Example.  $\langle \text{Expression} \rangle \rightarrow \langle \text{Expression} \rangle + \langle \text{Term} \rangle | \langle \text{Term} \rangle$   $\langle \text{Term} \rangle \rightarrow \langle \text{Term} \rangle \times \langle \text{Factor} \rangle | \langle \text{Factor} \rangle$   $\langle \text{Factor} \rangle \rightarrow \langle \langle \text{Expression} \rangle \rangle | a$   $I = \{a, +, \times, (,,)\}$ 

Example. <Sentence> -> <Noun-phrase> <Verb-phrase> <Noun-phrase> -> <Complex-Noun> | <Complex-Noun><Prep-Phrase> < complex-verb><prep-phrase> prep-phrase> -> <prep>< Complex-Noun> <complex-verb> -> < Verb> | < Verb> Noun-phrase> < complex-Noun> -> < Article> < Noun> <article> -> a | the < Noun > -> boy | girl | flower > -> with

## Definition.

A Context-free grammar is a 4-tuple  $\langle V, \Sigma, R.S \rangle$  Such that

1. Vis a finite set of variables.

2.  $\Sigma$  is a finite set of terminals.  $V \cap \Sigma = \phi$ 

3. R is a finite set of rules.

4. SEV is the Start variable.

Example

$$S \rightarrow OA1$$
 $A \rightarrow B$ 
 $B \rightarrow \lambda$ 

 $V=\{5,A,B\}, \Sigma=\{0,1,\lambda\}$  $V \cap \Sigma = \emptyset$