

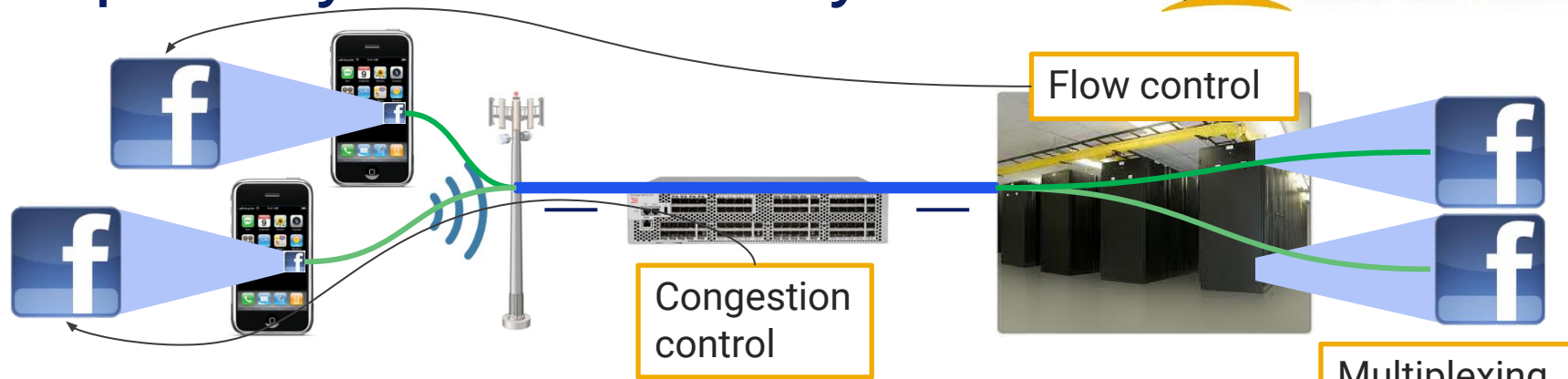


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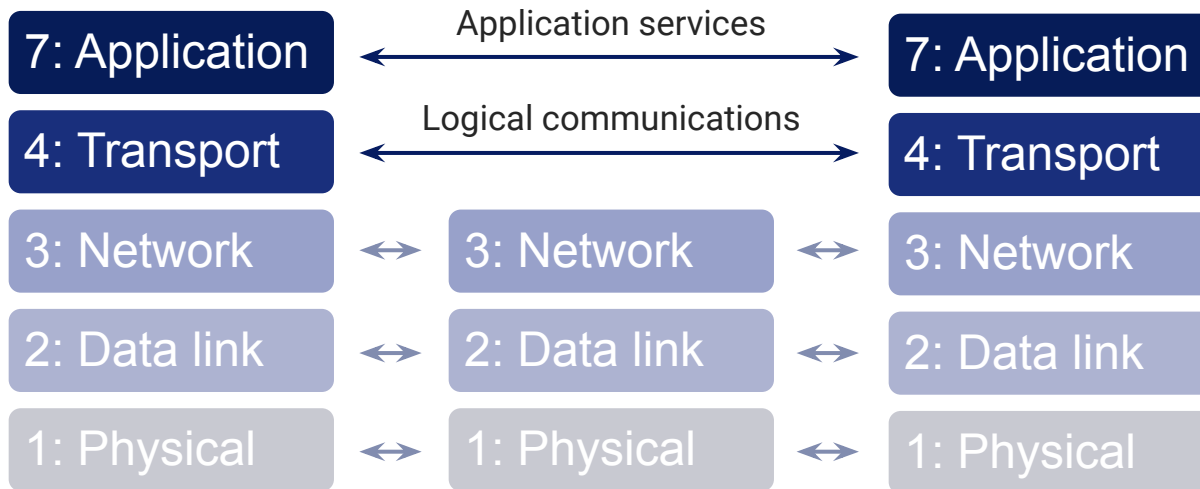
Chapter 3

RDT

Transport layer functionality



Multiplexing
via sockets

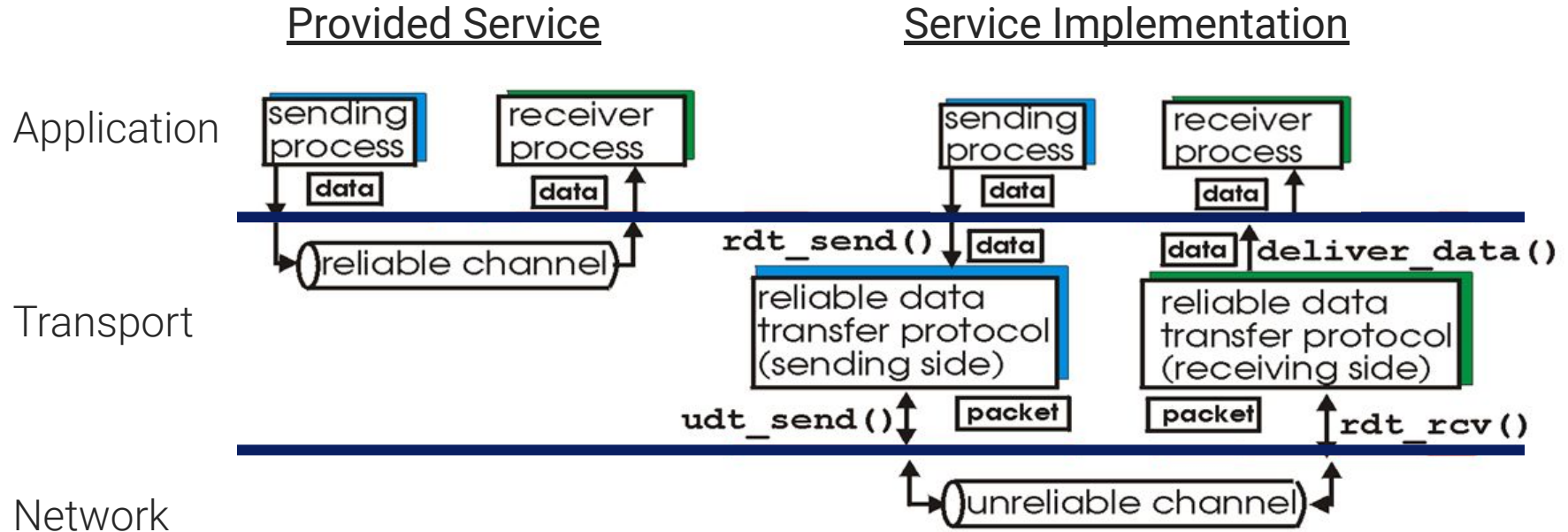


Inter-process comm.

- Segmentation and reassembly
- Error checking

- Reliability
- In-order delivery

Principles of reliable data transfer

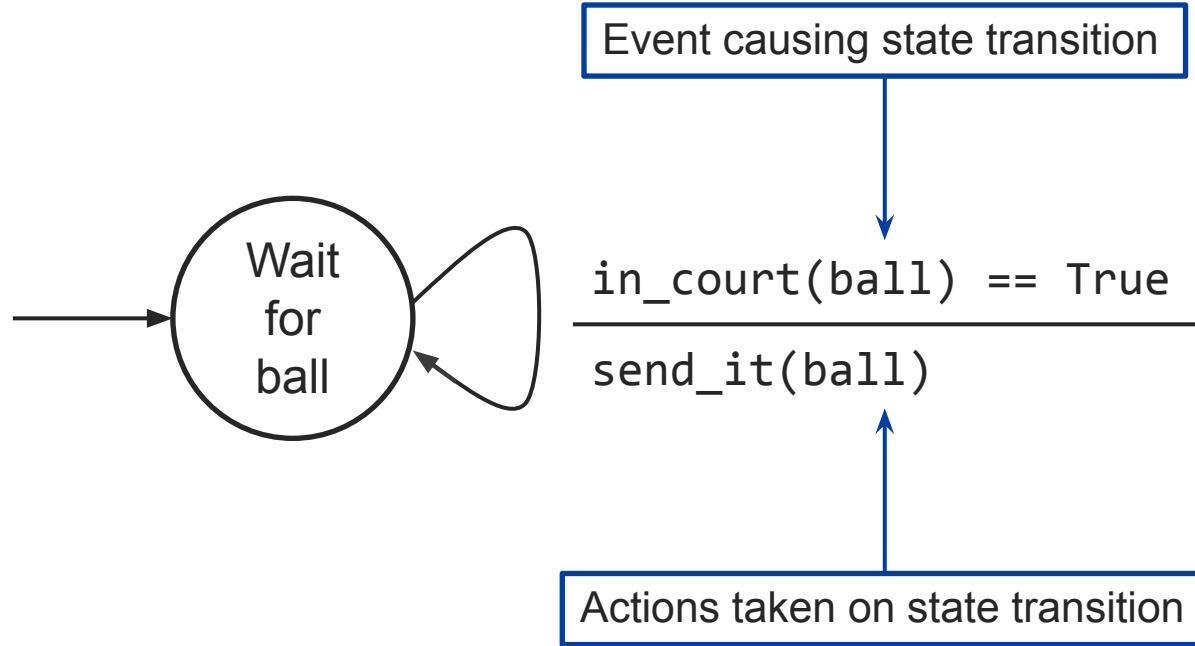


Characteristics of unreliable channel will determine complexity of reliable data transfer protocol (rdt)

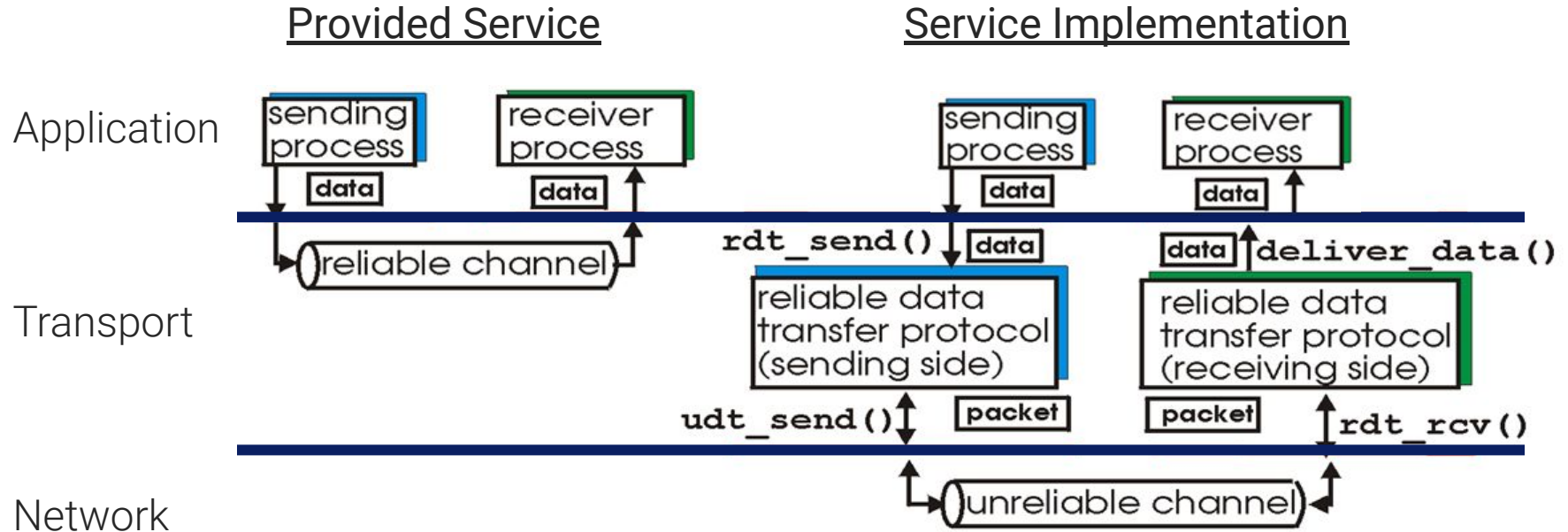
What are some ways in which the network channel can be unreliable?



Bruce Lee FSM



Principles of reliable data transfer

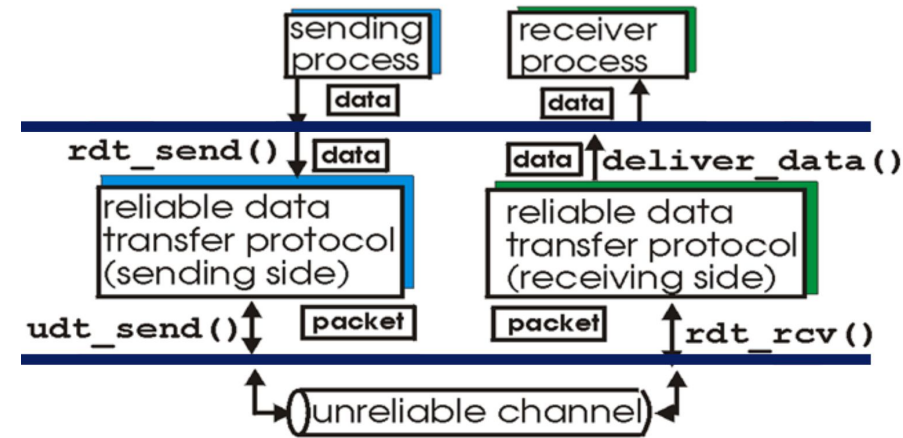


Characteristics of unreliable channel will determine complexity of reliable data transfer protocol (rdt)

What are some ways in which the network channel can be unreliable?

Reliable channel

- rdt1.0: reliable transfer over reliable channel
- Assumptions:
 - Unidirectional, long data flows
 - Perfectly reliable channel:
 - No bit errors
 - No packet loss
 - No packet reordering



Event causing state transition

Sender

Receiver



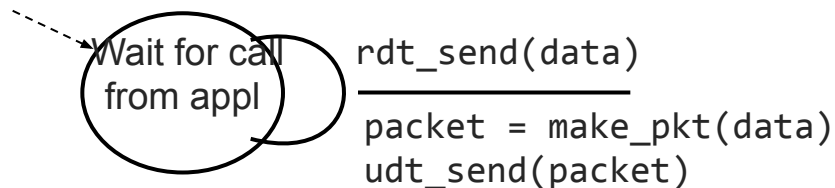
Actions taken on state transition

rdt2.0: Channel with bit errors

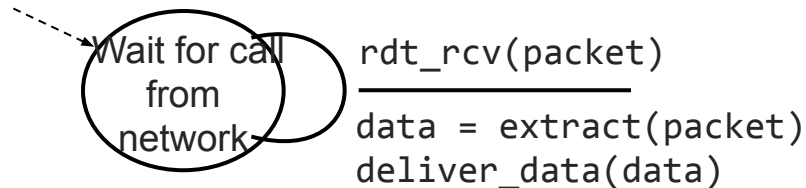
- How are errors detected?
 - Checksums: `make_pkt(data, checksum)`, `corrupt(rcvpkt)`
- How do humans recover from communication errors?
 - ACKs, NAKs, and retransmissions: `isACK(rcvpkt)`, `isNAK(rcvpkt)`
- Design sender and receiver FSMs for rdt2.0

rdt1.0:

Sender



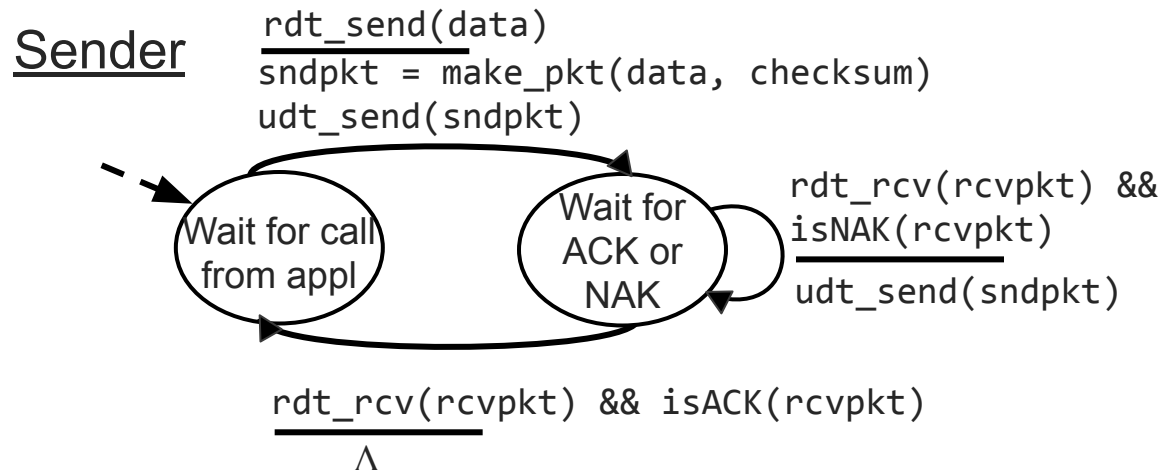
Receiver



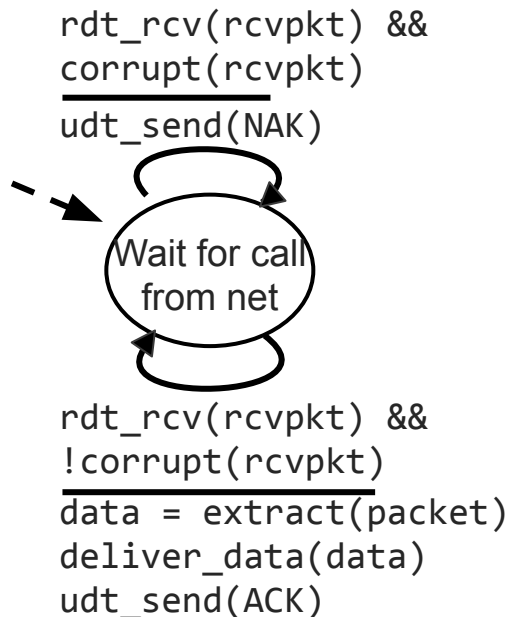
rdt2.0: Channel with bit errors

- How are errors detected?
 - Checksums
- How do humans recover from communication errors?
 - ACKs, NAKs, and retransmissions
- Design sender and receiver FSMs for rdt2.0

Stop-and-wait: sender sends one packet, then waits for receiver response



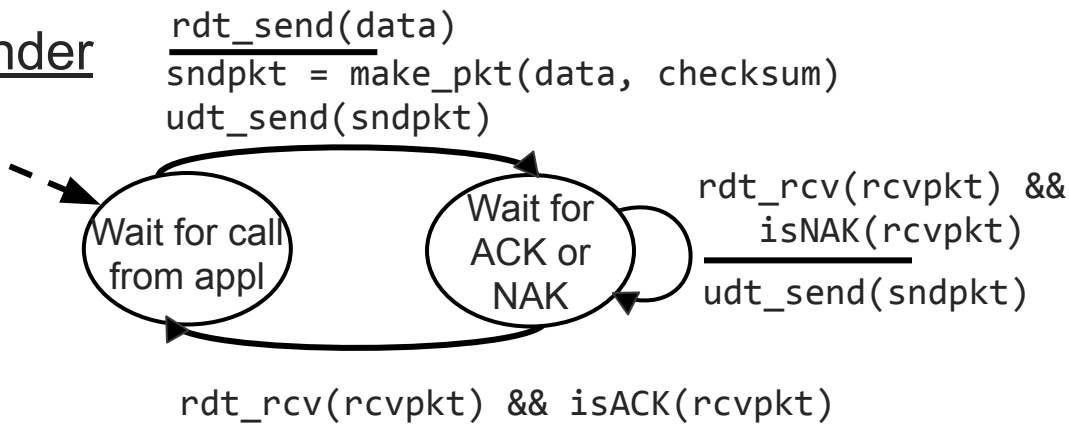
Receiver



rdt2.0: Channel with bit errors

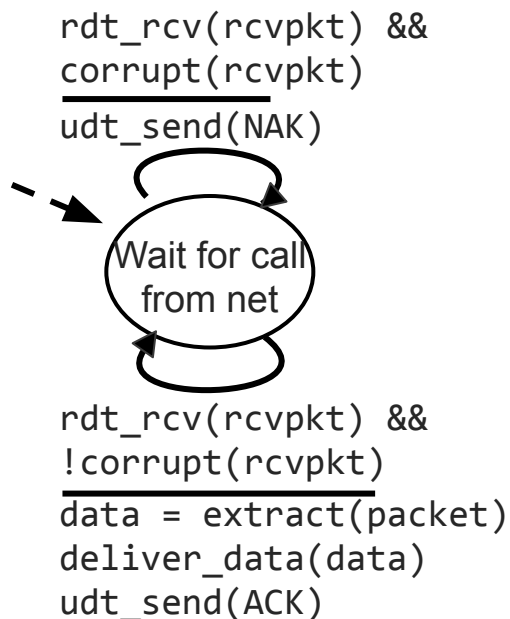
- What happens if ACK/NAK corrupted?
 - Duplicate delivery, or no retransmission when needed
- How can we deal with corrupted ACKs/NAKs?
 - Retransmission, but can get duplicate packets
- How to handle duplicate packets?
 - Embed sequence numbers in packets:
`make_pkt(seq_num, data, checksum), get_seq_num(rcvpkt)`
 - Only need 0 and 1 for `seq_num`. Why?
- Come up with sender and receiver FSMs for rdt2.0 with sequence numbers and retransmissions

Sender

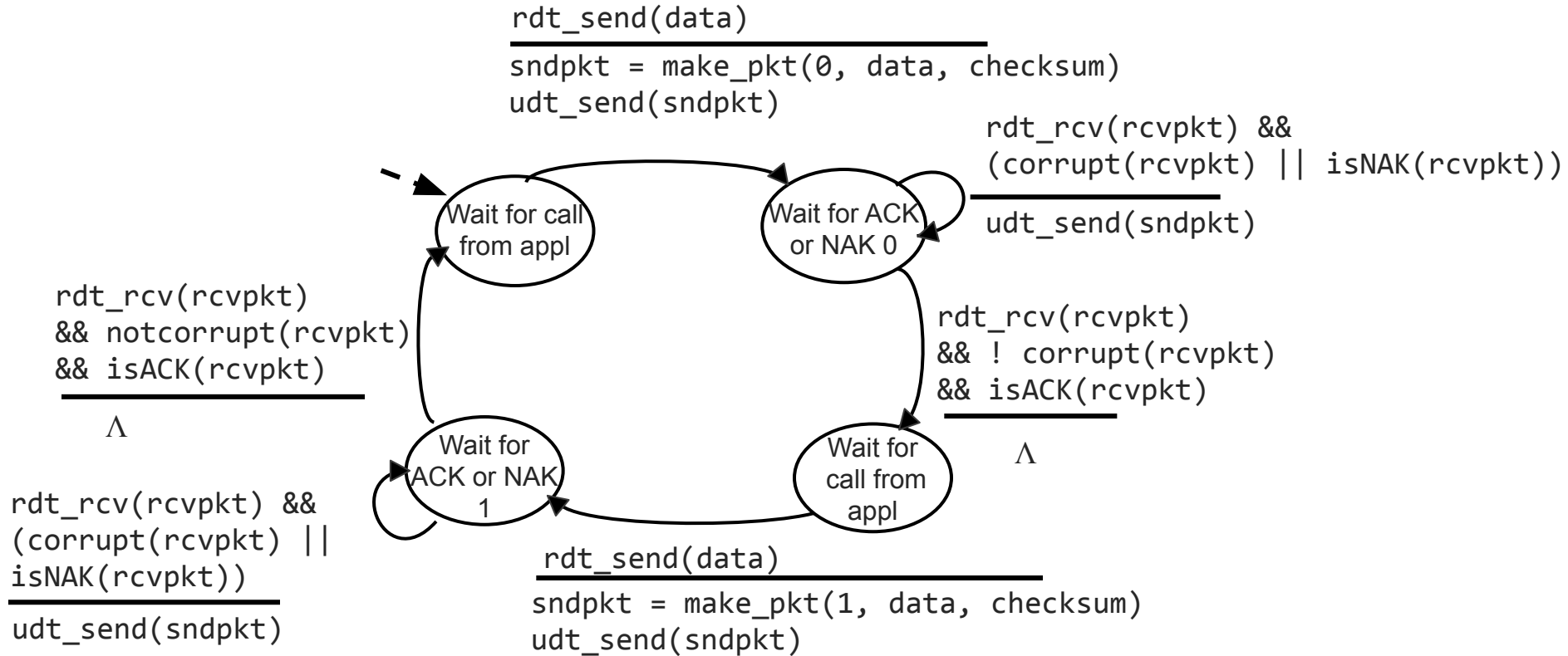


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Receiver



rdt2.1: sender, handles garbled ACKs



rdt2.1: receiver, handles garbled ACKs

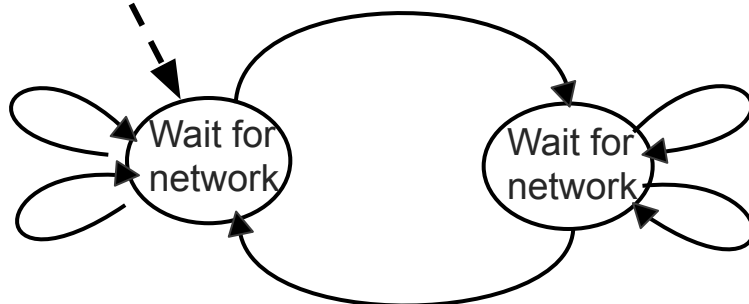
Deliver data only once

```
rdt_rcv(rcvpkt) &&  
corrupt(rcvpkt)  
-----  
sndpkt =  
    make_pkt(NAK,chksum)  
udt_send(sndpkt)
```

```
rdt_rcv(rcvpkt) &&  
not corrupt(rcvpkt) &&  
get_seq_num(rcvpkt) == 1  
-----  
sndpkt =  
    make_pkt(ACK,chksum)  
udt_send(sndpkt)
```

```
rdt_rcv(rcvpkt) && !corrupt(rcvpkt)  
    && get_seq_num(rcvpkt) == 0  
-----
```

```
data = extract(packet)  
deliver_data(data)  
sndpkt = make_pkt(ACK,chksum)  
udt_send(sndpkt)
```



```
rdt_rcv(rcvpkt) &&  
!corrupt(rcvpkt) &&  
get_seq_num(rcvpkt) == 1  
-----  
data = extract(packet)  
deliver_data(data)  
sndpkt = make_pkt(ACK,chksum)  
udt_send(sndpkt)
```

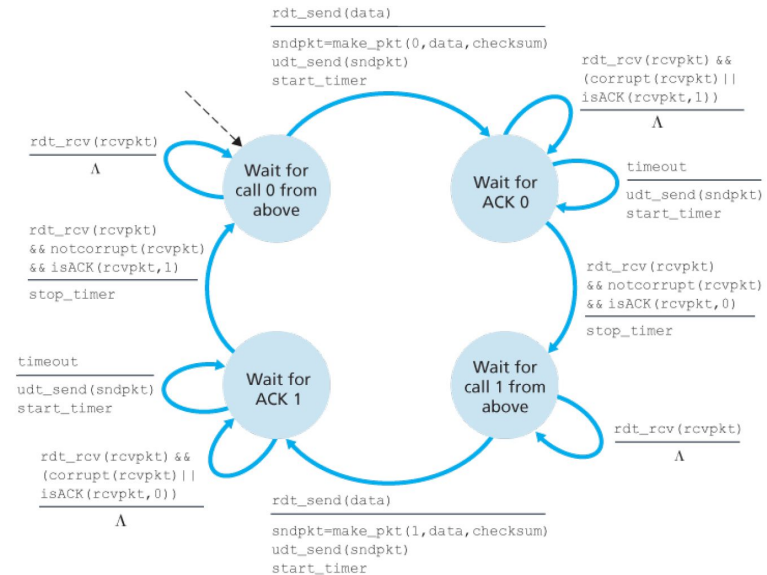
```
rdt_rcv(rcvpkt) &&  
corrupt(rcvpkt)  
-----  
sndpkt = make_pkt(NAK,chksum)  
udt_send(sndpkt)
```

```
rdt_rcv(rcvpkt) &&  
!corrupt(rcvpkt) &&  
get_seq_num(rcvpkt) == 0  
-----  
sndpkt = make_pkt(ACK,chksum)  
udt_send(sndpkt)
```

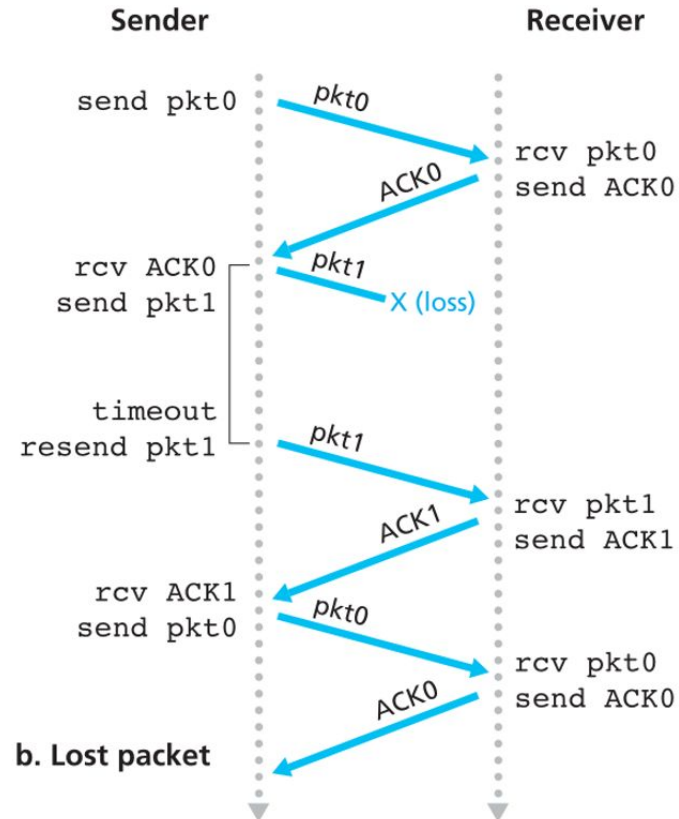
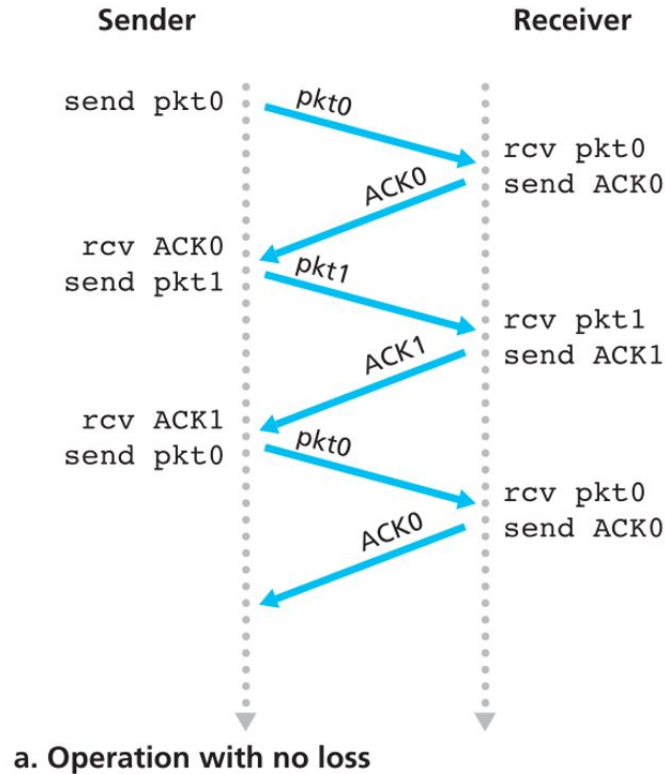
Loop on duplicates

rdt3.0: bit errors and loss

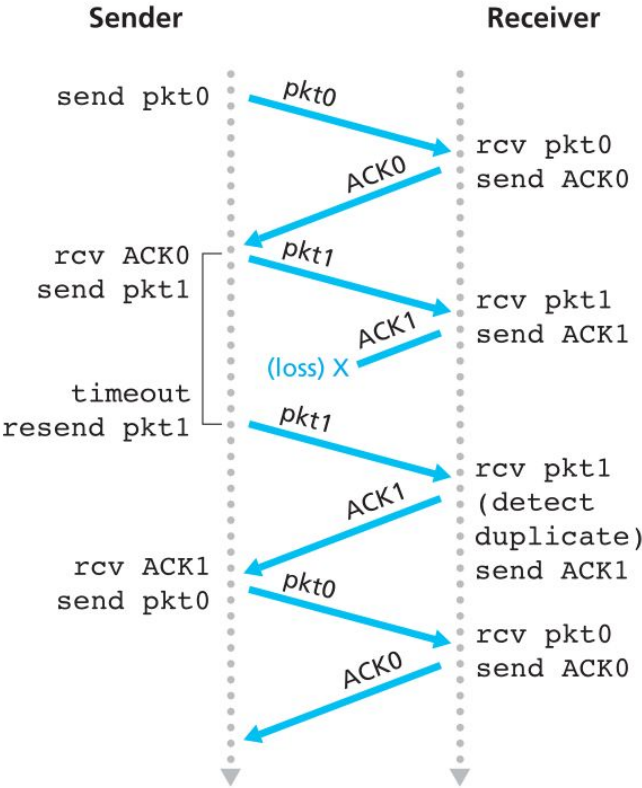
- New assumption:
 - packet loss
- How do we know a loss has occurred?
- What can we do about it?
- Approach:
 - Sender waits “reasonable” amount of time for ACK
 - Retransmits packet if no ACK received in this time
- What if packet/ACK only delayed?
 - Duplicate packets ignored at the receiver through sequence numbers
 - Receiver specifies sequences number of ACKed packet



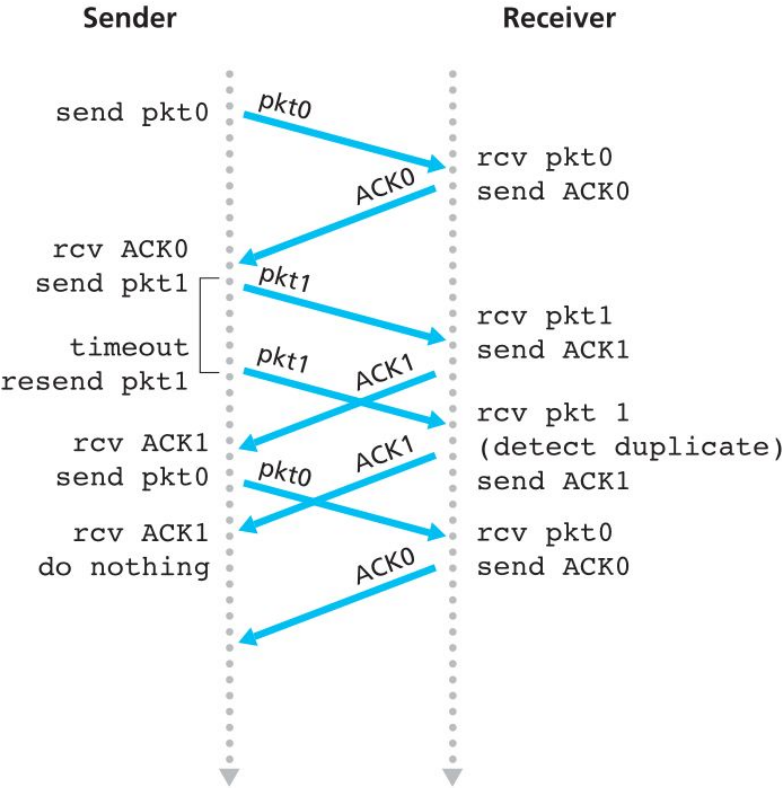
rdt3.0 in action



rdt3.0 in action



c. Lost ACK



d. Premature timeout



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