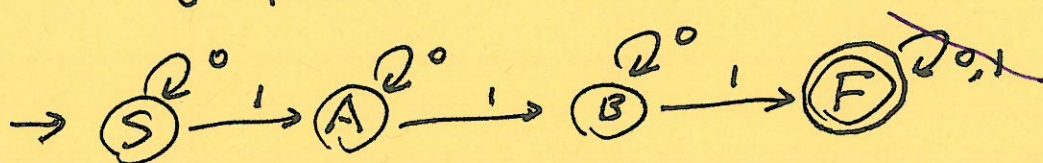


Jan 27

Questions from the previous lecture:

① EX1: $\{w \mid w \text{ contains at least three 1's}\}$ $\Sigma = \{0,1\}$.



What if I delete the transition at \textcircled{F} ?

Answer: You would miss some string like

1110000, 010100111111, ...

So, logically it would be an incomplete solution.

② $A \circ B = \{xy \mid x \in A, y \in B\}$.

How important is the order xy ?

Answer: We just use an example.

Let $A = \{0^*\}$, $B = \{1^+\}$.

$00011, 011111 \in A \circ B$.

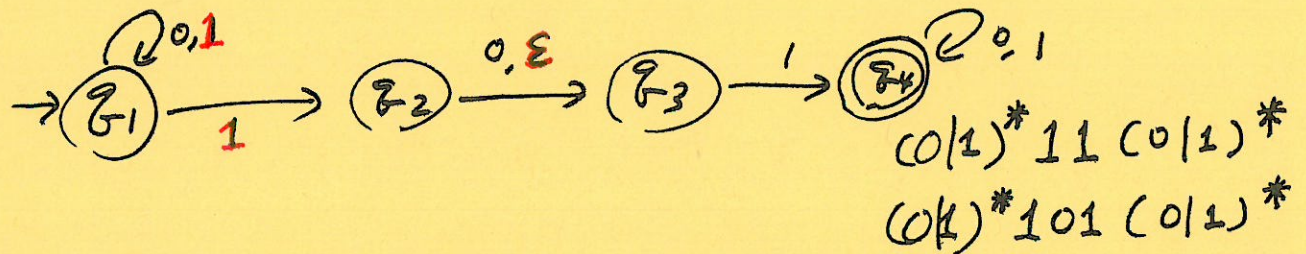
$11111 \in A \circ B \quad \parallel x = \varepsilon = 0^0$.

But $111100 \notin A \circ B$.

So, in essence, y can't be preceding x !

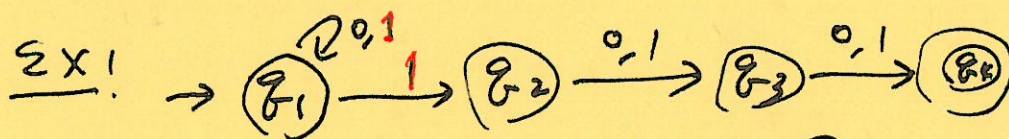
Deterministic — Given the current state and input, we know where to go next exactly.

Non-determinism —

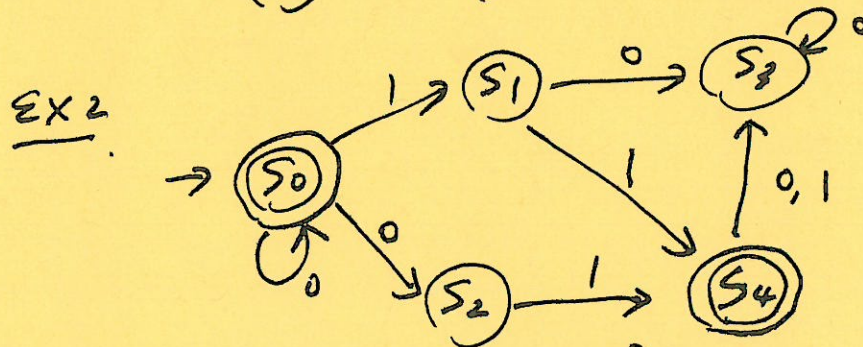


How does an NFA compute?

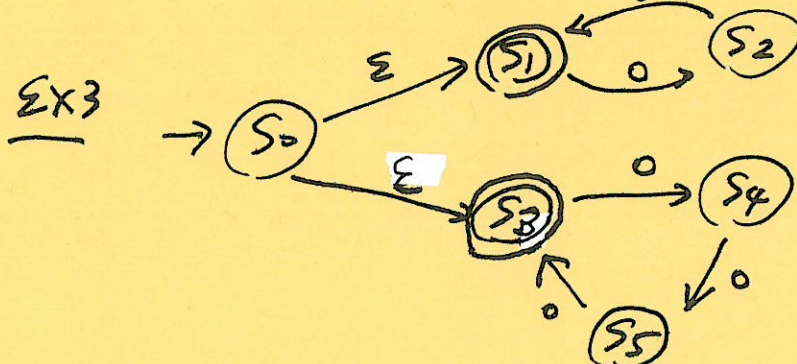
- If any one of the accept state can be reached by reading the input x in any way, we say the NFA accepts x .



$(0/1)^*1(1/0)(1/0)$
 $x^*1xx, x \in \{0,1\}$



$\{0^n, 0^n11, 0^n01\}, n \geq 0$



$\{0^k \mid k \text{ is a multiple of 2 or 3}\}$

- A nondeterministic finite automaton (NFA) is a 5-tuple $(Q, \Sigma, \delta, q_0, F)$, where

1. Q is the set of states
 2. Σ is a finite alphabet
 3. $\delta: Q \times \Sigma_\epsilon \rightarrow P(Q)$ is the transition function
 $\Sigma_\epsilon = \Sigma \cup \{\epsilon\}$. $P(Q)$ — power set of Q .
 4. $q_0 \in Q$ is the start state
 5. $F \subseteq Q$ is the set of accept states.
-

Ex 3 $\delta(s_0, \epsilon) = \{s_1, s_3\}$

Ex 1 $\delta(q_1, 1) = \{q_1, q_2\}$

- Let $N = (Q, \Sigma, \delta, q_0, F)$ be an NFA and $w = w_1 \dots w_m$
 $w_i \in \Sigma_\epsilon$

N accepts w if

$\exists r_0, r_1, \dots, r_m \in Q$, s.t.

① $r_0 = q_0$

② $r_{i+1} \in \delta(r_i, w_{i+1})$, $i = 0, \dots, m-1$

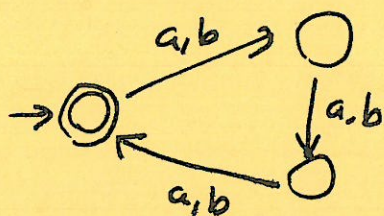
③ $r_m \in F$.

- Two machines are equivalent if they accept the same language.

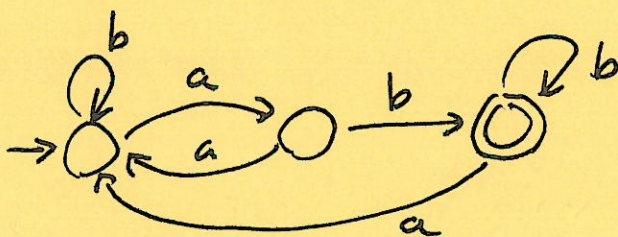
Exercises on DFA:

Let $\Sigma = \{a, b\}$, construct DFA's for the following languages.

- ① $A = \{w \mid \text{the length of } w \text{ is a multiple of } 3\}$



- ② $C = \{w \mid w \text{ has an odd number of } a\text{'s and ends with } b\}$



Exercises on NFA:

- ① $A = \{a^* b^* a^+\}$ →

An NFA with three states. The start state is a double circle. Transitions: from start state, 'a' loops back to start; 'ε' leads to state 1; from state 1, 'b' loops back to state 1; 'a' leads to state 2 (the final state); from state 2, 'a' loops back to state 2.

- ② $B = \{w \mid w \text{ doesn't contain } aab\}$

