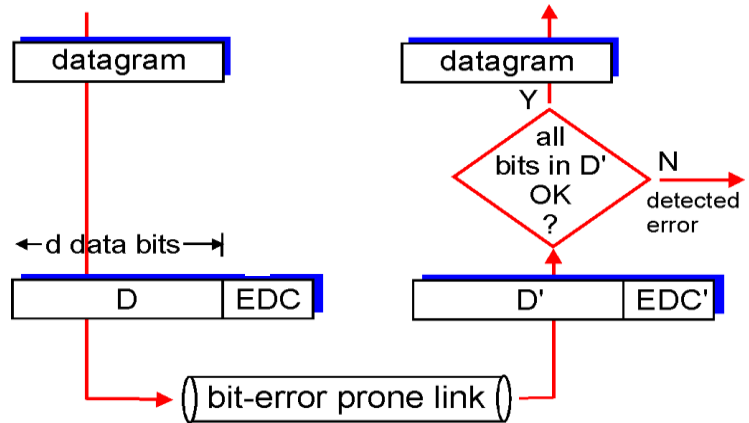


# Error detection



- EDC: Error Detection and Correction bits (redundancy)
- D: Data protected by error checking, may include header fields
- Error detection not 100% reliable!
  - Protocol may miss some errors, but rarely
  - Larger EDC field yields better detection and correction



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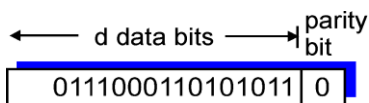
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# Parity checking



Single bit parity:

- Detects single bit error



What happens when errors occur in *bursts*? Say bits flip.

Two dimensional bit parity:

- Detects and **corrects** single bit error
- Detects two bit errors

		Row parity →			
Column parity ↓	$d_{1,1}$	...	$d_{1,j}$	$d_{1,j+1}$	
	$d_{2,1}$	...	$d_{2,j}$	$d_{2,j+1}$	
	...	...	...	...	
	$d_{i,1}$	...	$d_{i,j}$	$d_{i,j+1}$	
	$d_{i+1,1}$	...	$d_{i+1,j}$	$d_{i+1,j+1}$	
		No errors			
		Correctable single-bit error			
	1	0	1	0	1
	1	1	1	1	0
	0	1	1	1	0
	0	0	1	0	1
	0	0	1	0	1

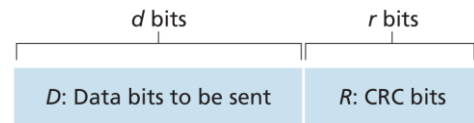
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# Cyclic Redundancy Check (CRC)



- Goal: Want a redundancy encoding that can tolerate **arbitrarily long** error bursts
- View data bits  $D$  as a binary number
- Agree on an  $r + 1$  bit generator  $G$ , with the leftmost digit 1
- Compute  $r$  bits  $R$  such that:
  - $D \ll r \oplus R = nG$  ( $nG$  – some multiple of  $G$ )
  - $D \ll r \oplus R = [D, R]$  (operation creates the packet)
- How to compute  $R$ 
  - $D \ll r \% G = R$
  - $D \ll r - R = D \ll r \oplus R = nG$   
( $\oplus$  equivalent to subtract mod 2 without carry)
- Destination uses known  $G$  to check if:
  - $(D \ll r \oplus R) \% G == 0$  (if not, then error tx error)



$1011 \text{ XOR } 0101 = 1110$   
 $1001 \text{ XOR } 1101 = 0100$   
 $1011 - 0101 = 1110$   
 $1001 - 1101 = 0100$

Can detect all burst errors up to  $r$  bits

$G_{\text{CRC-32}} = 10000010011000001000111011011011$

Widely used in practice (Ethernet, 802.11 WiFi)

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## Example



$$D = 101110 \quad G = 1011$$

$$R = D \ll r \% G = 101110\ 000 \% 1011 = 101$$

$$D \ll r - R = 101110\ 000 - 101 = 101101011$$

$$101101\ 011 \% 1011 = 0$$

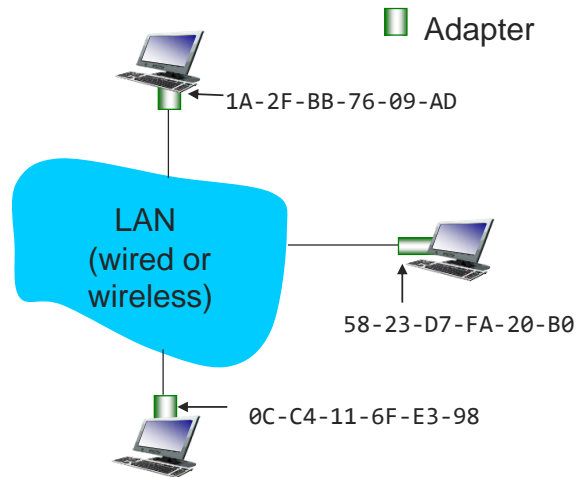
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# MAC Addresses



- MAC (or LAN or physical or Ethernet) address:
  - Addresses physical interfaces
  - Hardcoded 48 bit MAC address (from IEEE)
  - e.g.: 1A-2F-BB-76-09-AD
- Analogy:
  - MAC address: Social Security Number
  - IP address: postal address
- Flat MAC address space
  - Plus: portability between LANs
  - Minus: switch tables not hierarchical (lower scalability)



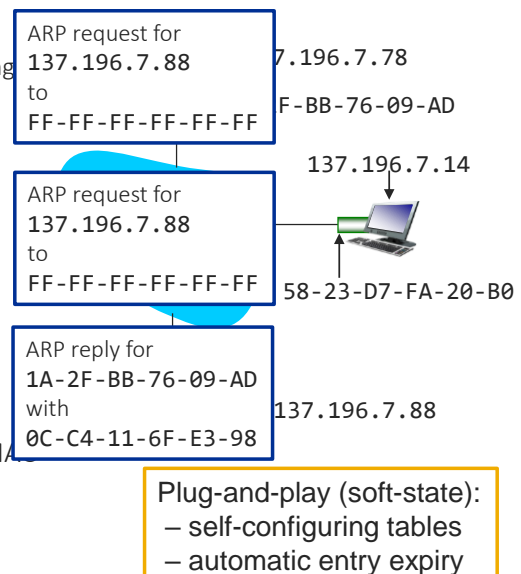
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# Address Resolution Protocol (ARP)



- Question
  - How to determine interface's MAC address, knowing IP address?
- ARP table
  - Each IP node (host, router) on LAN has table
  - < IP address; MAC address; TTL >
- A wants to send datagram to B
  - B's MAC address not in A's ARP table.
- A broadcasts ARP query packet, containing B's IP address
  - Broadcast address: FF-FF-FF-FF-FF-FF
  - All nodes on LAN receive ARP query
- B receives ARP packet, replies to A with its (B's) MAC address
  - Frame unicast to A's MAC address



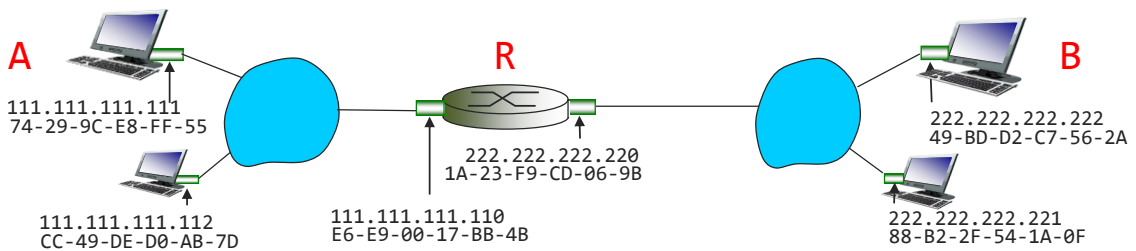
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# Routing to another LAN



- Walkthrough: send datagram from **A** to **B** via **R**
  - Focus on addressing – at IP (datagram) and MAC layer (frame)
  - Assume **A** knows **B**'s IP address. How? ← DNS
  - Assume **A** knows IP address of first hop router, **R**. How? ← DHCP
  - Assume **A** knows **R**'s MAC address. How? ← ARP



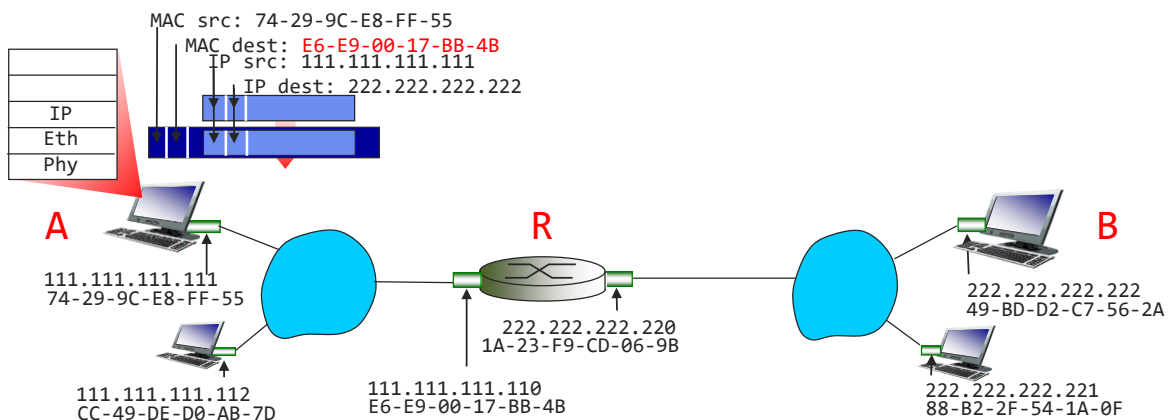
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# Routing to another LAN



- **A** creates IP datagram with IP source **A**, destination **B**
- **A** creates link-layer frame with **R**'s MAC address as dest, frame contains **A**-to-**B** IP datagram



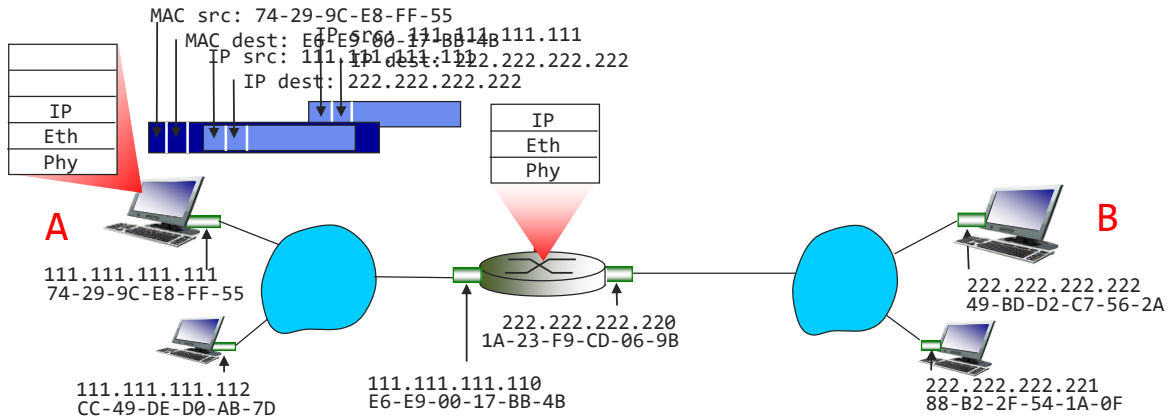
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# Routing to another LAN



- Frame sent from **A** to **R**
- Frame received at **R**, datagram removed, passed up to IP



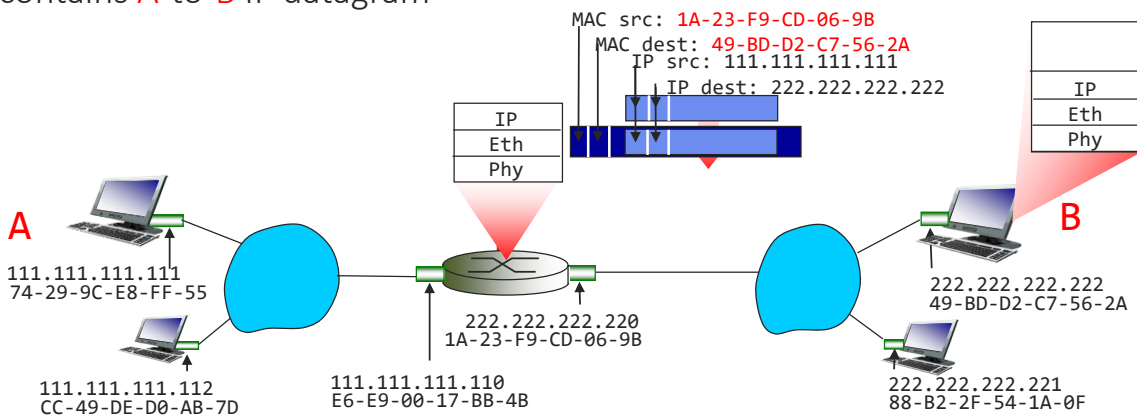
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# Routing to another LAN



- **R** forwards datagram with IP source **A**, destination **B**
- **R** creates link-layer frame with **B**'s MAC address as dest, frame contains **A-to-B** IP datagram



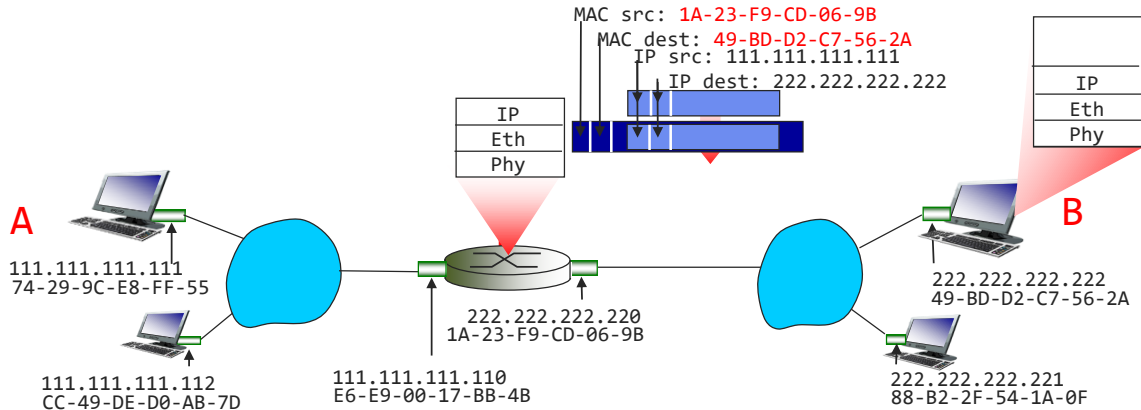
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# Routing to another LAN



- R forwards datagram with IP source **A**, destination **B**
- R creates link-layer frame with **B**'s MAC address as dest, frame contains **A**-to-**B** IP datagram



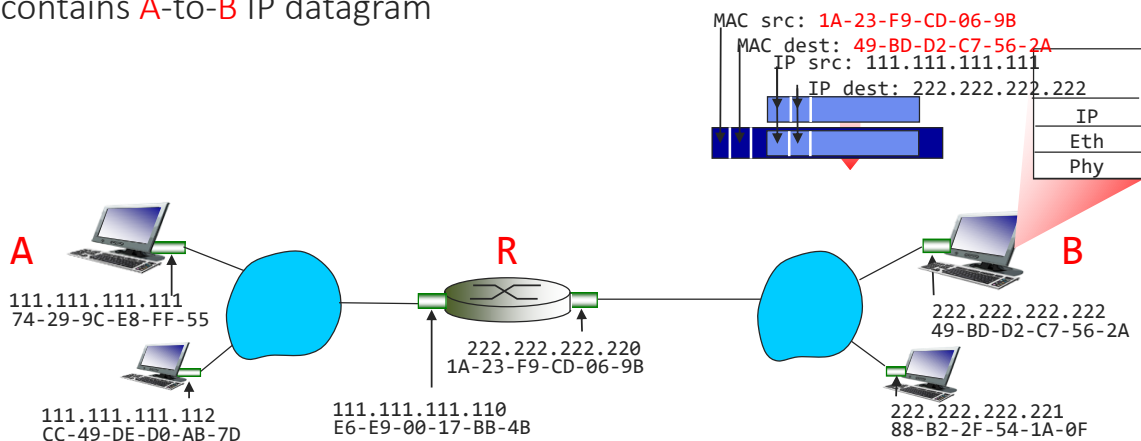
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# Routing to another LAN



- R forwards datagram with IP source **A**, destination **B**
- R creates link-layer frame with **B**'s MAC address as dest, frame contains **A**-to-**B** IP datagram



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# Exercise



Consider three LANs interconnected by two routers.

- Assign IP addresses to all of the interfaces
- Assign MAC addresses to all of the adapters
- Consider sending an IP datagram from Host A to Host E. Enumerate all the steps assuming that the ARP table in the sending host is empty (and the other tables are up to date)

