

## Definition: Natural Transformation

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A **natural transformation** provides a way to transform one **Functor** into another while respecting the internal structure of the categories involved.

#### Formal Definition

Let  $F, G : \mathcal{C} \rightarrow \mathcal{D}$  be two functors between categories  $\mathcal{C}$  and  $\mathcal{D}$ . A natural transformation  $\eta : F \Rightarrow G$  consists of:

- For each object  $X$  in  $\mathcal{C}$ , a **Morphism**  $\eta_X : F(X) \rightarrow G(X)$  in  $\mathcal{D}$

Such that for every morphism  $f : X \rightarrow Y$  in  $\mathcal{C}$ , the following diagram commutes:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(f)} & F(Y) \\ | & & | \\ X & & Y \\ \downarrow & & \downarrow \\ G(X) & \xrightarrow{G(f)} & G(Y) \end{array}$$

That is,  $\eta_Y \circ F(f) = G(f) \circ \eta_X$ .

#### Components and Naturality

- Each morphism  $\eta_X$  is called a **component** of the natural transformation
- The commutativity condition is called the **naturality condition**
- When this condition holds, we say that  $\eta$  is **natural in X**

#### Special Cases

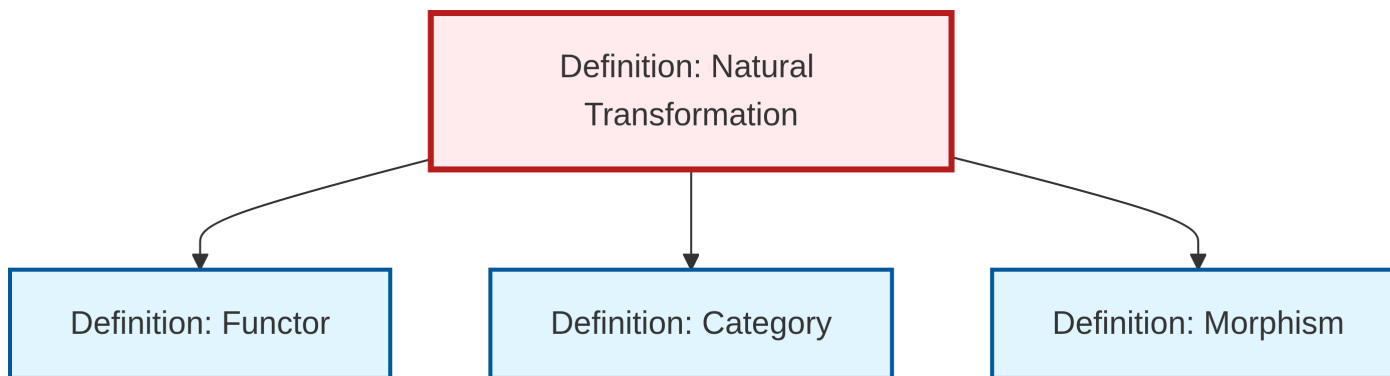
1. **Natural Isomorphism:** A natural transformation where every component  $\eta_X$  is an isomorphism
2. **Identity Natural Transformation:** For any functor  $F$ , the identity transformation  $\text{id}_F$  has components  $(\text{id}_F)_X = \text{id}_{F(X)}$

#### Composition

Natural transformations can be composed: - **Vertical Composition:** If  $\eta : F \Rightarrow G$  and  $\mu : G \Rightarrow H$ , then  $\mu \circ \eta : F \Rightarrow H$  - **Horizontal Composition:** Natural transformations can also be composed with functors



Dependency Graph



Local dependency graph