

March 3

Ex 4. $E = \{a^i b^j c^k \mid 0 \leq i \leq j \leq k\}$ is not CF.

Proof: Assume that E is CF, select $s = a^p b^p c^p$, where p is the pumping length. By the pumping lemma, s can be decomposed into $s = uvxyz$, s.t.

① $uv^i xy^i z \in E$ for $i \geq 0$,

② $|v| > 0$,

③ $|vxy| \leq p$.

a) When both v and y contain only one type of symbols (e.g., a 's and b 's),

a.1) a 's do not appear in v or y .

$uv^0 xy^0 z = uxz \notin E$ // # of a 's would be more than b 's or c 's

a.2) b 's do not appear

If a 's appear in v or y , $uv^2 xy^2 z \notin E$ // # of a 's \uparrow

If c 's appear in v or y , $uv^0 xy^0 z \notin E$ // as in a.1)

a.3) c 's do not appear

We pump up, $uv^2 xy^2 z \notin E$ // # of c 's $<$ # of a 's or b 's.

b) Either v or y contains more than one type of symbols, $uv^2 xy^2 z \notin E$ as we could have either $a-b-a-b$ or $b-c-b-c$ as subsequences (i.e., out of order).

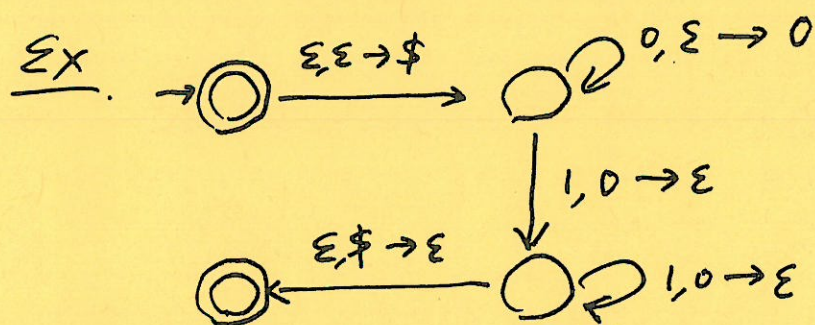
\therefore A contradiction to the pumping lemma!

$\therefore E$ is not CF.

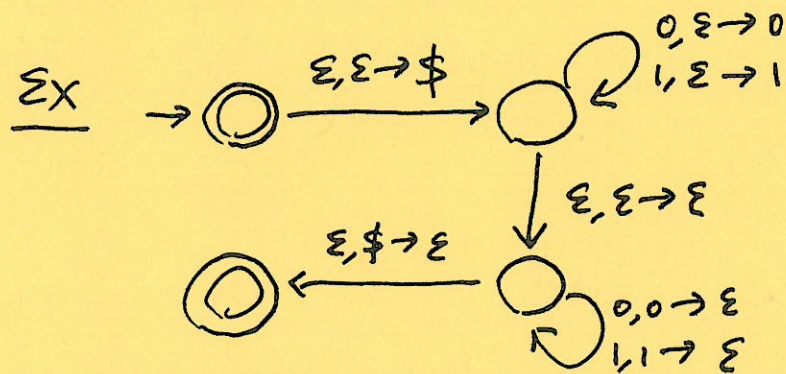
□

In Compilers, etc, we use deterministic CFL,
or deterministic PDA.

Exactly one of $\left. \begin{array}{l} \delta(q, a, x) \\ \delta(q, a, \epsilon) \\ \delta(q, \epsilon, x) \\ \delta(q, \epsilon, \epsilon) \end{array} \right\}$ is not ϕ



deterministic



non-deterministic

The diagonalization method (Cantor, 1873)

- How do we measure the size of a set?

Finite set — just count!

Infinite set — not trivial

$N = \{1, 2, 3, 4, \dots\}$

R — real numbers

- Def. $f: A \rightarrow B$

f is one to one if $f(a) \neq f(b)$ whenever $a \neq b$.

$$\text{abs}(x) = \begin{cases} x, & \text{if } x \geq 0 \\ -x, & \text{if } x < 0 \end{cases}$$

$$\text{abs}(-2) = 2$$

$$\text{abs}(2) = 2$$

$$\begin{array}{ccc} -2 & \xrightarrow{\quad} & 2 \\ 2 & \xrightarrow{\quad} & 2 \end{array}$$

f is onto if it hits every element of B .

$$\text{Square}(x) = x^2$$

- 3 can't be hit when the domain is N .

f is a correspondence if it is one-to-one and it is also onto.

A and B are of the same size if there is a correspondence between them.

- Def. A set A is countable if it is finite or it has the same size as N (natural numbers, or $\{1, 2, 3, \dots\}$).

Ex 1. $f(n) = 2n, n \in \mathbb{N}$ (set of ^{positive} even integers).

Ex 2. $Q = \{m/n \mid m, n \in \mathbb{N}\}$ is countable.

$m \backslash n$	1	2	3	4	...
1	$\frac{1}{1}$	$\frac{1}{2}$	$\frac{1}{3}$	$\frac{1}{4}$...
2	$\frac{2}{1}$	$\frac{2}{2}$	$\frac{2}{3}$	$\frac{2}{4}$...
3	$\frac{3}{1}$	$\frac{3}{2}$	$\frac{3}{3}$	$\frac{3}{4}$...
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Def. If an infinite set has no correspondence with \mathbb{N} , then it is called uncountable.

Def. \mathbb{R} - set of real numbers.

Thm 4.17 \mathbb{R} is uncountable.