



T.C. $O(n \log(n))$
S.C. $O(1)$

Leetcode Daily Challenge

24/03/2022



problem

Boats to Save People

pre-requisites

sorting, 2 pointer

difficulty
Medium

est. time
10-15 min

Let's build **Intuition**

can be asked in...



50%
Accuracy



Statement

Description

- You are given an array `people` where `people[i]` is the weight of the *i*th person, and an infinite number of boats where each boat can carry a maximum weight of `limit`. Each boat carries at most two people at the same time, provided the sum of the weight of those people is at most `limit`.
- Return the minimum number of boats to carry every given person.

i/p

```
people = [3,2,2,1],  
limit = 3
```

o/p

3

explanation

3 boats (1, 2), (2) and (3)



Observation

What are we given

- `people[]` (weights of `n` people)
- `limit` (limit of boat)
- we can put at max 2 people `i, j` in boat if $\text{people}[i] + \text{people}[j] \leq \text{limit}$
- return min boats required to carry all of them

Let's solve test cases...

consider this ex.

`people[] = {3,8,1,6}` `limit = 9`

min boats required = 2 (`{3,6}`, `{1,8}`)

I would encourage you to solve few more test cases before jumping to any solid conclusion

from above test case let's get some observation

next slide ----->



Observation

`people[] = {3,8,1,6} limit = 9` **test case #1**

there can be 3 ways with which we can pair 2 people in a boat

- 1 • 2 heavy weight people
- 2 • 2 light weight people
- 3 • 1 heavy + 1 light

1 2 heavy weight people

their combined weight can exceed boat limit.

from test case #1 $8 + 6 > \text{limit} (14 > 9)$

2 2 light weight people

their combined weight may be way less than boat limit & thus their would be vacant limit that we are just wasting(which we could have used).

from test case #1 $3 + 1 < \text{limit} (4 < 9)$

3 1 heavy + 1 light weight people

optimal

their combined weight would be closest to fill the max boat limit, hence the wastage of boat limit can be avoided.

from test case #1 $1 + 8 \leq \text{limit} (9 \leq 9)$



Observation

from previous example we concluded to pair

1 light + 1 heavy weight person.

how can we get 1 light + 1 heavy person ?

by **SORTING**

we can sort the people[] in asc. so left ones are light weight & right ones are heavy

now we want 1 person from left & 1 from right, what we can use ?

2 POINTER

1 pointer starts from extreme left & other from extreme right

final conclusion

we will sort the people[] array & use 2 pointer technique to traverse the array



Algoritihm

```
people[] = {3,8,1,6,4} limit = 9
```

```
int cnt = 0 (cnt = count of boats required)
```

- **sort people[]**

```
people[] = {1,3,4,6,8}
```

- **take 2 pointers**

```
l = 0, h = n-1      (n = people.size())
```

- **if(people[l] + people[h] <= limit)**

it means we can put these 2 people in a boat so

```
cnt++;
```

```
l++;
```

```
h--;
```

- **else**

it means their weight exceeds boat limit so heavy person(h) can't be paired with anyone, so he will occupy separate boat but light one can be paired so only

```
cnt++;
```

```
h--;
```

see here we haven't done l++ as l(th) person(light weight) still has chance to be paired up with someone, & we want to pair as many as possible to use least boats



```
class Solution {
public:
    int numRescueBoats(vector<int>& people, int limit) {

        sort(people.begin(), people.end());

        int cnt = 0;
        int l = 0;
        int h = people.size() - 1;

        while(l <= h) {

            if(people[l] + people[h] <= limit) {
                l++;
            }

            h--;
            cnt++;
        }

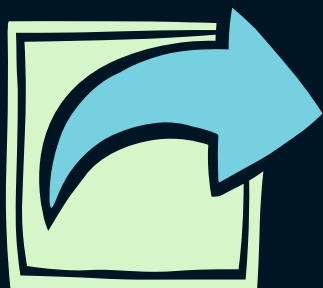
        return cnt;
    }
};
```



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