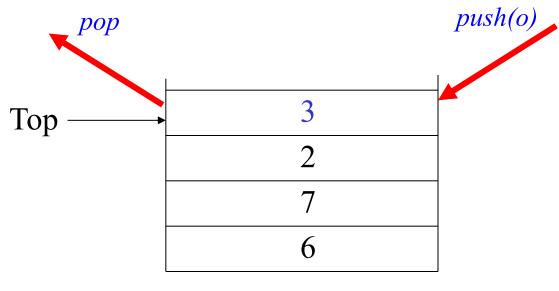
Stacks

What is a Stack?

- A stack is a list with the restriction that insertions and deletions can be performed in only one position, namely, the end of the list, called the top.
- The operations: push (insert) and pop (delete)



Designing and Building a Stack class

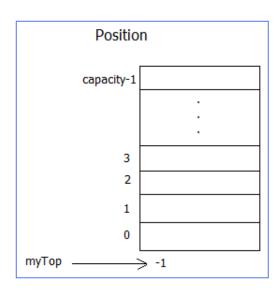
- The basic functions are:
 - Constructor: construct an empty stack
 - Empty(): Examines whether the stack is empty or not
 - Push(): Add a value at the top of the stack
 - Top(): Read the value at the top of the stack
 - Pop(): Remove the value at the top of the stack
 - Display(): Displays all the elements in the stack
 - Size(): Number of elements in stack

Constructor:

Create an array: (int) array[capacity]

Set myTop = -1

Empty():check if myTop == -1



Push(int x):

```
if array is not FULL (myTop < capacity-1)
   myTop++
   store the value x in array[myTop]
else
   output "out of space"</pre>
```

output "no elements in the stack"

Top():
 If the stack is not empty
 return the value in array[myTop]
 else:

```
    Pop():
        If the stack is not empty
        myTop -= 1
        else:
        output "no elements in the stack"
```

Further Considerations

- What if static array initially allocated for stack is too small?
 - Terminate execution?
 - Replace with larger array!

- Creating a larger array
 - Allocate larger array
 - Use loop to copy elements into new array
 - Delete old array

Growing the stack

- Double the stack each time
- Increase stack size by a constant k

Which one is preferred?

Double the Stack size

Initially size 1(phase1), then size 2 (in phase 2), size4(in phase 3) and so on In phase i, the size is 2ⁱ⁻¹

Cost of each phase: $2^{i-1} + 2^{i-2} + 2^{i-2} = 2^{i}$

- Construct array of size 2i-1
- Copy elements from previous array 2ⁱ⁻²
- Fill remaining 2ⁱ⁻² elements that can be accommodated.

For n elements we require log n phases.

Total cost: $2 + 4 + 8 + \dots + 2^{\log n} = O(n)$

Increase Stack size by c

Initially size c(phase1), then size 2c (in phase 2), size 3c (in phase 3) and so on In phase i, the size is i.c

Cost of each phase: i.c + (i-1).c+c=2.i.c

- Construct array of size i.c
- Copy elements from previous array (i-1).c
- Fill remaining elements that can be accommodated.: c

For n elements we require n/c phases.

Total cost:
$$2c + 2.2c + 2.3c + 2.4c + + n/c$$
. $2c = 2c(n/c)^2 = 2n^2/c$

Stack ADT Interface

• The main functions in the Stack ADT are (S is the stack)

```
boolean isEmpty(); // return true if empty
boolean isFull(S); // return true if full

void push(S, item); // insert item into stack

void pop(S); // remove most recent item

void clear(S); // remove all items from stack

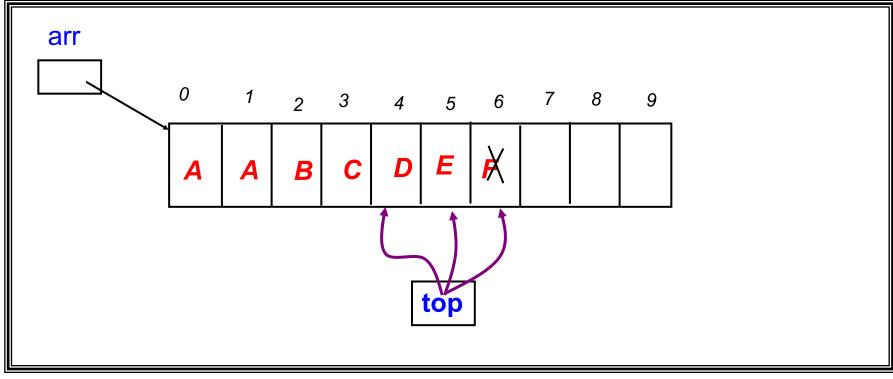
Item top(S); // retrieve most recent item

Item topAndPop(S); // return & remove most recent item
```

Implementation by Array

use Array with a top index pointer as an implementation of stack

StackAr



int top=-1,stack[MAX];

```
void pop()
  if(top==-1)
     printf("\nStack is empty!!");
  else
     printf("\nDeleted element is
%d",stack[top]);
     top=top-1;
```

```
void push()
  int val;
  if(top==MAX-1)
     printf("\nStack is full!!");
  else
     printf("\nEnter element to push:");
     scanf("%d",&val);
     top=top+1;
     stack[top]=val;
```

Applications

- Balanced Paranthesis
- Convert infix to postfix and vice versa
- Evaluating PostFix notation

- Infix to prefix and vice versa
- Evaluating Prefix notation

Infix to postfix

- If (push
- If) pop and print everything on stack until (
- If operand print
- If operator,
 - (a)-pop and print elements on stack of same or higher priority. If no such elements then do nothing.(b)-push operator

Priority order (, +-, */, ^

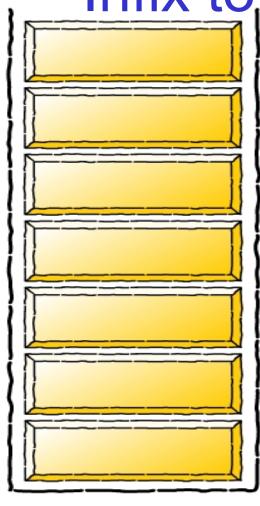
Try

• (a+b-c)*d-(e+f)

• $4^2*3-3+8/4/(1+1)$

a+b/c*(d+e)-f

Infix to postfix conversion

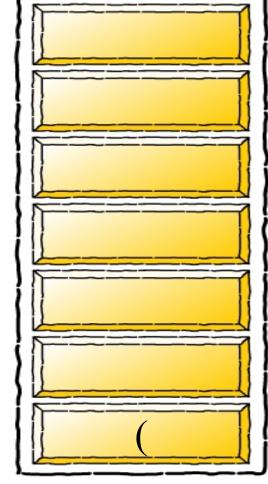


infix Vect

$$(a+b-c)*d-(e+f)$$



$$a + b - c$$
) * $d - (e + f)$



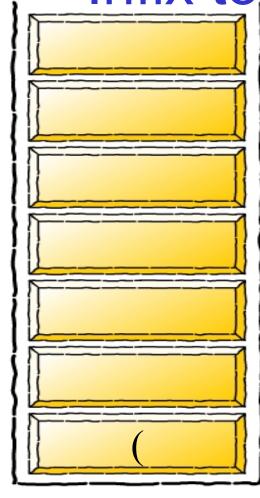




$$+ b - c) * d - (e + f)$$

postfix Vect

a

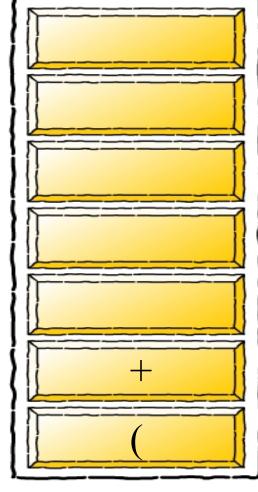




$$b - c) * d - (e + f)$$

postfix Vect

a

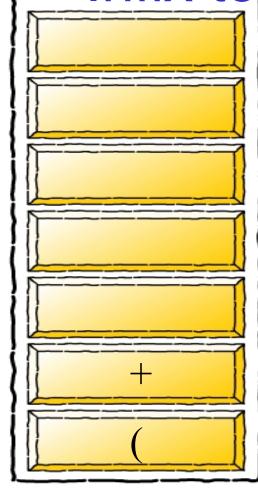




$$-c)*d-(e+f)$$

postfix Vect

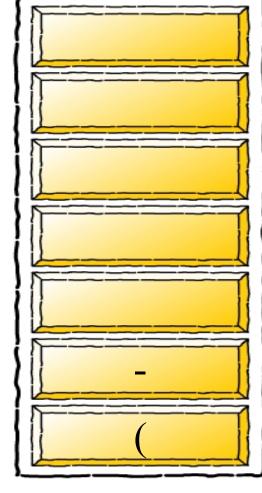
a b







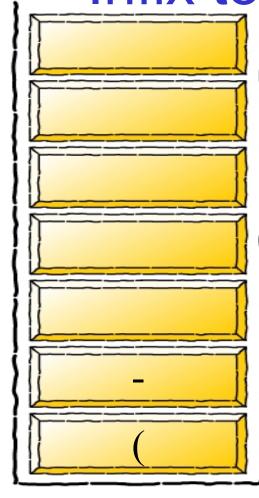
$$c)*d-(e+f)$$





$$) * d - (e + f)$$

$$ab+c$$

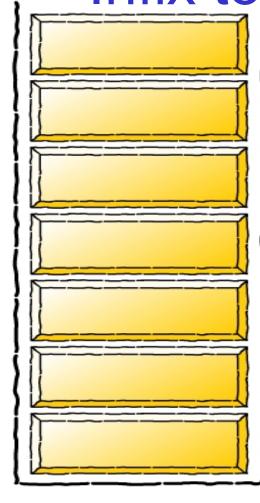






$$*d - (e + f)$$

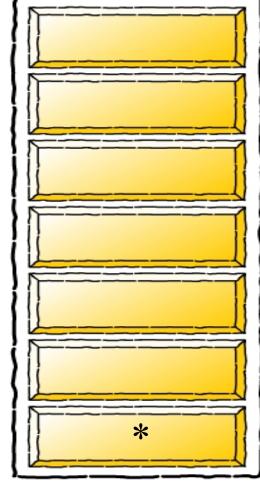
$$ab+c$$





$$d-(e+f)$$

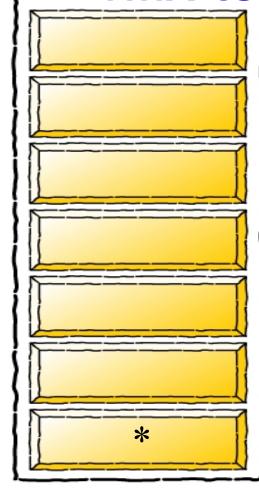
$$ab+c$$





$$-(e+f)$$

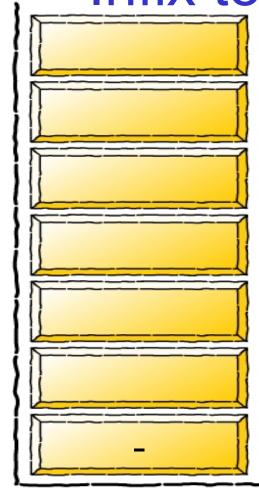
$$ab+c-d$$





$$(e+f)$$

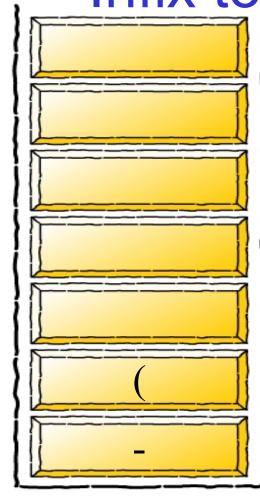
$$ab+c-d*$$





$$e + f$$

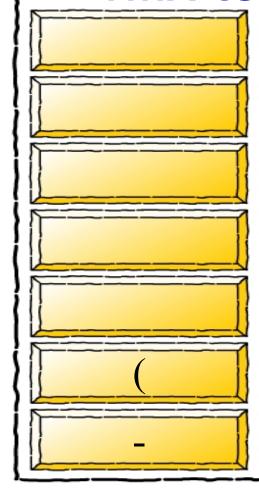
$$ab+c-d*$$

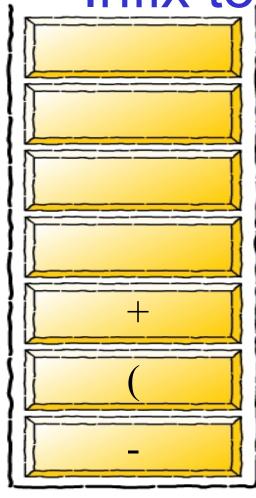




$$+f)$$

$$ab+c-d*e$$

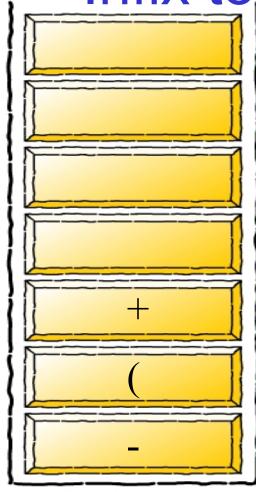




infixVect

f)

$$ab+c-d*e$$



infixVect

postfix Vect

ab+c-d*ef

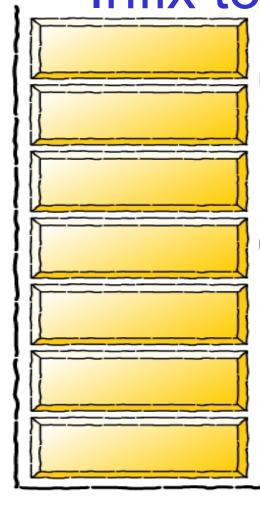




$$ab+c-d*ef+$$



$$a b + c - d * e f + -$$



- Scan the token list from left to right.
 - If the token is an operand, append it to the end of the output list.
 - If the token is a left parenthesis, push it on the opstack.
 - If the token is a right parenthesis, pop the opstack until the corresponding left parenthesis is removed. Append each operator to the end of the output list.
 - If the token is an operator, *, /, +, or -, push it on the opstack.
 However, first remove any operators already on the opstack that have higher or equal precedence and append them to the output list.
- When the input expression has been completely processed, check the opstack. Any operators still on the stack can be removed and appended to the end of the output list.

Try

- $4^2*3-3+8/4/(1+2)$
- A+b/c*(d+e)-f
- A^b^c [Note exponents evaluated right to left]. H.W
- Is $2^3^2 = 2^3 = 2^3 = 512$ or $(2^3)^2 = 64$?

Try

- 4²*3-3+8/4/(1+2)
 4 2 ³ * 3 8 4 / 1 2 + / +
- A+b/c*(d+e)-f
 A b c / d e + * + f -
- A^b^c [Note exponents evaluated right to left]. H.W
- Is $2^3^2 = 2^3(3^2) = 512$ or $(2^3)^2 = 64$?

Convert Postfix to Infix

- Abc/de+*+f-
- 42³*3-84/12+/+

Convert Postfix to Infix

- Abc/de+*+f (A +(b/c) * (d+e))-f
- 42³*3-84/12+/+

Evaluating Postfix

- Abc/de+*+f-
- 42³*3-84/12+/+

Try

- A^b^c [Note exponents evaluated right to left]. H.W
- In case of exponents rule changed
 - -From: pop same and higher priority elements
 - -To: pop higher priority elements.
- However, there are no higher priority operands and hence, we only push.
- a-b ^c ^d + f becomes a b c d ^ ^ f +

Notes

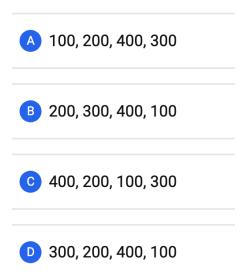
• Static array allocation O(1) but dynamic array allocation using malloc O(n) since n free chunks need to be checked for. These need not be contiguous

Questions

• Let S1 and S2 be two stacks. S1 has capacity of 4 elements. S2 has capacity of 2 elements. S1 already has 4 elements: 100, 200, 300, and 400(top), whereas S2 is empty.

Only the following three operations are available:

- PushToS2: Pop the top element from S1 and push it on S2.
 PushToS1: Pop the top element from S2 and push it on S1.
 GenerateOutput: Pop the top element from S1 and output it to the user.
- Note you cant pop an empty stack and push into a full stack
- Which of the following output sequences can be generated by using the above operations?



A single array A[1..MAXSIZE] is used to implement two stacks. The two stacks grow from opposite ends of the array. Variables topl and top2 (top1 < top2) point to the location of the topmost element in each of the stacks. If the space is to be used efficiently, the condition for "stack full" is ______.

•(GATE CSE 2004)

- 1. (top 1 = MAXSIZE/2) AND (top 2 = MAXSIZE/2 + 1)
- 2. top 1 + top 2 = MAXSIZE
- 3. (top 1 = MAXSIZE/2) or (top 2 = MAXSIZE)
- 4. top 1 = top 2 1

Homework

- Implement two stacks in an array by Dividing the space into two halves
- Sort a stack using a temporary stack Time:O(n^2), space:O(n)
- Design a stack that supports getMin() in O(1) time and O(1) extra space