1. Perform the pairwise disjointness test for the following grammar rules.

a.) S
$$\rightarrow$$
 aSb | bAA

Step 1.

$$FIRST(aSb) = \{a\}$$
 $FIRST(bAA) = \{b\}$

Step 2.

 $FIRST(aSb) \cap FIRST(bAA) = \emptyset$
Is the same as,
 $\{a\} \cap \{b\} = \emptyset$

Answer: Passes the test

b.
$$A \rightarrow B\{aB\} \mid a$$

Step 1.

$$FIRST(B\{aB\}) = \{b\}$$
 $FIRST(a) = \{a\}$

Step 2.

 $FIRST(B\{aB\}) \cap FIRST(a) = \emptyset$

Is the same as,
 $\{b\} \cap \{a\} = \emptyset$

Answer: Passes the test

c. $B \rightarrow aB \mid a$

Step 1.

$$FIRST(aB) = \{a\}$$
 $FIRST(a) = \{a\}$

Step 2.

 $FIRST(aB) \cap FIRST(a) = \{a\} \neq \emptyset$

Is the same as,
 $\{a\} \cap \{a\} = \{a\} \neq \emptyset$

Answer: Fails the test

2. Show a trace of the recursive descent parser given in Section 4.4.1 (in the book) for the string a + b * c .

Next token is: 11 Next lexeme is a

Enter <expr>
Enter <term>
Enter <factor>

Next token is: 21 Next lexeme is +

Exit <factor>
Exit <term>

Next token is: 11 Next lexeme is b

Enter <term>
Enter <factor>

Next token is: 23 Next lexeme is *

Exit <factor>

Next token is: 11 Next lexeme is c

Enter <factor>

Next token is: -1 Next lexeme is EOF

Exit <factor>
Exit <term>

Exit <expr>

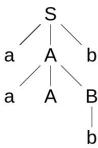
3. Given the following grammar and the right sentential form, draw a parse tree and show the phrases and simple phrases, as well as the handle.

a. aaAbb

Phrases: aaAbb, aaABb, aAb

Simple Phrases: b

Handle: b

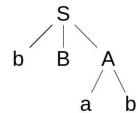


b. bBab

Phrases: bBab, bBA

Simple Phrases: ab

Handle: ab



c. aaAbBb

$$aaAbBb \rightarrow aSBb \rightarrow aSBB \rightarrow X$$

The last string cannot be derived from the given grammar, therefor the phrase, simple phrase or handle cannot be calculated.

4. Design a state diagram to recognize the floating-point literals of your favorite programming language. Similar to the state diagram from the book that I used numerous times in the lecture videos.

A state diagram to recognize the floating-point literals in C

