



# Core-Local Interrupt Controller (CLIC) RISC-V Privileged Architecture Extensions

Version 0.9-draft, 10/10/2023: This document is in the Development state. Assume anything can change. See <https://wiki.riscv.org/display/HOME/Specification+States>

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# Preamble

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# Chapter 1. Revision History

Date	Description
10/10/2023	issue #307 - NUM_INTERRUPT parameter text cleanup
10/10/2023	issue #355 - updated text clarifying clic interaction with rnmi spec
10/10/2023	issue #347/356/357/358 - updated mret/inhv pseudo-code
09/01/2023	issue #339 - typo changed xstatus.il to xintstatus.il
08/29/2023	issue #345 - stateen register required to block access to new ssclic CSRs?
08/29/2023	issue #350 - clarify which privilege modes will have mscratchcsw
08/29/2023	issue #351 - Clarify/Fix Smclic Ssclic Suclic memory map reserved areas
08/01/2023	issue #333 - xcause.xinhv - x is the faulting priv, i.e., when set, xinhv indicates is xepc is addr of a table entry.
06/20/2023	issue #339 - updated text describing mcause.pil and xret behavior to better match wording in priv spec
06/06/2023	issue #334 - removed text regarding debug changes to xintthresh as this does not apply #334
05/30/2023	spelling fixes.
05/09/2023	issue #317 - clarification that trampoline examples do not account for f or v registers
05/09/2023	issue #322 - xcliccfg xnlbits text cleanup
05/09/2023	issue #321 - intthresh clearing/zeroing text changed to setting to min value
05/09/2023	issue #320 - RNMI CSR CLIC details added
04/11/2023	issue #318 - use zero instead of x0 in assembly examples.
04/11/2023	issue #290 - force xepc to have table-entry alignment on xRET during inhv
03/28/2023	issue #311 - provided separate unlbits/snlbits/mnlbits in separate cliccfg registers per mode. This changes the cliccfg bit ordering.
03/28/2023	issue #304 - clean up VM page boundary alignment recommendations
03/28/2023	issue #295 - also include DRET in clearing intthresh
03/28/2023	issue #315 - typo fix in interrupt handler example
03/14/2023	issue #309 - Added discussion on NMI and RNMI handling.
03/14/2023	issue #295 - MRET/SRET clears current priv intthresh when going to a lower priv mode.
02/28/2023	issue #305 - clean up xideleg and xedeleg text
02/14/2023	pull #302 - make clicintrig, xnxti/xscratchcsw/xscratchcswl non-optional
02/14/2023	issue #293 - create separate cliccfg per priv mode
02/14/2023	issue #294 - Change addresses of mintstatus, sintstatus, uintstatus to 0xFB1, 0xDB1, and 0xCB1, respectively.
02/14/2023	issue #298 - typo cleanup - removed redundant note text in smclicshv section
01/31/2023	issues #75/#160 - reordered text into 5 extensions, smclic, ssclic, suclic, smclicshv, smclicconfig. No functional changes intended.

11/08/2022 issue #271 - text cleanup - updated inhv to xinhv

11/08/2022 issue #280 - clarified nxti CSR access types

11/08/2022 issue #235 - clarified inhv text for statement when a trap is taken all cause fields are updated.

11/08/2022 issue #88/#282 - moved xintstatus to read-only CSR addr range, moved clicbase, nvbits, clicinfo to parameters section

10/25/2022 issue #274 - clarify handling of horizontal traps on table fetch in pseudo-code

10/25/2022 issue #277 - be consistent in calling CLIC inputs local interrupts.

10/11/2022 issue #275 - clarify hwvector text

10/11/2022 issue #279 - Reserved use of all but simple csrrw access to scratchcsw/l, clarified operation of the instruction

10/11/2022 issue #277 - Added more clarification on csip

09/27/2022 issue #271/#272 - xinhv text cleanup. setting ucause\_inhv in pseudo-code when u-mode not implemented.

09/27/2022 issue #240/#255 - clarify CLIC vs CLINT mode settings

09/13/2022 issues #219/#222 - CLIC interrupt ordering text clarifications.

09/05/2022 issue #267 - update text from not defined to implementation-defined.

08/30/2022 clarify hw vectoring execute permissions,  
fixed text to say priv level instead of interrupt level.  
changed implicit read to implicit fetch.

08/30/2022 issue #219 - csip interrupt ordering clarification - CSIP interrupt ID was changed from 12 to 16

08/30/2022 issue #229 - clarify clicintattr.mode WARL behavior

08/30/2022 Define CLINT and replace references to "original basic local interrupts" with CLINT

08/30/2022 issue #191 - software vectoring read permission clarification

08/30/2022 issue #239/228 - clarification of breakpoints on hw vector table fetches, dpc

08/16/2022 issue #202 - 64-bit writes to  
{clicintctl,clicintattr,clicintie,clicintip} text clarification

08/02/2022 issue #250 - Clarified that not all specified CSRs are available in all privilege modes

08/02/2022 issue #248 - created a new xtvec submode field in clic mode

08/02/2022 issue #100 - reserving use of uimm bits in xnxti for future use

07/05/2022 Specified that xtvec.mode bits are writeable but hidden when in CLINT mode

06/21/2022 clarified "cleared" means set to 0 for interrupt pending bit

06/21/2022 issue #220 - reserved address space clarification (pull #243)

06/21/2022 issue #214 - xscratch pseudocode clarification (pull #215)

06/21/2022 issue #197 - Clarified xinhv pseudocode (pull #198)

06/21/2022 Made clear that mtvec[5:0]=000010 is still reserved

06/07/2022 pull #217 - allow implementing less than 8 bits for xintthresh

06/07/2022 issue #29/#155 (pull#190) - clarify clicinttrig details

06/07/2022 issue #212/pull#216 - fixed parameter value ranges for NUM\_INTERRUPT and CLICMTVECALIGN

06/07/2022 Pull #218 - typo fix. clicintattr regs are used to delegate interrupts

06/01/2022 Wording change in comparison with AIA features. Added reference to Bibliography.

05/10/2022 issue #235 - change “exception” to “trap” to match priv spec wording.

05/10/2022 issue #233 - mnxti pseudo-code clarification (added meaning of clic.priv,clic.level,clic.id)

05/10/2022 issue #225 - bounded time to respond to interrupts

04/26/2022 issue #191 - hw vector fetch permission changed to implicit read with execute permission required.

04/26/2022 issue #223/224 - mtval=0 allowed, hw vect xepc difference noted.

03/15/2022 issue #207 - further xret/inhv text clarification

03/06/2022 issue #210 - hw vector trap text clarification

02/15/2022 WFI text clarification

02/01/2022 issue #193 - xret/inhv text clarification

01/04/2022 issue #45 - remove new alignment constraint on CLINT mode when CLIC added

01/04/2022 issue #188 - clarification that writes to xcause affect xstatus

12/21/2021 issue #109 - add smclic arch string to spec

12/21/2021 issue #180 - change processor references to hart

11/09/2021 issue #48 - indicate when edge-triggered interrupts are cleared

11/09/2021 issue #179 - set interrupt bit during nxti access

10/28/2021 issue #154 - inhv clarification

10/28/2021 issue #31/#120 - wfi clarification

10/12/2021 issue #177 - Reduced mandatory reset requirements

09/29/2021 Added link to development states definition on top page

09/14/2021 pull #169 - nxti clarification

09/14/2021 pull #168 - only 0 or 8 level bits currently supported (other values reserved)

09/14/2021 issue #170 - clarified position of intthresh in CSR

08/31/2021 issue #86/#165 - Update mnxti pseudo-code to handle side-effects correctly.

08/31/2021 pull #164 - moved clicintattr.mode reset value to reset section of spec

08/17/2021 pull #163 - spec clarification that clicintie is held in bit 0 of byte.

07/20/2021 pull #161 - spec clarification that only writes to xnxti have side effects.

07/06/2021 issue #156,#77,#79 - more CLIC memory mapped text clarifications, clicintctl typo fixes

06/22/2021 issue #156 - reverted text and added clarification on CLIC memory mapped privilege regions.

05/25/2021 issue #149 - added text that 32-bit writes are legal but effects are not defined.

05/25/2021 issue #142 - added text that MPRV and SUM are obeyed on vector table accesses.

05/11/2021 issue #154 - added text that clarifies behavior when inhv is set when returning from a ret instruction.

04/27/2021 clicintip[i] state is undefined when switching from level to edge

triggered mode

04/22/2021 updated adoc format to align with risc-v template, added revision history

04/18/2021 Added Bibliography section

04/15/2021 issue #45 - for rev1.0 mtvec not xtvec controls enabling CLIC mode for all priv

04/13/2021 issue #141 - N-extension vs Bare S-mode note added.

04/13/2021 issue #117,#125 fix - change text to match table in M/S/U system if nmbits==1

04/12/2021 issue #47 fix - add CLIC reset behavior section

04/12/2021 issue #26 fix - modify wording that defined micro-architectural behavior of xINHV

04/12/2021 issue #91 - add DTS entry example

04/12/2021 added CLIC comparison to Advance Interrupt Architecture (AIA)

04/12/2021 issue #111,#105 fix - For hardware vectoring access exceptions, both {tval} and {epc} holds the faulting address

04/08/2021 issue #49, #79 - downplay M/S/U memory map requirements

03/30/2021 issue #29 - updated memory map table reserved section to give room for clicintrig

03/30/2021 issue #122 fix - remove wording referring to register

03/11/2021 issue #120 - update WFI wording

03/11/2021 typo fixing

03/11/2021 issue #51 - implementation of non CSRRW variants of xscratchcsw/xscratchcswl explicitly not defined/reserved.

03/11/2021 issue #58 - xintthresh was missing from table summarizing overall interrupt behavior

02/17/2021 issue #95 fix - removed N extension reference since not ratified.

02/17/2021 issue #90 fix - clarified that clicintip!=0 means interrupt pending

02/17/2021 issue #89 - updated CLIC interrupt ID ordering recommendations

02/17/2021 ihnv clarification - inhv bit has no effect except when returning from a trap using an {ret} instruction

02/17/2021 ihnv clarification - inhv only written by hw during table vector read. can be written by software.

02/02/2021 WFI wording change

01/19/2021 WFI wording change

01/07/2021 WFI section added

01/07/2021 Notes added clarifying clicintie and mstatus.xie

01/07/2021 interrupt priority clarification

12/17/2020 Added support for interrupt triggers

10/20/2020 clarified differences between level and priority

10/20/2020 fixed value range for CLICINTCTLBITS

10/20/2020 Clarified relationship among interrupt level, cliccfg.nlbits and CLICINTCTLBITS

09/08/2020 clarified description for interrupt level



## Chapter 2. Background and Motivation

The Core-Local Interrupt Controller (CLIC) Privileged Architecture Extensions are designed to provide low-latency, vectored, pre-emptive interrupts for RISC-V systems. When activated the CLIC subsumes and replaces the original RISC-V basic local interrupt scheme. The CLIC has a base design that requires minimal hardware, but supports additional extensions to provide hardware acceleration. The goal of the CLIC is to provide support for a variety of software ABI and interrupt models, without complex hardware that can impact high-performance implementations.

The CLIC also supports a Selective Hardware Vectoring extension that allow users to optimize each interrupt for either faster response or smaller code size.



While the current CLIC provides only hart-local interrupt control, future extensions might also support directing interrupts to harts within a core, hence the name (also CLIC sounds better than HLIC or HIC).



CLIC only replaces the original RISC-V basic local interrupt scheme. Exception behavior is unchanged.

### 2.1. Original RISC-V basic local Interrupts (CLINT mode)

The RISC-V Privileged Architecture specification defines CSRs such as `xip`, `xie` and interrupt behavior. A simple interrupt controller that provides inter-processor interrupts and timer functionalities for this RISC-V interrupt scheme has been called CLINT. This specification will use the term CLINT mode when `xtvec.mode` is set to either `00` or `01`.

CLINT mode supports interrupt preemption, but only based on privilege mode. At any point in time, a RISC-V hart is running with a current privilege mode. The global interrupt enable bits, `xstatus.xie`, control whether interrupts can be taken for the current or higher privilege modes; interrupts are always disabled for lower-privileged modes. Any enabled interrupt from a higher-privilege mode will stop execution at the current privilege mode, and enter the handler at the higher privilege mode. Each privilege mode has its own interrupt state registers, e.g. `mepc/mcause` for M-mode and `sepc/scause` for S-mode, to support preemption, or generically `xepc/xcause` for privilege mode `x`. Preemption by a higher-privilege-mode interrupt also pushes current privilege mode and interrupt enable status onto the `xpp` and `xpie` stacks in the `xstatus` register of the higher-privilege mode.

The `xtvec` register specifies both the interrupt mode and the base address of the interrupt vector table. The low bits of the WARL `xtvec` register indicate what interrupt model is supported. The CLINT mode settings of `xtvec` mode (`*00` and `*01`) indicate use of the basic interrupt model with either non-vectored or vectored transfer to a handler function, with the 4-byte (or greater) aligned table base address held in the upper bits of `xtvec`.



WARL means "Write Any, Read Legal" indicating that any value can be attempted to be written but only some supported values will actually be written.



The settings of `xtvec` mode with the value of `11` and a newly defined `xtvec` submode field with the value of `0000` indicate CLIC modes instead of CLINT modes. Refer to the `xtvec` section in this specification for details.

## 2.2. CLIC mode compared to CLINT mode

A CLINT mode interrupt controller is a small unit that provides local interrupts and manages the software, timer, and external interrupt signals (`xsip`/`xtip`/`xeip` signals in the `xip` register). This basic controller also allows additional custom fast local interrupt signals to be added in bits 16 and up of the `xip` register.

Priority for local interrupts is fixed. `xtvec` mode can be set so that all interrupts are direct and set the pc to the same vector base address. `xtvec` mode can also be set so that all interrupts are vectored using a vector table filled with jump instructions.

CLIC allows software to control interrupt mode, trigger type, priority, and a CLIC mode vectoring behavior for each individual interrupt. The CLIC mode vector table holds addresses so does not have the +/-1MiB jump instruction limitation. CLIC adds support for same privilege level interrupt preemption (horizontal interrupts) and additional support to reduce the number of memory or CSR accesses within an interrupt handler.

Platform profiles may require support for either or both of the CLINT and CLIC interrupt modes.

## 2.3. CLIC compared to PLIC

The standard RISC-V platform-level interrupt controller (PLIC) provides centralized interrupt prioritization and routes shared platform-level interrupts among multiple harts, but sends only a single external interrupt signal per privilege mode to each hart.

The PLIC routing scheme uses a notification/claim/response/completion sequence to route interrupts to individual harts which requires additional interrupt handler memory accesses.

The CLIC complements the PLIC. Smaller single-core systems might have only a CLIC, while multicore systems might have a CLIC per-core and a single shared PLIC. The PLIC `xeip` signals are treated as hart-local interrupt sources by the CLIC at each core.

## 2.4. CLIC compared to Advanced Interrupt Architecture

Advanced interrupt Architecture (AIA) supports message-signaled interrupts (MSIs) and an Advanced PLIC (APLIC) and targeted to support multiple harts, and support for virtualization. Like CLIC, the relative priority of all interrupts (not just external) can be configured. CLIC is targeted at CLIC per-core and has the option to give each interrupt source a separate trap entry address, preemption (nesting) of interrupts with adjustable priority threshold control, and support for reduced context switching with back-to-back interrupts.

## Chapter 3. CLIC Overview

This section gives an overview for the Core-Local Interrupt Controller (CLIC) that receives interrupt signals and presents the next interrupt to be processed by the hart.

The CLIC supports up to 4096 interrupt inputs per hart. Each interrupt input  $i$  has four 8-bit memory-mapped control registers: an interrupt-pending bit (`clicintip[i]`), an interrupt-enable bit (`clicintie[i]`), interrupt attributes (`clicintattr[i]`) to specify privilege mode and trigger type, and interrupt control bits to specify level and priority (`clicintctl[i]`).

When the first 16 interrupt inputs are reserved for the CLINT mode interrupts present in the low 16 bits of the `xip` and `xie` registers, up to 4080 local interrupts can be added.

### 3.1. Interrupt Preemption

The CLIC extends interrupt preemption to support up to 256 interrupt levels for each privilege mode, where higher-numbered interrupt levels can preempt lower-numbered interrupt levels. Interrupt level 0 corresponds to regular execution outside of an interrupt handler. Levels 1—255 correspond to interrupt handler levels. Platform profiles will dictate how many interrupt levels must be supported.

Incoming interrupts with a higher interrupt level can preempt an active interrupt handler running at a lower interrupt level in the same privilege mode, provided interrupts are globally enabled in this privilege mode.



Existing RISC-V interrupt behavior is retained, where incoming interrupts for a higher privilege mode can preempt an active interrupt handler running in a lower privilege mode, regardless of global interrupt enable in lower privilege mode.

### 3.2. CLIC Interaction with Other Local Interrupts

The CLIC subsumes the functionality of the basic local interrupts previously provided in bits 16 and up of `xip`/`xie`, so these are no longer visible in `xip`/`xie`.

The existing timer (`mtip`/`stip`), software (`msip`/`ssip`), and external interrupt inputs (`meip`/`seip`) are treated as additional local interrupt sources, where the privilege mode, interrupt level, and priority can be altered using memory-mapped `clicintattr[i]` and `clicintctl[i]` registers.



In CLIC mode, interrupt delegation for these signals is achieved via changing the interrupt's privilege mode in the CLIC Interrupt Attribute Register (`clicintattr`), as with any other CLIC interrupt input.

## Chapter 4. smcllc M-mode CLIC extension

### 4.1. CLIC Memory-Mapped Registers

#### 4.1.1. CLIC Memory Map

Each hart has a separate CLIC accessed by a separate address region. The M-mode CLIC memory map region must be made accessible to the M-mode software running on the hart.



A bus memory map or locked PMP entries could prevent M-mode software on a particular hart from reaching the CLIC memory map.



For reserved memory regions, specific trap behavior is not specified. Depending on system bus architecture, the system can ignore the access (e.g., read zero/write ignored) or cause a bus error (usually imprecise interrupt), or some other platform-specific behavior. The "reserved" annotation here implies that future standards might place additional standard registers in that space, and so using the space for non-standard features is inadvisable.

The base address of M-mode and the base addresses of any other privilege mode CLIC memory-mapped registers is specified via the general RISC-V discovery mechanism that is in development. See the CLIC Parameters section for additional detail.

The CLIC memory map supports up to 4096 total interrupt inputs.

## M-mode CLIC memory map

## Offset

###	0x0004-0x003F	reserved	###
###	0x00C0-0x07FF	reserved	###
###	0x0800-0x0FFF	custom	###

0x0000	4B	RW	reserved for smclicconfig extension
0x0040	4B	RW	clicinttrig[0]
0x0044	4B	RW	clicinttrig[1]
0x0048	4B	RW	clicinttrig[2]
...			
0x00B4	4B	RW	clicinttrig[29]
0x00B8	4B	RW	clicinttrig[30]
0x00BC	4B	RW	clicinttrig[31]
0x1000+4*i	1B/input	R or RW	clicintip[i]
0x1001+4*i	1B/input	RW	clicintie[i]
0x1002+4*i	1B/input	RW	clicintattr[i]
0x1003+4*i	1B/input	RW	clicintctl[i]
...			
0x4FFC	1B/input	R or RW	clicintip[4095]
0x4FFD	1B/input	RW	clicintie[4095]
0x4FFE	1B/input	RW	clicintattr[4095]
0x4FFF	1B/input	RW	clicintctl[4095]



Discovery mechanisms are still in development.

8b, 16b, and 32b stores to CLIC memory-mapped registers are atomic, however, there is no specified order in which the effects of the individual field updates take effect. For RV64, naturally aligned 64-bit memory accesses to the CLIC memory-mapped registers are additionally supported but 64b accesses can be broken into two 32b accesses in any order.

If an input  $i$  is not present in the hardware, the corresponding `clicintip[i]`, `clicintie[i]`, `clicintattr[i]`, `clicintctl[i]` memory locations appear hardwired to zero.

All CLIC-memory mapped registers are visible to M-mode.

The intent is that only the necessary address regions are made accessible to each privilege mode using the system's standard memory protection mechanisms. This can be done either using PMPs in microcontroller systems, or page tables (and/or PMPs) in harts with virtual memory support.

The location of the M-mode CLIC region is specified by the platform specification and made visible via the discovery mechanism for that platform.

The CLIC specification does not dictate how CLIC memory-mapped registers are split between privilege regions as well as the layout of multiple harts as this is generally a platform issue and each platform needs to define a discovery mechanism to determine the memory map locations. Some considerations for platforms to consider are

selecting regions that allow for efficient PMP and virtual memory configuration. For example, it may be desired that the base of each privilege mode CLIC region is naturally aligned to a virtual memory page (4KiB) so they can be mapped through the TLBs.

### 4.1.2. Specifying Interrupt Level

A configurable number of upper bits in `clicintctl[i]` are assigned to encode the interrupt level.

Only 0 or 8 level bits are currently supported, with other values currently reserved.



In effect, this switches the control bits from being used only for level or only for priority. The design supports a wider range of level-bit settings but this is not currently being standardized.

If the actual bits assigned or implemented are fewer than 8, then these bits are left-justified and appended with 1's for the lower missing bits. The following table shows how levels are encoded for these cases.

#bits	encoding	interrupt levels
0	.....	255
1	1.....	127, 255
2	11.....	63, 127, 191, 255
3	111.....	31, 63, 95, 127, 159, 191, 223, 255
4	1111....	15, 31, 47, 63, 79, 95, 111, 127, 143, 159, 175, 191, 207, 223, 239, 255

"1" bits are available variable bits in level specification

"." bits are non-existent bits for level encoding, assumed to be 1

### 4.1.3. Specifying Interrupt Priority

The least-significant bits in `clicintctl[i]` that are not configured to be part of the interrupt level are interrupt priority, which are used to prioritize among interrupts pending-and-enabled at the same privilege mode and interrupt level. The highest-priority interrupt at a given privilege mode and interrupt level is taken first. In case there are multiple pending-and-enabled interrupts at the same highest priority, the highest-numbered interrupt is taken first.



The highest numbered interrupt wins in a tie (when privilege mode, level and priority are all identical). This is the same as in CLINT interrupt mode, but different than the PLIC.

Notice that the 8-bit interrupt level is used to determine preemption (for nesting interrupts). In contrast, the 8-bit interrupt priority does not affect preemption but is only used as a tie-breaker when there are multiple pending interrupts with the same interrupt level.

Any implemented priority bits are treated as the most-significant bits of a 8-bit unsigned integer with lower unimplemented bits set to 1. For example, with one priority bit (`p111_1111`), interrupts can be set to have priorities 127 or 255, and with two priority bits (`pp11_1111`), interrupts can be set to have priorities 63, 127, 191, or 255.

## 4.2. CLIC Interrupt Pending (`clicintip`)

Each interrupt input has a dedicated interrupt pending bit (`clicintip[i]`) and occupies one byte in the memory map for ease of access. The pending bit is located in bit 0 of the byte. Software should assume

`clicintip[i]=0` means no interrupt pending, and `clicintip[i]!=0` indicates an interrupt is pending to accommodate possible future expansion of the `clicintip` field.

The conditions for an interrupt trap to occur must be evaluated in a bounded amount of time from when an interrupt becomes, or ceases to be, pending in `clicintip`, but unlike the MIP/MIE CSRs, there is no requirement that `clicintie` or `clicintip` are evaluated immediately following an explicit store to `clicintip` or `clicintie`.

When the input is configured for level-sensitive input, the `clicintip[i]` bit reflects the value of an input signal to the interrupt controller after any conditional inversion specified by the `clicintattr[i]` field, and software writes to the bit are ignored. Software clears the interrupt at the source device.

When the input is configured for edge-sensitive input, `clicintip[i]` is a read-write register that can be updated both by hardware interrupt inputs and by software. The bit is set by hardware after an edge of the appropriate polarity is observed on the interrupt input, as determined by the `clicintattr[i]` field. Software writes can set or clear edge-triggered pending bits directly by writes to the memory-mapped register. Edge-triggered pending bits can also be cleared when a CSR instruction that accesses `xnxti` includes a write.



Software is expected to use a CSR instruction that accesses `xnxti` that includes a write to clear an edge-triggered pending bit in non-vectored mode. Additional detail on this is described in the `xnxti` section.

The value in the `clicintip[i]` is undefined when switching from level-sensitive mode to edge-triggered mode in `clicintattr[i]`.



Software cannot rely on the underlying `clicintip[i]` register bits used in edge-triggered mode to hold state while in level-sensitive mode.

## 4.3. CLIC Interrupt Enable (`clicintie`)

Each interrupt input has a dedicated interrupt-enable bit (`clicintie[i]`) and occupies one byte in the memory map for ease of access. This control bit is read-write to enable/disable the corresponding interrupt. The enable bit is located in bit 0 of the byte. Software should assume `clicintie[i]=0` means no interrupt enabled, and `clicintie[i]!=0` indicates an interrupt is enabled to accommodate possible future expansion of the `clicintie` field.



`clicintie[i]` is the individual enable bit while `xstatus.xie` is the global enable bit for the current privilege mode. Therefore, for an interrupt `i` to be enabled in the current privilege mode, both `clicintie[i]` and `xstatus.xie` have to be set.



In contrast, since `xstatus.xie` only takes effect in the current privilege mode according to RISC-V convention, an interrupt `i` from a higher privilege mode is enabled as long as `clicintie[i]` is set (regardless of the setting of `xstatus.xie` in the higher privilege modes).

## 4.4. CLIC Interrupt Attribute (`clicintattr`)

This is an 8-bit WARL read-write register to specify various attributes for each interrupt.

`clicintattr` register layout

Bits	Field
7:6	mode
5:3	reserved (WPRI 0)
2:1	trig
0	reserved for smclicshv extension (WARL 0)

The 2-bit `trig` WARL field specifies the trigger type and polarity for each interrupt input. Bit 1, `trig[0]`, is defined as "edge-triggered" (0: level-triggered, 1: edge-triggered); while bit 2, `trig[1]`, is defined as "negative-edge" (0: positive-edge, 1: negative-edge). More specifically, there can be four possible combinations: positive level-triggered, negative level-triggered, positive edge-triggered, and negative edge-triggered.



Some implementations may want to save these bits so only certain trigger types are supported. In this case, these bits become hard-wired to fixed values (WARL).

The 2-bit `mode` WARL field specifies which privilege mode this interrupt operates in.

Encoding for RISC-V privilege levels (`mstatus.mpp`)

Level	Encoding	Name	Abbreviation
0	00	User/Application	U
1	01	Supervisor	S
2	10	Reserved	
3	11	Machine	M



For security purpose, the `mode` field can only be set to a privilege level that is equal to or lower than the currently running privilege level and if interrupts are supported at that privilege level (e.g. ssclic extension, suclic extension).

## 4.5. CLIC Interrupt Input Control (`clicintctl`)

`clicintctl[i]` is an 8-bit memory-mapped WARL control register to specify interrupt level and interrupt priority.

To select an interrupt to present to the core, the CLIC hardware combines the valid bits in `clicintattr.mode` and `clicintctl` to form an unsigned integer, then picks the global maximum across all pending-and-enabled interrupts based on this value. Next, the `smclicconfig` extension defines how to split the `clicintctl` value into interrupt level and interrupt priority. Finally, the interrupt level of this selected interrupt is compared with the interrupt-level threshold of the associated privilege mode to determine whether it is qualified or masked by the threshold (and thus no interrupt is presented).



Selecting an interrupt at a high privilege mode masks any interrupt at a lower privilege mode since the higher-privilege mode causes the interrupt signal to appear more urgent than any lower-privilege mode interrupt.



### 4.5.1. Interrupt Input Identification Number

The 4096 CLIC interrupt vectors are given unique identification numbers with `xcause` Exception Code (`exccode`) values. When maintaining backward compatibility is desired, the CLINT mode interrupts retain their original cause values, while the new interrupts are numbered starting at 16.



When upgrading from an earlier CLINT mode design that had local interrupts attached directly to bits 16 and above, these local interrupts can be now attached as CLIC inputs 16 and above to retain the same interrupt IDs.

## 4.6. CLIC Interrupt Trigger (`clicintrig`)

Optional interrupt triggers (`clicintrig[i]`) are used to generate a breakpoint exception, entry into Debug Mode, or a trace action. If these registers are not implemented, they appear as hard-wired zeros.

Each interrupt trigger is a 32-bit memory-mapped WARL register with the following layout:

`clicintrig` register layout

Bits	Field
31	<code>enable</code>
30:13	reserved (WARL 0)
12:0	<code>interrupt_number</code>

The `interrupt_number` field selects which number of interrupt input is used as the source for this interrupt trigger.

The `enable` control bit is read-write to enable/disable this interrupt trigger.

This logic is intended to be used with `tmexttrigger.intctl` as described in the RISC-V debug specification.

A trigger is signaled to the debug module if an interrupt is taken and the interrupt code matches a `clicintrig[i].interrupt_number` and the associated `clicintrig[i].enable` is set.

## 4.7. CLIC CSRs

This section describes the CLIC-related hart-specific Control and Status Registers (CSRs). When in CLINT interrupt mode, the behavior is intended to be software compatible with CLINT-mode-only systems.

Unless explicitly specified differently below, CSR state bits retain their value when switching between CLIC and CLINT modes. New CLIC CSRs and new CLIC CSR fields appear to be zero for both reads and implicit reads in CLINT mode.

The interrupt-handling CSRs are listed below, with changes and additions for CLIC mode described in the following sections.

	Number	Name	Description
	0x300	mstatus	Status register
	0x303	mideleg	Interrupt delegation register (INACTIVE IN CLIC MODE)
	0x304	mie	Interrupt-enable register (INACTIVE IN CLIC MODE)
	0x305	mtvec	Trap-handler base address / interrupt mode
(NEW)	0x307	mtvt	Trap-handler vector table base address
	0x340	mscratch	Scratch register for trap handlers
	0x341	mepc	Exception program counter
	0x342	mcause	Cause of trap
	0x343	mtval	Bad address or instruction
	0x344	mip	Interrupt-pending register (INACTIVE IN CLIC MODE)
(NEW)	0x345	mnxti	Interrupt handler address and enable modifier
(NEW)	0xFB1	mintstatus	Current interrupt levels
(NEW)	0x347	mintthresh	Interrupt-level threshold
(NEW)	0x348	mscratchcsw	Conditional scratch swap on priv mode change
(NEW)	0x349	mscratchcswl	Conditional scratch swap on level change

#### 4.7.1. Changes to `xstatus` CSRs

When in CLINT interrupt mode, the `xstatus` register behavior is unchanged (i.e., backwards-compatible with CLINT mode). When in CLIC mode, the `xpp` and `xpie` in `xstatus` are now accessible via fields in the `xcause` register.

#### 4.7.2. Changes to Delegation (`xideleg`) CSRs

In CLIC mode, the `mode` field in Interrupt Attribute Register (`clicintattr[i].mode`) specifies the privilege mode in which each interrupt should be taken. If `xideleg` exists, the `xideleg` CSR ceases to have effect in CLIC mode. If `xideleg` exists, the `xideleg` CSR is still accessible and state bits retain their values when switching between CLIC and CLINT interrupt modes.

#### 4.7.3. Changes to `xie/xip` CSRs

The `xie` CSR appears hardwired to zero in CLIC mode, replaced by separate memory-mapped interrupt enables (`clicintie[i]`).

The `xip` CSR appears hardwired to zero in CLIC mode, replaced by separate memory-mapped interrupt pendings (`clicintip[i]`).

Writes to `xie/xip` will be ignored and will not trap (i.e., no access faults). `xie/xip` always appear to be zero in CLIC mode.

In systems that support both CLINT and CLIC modes, the state bits in `xie` and `xip` retain their value when switching between modes.

#### 4.7.4. New `xtvec` CSR Mode for CLIC

The CLIC interrupt-handling mode is encoded as a new state in the existing `xtvec` WARL register, where `xtvec.mode` (the two least-significant bits) is 11, and bits `xtvec[5:2]` (`xtvec.submode`) are zero. The other encodings of `xtvec.submode` are reserved for future use. The trap vector base address is specified as the upper XLEN-6 bits of `xtvec(base)` with six lower zero bits appended, which constrains alignment on a 64-byte or larger

power-of-two boundary.

#### CLIC mode xtvec register layout

Bits	Field
XLEN-1:6	base (WARL)
5:2	submode (WARL)
1:0	mode (WARL)



Systems implementing both CLIC and CLINT mode may, but are not required to, limit alignment of `mtvec` to 64-byte boundaries in both modes.

If a system supports both modes, when `mtvec.mode` is set to `11` and `mtvec.submode` is set to `0000`, all privilege modes operate in CLIC mode. In CLIC mode, `xtvec.mode` and `xtvec.submode` in lower privilege modes are writeable but appear to be `11` and `0000` respectively when read or implicitly read in that mode.

If a system supports both modes, when `mtvec.mode` is set to a CLINT mode (`mtvec.mode=00` or `mtvec.mode=01`), all privilege modes operate in CLINT mode. In CLINT mode, both bits of `xtvec.mode` are writeable in lower-privilege modes but `xtvec.mode` bit 1 appears to be `0` when read or implicitly read in that mode. `xtvec` operates as before where each privilege mode can set their CLINT mode (direct or vectored) independently.



Although future CLIC versions may allow privileges to have different `xtvec.mode` settings, for now all privilege modes must run in either CLIC mode or all privilege modes must run in non-CLIC mode. These constraints might change if there are future additions to the CLIC or other new interrupt controller specs.

(`xtvec[5:0]`)

submode	mode	Action on Interrupt
---------	------	---------------------

aaaa	00	pc := OBASE (CLINT non-vectored basic mode)
------	----	---

aaaa	01	pc := OBASE + 4 * exccode (CLINT vectored basic mode)
------	----	---

0000	11	(CLIC mode)
------	----	-------------

(non-vectored)

pc := NBASE

0000	10	Reserved
------	----	----------

xxxx	1?	(xxxx!=0000) Reserved
------	----	-----------------------

OBASE = `xtvec[XLEN-1:2]<<2` # CLINT mode vector base is at least 4-byte aligned.

NBASE = `xtvec[XLEN-1:6]<<6` # CLIC mode vector base is at least 64-byte aligned.

TBASE = `xtvt[XLEN-1:6]<<6` # Software trap vector table base is at least 64-byte aligned.

In CLIC mode, if the `smclishv` extension is not supported, all interrupts are non-vectored, where the hart jumps to the trap handler address held in the upper XLEN-6 bits of `xtvec` for all exceptions and interrupts in privilege mode `x`.

Implementations might support only one of CLINT or CLIC mode. If only basic mode is supported, writes to bit 1 are ignored and it is always set to zero (current behavior). If only CLIC mode is supported, writes to bit 1 are also ignored and it is always set to one. CLIC mode hardwires `xtvec` bits 2-5 to zero (assuming no further CLIC extensions are supported).

In CLIC mode, synchronous exception traps always jump to NBASE.

#### 4.7.5. New `xtvt` CSRs

The `xtvt` WARL XLEN-bit CSR holds the base address of the trap vector table, aligned on a 64-byte or greater power-of-two boundary. The actual alignment can be determined by writing ones to the low-order bits then reading them back. Values other than 0 in the low 6 bits of `xtvt` are reserved.

In systems that support both CLINT and CLIC modes, the `xtvt` CSR is still accessible in basic mode (but does not have any effect).

#### 4.7.6. Changes to `xcause` CSRs

In both CLINT and CLIC modes, the `xcause` CSR is written at the time an interrupt or synchronous trap is taken, recording the reason for the interrupt or trap. For CLIC mode, `xcause` is also extended to record more information about the interrupted context, which is used to reduce the overhead to save and restore that context for an `xret` instruction. CLIC mode `xcause` also adds state to record progress through the trap handling process.

mcause		
Bits	Field	Description
XLEN-1	Interrupt	Interrupt=1, Exception=0
30	(reserved for smclishv extension)	
29:28	mpp[1:0]	Previous privilege mode, same as mstatus.mpp
27	mpie	Previous interrupt enable, same as mstatus.mpie
26:24	(reserved)	
23:16	mpil[7:0]	Previous interrupt level
15:12	(reserved)	
11:0	Exccode[11:0]	Exception/interrupt code

The `mcause.mpp` and `mcause.mpie` fields mirror the `mstatus.mpp` and `mstatus.mpie` fields, and are aliased into `mcause` to reduce context save/restore code.

Note: In a straightforward implementation, reading or writing mstatus fields mpp/mpie in mcause is equivalent to reading or writing the homonymous field in mstatus.

If the hart is currently running at some privilege mode (`pp`) at some interrupt level (`pil`) and an enabled interrupt becomes pending at any interrupt level in a higher privilege mode or if an interrupt at a higher interrupt level in the current privilege mode becomes pending and interrupts are globally enabled in this privilege mode, then execution is immediately transferred to a handler running with the new interrupt's privilege mode (`x`) and interrupt level (`il`).

As stated in the RISC-V privilege specification, when a trap is taken from privilege mode `y` into privilege mode `x`, `xPIE` is set to the value of `xIE`; `xIE` is set to 0; and `xPP` is set to `y`. `xepc` is written with the virtual address of the instruction that was interrupted or that encountered the exception. Additionally in CLIC mode, interrupt level (`xpil`) is set to `xintstatus.xil` and `xcause.exccode` is written with a code indicating the event (the id of the

interrupt or exception code) that caused the trap.

For backwards compatibility in systems supporting both CLINT and CLIC modes, when switching to CLINT mode the new CLIC `xcause` state field (`xpil`) is zeroed. The other new CLIC `xcause` fields, `xpp` and `xpie`, appear as zero in the `xcause` CSR but the corresponding state bits in the `mstatus` register are not cleared.

Note: For now all privilege modes must run in either CLIC mode or all privilege modes must run in non-CLIC mode so switching to CLINT mode from CLIC mode causes `xpil` in all privilege modes to be zeroed.

when not in CLIC mode, `xcause` has the CLINT mode format.

#### 4.7.7. Next Interrupt Handler Address and Interrupt-Enable CSRs (`xnxti`)

The `xnxti` CSR can be used by software to service the next horizontal interrupt for the same privilege mode when it has greater level than the saved interrupt context (held in `xcause`pil``) and greater level than the interrupt threshold of the corresponding privilege mode, without incurring the full cost of an interrupt pipeline flush and context save/restore. The `xnxti` CSR is designed to be accessed using CSRRSI/CSRRCI instructions, where the value read is a pointer to an entry in the trap handler table and the write back updates the interrupt-enable status. In addition, writes to the `xnxti` have side-effects that update the interrupt context state.



This is different than a regular CSR instruction as the value returned is different from the value used in the read-modify-write operation.

These CSRs are only designed to be used with the CSRR (CSRRS `rd,csr,x0`), CSRRSI, and CSRRCI instructions. Accessing the `xnxti` CSR using any other CSR instruction form (CSRRW/CSRRS,`rs1!=x0`/CSRRC/CSRRWI) is reserved. Note: Use of `xnxti` with CSRRSI with non-zero uimm values for bits 0, 2, and 4 are reserved for future use.

A read of the `xnxti` CSR using CSRR returns either zero, indicating there is no suitable interrupt to service or that the system is not in a CLIC mode, or returns a non-zero address of the entry in the trap handler table for software trap vectoring.

If the CSR instruction that accesses `xnxti` includes a write, the `xstatus` CSR is the one used for the read-modify-write portion of the operation, while the `xcause` register's `exccode` field and the `xintstatus` register's `xil` field can also be updated with the new interrupt id and level. If the interrupt is edge-triggered, then the pending bit is also zeroed.



Following the usual convention for CSR instructions, if the CSR instruction does not include write side effects (e.g., `csrr t0, xnxti`), then no state update on any CSR occurs. This can be used to determine if an interrupt could be taken without actually updating `xil` and `exccode`.

The `xnxti` CSR is intended to be used inside an interrupt handler after an initial interrupt has been taken and `xcause` and `xepc` registers updated with the interrupted context and the id of the interrupt.

If the pending interrupt is edge-triggered, hardware will automatically clear the corresponding pending bit when the CSR instruction that accesses `xnxti` includes a write. However, if the CSR instruction does not include write side effects (e.g., `csrr t0, xnxti`), then no state update on any CSR occurs and thus the interrupt pending bit is not zeroed. This behavior allows software to optimize the selection and execution of interrupts using `xnxti`.

```

// Pseudo-code for csrrsi rd, mnxti, uimm[4:0] in M mode.
// clic.priv, clic.level, clic.id represent the highest-ranked interrupt
currently present in the CLIC
mstatus |= uimm[4:0]; // Performed regardless of interrupt readiness.
if (clic.priv==M && clic.level > mcause.pil && clic.level > mintthresh.th) {
    // There is an available interrupt.
    if (uimm[4:0] != 0) { // Side-effects should occur.
        // Commit to servicing the available interrupt.
        mintstatus.mil = clic.level; // Update hart's interrupt level.
        mcause.exccode = clic.id;    // Update interrupt id in mcause.
        mcause.interrupt = 1;        // Set interrupt bit in mcause.
        if (clicintattr[clic.id][1] == 1) { // If edge interrupt,
            clicintip[clic.id] = 0;        // clear edge interrupt
        }
    }
    rd = TBASE + XLEN/8 * clic.id; // Return pointer to trap handler entry.
} else {
    // No interrupt or in non-CLIC mode.
    rd = 0;
}
// When a different CSR instruction is used, the update of mstatus and the test
// for whether side-effects should occur are modified accordingly.
// When a different privileges xnxti CSR is accessed then clic.priv is compared
with
// the corresponding privilege and xstatus, xintstatus.xil, xcause.exccode are
the
// corresponding privileges CSRs.

```



Vertical interrupts to different privilege modes will be taken preemptively by the hardware, so `xnxti` effectively only ever handles the next interrupt in the same privilege mode.

In CLINT mode, reads of `xnxti` return 0, updates to `xstatus` proceed as in CLIC mode, but updates to `xintstatus` and `xcause` do not take effect.

#### 4.7.8. New Interrupt Status (`xintstatus`) CSRs

A new M-mode CSR, `mintstatus`, holds the active interrupt level for each supported privilege mode. These fields are read-only. The primary reason to expose these fields is to support debug.

```

mintstatus fields
31:24 mil
23:16 (reserved) # To follow pattern of others.
15: 8 sil if sscllic is supported
7: 0 uil if uscllic is supported

```

The `xintstatus` registers are accessible in CLINT mode for system that support both modes.

### 4.7.9. New Interrupt-Level Threshold (`xintthresh`) CSRs

The interrupt-level threshold (`xintthresh`) is a new read-write WARL CSR, which holds an 8-bit field (`th`) for the threshold level of the associated privilege mode. The `th` field is held in the least-significant 8 bits of the CSR, and zero should be written to the upper bits.

A typical usage of the interrupt-level threshold is for implementing critical sections. The current handler can temporarily raise its effective interrupt level to implement a critical section among a subset of levels, while still allowing higher interrupt levels to preempt.

The current hart's effective interrupt level would then be: `effective_level = max(xintstatus.xil, xintthresh.th)`

The max is used to prevent a hart from dropping below its original level which would break assumptions in design, and also makes it simple for software to remove threshold without knowing its own level by simply setting `xintthresh` to the lowest supported `xintthresh` value.

The interrupt-level threshold is only valid when running in associated privilege mode and not in other modes. This is because interrupts for lower privilege modes are always disabled, whereas interrupts for higher privilege modes are always enabled.

If the hart is currently running at some privilege mode `x`, an MRET or SRET instruction that changes the privilege mode to a mode less privileged than `x` also sets `xintthresh` to the lowest supported `xintthresh` value. This helps software avoid a higher privilege mode from having a non-minimum threshold while a lower privilege mode is running.



The anticipated use of threshold is to provide critical sections within code running at one privilege level, not to selectively mask interrupts before running lower-privilege code. If desired, higher-privilege-mode interrupts can be selectively disabled using local interrupt enables before switching to a lower privilege mode.



This behavior significantly reduces the hardware cost because it only needs to select one global maximum interrupt and compare with the threshold of the associated privilege mode. If higher-privilege modes could have non-minimum thresholds, hardware would have to select multiple maximum interrupts (one for the current mode and one for each higher-privilege mode) qualified by the per-mode threshold, then pick a qualified maximum interrupt with the highest privilege mode.

### 4.7.10. Scratch Swap CSR (`xscratchsw`) for Multiple Privilege Modes

To accelerate interrupt handling with multiple privilege modes, a new CSR `xscratchsw` is defined for all but the lowest privilege mode supported in a given implementation to support conditional swapping of the `xscratch` register when transitioning between privilege modes. The CSR instruction is used once at the entry to a handler routine and once at handler exit, so only adds two instructions to the interrupt code path.

These CSRs are only designed to be used with the `csrrw` instruction with neither `rd` nor `rs1` set to `x0`. Accessing the `xscratchsw` register with the `csrrw` instruction with either `rd` or `rs1` set to `x0`, or using any other CSR instruction form (CSRRWI/CSRRS/CSRRC/CSRRSI/CSRRCI), is reserved.

When using `csrrw` to access `xscratchsw`, the value written into `rd` is either `xscratch` if `xpp` is different than the current privilege mode, or `rs1` if `xpp` is the same as the current privilege mode. The `xscratch` register is only written with the original value of `rs1` if there is a privilege mode difference.



This is different than a regular CSR instruction as the value returned is different from the value used in the read-modify-write operation.



The CSR instructions are defined to always copy a result (`xscratch` or `rs1`) to the `rd` destination to simplify implementations using register renaming, and in normal use the instructions set both `rs1 = sp` and `rd = sp`.

An example of normal usage of the `xscratchcsw` CSR is as follows:

```
csrrw sp, mscratchcsw, sp
# If mpp!=M-mode, swap mscratch and stack pointer (sp)
# otherwise sp copied to sp (i.e., no change) and mscratch unchanged
```

Formal description follows:

```
csrrw rd, mscratchcsw, rs1

match cur_privilege {
  Machine => match mstatus.MPP() {
    Machine => rd = rs1; // mscratch unchanged.
    _      => t = rs1; rd = mscratch; mscratch = t; /* default: for all other
priv modes*/
  }
}
```



To avoid virtualization holes, software cannot directly read the hart's current privilege mode. The swap instruction will trap if software tries to access a given mode's `xscratchcsw` CSR from a lesser-privileged mode, so the new CSR does not open a virtualization hole.

#### 4.7.11. Scratch Swap CSR (`xscratchcswl`) for Interrupt Levels

A new `xscratchcswl` CSR is added to support faster swapping of the stack pointer between interrupt and non-interrupt code running in the same privilege mode.

```
csrrw rd, mscratchcswl, rs1

// Pseudocode operation.
if ( (mcause.pil==0) != (mintstatus.mil==0) ) then {
  t = rs1; rd = mscratch; mscratch = t;
} else {
  rd = rs1; // mscratch unchanged.
}

// Usual use: csrrw sp, mscratchcswl, sp
```



This new CSR operates similarly to `xscratchcsw` except that the swap condition is true when the interrupter and interruptee are not both application tasks or not both interrupt handlers.

As with `xscratchcsw`, these CSRs are only designed to be used with the `csrrw` instruction with neither `rd` nor `rs1` set to `x0`. Accessing the `xscratchcsw1` register with the `csrrw` instruction with either `rd` or `rs1` set to `x0`, or using any other CSR instruction form (`CSRRWI/CSRRS/CSRRC/CSRRSI/CSRRCI`), is reserved.

## 4.8. CLIC Reset Behavior

In general in RISC-V, mandatory reset state is minimized but platform specifications or company policy might add additional reset requirements. Since the general privileged architecture states that `mstatus.mie` is reset to zero, interrupts will not be enabled coming out of reset.

### 4.8.1. CLIC mandatory reset state

`xintstatus.xil` fields reset to 0. Interrupt level 0 corresponds to regular execution outside of an interrupt handler.

The reset behavior of other fields is platform-specific.

## 4.9. CLIC Interrupt Operation

This section describes the operation of CLIC interrupts.

### 4.9.1. General Interrupt Overview

At any time, a hart is running in some privilege mode with some interrupt level. The hart's privilege mode is held internally but is not visible to software running on a hart (to avoid virtualization holes), but the current interrupt level is made visible in the `xintstatus` register.

Within a privilege mode `x`, if the associated global interrupt-enable `xie` is clear, then no interrupts will be taken in that privilege mode, but a pending-enabled interrupt in a higher privilege mode will preempt current execution. If `xie` is set, then pending-enabled interrupts at a higher interrupt level in the same privilege mode will preempt current execution and run the interrupt handler for the higher interrupt level.

As with the existing RISC-V mechanism, when an interrupt or synchronous exception is taken, the privilege mode and interrupt level are modified to reflect the new privilege mode and interrupt level. The global interrupt-enable bit of the handler's privilege mode is cleared, to prevent preemption by higher-level interrupts in the same privilege mode.

The overall behavior is summarized in the following table: the Current `p/ie/il` fields represent the current privilege mode `P` (not software visible), interrupt enable `ie = (xstatus.xie & clicintie[i])` and interrupt level `L = max(xintstatus.xil, xintthresh.th)`; the CLIC `priv,level`, and `id` fields represent the highest-ranked interrupt currently present in the CLIC with `nP` representing the new privilege mode, `nL` representing the new interrupt level, and `id` representing the interrupt's id; Current' shows the `p/ie/il` context in the handler's privilege mode; `pc` represents the program counter with `V` representing the result of any hardware vectoring; `cde` represents the `xcause exccode` field; while the Previous `pp/il/ie/epc` columns represent previous context fields in `xcause` and `xepc`.

Current	CLIC	Current'	Previous
p/ie/il	priv level id	p/ie/il pc cde	pp/il/ie epc
P ? ?	nP<P ? ?	- - - - -	- - - - - # Interrupt ignored
P 0 ?	nP=P ? ?	- - - - -	- - - - - # Interrupts disabled
P 1 ?	nP=P 0 ?	- - - - -	- - - - - # No interrupt
P 1 L	nP=P 0<nL<=L ?	- - - - -	- - - - - # Interrupt ignored
P 1 L	nP=P L<nL id	P 0 nL V id	P L 1 pc # Horizontal interrupt taken
P ? ?	nP>P 0 ?	- - - - -	- - - - - # No interrupt
P e L	nP>P 0<nL id	nP 0 nL V id	P L e pc # Vertical interrupt taken

### 4.9.2. Critical Sections in Interrupt Handlers

To implement a critical section between interrupt handlers at different levels in the same privilege mode, an interrupt handler at any interrupt level can temporarily raise the interrupt-level threshold (`mintthresh.th`) to mask a subset of levels, while still allowing higher interrupt levels to preempt. Alternatively, although not recommended due to worse system impacts, it can clear the mode's global interrupt-enable bit (`xie`) to prevent any interrupts with the same privilege mode from being taken.

### 4.9.3. CLIC events that cause the hart to resume execution after Wait for Interrupt (WFI) Instruction

As described in the privileged specification, the Wait for Interrupt instruction (WFI) provides a hint to the implementation that the current hart can be stalled. The hart may optionally resume execution anytime. This section describes CLIC events that must cause the hart to resume execution.



WFI can be a NOP and not actually pause hart execution. In addition, implementations can resume execution after a WFI for any other reason.

As in the privileged specification, if an interrupt is taken while the hart is stalled, the interrupt trap will be taken on the following instruction, i.e., execution resumes in the trap handler and  $mepc = pc + 4$ . If the event that causes the hart to resume execution does not cause an interrupt to be taken, execution will resume at  $pc + 4$ .

In CLIC mode, similar to CLINT mode, events causing the hart to resume execution after a Wait for Interrupt instruction (WFI) are unaffected by the global interrupt-enable bits in `xstatus.xie` but should honor `cl McIntie[i]` and `xintthresh`.

*A pending-and-enabled interrupt  $i$  causes the hart to resume execution if interrupt  $i$*

- has a higher privilege mode than the current privilege mode and
- the interrupt priority reduction tree selects interrupt  $i$  as the maximum across all pending-and-enabled interrupts and
- the interrupt  $i$  level is not equal to 0.

*A pending-and-enabled interrupt  $i$  causes the hart to resume execution if interrupt  $i$*

- has the same privilege mode as the current privilege mode and
- the interrupt priority reduction tree selects interrupt  $i$  as the maximum across all pending-and-enabled interrupts and
- the interrupt  $i$  level is greater than  $\max(xintstatus.xil, xintthresh.th)$

*A pending-and-enabled interrupt  $i$  causes the hart to resume execution if interrupt  $i$*

- has a lower privilege mode than the current privilege mode and
- the interrupt priority reduction tree selects interrupt  $i$  as the maximum across all pending-and-enabled interrupts and
- the interrupt  $i$  level is not equal to 0.



If an implementation allows setting an interrupts level to 0, level 0 will behave as a locally disabled interrupt but can still mask lower-mode interrupts. For example, if there is a non-zero level supervisor interrupt pending and a level-zero machine interrupt pending, the machine interrupt will be the global maximum across all pending-and-enabled interrupts but interrupt level 0 implies no interrupt. So programming an interrupt level to 0 should not be used to disable interrupts. `clicintie[i]` should be used instead.



`xintthresh` only applies to the current privilege mode. There is a proposal to add a new WFMI instruction ("wait for mode's interrupts") to the privilege specification. This instruction only has to wakeup for pending-and-enabled interrupts in the current mode, and is not required to wakeup for pending-and-enabled interrupts in lower privilege modes. Pending-enabled higher privilege-mode interrupts will interrupt/wakeup as usual.

#### 4.9.4. Synchronous Exception Handling

Horizontal synchronous exception traps, which stay within a privilege mode, are serviced with the same interrupt level as the instruction that raised the exception.

Vertical synchronous exception traps, which are serviced at a higher privilege mode, are taken at interrupt level 0 in the higher privilege mode.



Traps should be avoided at any time when `xepc/xcause` are live because these CSRs will be overwritten. Software should try to back them up if needed.

#### 4.9.5. Non-Resumable Non-Maskable Interrupts

The handling of NMIs is implementation-specific, but NMIs are always handled in M-mode and can overwrite `mepc` and `mcause` of an active M-mode CLIC interrupt handler.

#### 4.9.6. Resumable Non-Maskable Interrupts



This section describes the interaction of the CLIC with the proposed new RNMI specification.

The resumable NMI (RNMI) extension adds additional `mnepc`, `mncause`, and `mnstatus` CSRs. The RNMI CSRs are separate from the CLIC CSRs and so CLIC state does not need to be saved in RNMI CSRs. This makes the design of the RNMI extension orthogonal to the hart's general interrupt scheme.

#### 4.9.7. Returns from Handlers

As stated in the RISC-V Privilege Specification, when executing an `xRET` instruction, supposing `xPP` holds the value  $y$ , `xIE` is set to `xPIE`; the privilege mode is changed to  $y$ ; `xPIE` is set to 1; and `xPP` is set to the least-privileged supported mode. `xRET` sets the `pc` to the value stored in the `xepc` register. Additionally in CLIC mode, `xRET` sets `xintstatus.xil` to `xcause.xpil`. The `xret` instruction does not modify the `xcause.xpil` field in `xcause`.

## Chapter 5. ssclic S-mode CLIC extension

The ssclic extension depends on the smclic extension.

### 5.1. ssclic CLIC Memory-Mapped Registers

#### 5.1.1. ssclic CLIC Memory Map

Supervisor-mode CLIC regions only expose interrupts that have been configured to be supervisor-accessible via the M-mode CLIC region.

Offset			
###	0x0004-0x07FF	reserved	###
###	0x0800-0x0FFF	custom	###
Layout of Supervisor-mode CLIC regions			
0x0000	4B	RW	reserved for smclicconfig extension
0x1000+4*i	1B/input	R or RW	clicintip[i]
0x1001+4*i	1B/input	RW	clicintie[i]
0x1002+4*i	1B/input	RW	clicintattr[i]
0x1003+4*i	1B/input	RW	clicintctl[i]

The location of the S-mode CLIC regions are independent of the location of the M-mode CLIC region, and their base addresses are specified by the platform specification and made visible via the discovery mechanism for that platform.



It may desired that the base of each privilege mode CLIC region is aligned to a virtual memory page (4KiB) so they can be mapped through the TLBs.

Interrupt registers `clicintip[i]`, `clicintie[i]`, `clicintattr[i]`, `clicintctl[i]` configured as M-mode interrupts are not accessible to S-mode.

In S-mode, any interrupt *i* that is not accessible to S-mode appears as hard-wired zeros in `clicintip[i]`, `clicintie[i]`, `clicintattr[i]`, and `clicintctl[i]`.

It is not intended that the interconnect to the CLIC memory-mapped interrupt regions be required to carry the privilege mode of the initiator. A possible implementation of the CLIC memory map would be to alias the same physical CLIC memory-mapped registers to different address ranges, with each address range given different permissions for each privilege mode. Interrupts configured as M-mode interrupts appear as hard-wired zeros in the S-mode address range.

### 5.2. ssclic CLIC CSRs

The interrupt-handling CSRs are listed below, with changes and additions for CLIC mode described in the following sections.

	Number	Name	Description
	0x100	sstatus	Status register
	0x104	sie	Interrupt-enable register (INACTIVE IN CLIC MODE)
	0x105	stvec	Trap-handler base address / interrupt mode
(NEW)	0x107	stvt	Trap-handler vector table base address
	0x140	sscratch	Scratch register for trap handlers
	0x141	sepc	Exception program counter
	0x142	scause	Cause of trap
	0x143	stval	Bad address or instruction
	0x144	sip	Interrupt-pending register (INACTIVE IN CLIC MODE)
(NEW)	0x145	snxti	Interrupt handler address and enable modifier
(NEW)	0xDB1	sintstatus	Current interrupt levels
(NEW)	0x147	sintthresh	Interrupt-level threshold
(NEW)	0x148	sscratchcsw	Conditional scratch swap on priv mode change
(NEW)	0x149	sscratchcswl	Conditional scratch swap on level change

### 5.2.1. State Enable

If the Smstateen extension is implemented, then bit X in mstateen0 is implemented. If bit X of a controlling mstateen0 CSR is zero, then access to the new CSRs (stvt, snxti, sintstatus, sintthresh, sscratchcsw, sscratchcswl) by S-mode or a lower privilege level results in an illegal instruction exception, except if the hypervisor extension is implemented and the conditions for a virtual instruction exception apply, in which case a virtual instruction exception is raised when in VS or VU mode instead of an illegal instruction exception.

### 5.2.2. ssclic Changes to xcause CSRs

scause

Bits	Field	Description
XLEN-1	Interrupt	Interrupt=1, Exception=0
30	(reserved for smclishv extension)	
29	(reserved)	
28	spp	Previous privilege mode, same as sstatus.spp
27	spie	Previous interrupt enable, same as sstatus.spie
26:24	(reserved)	
23:16	spil[7:0]	Previous interrupt level
15:12	(reserved)	
11:0	exccode[11:0]	Exception/interrupt code

The supervisor **scause** register has only a single **spp** bit (to indicate user/supervisor) mirrored from **sstatus.spp**

### 5.2.3. ssclic New Interrupt Status (xintstatus) CSRs

A corresponding supervisor mode, **sintstatus** CSR provides restricted views of mintstatus.

```
sintstatus fields
31:16 (reserved)
15: 8 sil
7: 0 uil if uscllc is supported
```

#### 5.2.4. ssclic Scratch Swap CSR (`xscratchcsw`) for Multiple Privilege Modes

```
csrrw rd, sscratchcsw, rs1

match cur_privilege {
  Supervisor => if sstatus.SPP() then {
    rd = rs1; // sscratch unchanged.
  } else {
    t = rs1; rd = sscratch; sscratch = t;
  }
/* Although machine-mode access to sscratchcsw is not expected to be the normal
usage, */
/* it is specified in a way that simplifies hardware. */
  Machine    => match mstatus.MPP() {
    Supervisor => t = rs1; rd = sscratch; sscratch = t;
    Machine    => rd = rs1; // sscratch unchanged.
    _          => t = rs1; rd = sscratch; sscratch = t; /* default */
  }
}
```

### 5.3. ssclic CLIC Reset Behavior



For an S-mode execution environment, the EEI should specify that `status.sie` is also reset on entry. It is then responsibility of the execution environment to ensure that is true before beginning execution in S-mode.

## Chapter 6. suclic U-mode CLIC extension

The suclic extension depends on the smclic extension and the draft N-extension. Note: The proposed N-extension would add user-mode interrupts and traps, but has not been ratified and is not currently being advanced.

### 6.1. suclic CLIC Memory-Mapped Registers

#### 6.1.1. suclic CLIC Memory Map

User-mode CLIC regions only expose interrupts that have been configured to be user-accessible via the higher privilege mode CLIC regions.

Offset			
###	0x0004-0x07FF	reserved	###
###	0x0800-0x0FFF	custom	###
Layout of user-mode CLIC regions			
0x0000	4B	RW	reserved for smclicconfig extension
0x1000+4*i	1B/input	R or RW	clicintip[i]
0x1001+4*i	1B/input	RW	clicintie[i]
0x1002+4*i	1B/input	RW	clicintattr[i]
0x1003+4*i	1B/input	RW	clicintctl[i]

The location of the U-mode CLIC regions are independent of the location of other privilege mode CLIC regions, and their base addresses are specified by the platform specification and made visible via the discovery mechanism for that platform.

Note: It may be desired that the bases of each privilege mode CLIC region is VM page (4k) aligned so they can be mapped through the TLBs.

Interrupt registers `clicintip[i]`, `clicintie[i]`, `clicintattr[i]`, `clicintctl[i]` configured as higher privilege mode interrupts are not accessible to U-mode.

Likewise, in U-mode, any interrupt *i* that is not accessible to U-mode appears as hard-wired zeros in `clicintip[i]`, `clicintie[i]`, `clicintattr[i]`, and `clicintctl[i]`.

Interrupts configured as higher privilege modes would appear as hard-wired zeros in the U-mode address range.

### 6.2. suclic CLIC CSRs

The interrupt-handling CSRs are listed below, with changes and additions for CLIC mode described in the following sections.

Number	Name	Description
0x000	ustatus	Status register
0x004	uie	Interrupt-enable register (INACTIVE IN CLIC MODE)
0x005	utvec	Trap-handler base address / interrupt mode
(NEW) 0x007	utvt	Trap-handler vector table base address
0x040	uscratch	Scratch register for trap handlers
0x041	uepc	Exception program counter
0x042	ucause	Cause of trap
0x043	utval	Bad address or instruction
0x044	uip	Interrupt-pending register (INACTIVE IN CLIC MODE)
(NEW) 0x045	unxti	Interrupt handler address and enable modifier
(NEW) 0xCB1	uintstatus	Current interrupt levels
(NEW) 0x047	uintthresh	Interrupt-level threshold
(NEW) 0x049	uscratchcswl	Conditional scratch swap on level change

### 6.2.1. suclic Changes to `xcause` CSRs

ucause		
Bits	Field	Description
XLEN-1	Interrupt	Interrupt=1, Exception=0
30	(reserved for smclicshv extension)	
29:28	(reserved)	
27	upie	Previous interrupt enable, same as <code>ustatus.upie</code>
26:24	(reserved)	
23:16	upil[7:0]	Previous interrupt level
15:12	(reserved)	
11:0	exccode[11:0]	Exception/interrupt code

The user `ucause` register has no `upp` bit as interrupts can only have come from user mode.

### 6.2.2. suclic New Interrupt Status (`xintstatus`) CSRs

A corresponding user mode CSR, `uintstatus` provides restricted views of `mintstatus`.

```
uintstatus fields
31: 8 (reserved)
7: 0 uil
```

## 6.3. suclic CLIC Reset Behavior



For an U-mode execution environment, the EEI should specify that `status.uie` is also reset on entry. It is then responsibility of the execution environment to ensure that is true before beginning execution in U-mode.



## Chapter 7. smclishv CLIC selective hardware vectoring extension

The selective hardware vectoring extension gives users the flexibility to select the behavior for each interrupt: either hardware vectoring or non-vectoring. As a result, it allows users to optimize each interrupt and enjoy the benefits of both behaviors. More specifically, hardware vectoring has the advantage of faster interrupt response at the price of slightly increasing the code size (to save/restore contexts). On the other hand, non-vectoring has the advantage of smaller code size (by sharing and reusing one copy of common code to save/restore contexts) at the price of slightly slower interrupt response.

### 7.1. smclishv Changes to CLIC Memory-Mapped Registers

#### 7.1.1. smclishv Changes to CLIC Interrupt Pending (`clicintip`)

When the input is configured for edge-sensitive input, hardware clears the associated interrupt pending bit when an interrupt is serviced in vectored mode. See additional detail on hardware clearing in the `xtvec` section.



To improve performance, when a vectored interrupt is selected and serviced, the hardware will automatically clear a corresponding edge-triggered pending bit, so software doesn't need to clear the pending bit in the service routine.

In contrast, when a non-vectored (common code) interrupt is selected, the hardware will not automatically clear an edge-triggered pending bit.

#### 7.1.2. smclishv Changes to CLIC Interrupt Attribute (`clicintattr`)

This is an 8-bit WARL read-write register to specify various attributes for each interrupt.

`clicintattr` register layout

Bits	Field
7:6	mode
5:3	reserved (WPRI 0)
2:1	trig
0	shv

The 1-bit `shv` field is used for Selective Hardware Vectoring. If `shv` is 0, it assigns this interrupt to be non-vectored and thus it jumps to the common code at `xtvec`. If `shv` is 1, it assigns this interrupt to be hardware vectored and thus it automatically jumps to the trap-handler function pointer specified in `xtvt` CSR. This feature allows some interrupts to all jump to a common base address held in `xtvec`, while the others are vectored in hardware via a table pointed to by the additional `xtvt` CSR.

### 7.2. smclishv Changes to CLIC CSRs

### 7.2.1. smclishv Changes to `xtvec` CSR Mode for CLIC

```
(xtvec[5:0])
submode mode Action on Interrupt
aaaa 00 pc := OBASE (CLINT non-vectorized basic mode)
aaaa 01 pc := OBASE + 4 * exccode (CLINT vectorized basic mode)

0000 11 (CLIC mode)
      (non-vectorized)
      pc := NBASE if clicintattr[i].shv = 0
      (vectorized)
      pc := M[TBASE + XLEN/8 * exccode] & ~1 if clicintattr[i].shv = 1

0000 10 Reserved
xxxx 1? (xxxx!=0000) Reserved

OBASE = xtvec[XLEN-1:2]<<2 # CLINT mode vector base is at least 4-byte aligned.
NBASE = xtvec[XLEN-1:6]<<6 # CLIC mode vector base is at least 64-byte aligned.
TBASE = xtvt[XLEN-1:6]<<6 # Trap vector table base is at least 64-byte
aligned.
```

In CLIC mode, writing 0 to `clicintattr[i].shv` sets interrupt `i` to non-vectorized, where the hart jumps to the trap handler address held in the upper XLEN-6 bits of `xtvec` for all exceptions and interrupts in privilege mode `x`.

On the other hand, writing 1 to `clicintattr[i].shv` sets interrupt `i` to vectorized. When these interrupts are taken, the hart switches to the handler's privilege mode, and performs the trap side effects described in this and the privileged specification (e.g. update `xintstatus`, `xcause`, `xstatus` fields including clearing `xstatus.xie`). At this time, if the associated interrupt pending bit is configured for edge-sensitive input, it is cleared by hardware. The hart then fetches an XLEN-bit handler address with permissions corresponding to the handler's mode from the in-memory table whose base address (TBASE) is in `xtvt`. The trap handler function address is fetched from `TBASE+XLEN/8*exccode`. If the fetch is successful, the hart clears the low bit of the handler address and sets the PC to this handler address. If the trap handler function address fetch is unsuccessful and a exception trap occurs, the `xinhv` bit in `xcause` of the exception handler privilege mode is set indicating that `xepc` of the exception handler privilege mode contains a trap handler function address instead of the virtual address of an instruction.

The overall effect is:

```
pc := M[TBASE + XLEN/8 * exccode] & ~1
```

```

# Vector table layout for RV32 (4-byte function pointers)
mtvt -> 0x800000 # Interrupt 0 handler function pointer
        0x800004 # Interrupt 1 handler function pointer
        0x800008 # Interrupt 2 handler function pointer
        0x80000c # Interrupt 3 handler function pointer

# Vector table layout for RV64 (8-byte function pointers)
mtvt -> 0x800000 # Interrupt 0 handler function pointer
        0x800008 # Interrupt 1 handler function pointer
        0x800010 # Interrupt 2 handler function pointer
        0x800018 # Interrupt 3 handler function pointer

```



The CLINT vectored mode simply jumps to an address in the trap vector table, while the CLIC vectored mode reads a handler function address from the table, and jumps to it in hardware.



The vector table contains vector addresses rather than instructions because it simplifies static initialization in C. More specifically, the entries in the table are simple XLEN-bit function pointers.



The hardware vectoring bit `xinhv` is provided to allow resumable traps on fetches to the trap vector table.

When a trap is taken, the `xinhv` bit is set by hardware to indicate if `xepc` is the address of a table entry or cleared by hardware to indicate if `xepc` is the address of an instruction. The `xinhv` bits are only set by hardware if an exception occurs during the table vector read operation. The `xinhv` bits can be written by software, including when hardware vectoring is not in effect. The `xinhv` bit has no effect except when returning using an `xret` instruction.

When returning from an `xret` instruction, the `xinhv` bit modifies behavior as follows:

If the `xinhv` bit is set, the hart resumes the trap handler memory access to retrieve the function pointer for vectoring with permissions corresponding to the previous privilege mode. The trap handler function address is obtained from the current privilege mode's `xepc` with the low bits of the address cleared to force the access to be naturally aligned to an XLEN/8-byte table entry. If the fetch is successful, the hart clears the low bit of the handler address and sets the PC to this handler address.

```

/* MRET pseudo-code */
function exception_handler(cur_priv, xret, pc) {
    match (xret) {
    ...
    MRET => {
        /* standard MRET behavior */
        mstatus.MIE    = mstatus.MPIE;
        mstatus.MPIE   = 1;
        cur_priv       = mstatus.MPP;
        ... /* additional standard MRET behavior */

        if (minhv) {
            /* align mepc value to XLEN/8 byte boundary by ignoring low-order bits. */
            let table_addr = mepc & inhv_pc_alignment_mask();
            if (check_fetch_permissions(table_addr) = Addr_OK) {
                next_pc = mem_read(table_addr);
            } else {
                /* take table-fetch horizontal trap */
                mcause.minhv = 1;
                /* other side effects of taking a trap */
            }
        } else { /* Standard MRET behavior - mepc becomes next_pc */
            next_pc = mepc;
            /* mcause.minhv unchanged */
        }
    }
},

/* Note, next_pc is passed to fetch unit which will ignore 1 or 2
LSBs depending to IALIGN, so the masking is not shown above.
Similarly, permission checks on the ultimate instruction fetch
are not shown here. */
}

```



The `xinhv` bit when set at `xRET` informs hardware to repeat the table fetch using the address in `xEPC` to obtain the address of the trap handler that is then written to the PC instead of directly writing `xEPC` to the PC. One of the goals of this behavior is to avoid complicating the critical code paths for handling virtual memory in the more-privileged layer. The more-privileged layer does not have to distinguish CLIC vector table reads from other forms of page faults and can handle them using exactly the same code.

For permissions-checking purposes, the memory access to retrieve the function pointer for hardware vectoring is an *implicit* fetch at the privilege mode of the interrupt handler, and requires execute permission; read permission is irrelevant.



software vectoring will need vector table read permission.

If there is an access exception on the table fetch, `xepc` is written with the faulting address. `xtval` is either set to zero or written with the faulting address.



For simpler systems, we do not require that `xtval` is written with the faulting address. For systems with demand paging, `xtval` should be written with the faulting address to simplify page-fault handling code.



Interrupted context is lost on horizontal traps during table fetch where exceptions are the same privilege level as the interrupt handler. The interesting case is vertical traps, where a more privileged layer is handling page faults or other synchronous faults for the less privileged mode vector table access. The regular code path in more privileged layer will want to use `xtval` to determine what bad virtual address to page in, but will not normally restore `xtval` when returning to faulting context (potentially after some time and other contexts have run). However, it will restore `xepc` (using `x` for more privileged mode here) before using `xret` on normal code path. This is a rationale for why both `xtval` and `xepc` are recommended to be written with the faulting address in systems with demand paging.

Memory writes to the vector table require an instruction barrier (*fence.i*) to guarantee that they are visible to the instruction fetch.

### 7.2.2. smclishv Changes to `xepc` CSRs

The `xepc` CSRs behave the same in both modes, capturing the PC at which execution was interrupted. In CLIC mode, the `xepc` CSR additionally may hold the faulting address if there is an access exception on the table fetch during hardware vectoring.

### 7.2.3. smclishv Changes to `dpc` CSR

For implicit hardware vector table fetches, whether breakpoints trap on the table read is left as an implementation option. For explicit loads used in software vectoring, watchpoints operate as normal for any load. In CLIC mode, the `dpc` CSR additionally may hold the faulting address if breakpoints are allowed to trap on the table fetch during hardware vectoring. If breakpoints are allowed to trap on the table read, `dret` should honor `xinhv`.

### 7.2.4. smclishv Changes to `xcause` CSRs

mcause		
Bits	Field	Description
XLEN-1	Interrupt	Interrupt=1, Exception=0
30	minhv	When 1, indicates mepc is the address of a table entry. When 0, indicates mepc is the address of an instruction.
29:28	mpp[1:0]	Previous privilege mode, same as mstatus.mpp
27	mpie	Previous interrupt enable, same as mstatus.mpie
26:24	(reserved)	
23:16	mpil[7:0]	Previous interrupt level
15:12	(reserved)	
11:0	Exccode[11:0]	Exception/interrupt code

#### scause with ssclic extension

Bits	Field	Description
XLEN-1	Interrupt	Interrupt=1, Exception=0
30	sinhv	When 1, indicates sepc is the address of a table entry. When 0, indicates sepc is the address of an instruction.
29	(reserved)	
28	spp	Previous privilege mode, same as sstatus.spp
27	spie	Previous interrupt enable, same as sstatus.spie
26:24	(reserved)	
23:16	spil[7:0]	Previous interrupt level
15:12	(reserved)	
11:0	exccode[11:0]	Exception/interrupt code

#### ucause with suclic extension

Bits	Field	Description
XLEN-1	Interrupt	Interrupt=1, Exception=0
30	uinhv	When 1, indicates uepc is the address of a table entry. When 0, indicates uepc is the address of an instruction.
29:28	(reserved)	
27	upie	Previous interrupt enable, same as ustatus.upie
26:24	(reserved)	
23:16	upil[7:0]	Previous interrupt level
15:12	(reserved)	
11:0	exccode[11:0]	Exception/interrupt code

For backwards compatibility in systems supporting both CLINT and CLIC modes, when switching to CLINT mode the new CLIC `xcause` state fields (`xinhv` and `xpil`) are zeroed.

Note: For now all privilege modes must run in either CLIC mode or all privilege modes must run in non-CLIC mode so switching to CLINT mode from CLIC mode causes `xinhv` and `xpil` in all privilege modes to be zeroed.

In CLIC mode, when a trap is taken, `xcause` has the CLIC format and the `xcause` fields are updated. On the other hand, when not in CLIC mode, `xcause` has the CLINT mode format.

### 7.2.5. smclshv Changes to Next Interrupt Handler Address and Interrupt-Enable CSRs (`xnxti`)

A read of the `xnxti` CSR using CSRR returns either zero, indicating there is no suitable interrupt to service or that the highest ranked interrupt is SHV or that the system is not in a CLIC mode, or returns a non-zero address of the entry in the trap handler table for software trap vectoring.



The `xtvt` CSR could be set to memory addresses such that a table entry was at address zero, and this would be indistinguishable from the no-interrupt case.

```
// Pseudo-code for csrrsi rd, mnxti, uimm[4:0] in M mode.
// clic.priv, clic.level, clic.id represent the highest-ranked interrupt
// currently present in the CLIC
mstatus |= uimm[4:0]; // Performed regardless of interrupt readiness.
if (clic.priv==M && clic.level > mcause.pil && clic.level > mintthresh.th
    && (clicintattr.shv==0) ) {
    // There is an available, non-hardware-vectorized interrupt.
    if (uimm[4:0] != 0) { // Side-effects should occur.
        // Commit to servicing the available interrupt.
        mintstatus.mil = clic.level; // Update hart's interrupt level.
        mcause.exccode = clic.id;    // Update interrupt id in mcause.
        mcause.interrupt = 1;        // Set interrupt bit in mcause.
        if (clicintattr[clic.id][1] == 1) { // If edge interrupt,
            clicintip[clic.id] = 0;        // clear edge interrupt
        }
    }
    rd = TBASE + XLEN/8 * clic.id; // Return pointer to trap handler entry.
} else {
    // No interrupt, or a selectively hardware vectored interrupt, or in non-CLIC
    // mode.
    rd = 0;
}
// When a different CSR instruction is used, the update of mstatus and the test
// for whether side-effects should occur are modified accordingly.
// When a different privileges xnxti CSR is accessed then clic.priv is compared
// with
// the corresponding privilege and xstatus, xintstatus.xil, xcause.exccode are
// the
// corresponding privileges CSRs.
```

## Chapter 8. CLIC Parameters

### 8.1. MCLICBASE

The MCLICBASE parameter provides the base address of the m-mode CLIC memory mapped registers. Its value should be configured or set up at the platform level to indicate the starting address of m-mode CLIC memory mapped registers.

Since the CLIC memory map must be aligned at a 4KiB boundary, the MCLICBASE parameter has its 12 least-significant bits hardwired to zero. It is used to inform software about the location of CLIC m-mode memory mapped registers.

Systems with CLIC memory mapped registers for additional privilege modes will provide additional xCLICBASE parameters for each of those CLIC x-mode memory mapped register regions.

### 8.2. NVBITS Parameter - Specifying Support for smclicshv Selective Interrupt Hardware Vectoring Extension

The NVBITS Parameter specifies whether the smclicshv extension is implemented or not.

When NVBITS is 0, smclicshv extension is not implemented. In this case, all CLIC interrupts are non-vectorized and are directed to the common code at `xtvec` register.

When NVBITS is 1, smclicshv extension is implemented.

### 8.3. CLICINFO Parameters

The NUM\_INTERRUPT 13-bit parameter from 2-4096 that specifies the actual number of maximum interrupt inputs supported in this implementation. MSIP, MTIP are always included.

The VERSION 8-bit parameter specifies the implementation version of CLIC. The upper 4-bit specifies the architecture version, and the lower 4-bit specifies the implementation version.

The CLICINTCTLBITS 4-bit parameter specifies how many hardware bits are actually implemented in the `clicintctl` registers, with  $0 \leq \text{CLICINTCTLBITS} \leq 8$ . The implemented bits are kept left-justified in the most-significant bits of each 8-bit `clicintctl[i]` register, with the lower unimplemented bits treated as hardwired to 1.

The NUM\_TRIGGER 6-bit parameter specifies the number of maximum interrupt triggers supported in this implementation. Valid values are 0 to 32.

### 8.4. Additional CLIC Parameters



Name	Value Range	Description
CLICANDBASIC	0-1	Implements CLINT mode also?
CLICPRIVMODES	1-3	Number privilege modes: 1=M, 2=M/U, 3=M/S/U
CLICLEVELS including 0	2-256	Number of interrupt levels
CLICMAXID	12-4095	Largest interrupt ID
INTTHRESHBITS {intthresh}.th	1-8	Number of bits implemented in
CLICCFGMBITS	0- $\text{ceil}(\lg 2(\text{CLICPRIVMODES}))$	Number of bits implemented for cliccfg.nmbits
CLICCFGGBITS	0- $\text{ceil}(\lg 2(\text{CLICLEVELS}))$	Number of bits implemented for cliccfg.nlbits
CLICMTVECALIGN	$\geq 6$	Number of hardwired-zero least significant bits in mtvec address.



These parameters are likely to be available by the general discovery mechanism that is in development.

## Chapter 9. smcllicconfig CLIC configuration extension

Hardware implementations may wish to have a single implementation support different parameterizations of clic extensions. This extension defines that programmability.

### 9.1. CLIC Memory-Mapped Registers

#### 9.1.1. CLIC Configuration (`xcliccfg` registers)

The CLIC has a single memory-mapped 32-bit global configuration register per privilege mode, `xcliccfg`, that defines how the `clicintctl[i]` registers are subdivided into level and priority fields. `mcliccfg` has an additional field that defines interrupt privilege mode configuration.



WPRI means "Writes Preserve Values, Reads Ignore Values" indicating whole read/write fields are reserved for future use. Software should ignore the values read from these fields, and should preserve the values held in these fields when writing values to other fields of the same register. For forward compatibility, implementations that do not furnish these fields must hardwire them to zero.

`mcliccfg` register layout  
located at `MCLICBASE`

Bits	Field
31:28	reserved (WPRI 0)
27:24	<code>unlbits[3:0]</code> if <code>suclic</code> is supported, else reserved (WPRI 0)
23:20	reserved (WPRI 0)
19:16	<code>snlbits[3:0]</code> if <code>sscllic</code> is supported, else reserved (WPRI 0)
15:6	reserved (WPRI 0)
5:4	<code>nmbits[1:0]</code>
3:0	<code>mnlbits[3:0]</code>

`scliccfg` register layout - dependent on `sscllic` extension  
located at `SCLICBASE`

Bits	Field
31:28	reserved (WPRI 0)
27:24	<code>unlbits[3:0]</code> if <code>suclic</code> is supported, else reserved (WPRI 0)
23:20	reserved (WPRI 0)
19:16	<code>snlbits[3:0]</code> if <code>sscllic</code> is supported, else reserved (WPRI 0)
15:0	reserved (WPRI 0)

ucliccfg register layout - dependent on suclic extension  
located at UCLICBASE

Bits	Field
31:28	reserved (WPRI 0)
27:24	unlbits[3:0]
23:0	reserved (WPRI 0)

scliccfg and ucliccfg are subsets of the mcliccfg register.



In a straightforward implementation, reading or writing any field in ucliccfg or scliccfg is equivalent to reading or writing the homonymous field in mcliccfg.

Detailed explanation for each field are described in the following sections.

#### Interrupt Mode Table

priv-modes	nmbits	clicintattr[i].mode	Interpretation
M	0	xx	M-mode interrupt
M/U	0	xx	M-mode interrupt
M/U	1	0x	U-mode interrupt
M/U	1	1x	M-mode interrupt
M/S/U	0	xx	M-mode interrupt
M/S/U	1	0x	S-mode interrupt
M/S/U	1	1x	M-mode interrupt
M/S/U	2	00	U-mode interrupt
M/S/U	2	01	S-mode interrupt
M/S/U	2	10	Reserved (or extended S-mode)
M/S/U	2	11	M-mode interrupt
M/S/U	3	xx	Reserved

### 9.1.2. Specifying Interrupt Privilege Mode

The privilege mode of an interrupt is controlled by both `cliccfg.nmbits` and `clicintattr[i].mode` as described in the Specifying Interrupt Privilege Mode section below.

The 2-bit `cliccfg.nmbits` WARL field specifies how many bits are physically implemented in `clicintattr[i].mode` to represent an input *i*'s privilege mode. Although `cliccfg.nmbits` field is always 2-bit wide, the physically implemented bits in this field can be fewer than two (depending how many interrupt privilege-modes are supported).

For example, in M-mode-only systems, only M-mode exists so we do not need any extra bit to represent the supported privilege-modes. In this case, no physically implemented bits are needed in the `clicintattr.mode` and thus `cliccfg.nmbits` is 0 (i.e., `cliccfg.nmbits` can be hardwired to 0).

In M/U-mode systems with the `suclic` extension, `cllicfg.nmbits` can be set to 0 or 1. If `cllicfg.nmbits` = 0, then all interrupts are treated as M-mode interrupts. If the `cllicfg.nmbits` = 1, then a value of 1 in the most-significant bit (MSB) of a `clicintattr[i].mode` register indicates that interrupt input is taken in M-mode, while a value of 0 indicates that interrupt is taken in U-mode.

Similarly, in systems that support `sscl` and `suclic` extensions, `cllicfg.nmbits` can be set to 0, 1, or 2 bits to represent privilege-modes. `cllicfg.nmbits` = 0 indicates that all local interrupts are taken in M-mode. `cllicfg.nmbits` = 1 indicates that the MSB selects between M-mode (1) and S-mode (0). `cllicfg.nmbits` = 2 indicates that the two MSBs of each `clicintattr[i].mode` register encode the interrupt's privilege mode using the same encoding as the `mstatus.mpp` field.

`clicintattr[i].mode` field is writable and is unchanged by writes to `cllicfg.nmbits` but the read and implicit read value is the interpretation as specified in the Interrupt Mode Table above.



Bare S-mode (no MMU, `satp=0`) can be used in microcontrollers to allow hardware delegation of interrupts out of M-mode. Bare S-mode has already been ratified as part of privileged architecture. There are also proposals to add S-mode PMP support to allow an RTOS running in S-mode to isolate itself from tasks running in U-mode.

### 9.1.3. Specifying Interrupt Level

The 4-bit `mcliccfg.xnlbits` WARL fields indicate how many upper bits in `clicintctl[i]` are assigned to encode the interrupt level at that privilege level.

Although the interrupt level is an 8-bit unsigned integer, the number of bits actually assigned or implemented can be fewer than 8. As described above, the number of bits assigned is specified in `mcliccfg.xnlbits`. The number of bits actually implemented can be derived from `mcliccfg.xnlbits` and a fixed parameter `CLICINTCTLBITS` (with value between 0 to 8) which specifies bits implemented for both interrupt level and priority.



The number of available level bits can be determined by subtracting the number of mode bits from `CLICINTCTLBITS`.

For example, if the `xnlbits` > `CLICINTCTLBITS`, then the lower bits of the 8-bit interrupt level are assumed to be all 1s. Similarly, if `xnlbits` < 8, then the lower bits of the 8-bit interrupt level are assumed to be all 1s.

If `xnlbits` = 0, then all interrupts are treated as level 255.

Examples of `mcliccfg` settings:

CLICINTCTLBITS	mnbits	clicintctl[i]	interrupt levels
0	2	.....	255
1	2	1.....	127, 255
2	2	11.....	63, 127, 191, 255
3	3	111.....	31, 63, 95, 127, 159, 191, 223, 255
4	1	1ppp....	127, 255

"." bits are non-existent bits for level encoding, assumed to be 1

"1" bits are available variable bits in level specification

"p" bits are available variable bits in priority specification

The number of bits actually implemented in `clicintctl[i]` is specified by a fixed parameter `CLICINTCTLBITS`, which has a value between 0 to 8. The implemented bits are kept left-justified in the most-significant bits of each 8-bit `clicintctl[i]` register, with the lower unimplemented bits treated as hardwired to 1. These control bits are interpreted as level and priority according to the setting in the CLIC Configuration register (`mcliccfg.xnlbits`).

## 9.2. smclconfig Changes to CLIC CSRs

### 9.2.1. smclconfig Changes to Interrupt-Level Threshold (`xintthresh`) CSRs

If the number of bits actually implemented in the `th` field is less than 8 (e.g. an implementation option when `CLICINTCTLBITS` is less than 8), the number of implemented bits `INTTHRESHBITS` must be greater than `CLICINTCTLBITS` and the implemented bits should be kept left-justified in the most-significant bits of the 8-bit field, with the lower unimplemented bits treated as hardwired to 1. For example, if `CLICINTCTLBITS` is 1 and `INTTHRESHBITS` is 2, interrupts can be set to level 127 or 255 and `xintthresh.th` can be set to 63, 127, 191, or 255.

# Chapter 10. Interrupt Handling Software

## 10.1. Interrupt Stack Software Conventions

The CLIC supports multiple nested interrupt handlers, and each handler requires some working registers. To make registers available, each handler typically saves and restores registers from the interrupted context on a memory-resident stack. In addition, the memory-resident stack is used to hold other interrupted context information, such as `xepc` and `xcause`, which are required by the `xret` instruction.

The standard RISC-V ABI convention is that stacks grow downwards, and that memory addresses below the current stack pointer can be dynamically altered by another agent, such as an interrupt handler.

When interrupts are taken horizontally within the same privilege mode, the interrupt handler may be able to use the same stack as the interrupted thread, by allocating a new stack frame below the current stack pointer.

When interrupts are taken vertically into a higher privilege mode, the stack pointer must be swapped to a stack within the higher privilege mode to avoid a security hole. The `xscratch` registers can be used to hold the stack pointer of a higher-privilege mode while lower-privilege code is executing, or `xscratch` can be used to point to more extensive thread-local context that might contain a stack pointer.

## 10.2. Inline Interrupt Handlers and "Interrupt Attribute" for C

Inline interrupt handlers are small leaf functions that handle simple interrupts. To provide easy C coding for inline interrupt handlers, while reducing register save/restore overhead, we use standard interrupt attributes, which have the following syntax:

```
/* Small ISR to poke device to clear interrupt and increment in-memory counter.
*/
void __attribute__((interrupt))
foo (void)
{
    extern volatile int INTERRUPT_FLAG;
    INTERRUPT_FLAG = 0;
    extern volatile int COUNTER;
#ifdef __riscv_atomic
    __atomic_fetch_add (&COUNTER, 1, __ATOMIC_RELAXED);
#else
    COUNTER++;
#endif
}
```

The attribute tells the C compiler to use callee-save for all registers, so the handler has to "pay as it goes" to use registers, and only save the full caller-save set if it makes a nested regular C call. The attribute also tells the C compiler to align the function entry point on an 8-byte boundary.

```

.align 3
# Inline non-preemptible interrupt handler.
# Only safe for horizontal interrupts.
foo:
    addi sp, sp, -FRAMESIZE      # Create a frame on stack.
    sw a0, OFFSET(sp)           # Save working register.
    sw zero, INTERRUPT_FLAG, a0  # Clear interrupt flag.
    sw a1, OFFSET(sp)           # Save working register.
    la a0, COUNTER               # Get counter address.
    li a1, 1
    amoadd.w zero, (a0), a1      # Increment counter in memory.
    lw a1, OFFSET(sp)           # Restore registers.
    lw a0, OFFSET(sp)
    addi sp, sp, FRAMESIZE       # Free stack frame.
    mret                         # Return from handler using saved mepc.

```

With hardware vectoring, inline interrupt handlers can provide very rapid response for small tasks.



The above entire handler executes in 13 instructions. The `INTERRUPT_FLAG` store and the `la` require two instructions each to build up a global address. A simple pipeline would encounter two pipeline flushes (on entry and on exit), plus the cycles taken to fetch the hardware vector entry.

These inline handlers can be used with the CLINT mode as well as CLIC mode.

To take advantage of hardware preemption in CLIC mode, inline handlers must save and restore `xepc` and `xcause` before enabling interrupts:

```

.align 3
# Inline preemptible interrupt handler.
# Only safe for horizontal interrupts.
foo:
#----- Interrupts disabled on entry ---#
addi sp, sp, -FRAMESIZE      # Create a frame on stack.
sw a0, OFFSET(sp)           # Save working register.
csrr a0, mcause              # Read cause.
sw a1, OFFSET(sp)           # Save working register.
csrr a1, mepc                # Read epc.
csrrsi zero, mstatus, MIE    # Enable interrupts.
#----- Interrupts enabled -----#
sw a0, OFFSET(sp)           # Save cause on stack.
sw zero, INTERRUPT_FLAG, a0  # Clear interrupt flag.
sw a1, OFFSET(sp)           # Save epc on stack.
la a0, COUNTER               # Get counter address.
li a1, 1
amoadd.w zero, (a0), a1      # Increment counter in memory.
lw a1, OFFSET(sp)           # Restore epc
lw a0, OFFSET(sp)           # Restore cause
csrrci zero, mstatus, MIE    # Disable interrupts.
#----- Interrupts disabled -----#
csrw mepc, a1                # Put epc back.
lw a1, OFFSET(sp)           # Restore a1.
csrw mcause, a0              # Put cause back.
lw a0, OFFSET(sp)           # Restore a0.
addi sp, sp, FRAMESIZE       # Free stack frame.
mret                          # Return from handler.
#-----#

```



This version requires 10 more instructions, but reduces the time a preempting interrupt has to wait from a 13-instruction window to a 6-instruction window (the instruction that disables interrupts can be preempted before committing).



This form cannot be used with CLINT mode, unless the original interrupt pending signal is cleared before re-enabling interrupts.



## Chapter 11. Calling C-ABI Functions as Interrupt Handlers

An alternative model is where all interrupt handler routines use the standard C ABI. In this case, the CLIC would use no hardware vectoring for the C ABI handlers and instead use a common software trampoline, which uses the `xnxti` instruction to obtain the trap-handler address. The code sequence below is annotated with an explanation of its operation.



Example handlers in this specification do not account for the presence of `f` or `v` registers when saving registers.

### 11.1. C-ABI Trampoline Code

```
# Example Unix C ABI interrupt trampoline.
# Only safe for horizontal interrupts.
# FRAMESIZE should be defined appropriately to hold saved context with ABI-
# specified alignment.
# OFFSET should be replaced with individual stack frame locations.
# Register save/restore pseudo-code should be expanded to individual
# instructions.

irq_enter:
#----Interrupts disabled for 7 + SREGS instructions, where SREGS is number of
# registers saved. ①
    addi sp, sp, -FRAMESIZE # Allocate space on stack. ②
    sw a1, OFFSET(sp)      # Save a1.
    csrr a1, mcause         # Get mcause of interrupted context.
    sw a0, OFFSET(sp)      # Save a0.
    csrr a0, mepc           # Get mepc of interrupt context.
    bgez a1, handle_exc    # Handle synchronous exception. ③
    sw a0, OFFSET(sp)      # Save mepc.
    sw a1, OFFSET(sp)      # Save mcause of interrupted context.
    sw a2-a7, OFFSET(sp)   # Save other argument registers.
    sw t0-t6, OFFSET(sp)   # Save temporaries.
    sw ra, OFFSET(sp)      # 1 return address ⑤
    csrrsi a0, mnxti, MIE  # Get highest current interrupt and enable interrupts.
                          # Will return original interrupt if no others appear.
⑥
#----Interrupts enabled -----⑦
    beqz a0, exit          # Check if original interrupt vanished. ⑧

service_loop:             # 5 instructions in pending-interrupt service loop.
    lw a1, (a0)           # Indirect into handler vector table for function
# pointer. ⑨
    csrrsi zero, mstatus, MIE # Ensure interrupts enabled. ⑩
```

```

jalr a1                # Call C ABI Routine, a0 has interrupt ID encoded. ⑪
                        # Routine must clear down interrupt in CLIC.

csrrsi a0, mnxti, MIE  # Claim any pending interrupt at level > mcause.pil ⑫
bnez a0, service_loop  # Loop to service any interrupt. ⑬

#--- Restore ABI registers with interrupts enabled -⑭
lw ra, OFFSET(sp)      # Restore return address
lw t0-t6, OFFSET(sp)   # Restore temporaries.
lw a2-a7, OFFSET(sp)   # Restore other arguments.
lw a1, OFFSET(sp)      # Get saved mcause,
exit:                  # Fast exit point.
lw a0, OFFSET(sp)      # Get saved mepc.

csrrci zero, mstatus, MIE # Disable interrupts ⑮
#---- Critical section with interrupts disabled -----
csrw mcause, a1         # Restore previous context.

lw a1, OFFSET(sp)      # Restore original a1 value.
csrw mepc, a0           # Restore previous context.

csrrci a0, mnxti, MIE  # Claim highest current interrupt. ⑯
bnez a0, service_loop  # Go around if new interrupt.

lw a0, OFFSET(sp)      # Restore original a0 value.
addi sp, sp, FRAMESIZE # Reclaim stack space.
mret                   # Return from interrupt.
#-----
#-----
handle_exc:
# ...
# Perform exception processing with interrupts disabled ④
# ...
addi sp, sp, FRAMESIZE # Reclaim stack space.
mret # Return from exception
#-----

```

- ① An initial interrupt (II) causes entry to the handler with interrupts disabled, and `xepc` and `xcause` CSRs hold values representing the original interrupted context (OIC), including the PC in `xepc`, the privilege mode in `xpp` (visible in both `xcause` and `xstatus`), the interrupt level in `xpil` (in `xcause`) and the interrupt enable state in `xpie` (visible in both `xcause` and `xstatus`). The `xcause` CSR and the `xintstatus` CSRs additionally hold information on the interrupt to be handled, including `exccode` in `xcause` and `xil` in `xintstatus`.
- ② The interrupt trampoline needs sufficient space to store the OIC's caller-save registers as well as its `epc` and `cause` values, which are saved in a frame on the memory stack to support preemption. This routine is M-mode only so does not need to consider swapping stacks from other privilege modes. A simple constant bump of the stack pointer `sp` is sufficient to provide space to store the OIC.
- ③ The trap handler could have been entered by a synchronous exception instead of an interrupt, which can be determined by examining the sign bit of the returned `xcause` value. If the trap was for an exception (sign bit

zero), the code jumps to exception handler code while keeping interrupts disabled.

- ④ The exception handler code is located here out of line to reduce performance impact on interrupts. The main body of the trampoline only handles interrupts.
- ⑤ If this was an interrupt, the trampoline entry code continues to save all the caller-save registers to the stack. This is done with interrupts disabled, as even if an interrupt arrived with a higher interrupt level it would still require all registers to be saved.
- ⑥ When `xnxti` is read here, the interrupt inputs to the CLIC might have changed from the time the handler was initially entered. The return value of `xnxti`, which holds a pointer to an entry in the trap vector table, is saved in register `a0` so it can be passed as the first argument to the software-vectored interrupt handler, where it can be used to reconstruct the original interrupt id in the case where multiple vector entries use a common handler. There are multiple cases to consider, all of which are handled correctly by the definition of `xnxti`:
  - The II is still the ranking interrupt (no change). In this case, as the level of the II will still be higher than `pil` from the OIC, `xil` and `exccode` will be rewritten with the same value that they already had (effectively unchanged), and `xnxti` will return the table entry for the II.
  - The II has been superseded by a higher-level non-SHV interrupt. In this case, `xil` will be set to the new higher interrupt level, `exccode` will be updated to the new interrupt id, and `xnxti` will return the vector table entry for the new higher-level interrupt. The OIC is not disturbed, retaining the original `epc` and the original `pil`. This case reduces latency to service a more-important interrupt that arrives after the state-save sequence was begun for the less-important II. The II, if still pending-enabled, will be serviced sometime after the higher-level interrupt as described below.
  - The II has been superseded by a higher-priority non-SHV interrupt at the same level. This operates similarly to the previous case, with `exccode` updated to the new interrupt id. Because the lower-priority interrupt had not begun to run its service routine, this optimization preserves the property that interrupt handlers at the same interrupt level but different priorities execute atomically with respect to each other (i.e., they do not preempt each other).
  - The II has disappeared and a lower-ranked non-SHV interrupt, which has interrupt level greater than the OIC's `pil` is present in CLIC. In this case, the `xil` of the handler will be reduced to the lower-ranked interrupt's level, `exccode` will be updated with the new interrupt id, and `xnxti` will return a pointer to the appropriate handler in table. In this case, the new lower-ranked interrupt would still have caused the original context to have been interrupted to run the handler, and the disappearing II has simply caused the lower-ranked interrupt's entry and state-save sequence to begin earlier.
  - The II has disappeared and either there is no current interrupt from the CLIC, or the current ranking interrupt is a non-SHV interrupt with level lower than `xpil`. In this case, the `xil` and `exccode` are not updated, and 0 is returned by `xnxti`. The following trampoline code will then not fetch a vector from the table, and instead just restore the OIC context and `mret` back to it. This preserves the property that the OIC completes execution before servicing any new interrupt with a lower or equal interrupt level.
  - The II has been superseded by a higher-level SHV interrupt. In this case, the `xil` and `exccode` are not updated, and 0 is returned by `xnxti`. Once interrupts are reenabled for the following instruction, the hart will preempt the current handler and execute the vectored interrupt at a higher interrupt level using the function pointer stored in the vector table.
- ⑦ Interrupts are now enabled. If a higher-level SHV interrupt had arrived while interrupts were disabled, then the current handler will be preempted and execution starts at the SHV handler address. If a non-vectored higher-level interrupt arrives now, it will also preempt the current handler and begin a nested state-save sequence at the handler entry point `irq_enter`.
- ⑧ The branch checks if the II disappeared or if a higher priority SHV at the same level appeared, in which case the current handler returns to the OIC. As most registers have not been touched, the routine can skip past most of the register restore code. This preserves the property that interrupts (SHV or non-SHV) at the same

level do not preempt each other.

- ⑨ The value returned by `xnxti` is used to index the vector table and return the function pointer.
- ⑩ This `csrrsi` instruction enables interrupts and is redundant when proceeding sequentially from the first `xnxti` read (6) or if looping back from the end of the `service_loop` (13). However, it is required on the backward path from (16) to re-enable interrupts to allow preemption. It is scheduled after the table lookup to use what will often be a load-use delay slot.
- ⑪ The `jalr` instruction actually calls the C ABI function that implements the handler. Interrupts are enabled at this point, so the C function can be preempted at any time by an interrupt with a higher level than current `xil`.
- ⑫ Once the handler returns, another read of `xnxti` checks if there are any more interrupts to service. Interrupts remain enabled. The `csrrsi` includes a redundant set of the `xie` interrupt enable to force the CSR instruction to update CSR state. Only non-SHV interrupts with a level greater than `pil` will be serviced in this loop. Note that `xil` can decrease from its current value on the `xnxti` read. `xil` should not increase in this code, as interrupts are enabled here and if a higher-level interrupt was ready, it should have preempted this instruction.
- ⑬ If there was another appropriate interrupt to service, the code loops back to perform the next handler call. The `service_loop` only contains 5 instructions, allowing multiple back-back interrupts to be handled without saving and restoring contexts. On a simple pipeline with a one-cycle load-use penalty, single-cycle CSR access, and a one-cycle taken-branch penalty, the service loop can initiate a new interrupt service with only 7 clock cycles of overhead per handler call.
- ⑭ This instruction sequence restores the OIC. Interrupts are still enabled, so preemption is allowed during this restore.
- ⑮ Interrupts are disabled for the final steps of restoring the OIC, which requires loading `mcause` and `mepc` from the stacked values, and recovering the final register values from the OIC.
- ⑯ A final read of `xnxti` is performed before returning, to reduce the maximum interrupt latency. If a suitable interrupt arrives, it can be serviced without saving context. The `csrrci` instruction includes a redundant clear of the interrupt enable bit to ensure the CSR state updates occur. Interrupts must stay disabled until after the following branch to maintain the critical section used to restore the OIC in the case that there is no interrupt to service.

The following table summarizes the machine state changes that occur at the first `xnxti`:

IC	at entry  ->					at first nexti (6)								
il	CLIC					CLIC								
	level	id	V	->	mil	code		level	id	V	->	mil	code	rd
p	e<=p	?	?	->										# Shouldn't happen
p	e>p	i	0	->	e	i		f>p	j	0	->	f	j	T # Same or superseded interrupt
p	e>p	i	0	->	e	i		f>p	j	1	->	e	i	0 # Ignore vectored interrupt
p	e>p	i	0	->	e	i		f<=p	j	?	->	e	i	0 # Interrupt disappeared
p	e>p	i	1	->	e	i								# Won't be in trampoline

## 11.2. Revised C-ABI for Embedded RISC-V

The overhead to save and restore registers in the interrupt trampoline can be reduced with a new embedded ABI that reduces the number of caller-save registers. Work is underway to define such an ABI, but it is likely to require around 7 integer registers to be saved/restored instead of 16 in the standard Unix ABI.

This will result in 18 instructions executed in the trampoline code before arriving at the correct handler function,

of which 9 are stores (saving 7 registers plus 2 words for `xepc` and `xcause`).

## 11.3. Analysis of Worst-Case Interrupt Latencies for C-ABI Trampoline

The following analysis assumes a system with M-mode only and a new embedded ABI requiring 7 caller-save registers to be saved and restored. For cycle timings, we assume a simple 3-stage pipeline that has a one-cycle taken-branch or pipeline flush penalty, a one-cycle load-use delay, and single-cycle CSR access. This simple model ignores effects from contention in shared memory structures, or pipeline hazards from continuing long-latency operations in the interrupted code.

There are several cases to consider for the worst-case latency for a C-ABI higher-level interrupt handler that preempts lower-level code.

If an interrupt arrives while interrupts are enabled, either inside or outside of a current handler, the hart will jump directly to `irq_enter` at the new interrupt level. The system must flush the execution pipeline and then execute 18 instructions, the last of which is the `jalr` that calls the handler function. These 18 instructions execute in 20 cycles using the simple pipeline model.

When interrupts are disabled, the arriving preempting handler could be delayed. If the preempting interrupt arrives while interrupts are disabled during the initial entry sequence (1)–(6), there will be no additional delay as the first `xnxti` instruction (6) will cause the higher-level interrupt handler to be invoked, replacing the original interrupt cause.

If the preempting interrupt arrives after interrupts are disabled (15) but before `xnxti` is read (16), then the trampoline will observe the new interrupt during execution of the `xnxti` read (16), and take a short branch back to the `service_loop`, which is lower latency than the interrupt-disabled case.

If the preempting interrupt arrives after the read of `xnxti` commits (16), then the interrupt has to wait an additional 4 instructions until the `mret` reenables interrupts, at which point the interrupt will be taken and the handler entered at `irq_enter`. In the simple pipeline model, `mret` adds an additional pipeline flush cycle, so the preemption latency is 20+5 cycles, which represents the worst-case for a preempting C-ABI interrupt handler.

## Chapter 12. Interrupt-Driven C-ABI Model

For many embedded systems, after initialization, essentially all code is run in response to an interrupt, interrupt levels are used to prioritize execution of different tasks, and the hart should sleep in between interrupt events to save energy.

The following code can be used as the background code that runs at interrupt level 0 and which when there is no active work to do, puts the hart to sleep with no active context, waiting for an interrupt using the `wfi` instruction. The code is entered at the `enter_loop` location and never returns directly.

```
# Source code for interrupt-driven model background code.
sleep:
    csrrci zero, mstatus, MIE # Disable interrupts. ①
    wfi                      # Hart waits for next interrupt event.
    csrrsi a0, mnxti, MIE    # Gather interrupt details, and enable interrupts. ②
    beqz a0, sleep           # Go back to sleep if no interrupt (will be preempted
if SHV). ③

service_loop: ④
    lw a1, (a0)              # Get handler address.
    csrrsi zero, mstatus, MIE # Enable interrupts
    jalr a1                  # Call C-ABI handler routine
    csrrsi a0, mnxti, MIE    # Claim any pending interrupt at level > 0
    bnez a0, service_loop    # Loop to service any interrupt.

# This is also entry point to begin sleeping.
enter_sleep: ⑤
    la a0, sleep
    csrrci zero, mstatus, MIE # Disable interrupts.
    #--- Interrupts disabled
    csrw mepc, a0             # Initialize mepc to point to sleep
    li a0, (MMODE)<<PP|(0)<<PIL|(1)<<PIE
    csrw mcause, a0           # Initialize mcause to have pp=M, pil=0, pie=1
    mret                     # Jump to sleep at level 0 with interrupts enabled.
    #--- Interrupts enabled
```

- ① The `sleep` loop is used to stall the hart while waiting for work and is always entered at interrupt level 0. Interrupts are disabled, then a `wfi` is executed. The `wfi` will stall the hart until some event occurs. When an event, including an interrupt occurs, the `wfi` retires. Because interrupts are disabled, the hart does not jump to an interrupt handler but instead executes the next instruction, avoiding context save/restore overhead.
- ② The read of `mnxti` will determine if any non-SHV interrupt is available, and if so return a pointer to the table entry. Interrupts are enabled by this instruction to allow SHV interrupts to be taken via preemption.
- ③ The value in `a0` checked by the branch can be zero for two reasons. Either there was no interrupt detected or an SHV interrupt was detected. If there was no interrupt, the branch loops back to put the hart to sleep. Interrupts are enabled, so any SHV interrupt (which all have higher interrupt level than the current interrupt level of 0) will preempt the branch's execution and call the SHV handler. Once the SHV handler returns, the branch will resume and cause execution to return back to the `sleep_loop`.

- 
- ④ The service loop is identical to that in the C-ABI interrupt handler, except that the previous interrupt level is 0, so all pending interrupts will be serviced in the loop before the loop exits. Interrupts are enabled, so preemption is allowed for both C-ABI trampoline and SHV interrupts. When an SHV interrupt at the same or lower interrupt level is the next to be serviced, the `xnxti` instruction will return 0 causing execution to drop out of the loop. The following code will reinitialize the hart's interrupt level to 0, and disable interrupts for one instruction, to ensure the SHV interrupt will be taken.
  - ⑤ This code initializes `mepc` and `mcause` then uses an `mret` to jump to the `sleep` loop while simultaneously resetting interrupt level to 0 and enabling interrupts. This is also the entry point to initiate interrupt-driven execution. Interrupts are enabled to allow SHV interrupts to preempt execution on the first instruction in `sleep` (which disables interrupts again).

This code does not increase worst-case interrupt latency over that of the C-ABI trampoline.

## Chapter 13. Alternate Interrupt Models for Software Vectoring

Platforms may not implement the `sclicshv` extension, in which case, hardware vectoring can be emulated by a single software trampoline present at `NBASE` using the separate vector table address in `xtvt`. There are several different software approaches possible, depending on system requirements and constraints, as detailed in following subsections.

### 13.1. `gp` Trampoline to Inline Interrupt Handlers in Single Privilege Mode

Where interrupts are known to be generated and handled in a single privilege mode (i.e., M-mode only systems, or U-mode interrupt handlers), a three-instruction sequence using the `gp` register to hold the handler address can be used to indirect to an inline interrupt handler of the type described in [Inlines](#).

```
# Software-vectorized interrupt servicing.
# Only safe for horizontal interrupts.
# Must be placed three instructions back from gp.
irq_enter:
    csrrci gp, mnxti, MIE    # Overwrite gp, keep interrupts disabled.
    beqz gp, handle_exc      # Encountered exception.
    jalr gp, gp              # Recreate gp and jump to handler.
gp:                          # Must be right before system's gp location.
    # ... gp data section

    # Must be within range of beqz instruction.
handle_exc:
    # Has to recreate gp.
```

The three-instruction sequence relies on the `jalr` instruction recreating the value in the `gp` register, which is a known constant pointing into the middle of the global data area, by placing the `jalr` directly before the `gp` location in memory. The routine jumped to by the `jalr` does not return via a `j ra` but instead ends with an `mret`.



This constraint on memory layout might not always be possible, particularly if the system does not allow placing executable memory right next to read-write memory, for example if the system does not allow a protection boundary to be placed at 'gp' and if executable code must not be writeable.

The code can be used with preemptible inline interrupt handlers.

### 13.2. Trampoline for Preemptible Inline Handlers

This section describes a more general software-trampoline scheme for calling preemptible inline handlers, which factors out the `xepc/xcause` save code into the trampoline, and which uses a different interrupt handler calling convention.



The interrupt handlers for this scheme have a calling convention where there is one caller-save argument register **a0** that passes in the handler address to distinguish different interrupt inputs, and one temporary register **a1** that is also caller-save. These two registers had to be saved already by the trampoline. All other registers are callee-save, except for the return address **ra**. The handler normally returns with a regular **j ra**.

```
# Example handler with new calling convention.
# Only safe for horizontal interrupts.
# Handlers have two temporary registers available, a0, a1.
handler_example:
    sw zero, INTERRUPT_FLAG, a0      # Clear interrupt flag.
    la a0, COUNTER                   # Get counter address.
    li a1, 1                         # Increment value.
    amoad.w zero, (a0), a1           # Bump counter.
    j ra

# Interrupt trampoline code.
irq_enter:
    #----- Interrupts disabled on entry ---#
    addi sp, sp, -FRAMESIZE          # Create a frame on stack.
    sw a0, OFFSET(sp)               # Save working register.
    csrr a0, mcause                  # Read cause.
    bgez a0, handle_exc              # Handler exception.
    sw a1, OFFSET(sp)               # Save working register.
    csrr a1, mepc                    # Read epc.
    sw a0, OFFSET(sp)               # Save cause
    csrrsi a0, mnxti, MIE            # Get highest interrupt, enable interrupts.
    #----- Interrupts enabled -----#
    beqz a0, exit
    sw a1, OFFSET(sp)               # Save epc.
    sw ra, OFFSET(sp)               # Save return address.

irq_loop:
    lw a1, (a0)                     # Get function pointer.
    jalr a1                          # Call handler code.
    csrrsi a0, mnxti, MIE            # Get any next interrupt.
    bnez a0, irq_loop               # Service interrupt if any.

    lw ra, OFFSET(sp)               # Restore ra.
    lw a1, OFFSET(sp)               # Get epc.
exit:
    lw a0, OFFSET(sp)               # Get cause.
    csrrci zero, mstatus, MIE        # Disable interrupts.
    #----- Interrupts disabled -----#
    csrw mepc, a1                   # Put epc back.
    lw a1, OFFSET(sp)               # Restore a1.
    csrw mcause, a0                 # Put cause back.
    lw a0, OFFSET(sp)               # Restore a0.
```

```

    addi sp, sp, FRAMESIZE      # Free stack frame.
    mret                        # Return from handler.
    #-----#

handle_exc:
    # ...
    # Handle exception with interrupts disabled.
    # ...
    addi sp, sp, FRAMESIZE     # Deallocate stack space
    mret                        # Return from handler.
    #-----#

```

This interrupt handler can be used together with the `wfi` sleep background routine shown above.

## Chapter 14. Managing Interrupt Stacks Across Privilege Modes

Interrupt handlers need to have a place to save the previous context's state to provide working registers for the handler code. If a handler can be entered from a lower-privilege mode, a pointer to some safe memory for the context save must be swapped in at entry to the higher-privileged handler to avoid security holes. The RISC-V privileged architecture provides the `xscratch` register to hold this information for a higher-privilege mode while executing in a lower-privilege mode. For the following discussion and code examples, the assumption is that `xscratch` is used to hold the higher-privilege-mode stack pointer but other software conventions are possible (e.g., `xscratch` points to a thread context block).

Existing RISC-V ABIs allow addresses immediately below the stack pointer to be overwritten by interrupt service routines. The current stack pointer in `sp` (`x2`) should be swapped with `xscratch` whenever a handler is entered from a lower-privilege mode, but should not be swapped if entered from another handler in the same privilege mode, including when preempting an existing interrupt handler. At exit from a handler, the lower-privilege stack pointer should be swapped back in if transitioning back to the lower-privilege mode.

### 14.1. Software Privileged Stack Swap

In this convention, when code is running in a lower privilege mode, `xscratch` holds the stack pointer for the higher-privilege mode. When the higher-privilege mode is entered, `xscratch` is set to zero to signal to any preempting handlers that the stack pointer has already been swapped.

The old stack pointer is saved to new stack frame before new frame is created by bumping stack pointer, but this is done with interrupts disabled.

```

# This code is out of line to reduce worst-case preemption latency.
enter_M:
    sw sp, OFFSET-FRAMESIZE(sp) # Save previous mscratch (M-mode sp)
    addi sp, sp, -FRAMESIZE      # Create a frame on stack.
    sw a0, OFFSET(sp)           # Save a register.
    csrrw a0, mscratch, 0       # Get previous sp, and zero mscratch.
    sw a0, OFFSET(sp)           # Save previous sp (U-mode sp)
    j continue                  # Jump back into handler

irq_enter:
    #----- Interrupts disabled on entry ---#
    csrrw sp, mscratch, sp      # Swap stack pointer and scratch.
    bnez sp, enter_M            # Check if entering M-mode
    csrrw sp, mscratch, sp      # Already in M-mode, so swap sp back.
    sw sp, OFFSET-FRAMESIZE(sp) # Save previous sp to stack.
    addi sp, sp, -FRAMESIZE      # Create a frame on stack.
    sw zero, OFFSET(sp)         # Save previous mscratch to stack (was zero).
    sw a0, OFFSET(sp)           # Save a register.
continue:
    csrr a0, mcause              # Read cause.
    bgez a0, handle_exc         # Handle exception.
    sw a1, OFFSET(sp)           # Save working register.
    csrr a1, mepc                # Read epc.
    sw a0, OFFSET(sp)           # Save cause
    csrrsi a0, mnxti, MIE       # Get highest interrupt, enable interrupts.
    #----- Interrupts enabled -----#
    beqz a0, exit
    ...

#---- Critical section with interrupts disabled -----
    ...

    lw a0, OFFSET(sp)           # Get previous mscratch.
    csrw mscratch, a0           # Put back in mscratch.
    lw a0, OFFSET(sp)           # Restore original a0 value.
    lw sp, OFFSET(sp)           # Restore previous sp
    mret                        # Return from interrupt.
#-----

```

This code can be used in a secure model where user-level code has one stack, and all interrupts and exceptions are handled on a second M-mode-only stack. In addition, background non-handler code in M-mode can either use the same M-mode stack as the interrupt handler, or a separate M-mode stack. The only difference is in the value held in `xscratch` while the M-mode background thread is running (either 0 to indicate use the existing stack pointer in `sp` or non-zero to indicate this stack pointer should be used in the handler).

The above software scheme adds 7 instructions to the interrupt code path when preempting the same privilege mode, and adds an additional 6 instructions (13 total including two taken branches) for interrupts from a lower-

privilege mode into a higher-privileged mode.

### 14.1.1. Stack Swap Example Code

Interrupt handlers running in the lowest privilege mode do not need to swap stack pointers, as they will only be entered by a horizontal interrupt from the same privilege mode. In systems with multiple privilege modes, handlers running in higher privilege modes must account for vertical interrupts taken from a lower privilege mode (in which case the stack pointer must be swapped) as well as horizontal interrupts from the same privilege mode.

```
# Example of inline interrupt with stack swapping.
.align 3
foo:
    csrrw sp, mscratchcsw, sp    # Conditionally swap in stack pointer.
    addi sp, sp, -FRAMESIZE      # Create a frame on stack.
    sw s0, OFFSET(sp)           # Save working register.
    sw zero, INTERRUPT_FLAG, s0  # Clear interrupt flag.
    sw s1, OFFSET(sp)           # Save working register.
    la s0, COUNTER               # Get counter address.
    li s1, 1
    amoadd.w zero, (s0), s1      # Increment counter in memory.
    lw s1, OFFSET(sp)           # Restore registers.
    lw s0, OFFSET(sp)
    addi sp, sp, FRAMESIZE       # Free stack frame.
    csrrw sp, mscratchcsw, sp    # Conditionally swap out stack pointer.
    mret                         # Return from handler using saved mepc.
```

```

# Example of inline preemptible interrupt with stack swapping.
.align 3
foo:
    #----- Interrupts disabled on entry ---#
    csrrw sp, mscratchcsw, sp    # Conditionally swap in stack pointer.
    addi sp, sp, -FRAMESIZE      # Create a frame on stack.
    sw s0, OFFSET(sp)           # Save working register.
    sw s1, OFFSET(sp)           # Save working register.
    csrr s0, mcause              # Read cause.
    csrr s1, mepc                # Read epc.
    csrrsi zero, mstatus, MIE    # Enable interrupts.
    #----- Interrupts enabled -----#
    sw s0, OFFSET(sp)           # Save cause on stack.
    sw zero, INTERRUPT_FLAG, s0  # Clear interrupt flag.
    sw s1, OFFSET(sp)           # Save epc on stack.
    la s0, COUNTER              # Get counter address.
    li s1, 1
    amoad.w zero, (s0), s1       # Increment counter in memory.
    lw s1, OFFSET(sp)           # Restore epc
    lw s0, OFFSET(sp)           # Restore cause
    #----- Interrupts disabled -----#
    csrrci zero, mstatus, MIE    # Disable interrupts.
    csr w mepc, s1               # Put epc back.
    csr w mcause, s0             # Put cause back.
    lw s1, OFFSET(sp)           # Restore s1.
    lw s0, OFFSET(sp)           # Restore s0.
    addi sp, sp, FRAMESIZE       # Free stack frame.
    csrrw sp, mscratchcsw, sp    # Conditionally swap out stack pointer.
    mret                         # Return from handler.
    #-----#

```

```

# Example C-ABI interrupt trampoline with stack swapping.

irq_enter:
#----
    csrrw sp, mscratchcsw, sp # Conditionally swap in stack pointer.
    addi sp, sp, -FRAMESIZE    # Allocate space on stack.
    # ...
    # Everything else same as above.
    # ...
    addi sp, sp, FRAMESIZE     # Reclaim stack space.
    csrrw sp, mscratchcsw, sp # Conditionally swap back stack pointer.
    mret                      # Return from interrupt.
#-----
#-----

handle_exc:
    # ...
    # Perform exception processing with interrupts disabled
    # ...
    addi sp, sp, FRAMESIZE     # Reclaim stack space.
    csrrw sp, mscratchcsw, sp # Conditionally swap back stack pointer.
    mret                      # Return from exception
#-----

```

In all cases, conditionally swapping the stack to account for potential privilege-mode changes adds two extra instructions to all interrupt handlers.

## Chapter 15. Separating Stack per Interrupt Level

Within a single privilege mode, it can be useful to separate interrupt handler tasks from application tasks to increase robustness, reduce space usage, and aid in system debugging. Interrupt handler tasks have non-zero interrupt levels, while application tasks have an interrupt level of zero.



## Chapter 16. CLIC Interrupt ID ordering recommendations

The specific numbering of interrupts is defined by the platform. A platform definition will often be based on a specific RISC-V ISA profile, where RISC-V ISA profiles specify a common set of ISA choices that capture the most value for most users to enable software compatibility.

The current RVA profiles (RVA20/RVA22/RVA23) include the CLINT interrupt scheme.

Certain RISC-V profiles may include the CLIC as an option.



The CLIC is not included in RVA series of profiles as mandatory or supported optional.

Four different CLIC interrupt ID orderings are enumerated below for ease of reference in profile specifications.

### 16.1. CLIC-mode interrupt-map for systems retaining interrupt ID compatible with CLINT mode:

The CLINT-mode interrupts retain their interrupt ID in CLIC mode. The `clicintattr` settings are now used to delegate these interrupts as required.

We recommend allocating an edge-triggered CLIC software interrupt bit (`csip`) with the lowest number (16) among the local interrupts to provide support for a local background interrupt thread, e.g. context switching. By placing the `csip` interrupt at the lowest number among local interrupts, we avoid requiring another interrupt level to ensure this only runs when there are no other interrupts.

The existing CLINT software interrupt bits are primarily intended for inter-hart interrupt signaling, and so are retained for that purpose. The software interrupts can have their priority changed using their `clicintctl` setting.

CLIC interrupt inputs are allocated IDs beginning at interrupt ID 17. Any fast local interrupts that would have been connected at interrupt ID 16 and above should now be mapped into corresponding inputs of the CLIC.

ID	Interrupt	Note
0	usip	User software Interrupt
1	ssip	Supervisor software Interrupt
2	reserved	
3	msip	Machine software interrupt
4	utip	User timer interrupt
5	stip	Supervisor timer interrupt
6	reserved	
7	mtip	Machine timer interrupt
8	ueip	User external (PLIC/APLIC) interrupt
9	seip	Supervisor external (PLIC/APLIC) interrupt
10	reserved	
11	meip	Machine external (PLIC/APLIC) interrupt
12	reserved	
13	reserved	
14	reserved	
15	reserved	
16	csip	CLIC software interrupt
17+	inputs	CLIC local inputs

In CLINT mode, as stated in the RISC-V privilege specification, each individual bit in CSR register `xip` may be software writable or may be software read-only. When bit  $i$  in `xip` is writable, a pending interrupt  $i$  can be cleared by writing 0 to this bit. STIP is writable in `mip`, and may be written by M-mode software to deliver timer interrupts to S-mode.

However, in CLIC mode, `xip` bits cannot be updated in software when they are in level-sensitive mode. One possible equivalent implementation to deliver timer interrupts would be to hardwire the STIP input to 0 and then M-mode software can deliver timer interrupts to S-mode by setting `clicintattr[stip].trig` to 00 (positive-edge-triggered) and writing `clicintip[stip]` to 1.

The following optional recommendations remove unused interrupt inputs and compress the map to simplify hardware implementations.

## 16.2. CLIC-mode interrupt-map recommendation for single-hart systems with PLIC/APLIC:

ID	Interrupt
0	S-mode software interrupt
1	S-mode timer interrupt
2	S-mode external (PLIC/APLIC) interrupt
3	M-mode software interrupt
4	M-mode timer interrupt
5	M-mode external (PLIC/APLIC) interrupt
6+	local

## 16.3. CLIC-mode interrupt-map recommendation for single-hart M/S/U systems without N-extension with no PLIC/APLIC:

ID	Interrupt
0	S-mode software interrupt
1	S-mode timer interrupt
2	M-mode software interrupt
3	M-mode timer interrupt
4+	other

## 16.4. CLIC-mode interrupt-map recommendation for single-hart M-mode only or M/U mode

systems without N extension and no PLIC/APLIC:

ID	Interrupt
0	M-mode software interrupt
1	M-mode timer interrupt
2+	other

# Appendix A: Appendix

## A.1. Prototype DTS Entry

Modified from Example at [elixir.bootlin.com/linux/v4.14.0/source/arch/arm/boot/dts/example.dtsi](https://elixir.bootlin.com/linux/v4.14.0/source/arch/arm/boot/dts/example.dtsi)

```
/dts-v1/;

/ {
    ...

    interrupt-parent = <&intc>;

    ...

    intc: interrupt-controller@MCLICBASE {
        compatible = "xxx,xxx";
        reg = <MCLICBASE 0x5000 >;
        interrupt-controller;
        #interrupt-cells = <NUM_INTERRUPT>;
    };

    ...
};
```

## Chapter 17. Bibliography

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