CSCC37 Linear Systems Notes

1. Linear Algebra Review:

a) Terminology:

Let A be a mxn matrix:

- This means that A has m rows and n colns.

- If m=n, then A is a square matrix.

Let B be a nxn matrix:

- B is a square matrix.

- B is said to be singular if it has I of the following equivalent properties:

1. B has no inverse.

2. det (B) =0

3. $B\bar{z}=\bar{o}$ for some vector $\bar{z}\neq\bar{o}$ Otherwise, B is non-singular. If B is non-singular, then B^{-1} exists and the system $B\bar{x}=\bar{b}$ always has a unique soln $\bar{x}=B^{-1}\bar{b}$ regardless of the value of \bar{b} . If B is singular, it will either have no solns or infinitely many solns.

The main diagonal of B is the values B_{II} , Bzz, ..., Bnn.

Eig. Let $B=\begin{bmatrix} z\\ z\\ z\\ z\end{bmatrix}$

The main diagonal is circled in red.

General Terminology:

- The transpose of matrix Am,n, denoted as AT, is created when you switch the row and coln indices of each element in A.

E.g. $A = \begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix}$ $AT = \begin{bmatrix} 1 & 3 & 5 \\ 2 & 4 & 6 \end{bmatrix}$

If A is a mxn matrix, AT is a nxm matrix

To create AT from A, write the rows of A as the colors of AT.

A square matrix whose transpose is equal to itself is a symmetric matrix.

I.e. If $A^{T} = A$, then A is a symmetric matrix.

A square matrix whose transpose is equal to its negative is a skew-symmetric matrix.

I.e. If $A^{T} = -A$, then A is a skew-symmetric metrix.

- The identity matrix, denoted as I, is a square matrix with 1's along the main diagonal and 0 elsewhere.

E.g. [1 0] [1 0 0] are identity matrices.

Note: The identity matrix is a symmetric matrix. Note: The product of 2 inverse matrices is always the identity matrix.

I.e Let $B=A^{-1}$. Then, AB=BA=I

- A lower triangular matrix is a square matrix if all entries above the main diagonal is 0.

 Fig. A = [1 0]

 [-3 2] is a lower triangular matrix.
- An upper triangular matrix is a square matrix if all entries below the main diagonal is 0.

 E.g. A = [1 2]

 [0 3] is an upper triangular matrix.

- A permutation matrix, denoted as P, is a square matrix having exactly one 1 in each row and coln and 0 elsewhere. It is used if you want to swap 2 rows.

Eig. Let $A = \begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix}$

Suppose you want to swap rows I and 2.
Your permutation matrix, P, would be [0 1]

 $\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix} = \begin{bmatrix} 3 & 4 \\ 1 & 2 \end{bmatrix}$

E.g. Let B = [123] 456 789]

Suppose you want to swap rows I and 3.
Your permutation matrix P would be [0 0 1]

0 0 1 1 2 3 7 8 9 0 1 0 4 5 6 = 4 5 6 1 0 0 7 8 9 1 2 3

Now suppose you want to swap rows 2 and 3 of the original matrix. Your P would be 1 0 0

Suppose you want to swap rows i and j, i is.

To create determine your permutation matrix, start with the identity matrix. Then, move the 1 at position (i,i) to (i,j) and move the 1 at position (j,i) to position (j,i).

In the first example, when I wanted to swap rows I and 2, i=1 and j=2. The I at (1,1) got moved to (1,2) and the 1 at (2,2) got moved to (2,1).

b) Calculations:

Matrix Addition and Subtraction:

- Two matrices, A and B can only be added or subtracted if they have the same number of rows and colors.

E.g. 1 W
$$A = \begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix}$$
, $B = \begin{bmatrix} 3 & 5 \\ 1 & 5 \end{bmatrix}$

Find A+B and A-B

Soln:

$$A+B=\begin{bmatrix} 1+3 & 2+5 \\ 3+1 & 4+5 \end{bmatrix}$$
 $A-B=\begin{bmatrix} 1-3 & 2-5 \\ 3-1 & 4-5 \end{bmatrix}$
 $=\begin{bmatrix} 4 & 7 \\ 4 & 9 \end{bmatrix}$ $=\begin{bmatrix} -2 & -3 \\ 2 & -1 \end{bmatrix}$

Matrix Multiplication:

- We can only multiply 2 matrices. A and B, if the number of colns of A = the num of rows of B. The resulting matrix will have the same number of rows as A and the same number of colns as B.

E.g. 2 Let
$$A = \begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix}$$
, $B = \begin{bmatrix} 3 & 1 & 2 \\ 5 & 2 & 3 \end{bmatrix}$

Find AXB

Soln:

$$A \times B = (1)(3) + (2)(5) (1)(1) + (2)(2) (1)(2) + (2)(3)$$

 $(3)(3) + (4)(5) (3)(1) + (4)(2) (3)(2) + (4)(3)$
 $(5)(3) + (6)(5) (5)(1) + (6)(2) (5)(2) + (6)(3)$

2. Linear Systems:

- $A\bar{x} = \bar{b}$, where $A \in \mathbb{R}^{n \times n}$, $\bar{x}, \bar{b} \in \mathbb{R}^{n}$ We're given A and \bar{b} and have to solve for \bar{x} .

- General Soln Technique:
- 1. Reduce the problem to an equivalent one that's easier to solve.
- 2. Solve the reduced problem.

Soln:

$$\begin{bmatrix} 3 & -5 \\ 6 & -7 \end{bmatrix} \begin{bmatrix} \chi_1 \\ \chi_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 5 \end{bmatrix} \underbrace{ \begin{cases} \zeta_1 \\ \zeta_2 \end{cases}}$$

$$\begin{bmatrix} 3 & -5 \\ 0 & 3 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 3 \end{bmatrix}$$

$$3x_2 = 3 - 5x_2 = 1$$

 $6x_1 - 7x_2 = 5$
 $6x_1 - 7 = 5$
 $6x_1 = 12$
 $x_1 = 2$

X1=2, X2=1

Now, we'll generalize this to n unknowns and n eqns. Ut matrix $A = \mathbb{Z}a_{ij}\mathbb{Z}$ where a_{ii} is the top left element.

Eqn 1: an X1 + an X2 + ... + an Xn = b1

Egn n: anix, + anixz + ... + ann xn = bn

We will use the following steps to reduce this system to triangular form.

Recall: An upper triangular matrix is a square matrix with 0's below the main diagonal.

Step 1:

- Assume that Q = 70
- Multiply Eqn 1 by azı and subtract from Eqn 2.
- Multiply Eqn 1 by asi and subtract from Eqn 3.
- Repeat for all remaining rows.

 I.e. Multiply Eqn 1 by air and subtract from Eqn i

where 4 & i & n.

- We now have an equivalent system where XI has been eliminated from eqns 2 to n.

I.e Now we have:

Eqn 1: a., x, + a., x + ... + a., xn = b. Eqn 2: 0 + a., x + ... + a., xn = b.

Egn n: 0 + ânz xz + ... + ânn xn = bn

Note: The small had, , is used to note that these values changed.

Step 2:

- Assume that azz \$0.
- Multiply eqn 2 by a32 and subtract from Eqn 3.
- Repeat for all remaining rows.
- We now have an equivalent system where Xz has been eliminated from eqns 3 to n.

Repeat this pattern up to and including eqn n-1.
Afterwards, we will have an upper triangular system.

- Another way of looking at this is using vector notation. $(A\bar{x}=\bar{b})$

Step 1: Eliminate the First caln of A using an.

Li
$$A\bar{x} = Li\bar{b}$$
, where $L_i = \begin{bmatrix} 1 & 0 & ... & 0 \\ -\frac{\alpha z_i}{\alpha_{i,i}} & 1 \\ \vdots & \ddots & \vdots \\ -\frac{\alpha n_i}{\alpha_{i,i}} & 1 \end{bmatrix}$

Note: Li is very similar to the identity matrix except the first coln is filled with multipliers used in the first step.

Step 2: Eliminate the second coln of LiA using azz.

$$L_{2}(L_{1}A)_{\overline{x}} = L_{2}(L_{1}\overline{b}) \text{ where } l_{2} = \begin{bmatrix} 1 & 0 & \dots & 0 \\ 0 & 1 & \vdots \\ \vdots & -\hat{a}_{3z} & \vdots \\ 0 & \vdots & 1 \end{bmatrix}$$

$$\frac{-\hat{a}_{0}z}{\hat{a}_{2}z}$$

We continue until we have Ln-, Ln-z ... Lz L, Ax = Ln-, Ln-z ... Lz L, B.

We let Ln-1 Ln-2 ... L. A = U where U is an upper triangular matrix. This becomes very easy to solve.

Fig. 4 Solve the following system of eqns using the technique we just learned.

$$2X_1 + 4X_2 - 2X_3 = 2$$

 $4X_1 + 9X_2 - 3X_3 = 8$
 $-2X_1 - 3X_2 + 7X_3 = 10$

Soln:

$$\begin{bmatrix} 2 & 4 & -2 \\ 4 & 9 & -3 \\ -2 & -3 & 7 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 2 \\ 8 \\ x_4 \end{bmatrix}$$

$$L_{1}(A) = \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \end{bmatrix} \begin{bmatrix} 2 & 4 & -2 \\ 4 & 9 & -3 \\ -2 & -3 & 7 \end{bmatrix}$$

$$= \begin{bmatrix} 2 & 4 & -2 \\ 0 & 1 & 1 \\ 0 & 1 & 5 \end{bmatrix}$$

$$L_2(L,A) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -1 & 1 \end{bmatrix} \begin{bmatrix} 2 & 4 & -2 \\ 0 & 1 & 5 \end{bmatrix}$$

$$L_2(L, B) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -1 & 1 \end{bmatrix} \begin{bmatrix} 2 \\ 4 \\ 12 \end{bmatrix}$$

We now have

$$\begin{bmatrix} 2 & 4 & -2 \\ 0 & 1 & 1 \\ 0 & 0 & 4 \end{bmatrix} \begin{bmatrix} \chi_1 \\ \chi_2 \\ \chi_3 \end{bmatrix} = \begin{bmatrix} 2 \\ 4 \\ 8 \end{bmatrix}$$

This makes finding x much easier.

3. LU Factorization:

- We have Ln-, Ln-z ... L, A = U => A= Li' Lz' ... Ln-, U Lemma 1: If Li is a Gauss Transformation, then Li' exists and is also a Gauss Transformation.

Lemma 2: If Li and Li are Gauss Transformations and joi, then Li Lj = Li + Lj - I

- A= Li' Lz' ... Ln-1' U C A= LU

- Using A = LU to solve $A\bar{x} = \bar{b}$, we can convert $A\bar{x} = \bar{b}$ into $(LU)\bar{x} = \bar{b}$. Then, let $\bar{d} = U\bar{x}$ where $L\bar{d}$ is a lower triangular matrix. We now have $L\bar{d} = \bar{b}$.

I.e. $A\bar{x} = \bar{b}$ C LU $\bar{x} = \bar{b}$ C L $\bar{d} = \bar{b}$ where $\bar{d} = U\bar{x}$.

while Ld is a lower triangular matrix and Ux is an upper triangular matrix.

- We use LU factorization because if we have the same coefficient matrix A but different RHS, we can use the same LU.

E.g. 5

Let
$$A = \begin{bmatrix} 2 & -1 & 1 \\ 2 & 2 & 2 \\ -2 & 4 & 1 \end{bmatrix}$$

Find L1 and L2

Soln:

$$\begin{bmatrix}
 1 & 0 & 0 \\
 -1 & 1 & 0
 \end{bmatrix}
 \begin{bmatrix}
 2 & -1 & 1 \\
 2 & 2 & 2 \\
 -2 & 4 & 1
 \end{bmatrix}$$

$$= \begin{bmatrix} 2 & -1 & 1 \\ 0 & 3 & 1 \\ 0 & 3 & 2 \end{bmatrix}$$

Now, let's show Lemma 2.

Recall that $L = L_i' L_i' L_i' \ldots L_{n-i}'$ and that $L_i L_i' = L_i' L_i' - L_i' = L_i' L_i' - L_i' = L_i' + L_i' + L_i' - L_i' = L_i' + L_i' +$

To compute Li', simply take Li and toggle /switch the sign of the multipliers.

$$L_1 = \begin{bmatrix} 1 & 0 & 0 \\ -1 & 1 & 0 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 1 & 0 \\ -1 & 0 & 1 \end{bmatrix}$$

$$L = L_{1}^{-1} \cdot L_{2}^{-1}$$

$$= \begin{bmatrix} 1 & 0 & 0 & \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$$

Another way to compute Li'Lz' is Li' + Lz' - I

$$= \begin{bmatrix} 2 & 0 & 0 \\ 1 & 2 & 0 \\ -1 & 1 & 2 \end{bmatrix} - \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

Now that we have I and U, let's see what LU is.

$$\begin{bmatrix} 1 & 0 & 0 & 0 & 2 & -1 & 1 \\ 1 & 1 & 0 & 0 & 3 & 1 \end{bmatrix} = \begin{bmatrix} 2 & -1 & 1 \\ 2 & 2 & 2 \\ -1 & 1 & 1 & 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} 2 & -1 & 1 \\ 2 & 2 & 2 \\ -2 & 4 & 1 \end{bmatrix}$$

find L and U.

$$L2(L, A) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -\frac{1}{16} & 1 \end{bmatrix} \begin{bmatrix} 8 & 2 & 9 \\ 0 & 8 & -\frac{1}{12} \\ 0 & \frac{2^2}{4} & \frac{9}{14} \end{bmatrix}$$

$$L = L'' + L'' - I$$

$$= \begin{bmatrix} 1 & 0 & 0 & & 1 & 0 & 0 \\ 1/2 & 1 & 0 & & 4 & 0 & 1 & 0 \\ 3/4 & 0 & 1 & & 0 & \frac{11}{16} & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

4. GE with Pivoting:
- If you go back to page 7, you'll see that we have "Assume that a " = 0" and = Assume that azz = 0." But what happens if at some step, a :: = 0?

E.g.
$$A = \begin{bmatrix} 1 & 2 & 1 \\ 1 & 2 & 4 \\ 1 & 3 & 3 \end{bmatrix}$$

A possible soln, in this case, is to swap rows 2 and 3.

- In general, if aii = 0, go down coln i, starting from row i+1, and find a suitable row to swap with row i. Note: If we want to find a row to swap with row i, we can only choose rows that are below row i. That is, we can only choose row j to swap with row i if j>i.
- A similar problem arises/occurs if one of the elements along the main diagonal is very small.

 This is a more common scenario than the previous one.

$$L_2 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -10^{16} & 1 \end{bmatrix}$$

$$L_{2}(L, A) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 10^{16} & 3 \end{bmatrix}$$

$$\begin{bmatrix} 0 & 10^{16} & 3 \\ 0 & 10^{16} & 1 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 2 & 1 \\ 0 & 10^{16} & 3 \\ 0 & 0 & 2-3 \cdot 10^{16} \end{bmatrix}$$

This term will cause a lot of issues.

A possible soln is to go down the coln, starting from the very small element, and to swap rows where the second row has a bigger element in that coln. This is called partial row pivoting.

For this example, swap rows 2 and 3.

Fig. 7 Solve
$$\begin{bmatrix} 2 & 6 & 6 \\ 3 & 5 & 12 \\ 6 & 6 & 12 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 20 \\ 25 \\ 30 \end{bmatrix}$$

using row pivoting.

Soln:

Step 1: We want the big gest value in cal 1 to be on the main diagonal. Hence, we swap row 1 and 3.

$$P_1 A \bar{\chi} = P_1 \bar{b}$$
, where $P_1 = \begin{bmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}$ Recall: P_1 is a $\begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \end{bmatrix}$ permutation matrix.

$$P. A = \begin{bmatrix} 6 & 6 & 12 \\ 3 & 5 & 12 \\ 2 & 6 & 6 \end{bmatrix} P. \overline{b} = \begin{bmatrix} 30 \\ 25 \\ 20 \end{bmatrix}$$

Step 2: Find Li

$$L_{1}(P, A) = \begin{bmatrix} 1 & 0 & 0 \\ -\frac{1}{2} & 1 & 0 \\ -\frac{1}{3} & 0 & 1 \end{bmatrix} \begin{bmatrix} 6 & 6 & 12 \\ 3 & 5 & 12 \\ 2 & 6 & 6 \end{bmatrix}$$

$$= \begin{bmatrix} 6 & 6 & 12 \\ 0 & 2 & 6 \\ 0 & 4 & 2 \end{bmatrix}$$

$$L_1(P_1 - \overline{b}) = \begin{bmatrix} 1 & 0 & 0 & \end{bmatrix} \begin{bmatrix} 30 \\ -\frac{1}{2} & 1 & 0 & \end{bmatrix} \begin{bmatrix} 30 \\ 25 \end{bmatrix} = \begin{bmatrix} 10 \\ 10 \end{bmatrix}$$

Step 3: We want the biggest value in col 2, Starting, from row 2 to be on the main diagonal. Hence, we swap rows 2 and 3.

$$P_2 L_1 P_1 A \bar{\chi} = P_2 L_1 P_1 \bar{b}$$
, where $P_2 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}$

Step 4: Find Lz

Now, we have
$$\begin{bmatrix} 6 & 6 & 12 \\ 0 & 4 & 2 \\ 0 & 0 & 5 \end{bmatrix} \begin{bmatrix} X_1 \\ X_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 10 \\ \hline 5 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

- We have LzPzL,P,A=U. LzPzL,P,A (1) LzPzL,P,A (2)

Note: The inverse of a permutation matrix is itself. Hence, that's why we can multiply LzPz LiPiA by Pz. Pz in (1).

Note: 2. = Pz L.Pz is a modified Gauss Transformation.

Solve PZLIPZ

Soln:
P2L.P2 = [100 | 100 | 100 |
001 | l21 | 0 | 001 |
010 | l3101 | 010]

Note: When you pre-multiply a matrix with a permutation matrix, you switch 2 rows. However, when you post-multiply a matrix with a permutation matrix, you switch 2 columns.

I.e. PL -> Changes 2 rows of L.
LP -> Changes 2 cols of L.

Hence, I, is Li with its 2 multipliers swapped.

Now we have $L_2 \hat{L}$, $P_2 P_1 A = U$ \iff $P_2 P_1 A = \hat{L}_1^{-1} \hat{L}_2^{-1} U$ \iff $P_4 = LU$ where $P = P_2 P_1$ and $L = \hat{L}_1^{-1} L_2^{-1}$

- Now, we have to solve $A\bar{x} = \bar{b}$ given PA = LU. $A\bar{x} = \bar{b}$ $PA\bar{x} = P\bar{b}$ $PA\bar{x} = P\bar{b}$ $DLU\bar{x} = \bar{b}$ where $\bar{b} = P\bar{b}$

Let $\overline{d} = 0\overline{x}$.

Hence, we solve:

1. $L\overline{d} = \overline{b}$ for \overline{d} (Forward Solve)

2. $U\overline{x} = \overline{d}$ for \overline{x} (Backward Solve)

- What happens if at the kth step, one element along the main diagonal, akk, and everything below it is 0?

Remember that our goal is to make every element in the kth coln under k 0, so we just continue.

However, this will result in a matrix U with a O for one of the entries along the main diagonal. Then, we will have a singular matrix U. If U is singular, when we do $U\bar{x}=\bar{d}$, we could have either O solns or infinitely many solns

E.g. $0x = \overline{d}$, where $0 = \begin{bmatrix} 2 & 5 & 4 \\ 0 & 0 & 1 \\ 0 & 0 & 2 \end{bmatrix}$

1	2	5	4	X		[4,]
Ì	0	0	1	72	n	95
	0	0	2	X3		L 93

So, we have X3 = dz 2X3 = d3

If d3 is twice dz, then Xz is a free variable and we have infinitely many soln.

If d3 is not twice dz, then we have no soln.

Note: While it's possible for U to be a singular matrix. L cannot be a singular matrix.

- Now suppose that at the kth step, if all elements below akk and the element akk have a magnitude of eps. max lujil. We call this numerical singularity

or near singularity.

- 5. Complexity of GE:
 Let A be a nxn matrix.

 - We will count additional/multiplication pairs, i.e. mx+b, as Floating Point Operation or FLOP.
 - a) Computing the complexity of LU Factorization:
 - col after a... Hence, we have (n-1)2 FLOPs.
 - Our second step is zeroing out the second col after azz. Hence, we have (n-z)2 FLOPs.
 - Our last step is zeroing out the (n-1)th col after a (n-1)(n-1). Hence, we have (n-(n-1))2 or 1 FLOP.
 - In total, we have (n-1)2 + (n-2)2 + ... +1 FLOPs.

 $\sum_{i=1}^{n} i^2 = n(n+i)(2n+i)$

However, since we go up to Cn-D2, we have n(n-1)(2n) or n3 + O(n2) FLOPs for computing the complexity of LU Factorization.

6) Computing the Complexity of Forward and Backward Solve:

- Consider forward solve:

TI	1 31		bi
l21 1 0	95		bz
l3, 1	1	=	1
1 14	1		-
ln, 1	90		[bn]

Recall forward solve is LJ=b where L is a lower triangular matrix.

$$d_1 = p_1$$

$$d_2 = p_2 - p_3 - p_3$$

dn = cn-o flops

Hence, we have 0 +1+2+...+ (n-1) FLOPS.

$$\sum_{i=1}^{n} i = n(n+i)$$

However, since we go up to (n-1), we get $\frac{n(n-1)}{2}$ or $\frac{n^2}{2}$ + O(n) FLOPs.

- Backward Solve is similar, resulting in another n2 + o(n) FLOPs.
 - In total, for forward and backward solve, the complexity is n2 + O(n) FLOPs.
 - 6. Round Off Error:
 Recall that we have PA = Lu computed in a floating point system (FPS).
 - Because of machine round off error, we actually get $\hat{\rho}(A+E)=\hat{L}\hat{U}$ where $\hat{\rho},\hat{L},\hat{u}$ are the computed factors and E is the error that occurs during factorization process.
 - We hope that E is small relative to A.
 - Now, solving $A\bar{x} = \bar{b}$ becomes $(A+E)\bar{x} = \bar{b}$ where \bar{x} is the computed soln.
 - Equivalently, let $E\hat{x} = \bar{x}$. Then, $(A+E)\hat{x} = \bar{b}$ \iff $A\hat{x} + \bar{r} = \bar{b}$ \iff $\bar{r} = \bar{b} - A\hat{x}$, where \bar{r} is the residual. We would like \bar{r} to be \bar{o} .
 - If we use row partial pivot, we can show that

 a) II EII & k. eps. II All where k is not too large and

 depends on n.
 - b) ||7|| € k.eps.||5|| € k.eps

11711 is called the relative residual.

Note: This does not mean that $1|\bar{x} - \hat{x}|$ or $\frac{1|\bar{x} - \hat{x}|}{|\bar{x}|}$ is small.

11x - x 11 is called the absolute error.

 $11\bar{x} - \hat{x}11$ is called the relative error.

X is the true solution.

E.g. 9. Given
$$\begin{bmatrix} 0.780 & 0.563 \\ 0.913 & 0.656 \end{bmatrix} = \begin{bmatrix} 0.217 \\ 0.254 \end{bmatrix}$$
 where

$$\bar{\chi} = \begin{bmatrix} 1 \\ -1 \end{bmatrix}$$
 and 2 computed solns $\hat{\chi_i} = \begin{bmatrix} 0.999 \\ -1.001 \end{bmatrix}$

and
$$\hat{\chi}_2 = \begin{bmatrix} 0.341 \\ -0.087 \end{bmatrix}$$
, find $\bar{\tau}_1$ and $\bar{\tau}_2$.

We see that <u>IIFIII</u> is much smaller than <u>IIFIII</u>, yet we see that $\hat{X}z$ is a terrible soln.

Furthermore, why is $11\bar{x} - \hat{x}$, 11 so much smaller $11\bar{x}11$

than 11x-x211?

- We need a relationship between relative error and relative residual.

Note: A small residual does not always mean a small error.

We have regns to start off: 1. Ar = b-F 2. Ar = B

We will now Subtract (1) from (2) We get: $A(\bar{x} - \hat{x}) = \bar{r}$ (3)

Rearranging (3) by multiplying both sides by A^{-1} gets us: $\bar{x} - \bar{x} = A^{-1} \bar{\tau}$ (4)

Taking the norm of both sides of (4) gets us: $||\bar{x} - \bar{x}|| = ||A^{-1}\bar{r}||$ $\leq ||A^{-1}|| ||\bar{r}|| (5)$

Taking the norm of $\overline{b} = A\overline{x}$, we get $||\overline{b}|| = ||A\overline{x}||$ $\leq ||A|| ||\overline{x}||$ (6)

Combining (5) and (6), we get $\frac{||\bar{x} - \bar{x}||}{||\bar{x}||} \leq \frac{||\bar{x}||}{||\bar{b}||}$

Relative = [IAII IIA-'II] IITII Note: 1= IIIII

error

Cond (A)

Relative residue = cond (A)

Relative residue = cond (A)

Cond (A) = 11A11 11A-111, Cond (A) 21 always.

If cond(A) is very large, the problem is poorly conditioned and small relative residuals do not mean small relative errors.

If cond(A) is not too large, the problem is well conditioned and a small relative residual is a reliable indicator of small relative error.

Note: Conditioning is a continuous spectrum. How large is "very large" depends on context.

Groing back to example 9, $A = \begin{bmatrix} 0.780 & 0.563 \\ 0.913 & 0.656 \end{bmatrix}$. Hence, $A^{-1} = \begin{bmatrix} 1 & 0.656 & -0.563 \\ det(A) & -0.913 & 0.780 \end{bmatrix}$

 $= 10^{6} \begin{bmatrix} 0.656 & -0.563 \\ -0.913 & 0.780 \end{bmatrix}$

Now, let's find cond(A). 11 A11 = 1,572 = 0.66.106 = 2.66.106 = 2.66.106

11x-x11 < 2.66x106 11711 meaning that the relative
11x11 error in x could be as
big as 2.66.106 times the
relative residual.

This means that A is a poorly conditioned matrix and relative residual is not a reliable indicator of relative error.

7. Iterative Refinement:

- One way to improve \bar{x} is to improve the mantissa length.

- Suppose that you've already solved (AtE) $\bar{x} = \bar{b}$ and you want to solve $A\bar{x} = \bar{b}$.

$$\vec{d}=\vec{5}+\vec{\hat{\chi}}A \iff \vec{d}=\vec{\hat{\chi}}(\exists +A)$$

 $\vec{5}-\vec{d}=\vec{\hat{\chi}}A \iff$

Now we have $A\bar{x} = \bar{b} - \bar{r}$ (2)

If we do (1) - (2), we get $A(\bar{x} - \bar{x}) = \bar{r}$. Let $\bar{z} = \bar{x} - \bar{x}$

Now, I'll solve $A\bar{z}=\bar{x}$. Furthermore, $\bar{x}=\bar{x}+\bar{z}$. However, this is a fallacy because we can't get \bar{z} . Instead, we get \bar{z} .

- Algorithim:
- 1. Compute & (0) by solving Ax = b in a FPS.
- 2. For i = 0,1,2, ... until the soln is good enough:
- 3. Compute 7(1) = b A x (1)
- 4. Solve AZ(i) = F(i) for some Z(i)
- 5. Update \$ (i+1) = \$ (i) + \(\frac{2}{5}(ii)\)

a.
$$1|\overline{x}||_{i} = n$$

$$\sum_{i=1}^{n} |x_{i}|$$

b)
$$||\bar{x}||_2 = \left(\frac{n}{\sum_{i=1}^{n} |x_i|^2}\right)^{1/2}$$

c)
$$11 \times 11 \infty = \max_{1 \leq i \leq n} (1 \times i1)$$

In general, for
$$p>0$$
, $||\bar{x}||_p = \left(\frac{n}{\sum_{i=1}^{n} |x_i|^p}\right)^{1/p}$

- Properties:

= 23

E.g. 10 Let
$$\bar{x} = [3, 5, -7, 8]$$

 $||\bar{x}||_1 = n$
 $|\bar{x}||_1 = n$

11x1100 = 8

9. Matrix Norms:

- Let A E IR nxm (I.e. A is a nxm matrix).

 $- \|A\|_{1} = \max_{1 \leq j \leq m} \left(\sum_{i=1}^{n} |a_{ij}| \right)$

= Max absolute col sum

 $- \|A\|_{\infty} = \max_{1 \le i \le n} \left(\sum_{j=1}^{m} |a_{ij}| \right)$

= Max absolute row sum

E.g. 11 Let A = [1 -7]. Find 11A11, and 11A1100

Soln:

| IAII, = max (1+1-21, 1-71+1-31) = max (3, 10)

=10

11A11 = max (1+1-71, 1-21+1-31) = max (8, 5) - Properties:

b) IIdAII = Id1 IIAII for any scalar &

d) IIABII & IIAII.IIBII

e) IIAXII ≤ IIAII. IIXII for any vector X. f) In general, IIAII = max IIAXII x≠0 IIXII

g) IIAIIz = Omax (A), where Omax (A) is the largest Singular value of matrix A.

h) $||A||_2 \le ||A||_F$, where $||A||_F = \left(\sum_{i=1}^m \sum_{j=i}^n |a_{ij}|^2\right)^{1/2}$

and is called the Frobenius norm.

10. Tutorial Examples: E.g. 12 Given $A = \begin{bmatrix} 2 & 6 & 6 \\ 3 & 5 & 12 \end{bmatrix}$ and $\overline{b} = \begin{bmatrix} 20 \\ 25 \\ 6 & 6 & 12 \end{bmatrix}$ 30

use LU factorization to solve Ax = b.

Soln:

We want A = LU. Then, Ax=b (LU)x=b ED LJ=b, where J=UX

$$L_{1}A = \begin{bmatrix} 1 & 0 & 0 \\ -3/2 & 1 & 0 \end{bmatrix} \begin{bmatrix} 2 & 6 & 6 \\ 3 & 5 & 12 \\ 6 & 6 & 12 \end{bmatrix}$$

$$= \begin{bmatrix} 2 & 6 & 6 \\ 0 & -4 & 3 \\ 0 & -12 & -6 \end{bmatrix}$$

$$L_{2}(L,A) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} \begin{bmatrix} 2 & 6 & 6 \\ 0 & -4 & 3 \\ 0 & -3 & 1 \end{bmatrix} \begin{bmatrix} 0 & -12 & -6 \end{bmatrix}$$

$$= \begin{bmatrix} 2 & 6 & 6 \\ 0 & -4 & 3 \end{bmatrix} \leftarrow u$$

$$\begin{bmatrix} 0 & 6 & -15 \end{bmatrix}$$

$$u = \begin{bmatrix} 2 & 6 & 6 \\ 0 & -4 & 3 \\ 0 & 0 & -15 \end{bmatrix}$$

$$L = L_1' \cdot L_2'$$

$$= L_1' + L_2' - I$$

$$= \begin{bmatrix} 1 & 0 & 0 \\ & & \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ & & & \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ & & & \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 0 & 0 \\ & & & \end{bmatrix} + \begin{bmatrix} 0 & 0 & 0 \\ & & & \end{bmatrix} - \begin{bmatrix} 0 & 0 & 0 \\ & & & \end{bmatrix}$$

We now have L and U.

I will define d = Ux.

LJ = b, solve for J

	1	0	0] [],		T20
1	3/2	1	0	95	2	25
	3	3	١	1 33		30

 $d_1 = 20$ 30 $d_1 + d_2 = 25 - 3d_2 = -5$

3d, + 3d2 + d3 = 30 -> 60 + (-15) +d3 = 30 -> d3 = -15

Now, solve for x in Ux = d

$$\begin{bmatrix} 2 & 6 & 6 & 1 \\ 0 & -4 & 3 & 1 \\ 0 & 0 & -15 & 1 \\ \end{bmatrix} \begin{array}{c|c} X_1 & 20 \\ X_2 & = -5 \\ \hline -15 \\ \end{bmatrix}$$

 $X_3 = 1$ $-4X_2 + 3X_3 = -5 \rightarrow -4X_2 + 3 = -5 \rightarrow -4X_2 = -8 \rightarrow X_2 = 2$ $2X_1 + 6X_2 + 6X_3 = 20 \rightarrow 2X_1 + 12 + 6 = 20 \rightarrow 2X_1 = 2 \rightarrow X_1 = 1$

E.g. 13 Do the same example as example 12, but this time, also do pivoting.

Soln:

Recall that we want the elements along the main diagonal, the pivot, to be the largest value in that column where the entries are chosen from and below the pivot.

I.e. The position of a pivot is akk, I = k = n. We want akk to be the largest value in colk starting from row k and going down.

Looking at the first col of A, we see that the largest value of col 1 starting from row 1 is the 6 on row 3. Hence, we swap rows I and 3.

$$P, A = \begin{bmatrix} 6 & 6 & 12 \\ 3 & 5 & 12 \\ 2 & 6 & 6 \end{bmatrix}$$

Now, looking at col 2 of L. (P. A), we see that the highest value of col 2 starting from row 2 is 4. So, we swap rows 2 and 3.

Lz Pz L, P, A => Lz Pz L, Pz Pz P, A, because P; Pi = I

1 0 6 | Recall: If you do P2M, you Switch the 2nd and 3rd rows J of M. If you do MPz, you switch the 2nd and 3rd cols of M.

	Lz L, F	PZ P, A=U
(-)	Pz P. A	= ~ Li' Lz' U
	P	L

$$PA = LU$$
= 1 0 0 6 6 12

'3 1 0 6 4 2

1/2 1/2 1 0 0 5

Recall that we started off with $A\bar{x} = \bar{b}$. Now, we have $PA\bar{x} = P\bar{b}$, where $P = PzP_1$

$$P\bar{b} = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} \begin{bmatrix} 20 \\ 30 \\ 30 \end{bmatrix}$$

LUX = PB Let $\overline{J} = U\overline{X}$ I will solve $L\overline{J} = P\overline{D}$ for \overline{J} . Then, I will solve $U\overline{X} = \overline{J}$ for \overline{X} .

$$\frac{3}{q^{1}} + qs = 50 \longrightarrow 10 + qs = 50 \longrightarrow qs = 10$$

$$\frac{s}{d_1} + \frac{s}{ds} + 43 = 52 - 312 + 2443 = 52 - 343 = 2$$

$$\begin{bmatrix} 6 & 6 & 12 \\ 0 & 4 & 2 \\ 0 & 0 & 5 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 10 \\ 5 \end{bmatrix}$$

73=1 $4x_{2}+2x_{3}=10 \rightarrow 4x_{2}+2=10 \rightarrow 4x_{2}=8 \rightarrow 5x_{2}=2$ $6x_{1}+6x_{2}+12x_{3}=10 \rightarrow 6x_{1}+12+12=30 \rightarrow 6x_{1}=6-3x_{1}=1$

$$\bar{\chi} = \begin{bmatrix} 1 \\ 2 \end{bmatrix}$$

which is what we got in the previous 2 tries.