower is to determine the set  $Y_r$  of  $r$  locations that maximizes its market share  $W_B(X_p, Y_r)$ . An optimal solution to this problem,  $Y_r^*(X_p)$ , is an ( $r|X_p$ )-medianoid. The problem of the leader is to determine the set  $X_p$  that minimizes  $W_B(X_p, Y_r^*(X_p))$ , that is, the set  $X_p$  which minimizes the maximum market share that the follower could achieve. An optimal solution to the problem of the leader is an ( $r|p$ )-centroid. Formally, the ( $r|p$ )-centroid problem or leader's problem is the following minimax problem

$$\min_{X \subset L, |X|=p} \max_{Y \subset L, |Y|=r} W_B(X, Y).$$

That is,

$$\min_{X \subset L, |X|=p} S(X)$$

where

$$S(X) = \max_{Y \subset L, |Y|=r} W_B(X, Y)$$

This is a bi-level problem where the lower level problem is the ( $r|X_p$ )-medianoid problem and the upper level problem is the ( $r|p$ )-centroid problem.

## Linear formulations

### The ( $r|X_p$ )-medianoid problem:

$$\max \sum_{i=1}^m \sum_{k=1}^n h_{ki} z_{ki}$$

$$\sum_{i=1}^m y_i = r$$

$$\sum_{i=1}^m z_{ki} \leq 1, \quad k \in [1..n]$$

$$z_{ki} \leq y_i, \quad i \in [1..m], k \in [1..n]$$

where  $y_i=1$  if the follower opens a facility at point  $l_i$  and  $y_i=0$  otherwise, and  $z_{ki}=1$  if client  $c_k$  visits a facility of the follower at point  $l_i$  and  $z_{ki}=0$  otherwise. The coefficient  $h_{ki}$  in the objective function is

$$h_{ki} = w_k (1 - f_k(d_k)).$$

### The ( $r|p$ )-centroid problem:

$$\min W$$

$$\sum_{i=1}^m x_i = p$$

$$\sum_{i=1}^m \sum_{k=1}^n h_{ki} u_{ki} \leq W, \quad j \in [1..(m \choose r)]$$

$$\sum_{i=1}^m u_{ki} = 1, \quad k \in [1..n]$$

$$u_{ki} \leq x_i, \quad i \in [1..m], k \in [1..n]$$

$$u_{ki}, x_i \in \{0,1\}, \quad i \in [1..m], k \in [1..n]$$

where  $x_i=1$  if the leader opens a facility at point  $l_i$  and  $x_i=0$  otherwise,  $u_{ki}=1$  if client  $c_k$  visits a facility of the leader at point  $l_i$  and  $u_{ki}=0$  otherwise, and

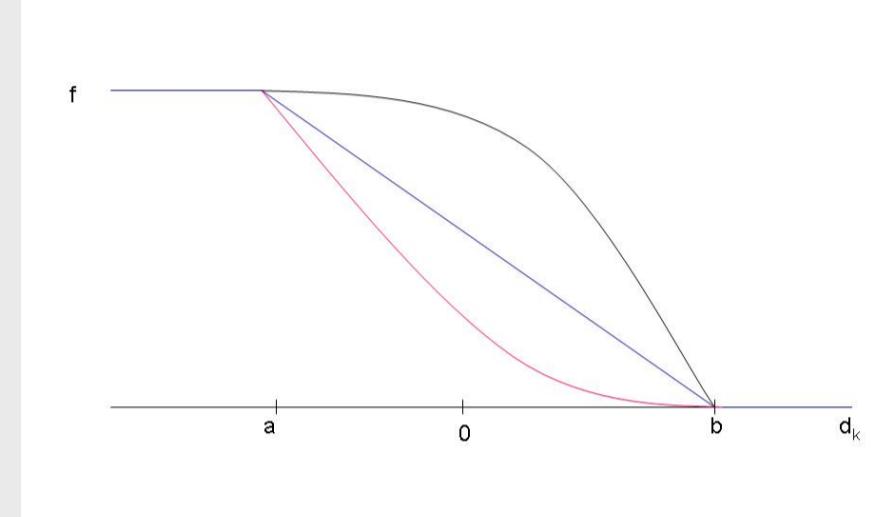
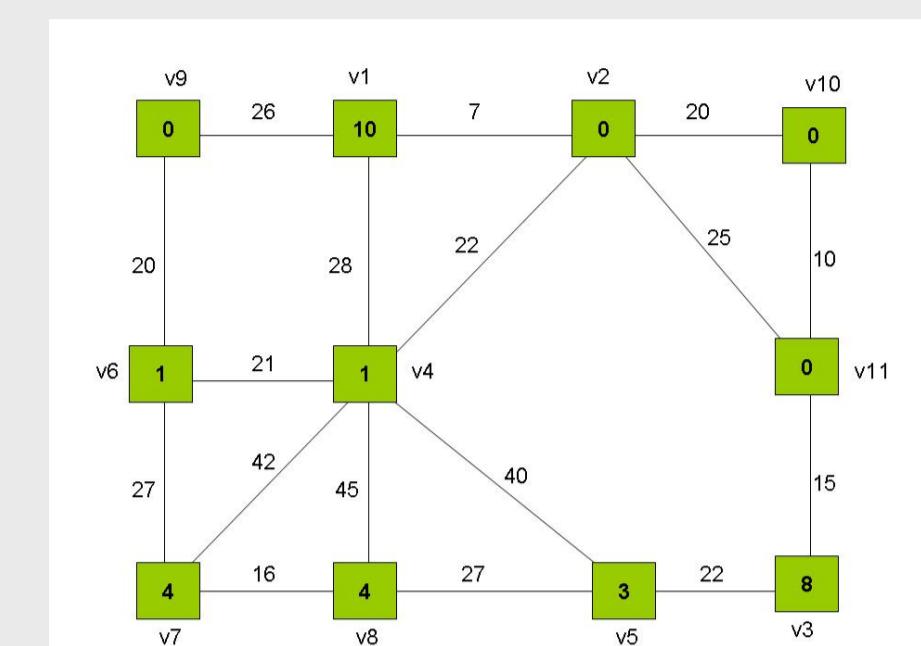
$$h_{kj} = w_k f_k(d_k(y_j) - d_{kj})$$

$$Y_j \in L_r = \{Y \subset L: |Y| = r\}, j \in [1..(m \choose r)]$$

$L_r$  is the set of feasible solutions for the follower.

The value  $h_{kj}$  represents the demand at  $c_k$  captured by the follower if he/she opens facilities at  $Y_j$  and the closest leader's facility is located at  $l_i$ .

## An example



Rule ( $p=r=2$ )	Coincidences allowed		Coincidences not allowed		X=Y
	X,Y	W=S(X)	X,Y	W=S(X)	W=S(X)
Binary ( $\mu=0.5$ )	1,7; 1,3	16,50	1,3; 5,6	13	15,50
$a+b=2,86$					
Linear	1,7; 1,11	16,50	1,3; 5,6	13	15,50
Concave	1,3; 1,3	23,25	1,3; 5,6	13	23,25
Convex	1,3; 1,5	14	1,3; 5,6	13	7,75
$a=-2,86, b=40$					
Linear	1,3; 1,3	28,93	1,5; 2,8	22,72	28,93
Concave	1,3; 1,3	30,88	1,2; 5,8	26,23	30,88
Convex	1,3; 1,3	27,00	1,5; 2,8	18,31	27,00
$a=-40, b=2,86$					
Linear	1,3; 5,8	9,88	1,3; 5,8	9,88	2,07
Concave	1,3; 5,6	11,16	1,3; 5,6	11,16	0,14
Convex	2,5; 3,7	7,30	2,5; 3,7	7,30	4,00

## A solution procedure

### Step 1. Initialization

1.1. Select  $s$  feasible leader's solutions  $X_i$ ,  $i=1, \dots, s$ . Solve the follower's problem for  $X_i$ ,  $i=1, \dots, s$ . An upper bound of the optimum  $W^*$  is  $W_{UP} = \min_i S(X_i)$ . Let  $X^*=X$  with  $S(X)=W_{UP}$ .

1.2. Let  $F$  be the selected family of good follower candidates,  $F=\{Y_j\}_{j \in L_s}$ . Set  $W_{LO}=0$ .

1.3. MaxIte= maximum number of iterations. Let  $i=0$ .

### Step 2. Iterations

Repeat, until  $W_{UP}=W_{LO}$  or until  $i=MaxIte$ , the steps:

2.1. Do  $i=i+1$ .

Solve the leader's problem constrained to  $F$ . Let  $X$  be the optimal solution obtained. If the optimal value obtained  $S_c(X)$  verifies  $S_c(X)>W_{LO}$  then do  $W_{LO}=S_c(X)$ . If  $W_{LO}=W_{UP}$ , then  $W^*=W_{LO}=W_{UP}$  is the optimal value and  $X^*=X$  is the optimal location set for the leader.

2.2. Solve the follower's problem for  $X$ . If  $S(X)<W_{UP}$  then set  $W_{UP}=S(X)$  and  $X^*=X$ . If  $W_{LO}=W_{UP}$ , then  $W^*=W_{LO}=W_{UP}$  is the optimal value and  $X^*$  is the optimal location set for the leader. Add  $Y(X)$  to  $F$ .

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