

Linear Programming Examples

1. Given

$$\begin{aligned} \min \quad & 4x + 3y - 6z \\ \text{s.t.} \quad & y - 3z \geq 2x + 2 \\ & 3x + 2y + 5z = 10 \\ & x, z \geq 0 \end{aligned}$$

a) Convert this LP into Standard Form

Soln:

Recall that in Standard Form, we want

$$\begin{aligned} \text{Max} \quad & z = c^T x \\ \text{s.t.} \quad & a_1^T x \leq b_1 \\ & \vdots \\ & a_m^T x \leq b_m \\ & x \geq 0 \end{aligned}$$

Here are a couple of issues:

1. We're minimizing the obj func instead of maximizing it.
2. The first constraint has \geq instead of \leq and $2x$ is on the RHS instead of the LHS.
3. The second constraint has $=$ instead of \leq .
4. $y \geq 0$ is missing.

To fix #4, replace all instances of y with $y' - y''$ s.t. $y' \geq 0$, $y'' \geq 0$.

Now we have:

$$\begin{aligned} \text{Min} \quad & 4x + 3y' - 3y'' - 6z \\ \text{s.t.} \quad & y' - y'' - 3z \geq 2x + 2 \\ & 3x + 2y' - 2y'' + 5z = 10 \\ & x, y', y'', z \geq 0 \end{aligned}$$

To fix #3, change the $=$ to LHS and RHS and then multiply the constraint that has \geq by -1 to change it to \leq .

Now we have:

$$\begin{array}{ll} \text{Min} & 4x + 3y' - 3y'' - 6z \\ \text{s.t.} & y' - y'' - 3z \geq 2x + 2 \\ & 3x + 2y' - 2y'' + 5z \leq 10 \\ & -3x - 2y' + 2y'' - 5z \leq -10 \\ & x, y', y'', z \geq 0 \end{array}$$

To fix #2, move $2x$ to the LHS and multiply the row by -1 to change \geq to \leq .

Now we have:

$$\begin{array}{ll} \text{Min} & 4x + 3y' - 3y'' - 6z \\ \text{s.t.} & 2x - y' + y'' + 3z \leq -2 \\ & 3x + 2y' - 2y'' + 5z \leq 10 \\ & -3x - 2y' + 2y'' - 5z \leq -10 \\ & x, y', y'', z \geq 0 \end{array}$$

Lastly, to change #1, we max $-(\text{obj Func})$.

Now we have:

$$\begin{array}{ll} \text{Max} & -4x - 3y' + 3y'' + 6z \\ \text{s.t.} & 2x - y' + y'' + 3z \leq -2 \\ & 3x + 2y' - 2y'' + 5z \leq 10 \\ & -3x - 2y' + 2y'' - 5z \leq -10 \\ & x, y', y'', z \geq 0 \end{array}$$

Final
Solution

b) Write the dual LP of the result from part a)

Soln:

Recall that

Primal LP

$$\begin{aligned} \max \quad & c^T x \\ \text{s.t.} \quad & Ax \leq b \\ & x \geq 0 \end{aligned}$$

Dual LP

$$\begin{aligned} \min \quad & b^T y \\ \text{s.t.} \quad & A^T y \geq c \\ & y \geq 0 \end{aligned}$$

This was what we got in part a)

$$\begin{aligned} \text{Max} \quad & -4x - 3y' + 3y'' + 6z \\ \text{s.t.} \quad & 2x - y' + y'' + 3z \leq -2 \\ & 3x + 2y' - 2y'' + 5z \leq 10 \\ & -3x - 2y' + 2y'' - 5z \leq -10 \\ & x, y', y'', z \geq 0 \end{aligned}$$

$$A = \begin{bmatrix} 2 & -1 & 1 & 3 \\ 3 & 2 & -2 & 5 \\ -3 & -2 & 2 & -5 \end{bmatrix}$$

$$c = [-4 \quad -3 \quad 3 \quad 6]$$

$$b = [-2 \quad 10 \quad -10]$$

Hence, our dual LP is:

$$\begin{aligned} \text{Min} \quad & -2y_1 + 10y_2 - 10y_3 \\ \text{s.t.} \quad & 2y_1 + 3y_2 - 3y_3 \geq -4 \\ & -y_1 + 2y_2 - 2y_3 \geq -3 \\ & y_1 - 2y_2 + 2y_3 \geq 3 \\ & 3y_1 + 5y_2 - 5y_3 \geq 6 \\ & y_1, y_2, y_3 \geq 0 \end{aligned}$$

Note: Since the primal LP has 3

constraints and 4

vars, the dual LP

has 3 vars and 4

constraints.

2. Convert the following LP to Standard Form:

$$\text{Min} \quad 2x_1 + 7x_2 + x_3$$

$$\text{s.t.} \quad x_1 - x_3 = 7$$

$$3x_1 + x_2 \geq 24$$

$$x_2 \geq 0$$

$$x_3 \leq 0$$

Soln:

$$\text{Let } x_1 = x_1' - x_1'', \quad x_1' \geq 0, \quad x_1'' \geq 0$$

$$\text{Max} \quad -2x_1' + 2x_1'' - 7x_2 - x_3$$

$$\text{s.t.} \quad x_1' - x_1'' - x_3 \leq 7$$

$$-x_1' + x_1'' + x_3 \leq -7$$

$$x_1', x_1'', x_2, -x_3 \geq 0$$

3. Convert the following LP to slack form.

$$\text{Max} \quad 2x_1 - 6x_3$$

$$\text{s.t.} \quad x_1 + x_2 - x_3 \leq 7$$

$$3x_1 - x_2 \geq 8$$

$$-x_1 + 2x_2 + 2x_3 \geq 0$$

$$x_1, x_2, x_3 \geq 0$$

Soln:

First, convert it to standard form

$$\text{Max} \quad 2x_1 - 6x_3$$

$$\text{s.t.} \quad x_1 + x_2 - x_3 \leq 7$$

$$-3x_1 + x_2 \leq -8$$

$$x_1 - 2x_2 - 2x_3 \leq 0$$

$$x_1, x_2, x_3 \geq 0$$

Recall:

Standard form

$$\begin{aligned} \text{Max } & c^T x \\ \text{s.t. } & Ax \leq b \\ & x \geq 0 \end{aligned}$$

Slack form

$$\begin{aligned} z &= c^T x \\ s &= b - Ax \\ x, s &\geq 0 \end{aligned}$$

Slack form:

$$z = 2x_1 - 6x_3$$

$$x_4 = 7 - x_1 - x_2 + x_3$$

$$x_5 = -8 + 3x_1 - x_2$$

$$x_6 = -x_1 + 2x_2 + 2x_3$$

$$x_1, \dots, x_6 \geq 0$$

x_4, x_5, x_6

Basic vars:

Nonbasic vars: x_1, x_2, x_3

4. Show that the following LP problem is infeasible:

$$\text{Max } 3x_1 - 2x_2$$

$$\begin{aligned} \text{s.t. } & x_1 + x_2 \leq 2 \\ & -2x_1 - 2x_2 \leq -10 \\ & x_1, x_2 \geq 0 \end{aligned}$$

Soln:

Rewriting the second constraint, we get
 $x_1 + x_2 \geq 5$.

This is at odds with the first constraint.

5. Solve the following LP with the simplex algo.

$$\begin{aligned}
 \text{Max} \quad & 18x_1 + 12.5x_2 \\
 \text{s.t.} \quad & x_1 + x_2 \leq 20 \\
 & x_1 \leq 12 \\
 & x_2 \leq 16 \\
 & x_1, x_2 \geq 0
 \end{aligned}$$

Soln:

First, rewrite the LP into slack form.

$$Z = 18x_1 + 12.5x_2$$

$$x_3 = 20 - x_1 - x_2$$

$$x_4 = 12 - x_1$$

$$x_5 = 16 - x_2$$

$$x_1, x_2, x_3, x_4, x_5 \geq 0$$

Then, set all non-basic vars (x_1, x_2) to 0.

Next, choose a non-basic var with a positive coefficient. This will be our entering var.

I'll choose x_1 . I'll find the tightest bound for x_1 .

$$\begin{aligned}
 x_1 &= 20 - x_2 - x_3 \\
 &= 20 - x_3 \\
 &\leq 20
 \end{aligned}$$

$$\begin{aligned}
 x_1 &= 12 - x_4 \\
 &\leq 12 \leftarrow \text{Tightest bound}
 \end{aligned}$$

Replace all instances of x_1 with $12 - x_4$.

$$Z = 18(12 - X_4) + 12.5X_2$$

$$= 216 - 18X_4 + 12.5X_2$$

$$X_3 = 20 - (12 - X_4) - X_2$$

$$= 8 + X_4 - X_2$$

$$X_1 = 12 - X_4$$

$$X_5 = 16 - X_2$$

$$X_1, \dots, X_5 \geq 0$$

I'll choose X_2 as the entering var and find its tightest bound.

$$X_2 = 8 + X_4 - X_3$$

$$= 8 - X_3$$

Since X_4 is a non-basic var, its value is 0.

Tightest
bound \rightarrow

$$\leq 8$$

$$X_2 = 16 - X_5$$

$$\leq 16$$

Replace all instances of X_2 with $8 + X_4 - X_3$

$$Z = 216 - 18X_4 + 12.5(8 + X_4 - X_3)$$

$$= 316 - 5.5X_4 - 12.5X_3$$

$$X_1 = 12 - X_4$$

$$X_2 = 8 + X_4 - X_3$$

$$X_5 = 16 - (8 + X_4 - X_3)$$

$$= 8 - X_4 + X_3$$

Since there are no more non-basic vars,
we stop.

$$Z = 316$$

$$x_1 = 12$$

$$x_2 = 8$$

$$x_3 = 0$$

$$x_4 = 0$$

$$x_5 = 8$$

6. Solve the following LP using Simplex.

$$\text{Max } 5x_1 - 3x_2$$

$$\text{s.t. } x_1 - x_2 \leq 1$$

$$2x_1 + x_2 \leq 2$$

$$x_1, x_2 \geq 0$$

Soln:

First, write the LP in slack form.

$$Z = 5x_1 - 3x_2$$

$$x_3 = 1 - x_1 + x_2$$

$$x_4 = 2 - 2x_1 - x_2$$

$$x_1, \dots, x_4 \geq 0$$

Then, set all non-basic vars as 0.

I'll choose x_1 as my entering var.

$$x_1 = 1 + x_2 - x_3$$

$$= 1 - x_3$$

$$\leq 1$$

$$2x_1 = 2 - x_2 - x_4$$

$$x_1 = 1 - \frac{x_2}{2}$$

$$\leq 1$$

We see that both eqns give us $X_1 \leq 1$.
I'll choose the first eqn.

$$\begin{aligned} Z &= 5(1 + X_2 - X_3) - 3X_2 \\ &= 5 + 2X_2 - 5X_3 \end{aligned}$$

$$X_1 = 1 + X_2 - X_3$$

$$\begin{aligned} X_4 &= 2 - 2(1 + X_2 - X_3) - X_2 \\ &= 2 - 2 - 2X_2 + 2X_3 - X_2 \\ &= 2X_3 - 3X_2 \end{aligned}$$

Now, I'll choose X_2 as my entering var.

$$\begin{aligned} X_2 &= 1 - X_1 - X_3 \\ &\leq 1 \end{aligned}$$

$$\begin{aligned} 3X_2 &= 2X_3 - X_4 \\ &= 0 - X_4 \\ &\leq 0 \leftarrow \text{Tightest bound} \end{aligned}$$

$$X_2 = \frac{2}{3}X_3 - \frac{1}{3}X_4$$

$$\begin{aligned} Z &= 5 + 2\left(\frac{2}{3}X_3 - \frac{1}{3}X_4\right) - 5X_3 \\ &= 5 + \frac{4}{3}X_3 - \frac{2}{3}X_4 - 5X_3 \\ &= 5 - \frac{11}{3}X_3 - \frac{2}{3}X_4 \leftarrow \text{No more entering var} \end{aligned}$$

$$X_1 = 1 + \frac{1}{3}X_3 - \frac{1}{3}X_4$$

$$X_2 = \frac{2}{3}X_3 - \frac{1}{3}X_4$$

Final Soln:

$$Z = 5$$

$$X_1 = 1, X_2 = 0, X_3 = 0, X_4 = 0$$

7. You are given n jobs with a list of durations d_1, d_2, \dots, d_n . For every pair of jobs (i, j) , you are given a boolean $P_{i,j}$.

If this is true, then job i must finish before job j can begin.

We want to find start times s_1, s_2, \dots, s_n for the jobs (no job can start before time 0) s.t. the total time to complete all jobs is minimized while ensuring that the pre-requisite constraints are met. Write a LP to solve this problem.

Soln:

I'll let s_i denote the start time of job i .

I'll let T be an upper bound on the time completion.

Our LP:

$$\begin{aligned} \text{Min} \quad & T \\ \text{s.t.} \quad & \sum_i s_i + d_i \leq T \\ & d_i + s_i \leq s_j, \quad i \neq j \text{ and } P_{i,j} = 1 \text{ for } i, j \in \{1, \dots, n\} \\ & s_i \geq 0, \text{ for } i \in \{1, \dots, n\} \end{aligned}$$

8. Suppose you're writing down a binary integer linear program. You have 3 vars x, y , and z and the constraint $x, y, z \in \{0, 1\}$. For each relationship below, write a LP to solve it.

a) $z = x \wedge y$. We want $z = 1$ when x and $y = 1$ and 0 otherwise

Soln:

$$z \leq x$$

$$z \leq y$$

$$z \geq x + y - 1$$

b) $Z = x \vee y$

Soln:

$$Z \geq x$$

$$Z \geq y$$

$$Z \leq x + y$$

c) $Z = 7x$ (Not x)

Soln:

$$Z = 1 - x$$

9. A farmer has 110 hectares piece of land. He plans to grow wheat and barley. He has data regarding the cost, profit and man-days for growing each variety (shown below)

Variety	Cost (Price/Hec)	Net Profit (Price/Hec)	Man-days/Hec
Wheat	100	50	10
Barley	200	120	30

The farmer has a budget of \$10,000 and 1200 man-days. Find the opt value and opt soln.

Soln:

Let x be the total amount of hecs used to grow wheat.

Let y be the total amount of hecs used to grow barley.

LP:

$$\text{Max } 50x + 120y$$

$$\text{s.t. } 100x + 200y \leq 10,000 \rightarrow x + 2y \leq 100$$

$$10x + 30y \leq 1,200 \rightarrow x + 3y \leq 120$$

$$x + y \leq 110$$

$$x, y \geq 0$$

I'll use the simplex algo to solve it.

First, I'll change it to slack form.

$$Z = 50x + 120y$$

$$S_1 = 100 - x - 2y$$

$$S_2 = 120 - x - 3y$$

$$S_3 = 110 - x - y$$

$$x, y, S_1, S_2, S_3 \geq 0$$

I'll pick x as my entering var.

$$x = 100 - 2y - S_1 \quad | \quad x = 120 - 3y - S_2 \quad | \quad x = 110 - y - S_3$$

Tightest Bound $\rightarrow \leq 100$

≤ 120

≤ 110

$$x = 100 - 2y - S_1$$

Now, we have

$$Z = 5000 + 20y - 50S_1$$

$$x = 100 - 2y - S_1$$

$$S_2 = 20 - y + S_1$$

$$S_3 = 10 + y + S_1$$

$$x, y, S_1, S_2, S_3 \geq 0$$

Now, I'll pick y as my entering var.

$$y = 50 - \frac{1}{2}x \quad | \quad y = 20 + S_1 - S_2 \quad | \quad y = S_3 - 10 - S_1$$

$$\leq 50 \quad | \quad \leq 20 \quad | \quad \uparrow \text{Unusable}$$

Tightest Bound

$$y = 20 + S_1 - S_2$$

Now we have

$$Z = 5400 - 30S_1 - 20S_2$$

$$X = 100 - 2(20 + S_1 - S_2) - S_1$$

$$= 60 - 3S_1 + 2S_2$$

$$y = 20 + S_1 - S_2$$

$$S_3 = 10 - (20 + S_1 - S_2) + S_1$$

$$= -10 + S_2$$

$$X, y, S_1, S_2, S_3 \geq 0$$

Max profit: \$5400

Should produce wheat in 60 hecs.

Should produce barley in 20 hecs.