Entropy and Information Gain Notes

Entropy:
- Entropy provides a measurement of uncertainity.

associated with a random variable or random process.

Note: This is the definition of entropy under in the context of information theory.

- For a discrete r.v. X with possible outcomes X1, X2, 1..., Xn which occur with probability P(X1), P(X2), 1..., P(Xn), the entropy of X, denoted as H(X), is defined as:

 $H(x) = -\sum_{i=1}^{n} P(x_i) \log_2 P(x_i)$

Note: $H(x) = E[-log_2P(x)]$, where E is expected value

- Note: When P(Xi)=0, for some Xi, we take P(Xi) log 2 P(Xi) to be 0, which is consistent with its limit.

I.e. lim plog(p)=0
P=0

- Eig. Suppose we flip a fair coin. $P_1 = P_2 = \frac{1}{2}$ $H(x) = -(2)(\frac{1}{2}\log_2(\frac{1}{2}))$ $= -\log_2(\frac{1}{2})$ = -(-1)= -(-1) - E.g. Say we toss an unfair coin now.

Suppose the probability of getting heads is 70%.

Then, the entropy becomes:

$$H(x) = -\sum_{i=1}^{n} P(x_i) \log_z P(x_i)$$

 $=-(0.7 \cdot \log_2(0.7) + 0.3 \cdot \log_2(0.3))$

≈ 0.8816 € Since one side comes up more frequently, there is reduced
 uncertainty and hence entropy.
 - E.g. The entropy of rolling a fair die is:

$$H(x) = -\sum_{i=1}^{n} P(x_i) - \log_2 P(x_i)$$

 $= -(6)(6)(\log_2(6))$ $= -(\log_2(1) - \log_2(6))$

= log 2 (6)

= 2.58
Since the probability of rolling a die (1/6) is smaller than the prob of flipping a coin (1/2), its entropy will be higher.

- We also have conditional entropy.

- $H(x|y) = -\sum_{i,j} P(x_i, y_i) \log_2 P(x_i|y_j)$ = $-\sum_{j} P(y_j) \sum_{i} P(x_i|y_j) \log_2 P(x_i|y_j)$ = $\sum_{j} P(y_j) H(x|y_j)$

Mutual Information:

- Mutual information is a measure of the into shared by 2 r.v.'s.

I.e. It is a measure of how much about the state of one such var is known when it is conditioned on the state of the other.

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$$I(x;y) = H(x) - H(x|y)$$

= $H(y) - H(y|x)$

- Also called information gain.