Principle of Extremal Example

Consider n points on a plane s.t. every point is connected via edges with at least 3 other points. Show that there must exist a cycle with an even number of edges.

Soln

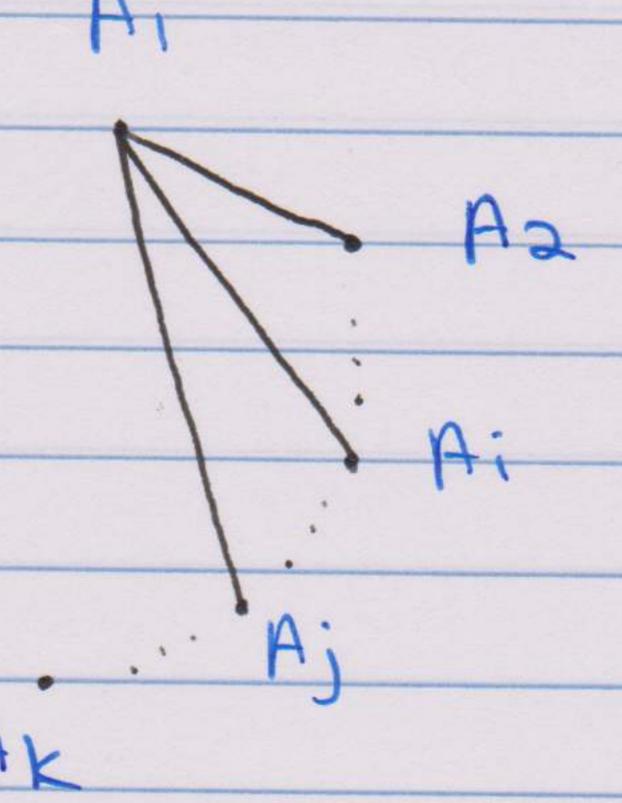
Consider the longest path which does not revisit a vertex twice in our finite graph:

A, \rightarrow A2 \rightarrow ... Ak.

By assumption, the point A1 must be

Connected to at least 3 attarrow into

connected to at least 3 other points, one of which is A2. The other 2 points must be points of the path or otherwise we would be able to create an even longer path, contradicting the fact that we considered the longest path. Let's assume that A1 is connected with A1 and A3 s.t. ic;



We have the following 3 cycles: Cycle 1: $A_1 \rightarrow A_2 \rightarrow ... A_i \rightarrow ... A_j \rightarrow A_1$ Cycle 2: $A_1 \rightarrow A_2 \rightarrow ... A_i \rightarrow A_1$ Cycle 3: $A_1 \rightarrow A_1 \rightarrow ... A_j \rightarrow A_1$

If i is even, then cycle 2 is even. This is because from $A_1 \rightarrow A_2 \rightarrow ... Ai$, there are i-1 edges. However, since there is another edge from A_i to A_1 , there are i edges in total.

Remark: To show why there are i-1 edges from A, -> A2 -> ... Ai, consider the following diagram below:

From $A_1 \rightarrow A_2$, there is ledge.

From $A_1 \rightarrow A_2 \rightarrow A_3$, there are $A_2 \rightarrow A_3 \rightarrow A_4$ From $A_1 \rightarrow A_2 \rightarrow A_3 \rightarrow A_4$ From $A_1 \rightarrow A_2 \rightarrow A_3 \rightarrow A_4$ there are 3 edges.

If j is even, then cycle I is even based on the logic applied above.

If both i and j are odd, then cycle 3 is even since its length is j-ita.

Remark: To explain why cycle 3 has length, consider the following below: a) Ai has length 1. b) Ai ->... Aj has length j-i.
To prove this, consider the following:
Ai ->... Ai = i-1 A, -> ... A; = j-1 Ai -> ... Aj = (A, -> ... Aj) - (A, -> ... Ai) = (j-1) - (i-1) c) Aj-> A, has length 1. Putting it all together, we get j-it2.