### Queue

#### Queue -Basics

- Queue List with access restrictions
  - ☐ FIFO First-In First-Out
  - Double ended- Head and Tail (front and rear)
  - Insertion always to the tail
  - Deletion always from the head

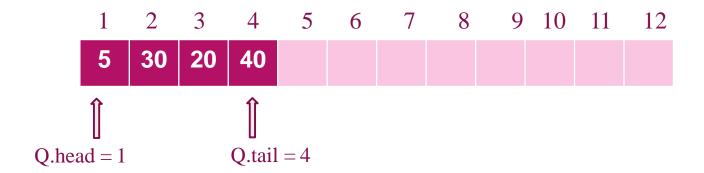
#### Queue - Implementations

- Array Based
  - □ Q [1..n]
  - □ Head, Tail array indices
- Pointer Based
  - As a linked list
  - □ Head, Tail pointers to the nodes at front and rear respectively

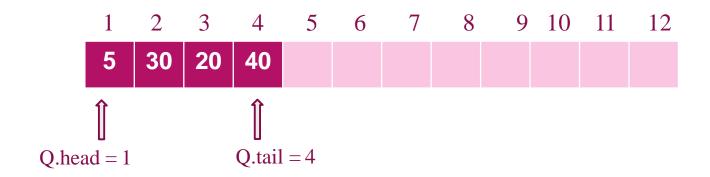
#### Queue - Array Based Implementation #1

- Array Based
  - ☐ Array **Q[1..n]** an array of at most n elements
  - ☐ An attribute **Q.head**—index of the head element
  - ☐ An attribute **Q.tail** index of the tail element
  - ☐ Elements from *Q[Q.head..Q.tail]*

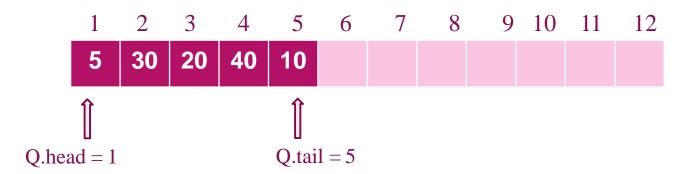
#### Queue - Implementation using Array #1

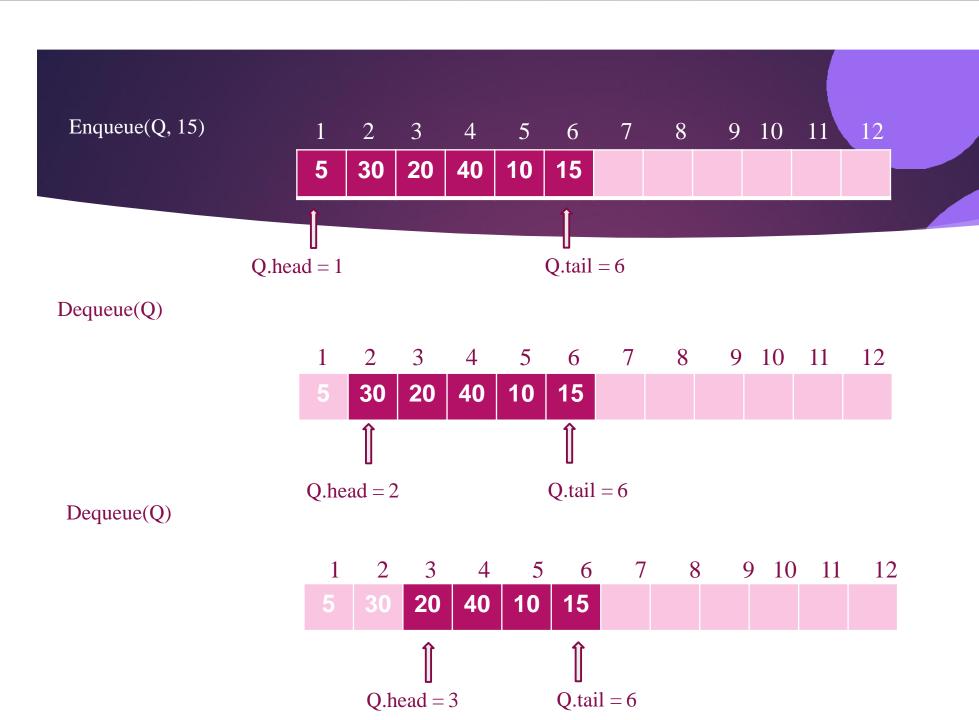


#### Queue - Implementation using Array #1



Enqueue(Q, 10)





**ENQUEUE (Q, x) // check the correctness** 

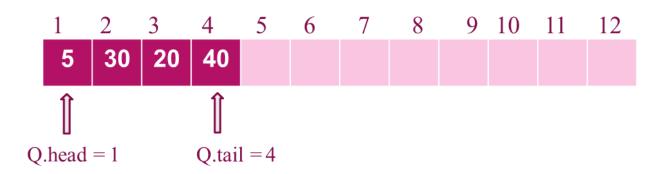
if (QUEUE-FULL(Q))

error "overflow"

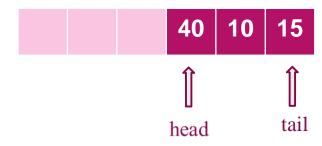
else

$$Q. tail = Q. tail + 1$$

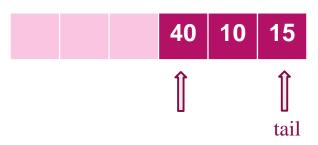
$$Q[Q.tail] = x$$



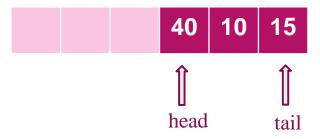
#### QUEUE-FULL(Q)?



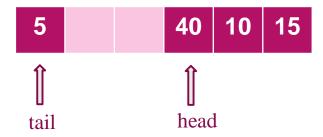
No further EnQueue possible even though the queue is not full

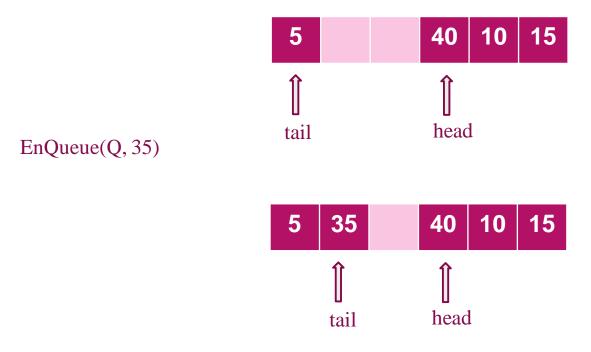


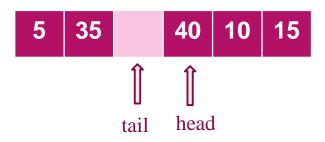
- No further EnQueue possible even though the queue is not full
  - > Left Shift the element
    - > Takes linear time



#### **Start adding from left**



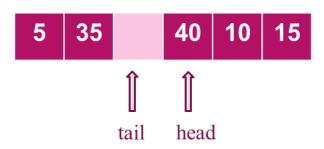


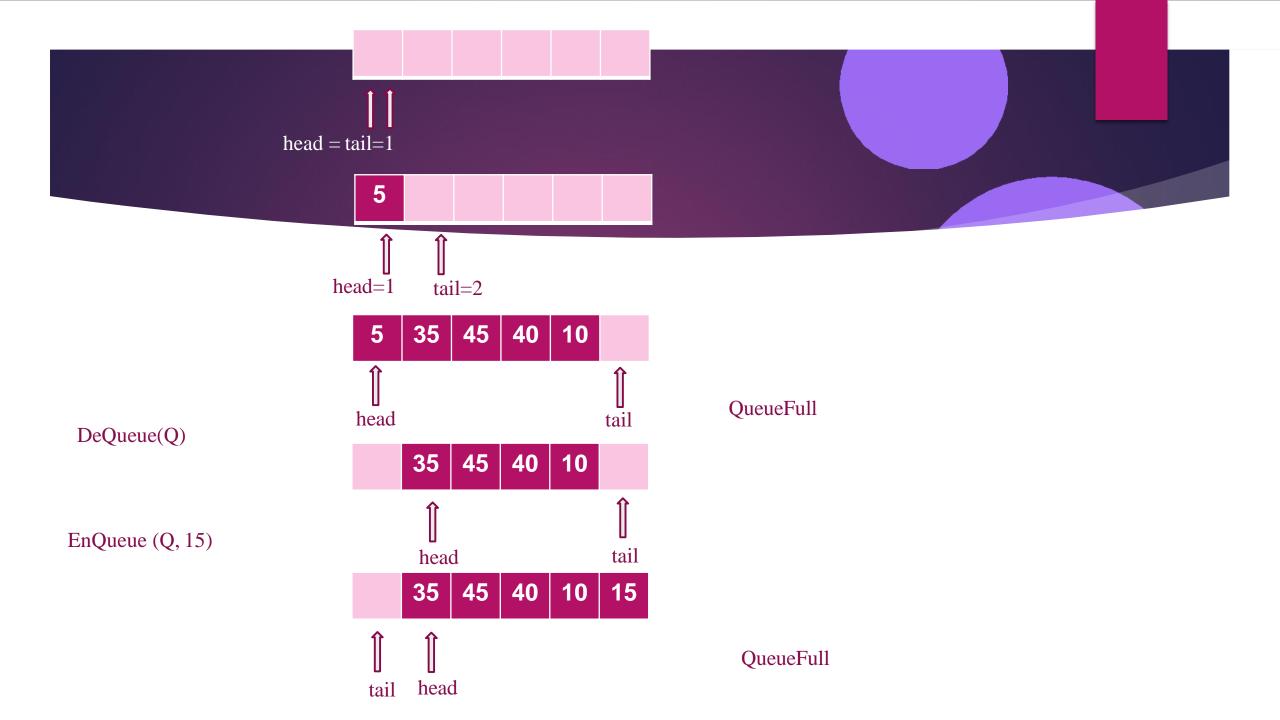


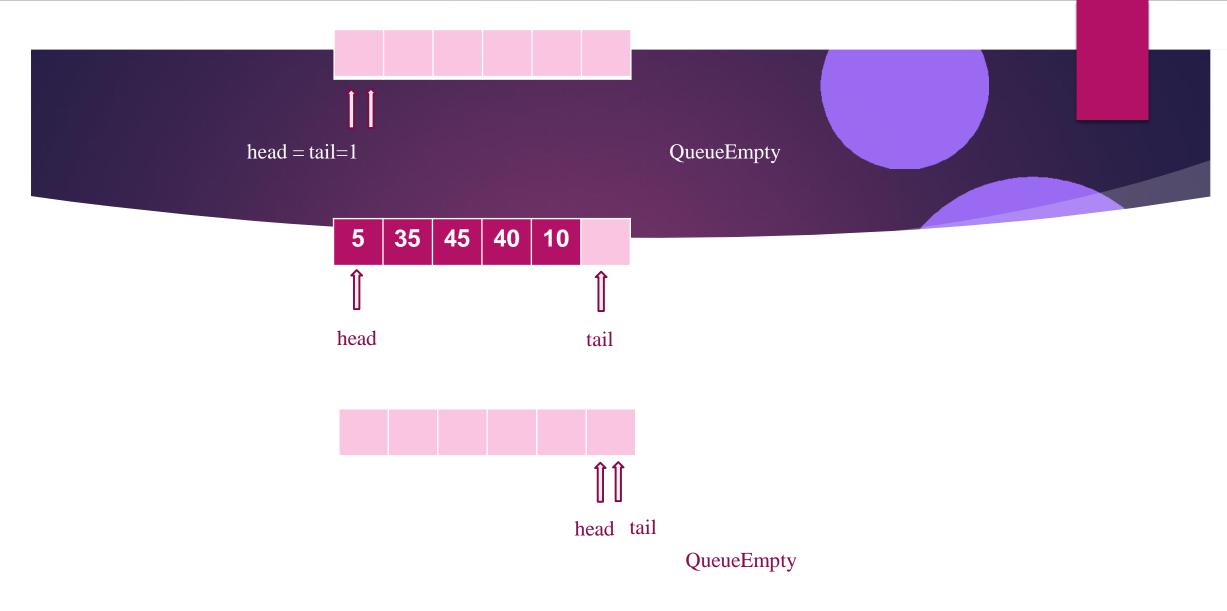
- One slot left vacant (maximum n-1 elements can be queued in an n element array)
- Q.tail points to the next location where a new element can be added
- Makes checking QueueFull() / QueueEmpty() easier

#### Queue - Array Based Implementation

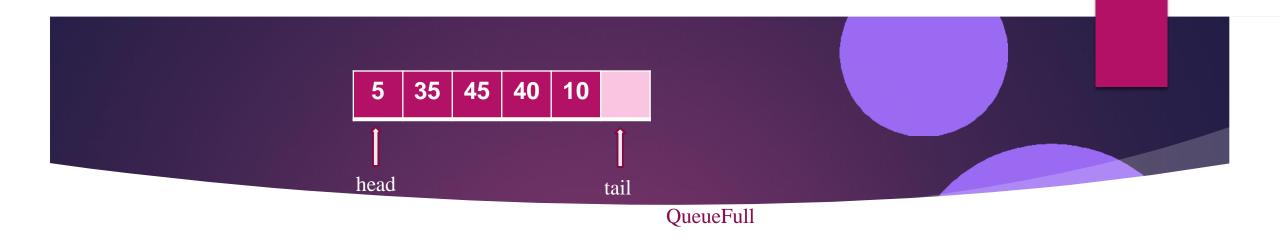
- Array Q[1..n]
- Q.head points to actual head
- Q.tail points to the next location for insertion
- □ Initially Q.head =Q.tail =1

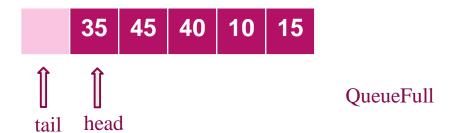






**QueueEmpty:** Q.head = Q.tail





Think about how do we write QueueFull: Q.head = (..... Q.tail....)?

#### Example

n=6, Q is initially empty

ENQUEUE (Q,8)

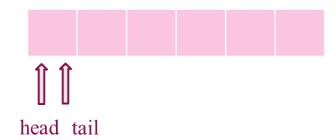
ENQUEUE (Q, 2)

ENQUEUE (Q, 3)

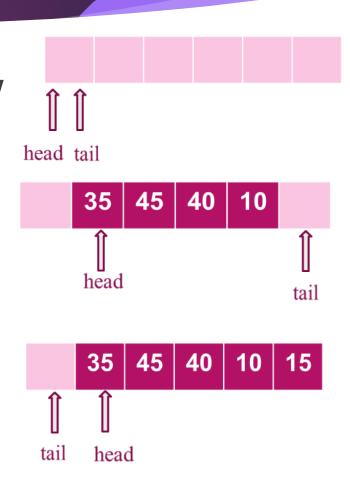
DEQUEUE(Q)

ENQUEUE (Q, 16)

DEQUEUE(Q)



ENQUEUE (Q, x) //CLRS: ignoring oveflow
Q[Q.tail]=x
if Q.tail == Q.length
Q.tail =1
else Q.tail = Q.tail+1



DEQUEUE (Q) //CLRS: ignoring underflow

```
x = Q[Q. head]
```

if (Q.head == Q.length)

Q.head=1

else Q.head = Q.head+1

return x



### Overflow/Underflow

- □ Write Algorithms for
  - □ QUEUE-EMPTY()
  - □ QUEUE-FULL()

Stack, Queue - Pointer based implementation

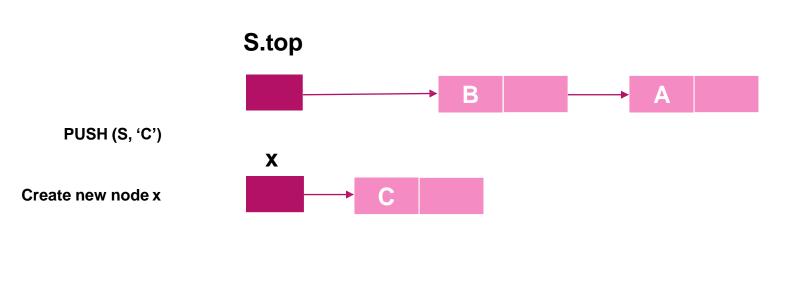
#### Stack - Basics

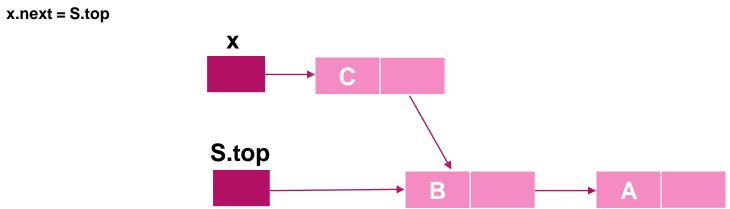
- List of elements, INSERT / DELETE only at one end of the list
  - Access restriction
  - Last-in, First-out (LIFO)
  - The last inserted element is the first one to be removed

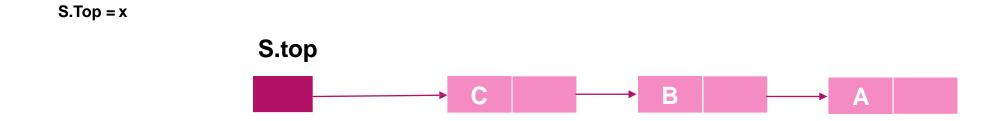
- As a linked list of nodes
- □ Top pointer to the front node
- □ PUSH() Insert to front of the list
- POP()- Delete the front node

- Stack S is a linked list
- Attribute S.top
  - points to the top element (node at the front of the list)
- □ PUSH(S, x)
  - □ Insert node x as the new top node
- □ POP(S)
  - □ Removes the node pointed to by S.top

## S.top NIL **PUSH (S, 'A')** PUSH (S, 'B') PUSH (S, 'C')

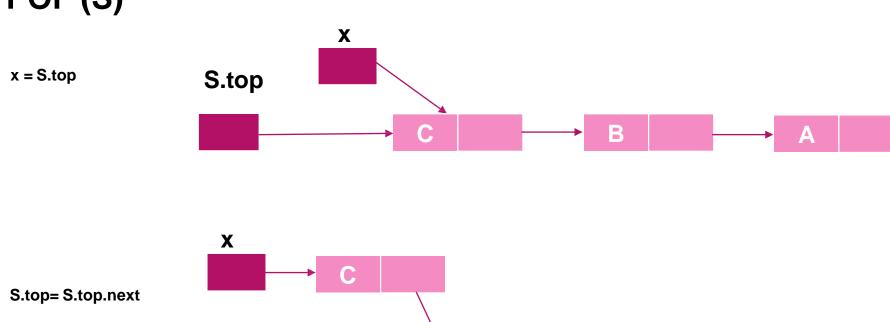


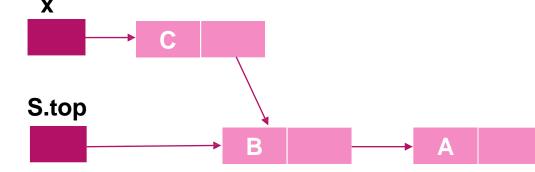




Think about what has to be done, If S.top is NIL

### S.top POP(S) X S.top







```
POP (S)

if (STACK-EMPTY(S))

error "underflow"

x = S.top

S.top = S.top.next

return x;
```

STACK-EMPTY (S)

if S.top == NIL

return true

return false

- > Time Complexity of operations
  - > PUSH ?
  - > POP ?
  - > STACKEMPTY() ?

- > Time Complexity of operations
  - > PUSH O(1)
  - > POP O(1)
  - > STACKEMPTY() O(1)

#### Implementation Details

```
struct node { //a node in the stack
     ElemType elem;
     struct node *next
     };

struct stack {
     struct node * top;
     };
```

#### Queue - Basics

- Queue List with access restrictions
  - ☐ FIFO First-In First-Out
  - Double ended- Head and Tail (front and rear)
  - Insertion always to the tail
  - Deletion always from the head

## Queue – Pointer Based Implementation

- As a linked list of nodes
- head pointer to the first node (at the front end)
- tail pointer to the last node (at the rear end)
- □ ENQUEUE() insert to end of the list
- DEQUEUE()- delete the first node

### Queue – Pointer Based Implementation

- Queue Q as a linked list
- Attribute Q.head
  - points to the first element (node at the front of the list)
- Attribute Q.tail
  - points to the last element (node at the rear of the list)
- $\Box$  ENQUEUE(Q, x)
  - Insert node x at the tail
- DEQUEUE(Q)
  - Removes the node at the head

# Queue - Pointer Based Implementation

```
ENQUEUE (Q, x)
if (QUEUE-EMPTY (Q))
  Q.front = x
else
  Q.tail.next = x
Q.tail = x
```

## Queue - Pointer Based Implementation

```
DEQUEUE (Q)
  if (QUEUE-EMPTY (Q))
     error "underflow"
  else
     x = Q. head
     if (Q.head == Q.tail)
        Q.tail = NIL
      Q.head = Q.head.next
     return x;
```

## Queue - Pointer Based Implementation

QUEUE-EMPTY (Q)

if Q. head == NIL

return true

return false

#### Implementation Details

```
struct node { //a node in the queue
    ElemType elem;
    struct node *next
    };

struct Queue {
    struct node * head;
    struct node * tail;
    };
```

- > Time Complexity of operations
  - > ENQUEUE ?
  - > DEQUEUE ?
  - > QUEUE-EMPTY() ?

- > Time Complexity of operations
  - > ENQUEUE O(1)
  - DEQUEUE O(1)
  - ➤ QUEUE-EMPTY() O(1)

#### Reference

T H Cormen, C E Leiserson, R L Rivest, C Stein *Introduction to Algorithms*, 3<sup>rd</sup> ed., PHI, 2010