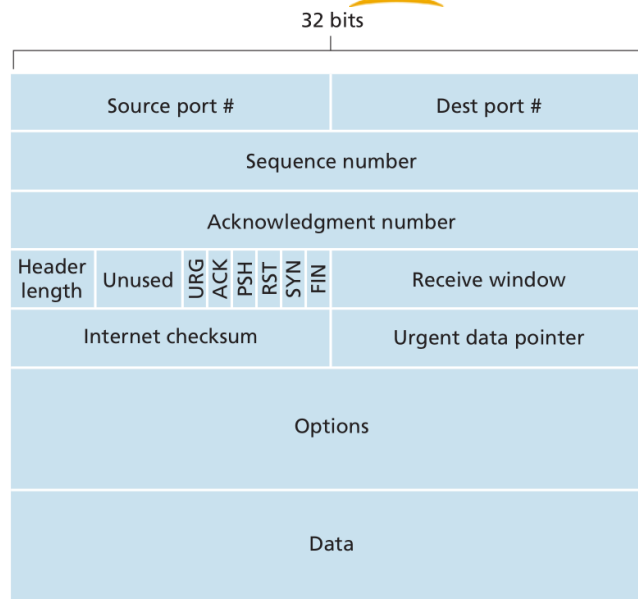


TCP packet structure



How is the TCP packet different from a UDP packet?

- UDP
 - Source port, destination port, length, checksum, data
- TCP
 - Sequence numbers
 - Flags
 - Receive window
 - Urgent data



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TCP 3-way handshake



Before exchanging data, sender/receiver “handshake”:

- Agree to establish connection (each knowing the other willing to establish connection)
- Agree on connection parameters
- Allocate resources

client state

LISTEN

SYNSENT

ESTAB

choose init seq num, x
send TCP SYN msg

SYNbit=1, Seq=x

SYNbit=1, Seq=y
ACKbit=1; ACKnum=x+1

received SYNACK(x)
indicates server is live;
send ACK for SYNACK;
this segment may contain
client-to-server data

ACKbit=1, ACKnum=y+1



server state

LISTEN

SYN RCVD

ESTAB

choose init seq num, y
send TCP SYNACK
msg, acking SYN

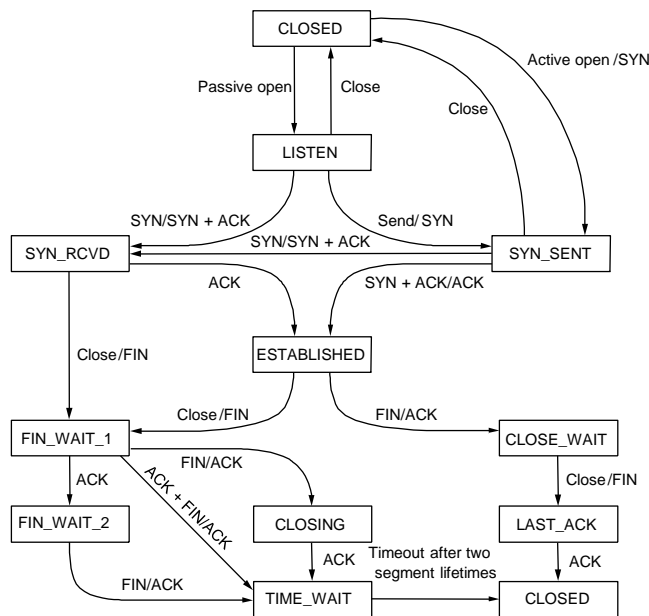
received ACK(y)
indicates client is live

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State Transition Diagram



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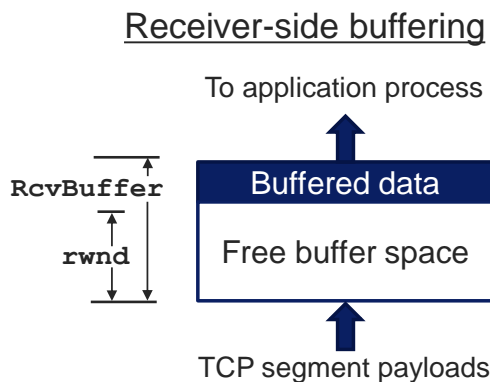
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Flow Control



- Receiver throttles sender to avoid receiver buffer overflow
- Receiver “advertises” free buffer space by including **rwnd** value in TCP header of receiver-to-sender segments
 - RcvBuffer** size set via socket options (typical default is 4096 bytes)
 - Many operating systems autoadjust **RcvBuffer**
- Sender limits amount of unacked (“in-flight”) data to receiver’s **rwnd** value
- Guarantees receive buffer will not overflow



What happens after **rwnd** set to 0?

- Sender stops sending and is not notified when buffer empties!
- Sender will continue to send 1B packets to get **rwnd** updates

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Nagle's Algorithm



- How long does sender delay sending data?
 - too long: hurts interactive applications
 - too short: poor network utilization – a packet may be mostly empty
 - strategies: timer-based vs self-clocking
- When application generates additional data
 - if fills a max segment (and window open): send it
 - else
 - if there is unack'ed data in transit: buffer it until ACK arrives
 - else: send it
- Nagle's algorithm can interact poorly with TCP's 'delayed acknowledgement' algorithm, so a way to disable it is quite commonly provided

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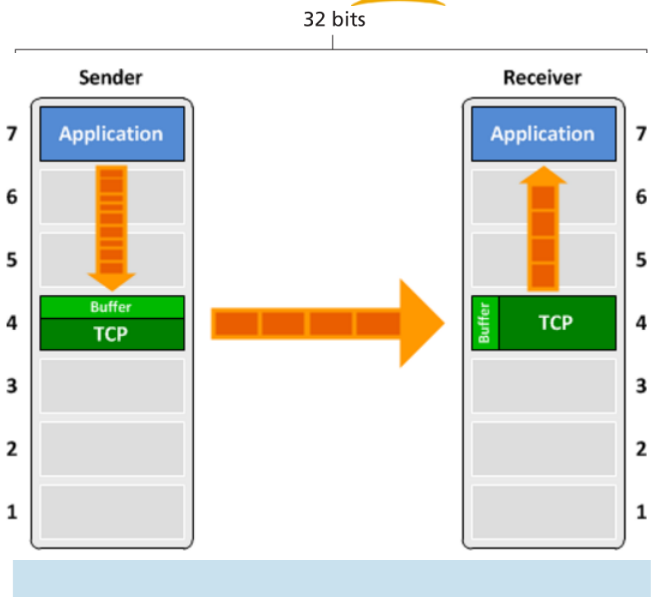
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TCP: Push vs Urgent



- PSH
 - TCP sender doesn't wait for more data (Nagle's algo)
 - Receiver forwards buffer content to application immediately
- URG
 - Receiver gets offset of urgent bytes (could be the whole packet)
 - Read by application first (potentially out of order)



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