

Semantics for Linear Logic

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some date

a work in progress... draft version 25 August 2025

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1 The syntax of linear logic

When one gives up the law of excluded middle, one creates the space for a number of strange phenomena to arise. Given any formula φ , we know that $\varphi \rightarrow \neg\neg\varphi$ holds intuitionistically. That is, it can be proven in a usual syntactic system for propositional (or first-order) logic without the use of the law of excluded middle. The converse $\neg\neg\varphi \rightarrow \varphi$, however, is not intuitionistically valid; one must appeal to the law of excluded middle or one of its equivalents to establish double negation elimination. As another example, given formulas φ and ψ , the following three de Morgan's laws

1. $(\neg\varphi \wedge \neg\psi) \rightarrow \neg(\varphi \vee \psi)$
2. $(\neg\varphi \vee \neg\psi) \rightarrow \neg(\varphi \wedge \psi)$
3. $\neg(\varphi \vee \psi) \rightarrow (\neg\varphi \wedge \neg\psi)$

can all be proven intuitionistically. However, the remaining de Morgan's law

4. $\neg(\varphi \wedge \psi) \rightarrow (\neg\varphi \vee \neg\psi)$

is not intuitionistically valid. Furthermore, the formula $(\neg\varphi \vee \psi) \rightarrow (\varphi \rightarrow \psi)$ is intuitionistically valid, while the formula $(\varphi \rightarrow \psi) \rightarrow (\neg\varphi \vee \psi)$ is not.

Must we settle for these asymmetries if we wish to have intuitionistic aspects in our logic?

1.1 Connectives

We will mainly concern ourselves with *propositional* linear logic, for there is already a stark enough difference in the propositional fragments of linear logic in contrast to intuitionistic or classical logic.

Definition 1.1.1

Let P be a countable set of propositional constants. The set \mathcal{F} of formulas of propositional linear logic is generated by the grammar

$$\mathcal{F} ::= P \mid P^\perp \mid 0 \mid \top \mid \perp \mid 1 \mid \mathcal{F} \& \mathcal{F} \mid \mathcal{F} \oplus \mathcal{F} \mid \mathcal{F} \otimes \mathcal{F} \mid \mathcal{F} \wp \mathcal{F} \mid ?\mathcal{F} \mid !\mathcal{F}.$$

These constants and connectives are split into three different classifications.

The additive constants consist of 0 and \top , with 0 representing the additive falsity and \top representing the additive truth. The additive conjunction is $\&$ and is called with. The additive disjunction is \oplus and is called plus.

The multiplicative constants consist of \perp and 1, with \perp representing the multiplicative falsity and 1 representing the multiplicative truth. The symbol \otimes is called tensor, and is to be interpreted as multiplicative conjunction. The symbol \wp is called par, and is to be interpreted as multiplicative disjunction.

Finally, the symbols $?$ and $!$ make up the exponential class of connectives. We will later see that these allow us to have an interaction between the additive and the multiplicative connectives. Much like with real numbers, where we have the exponential law $a^{b+c} = a^b \cdot a^c$ (whenever $a \neq 0$), we will obtain, for instance, an equivalence between $?(A \& B)$ and $(?A) \otimes (?B)$ for any formulas A and B .

We remark that negation of formulas was only defined on the propositional constants. Our syntax for linear logic cannot make sense of $(p \& q)^\perp$ whenever p and q are propositional constants. Instead, we define negation as a meta-operation on formulas.

Definition 1.1.2

We define the meta-operation $(-)^\perp$ on formulas inductively as follows.

1. $(p)^\perp := p^\perp$ for all propositional constants p .
2. $(p^\perp)^\perp := p$ for all propositional constants p .
3. $\perp^\perp := 1$.
4. $1^\perp := \perp$.
5. $0^\perp := \top$.
6. $\top^\perp := 0$.
7. If A and B are formulas of linear logic, then $(A \& B)^\perp := A^\perp \oplus B^\perp$.
8. If A and B are formulas of linear logic, then $(A \oplus B)^\perp := A^\perp \& B^\perp$.
9. If A and B are formulas of linear logic, then $(A \otimes B)^\perp := A^\perp \wp B^\perp$.
10. If A and B are formulas of linear logic, then $(A \wp B)^\perp := A^\perp \otimes B^\perp$.
11. If A is a formula of linear logic, then $(?A)^\perp := !(A^\perp)$.
12. If A is a formula of linear logic, then $(!A)^\perp := ?(A^\perp)$.

We also define $A \multimap B := A^\perp \wp B$ for any two formulas A and B .

The symbol \multimap is called linear implication or lollipop.

Exercise 1.1.3

Show that $(A^\perp)^\perp = A$ for all formulas A .

1.2 Inference rules

We will adopt a sequent calculus approach for proofs in linear logic. Capital Latin letters A, B, C, \dots will be used to denote individual formulas, whereas capital Greek letters $\Gamma, \Delta, \Theta, \dots$ will be used to denote (possibly empty) sets of formulas. Those familiar with the sequent calculus would expect two rules for each connective, for example

$$\frac{\Gamma, A \vdash C, \Delta}{\Gamma, A \& B \vdash C, \Delta} \&L \qquad \frac{\Gamma \vdash A, \Delta \quad \Gamma \vdash B, \Delta}{\Gamma \vdash A \& B, \Delta} \&R$$

Notice, however, that there are *six* connectives and *four* propositional constants, making for an expected number of *twenty* inference rules if we had a left and right rule for each connective and constant. While manageable, this is far too many inference rules for pedagogical purposes. Instead, we shall identify the statement $A \vdash B$ with $\vdash A^\perp, B$. In general, $\Gamma \vdash \Delta$ is identified with $\vdash \Gamma^\perp, \Delta$, where $\Gamma^\perp := \{A^\perp : A \in \Gamma\}$, recalling that $(-)^{\perp}$ is a meta-operation. This makes all our contexts empty, and we only need a single rule for each connective and constant for the introduction of that symbol into the right-hand side of the turnstile. Henceforth, when we write

$$\overline{\Gamma}$$

we mean that we can prove Γ from the empty context, i.e.

$$\overline{\vdash \Gamma}$$

Let us start with the inference rules. The first is the identity rule, inspired from the rule $A \vdash A$ from the sequent calculus.

$$\frac{}{A, A^\perp} \text{id}$$

We will also adopt the exchange rule, so that the order of the formulas on either side of the turnstile does not matter.

$$\frac{\Gamma, A, B, \Delta}{\Gamma, B, A, \Delta} \text{ex}$$

We will often use this rule implicitly, for the sake of space.

The (infamous) cut rule will also be adopted.

$$\frac{\Gamma, A \quad \Delta, A^\perp}{\Gamma, \Delta} \text{cut}$$

It can be shown that the calculus we are developing will allow for the omission of the cut rule.

We now move on to the additive rules. The additive conjunction rule is reminiscent of the $\wedge R$ rule from the LK sequent calculus.

$$\frac{\Gamma, A \quad \Gamma, B}{\Gamma, A \& B} \&$$

The additive disjunction rules are reminiscent of the $\vee R$ rules from the LK sequent calculus.

$$\frac{\Gamma, A}{\Gamma, A \oplus B} \oplus_0 \quad \frac{\Gamma, B}{\Gamma, A \oplus B} \oplus_1$$

We will often not distinguish between the \oplus_0 and the \oplus_1 rule when writing down a proof involving the additive disjunction rules.

The additive truth rule is inspired from the $\top R$ rule from the LK sequent calculus.

$$\frac{}{\Gamma, \top} \top$$

In contrast, much like the LK sequent calculus having no $\perp R$ rule, we will have no inference rule for the additive falsity 0. This does *not* mean that the constant 0 can never appear in proofs.

We will have a rule for each multiplicative symbol. The multiplicative conjunction rule is

$$\frac{\Gamma, A \quad \Delta, B}{\Gamma, \Delta, A \otimes B} \otimes$$

Note the difference between the multiplicative conjunction rule and the additive conjunction rule. In-

tuitively, the inference $\frac{\Gamma^\perp, A \quad \Delta^\perp, B}{\Gamma^\perp, \Delta^\perp, A \otimes B} \otimes$ says that from $\Gamma \vdash A$ and $\Delta \vdash B$, we can conclude $\Gamma, \Delta \vdash A \otimes B$. It is perhaps most instructive to observe the case when $\Gamma = \Delta$. We would have $\frac{\Gamma^\perp, A \quad \Gamma^\perp, B}{\Gamma^\perp, \Gamma^\perp, A \otimes B} \otimes$, whereas $\frac{\Gamma^\perp, A \quad \Gamma^\perp, B}{\Gamma^\perp, A \& B} \&$. That is, if $\Gamma \vdash A$ and $\Gamma \vdash B$, we would need two copies of Γ to prove $A \otimes B$, whereas we would only need one copy of Γ to prove $A \& B$.

The multiplicative disjunction rule is

$$\frac{\Gamma, A, B}{\Gamma, A \wp B} \wp$$

Observe that, because $\vdash A, A^\perp$ for any formula A using the identity rule, the multiplicative disjunction rule gives a law of excluded middle $\vdash A \wp A^\perp$.

The multiplicative truth rule only allows us to instantiate 1 from no other premises.

$$\frac{}{1} 1$$

Unlike the additive false 0, we *do* have a multiplicative falsity rule.

$$\frac{\Gamma}{\Gamma, \perp} \perp$$

We now move on to the exponential rules. Note that, so far, we have had no other structural rules other than the exchange rule; we do not have the weakening rule or the contraction rule for arbitrary formulas like in the LK sequent calculus. We will only have these structural rules for formulas with the question mark as their main connective.

$$\frac{\Gamma}{\Gamma, ?A} \text{wk} \quad \frac{\Gamma, ?A, ?A}{\Gamma, ?A} \text{ctr}$$

We have two more rules associated with exponentials, which are rather unique to linear logic. The first is the dereliction rule, whose choice of name can be justified upon seeing the inference rule.

$$\frac{\Gamma, A}{\Gamma, ?A} \text{der}$$

Finally, we have the promotion rule for introducing the ! symbol.

$$\frac{?\Gamma, A}{?\Gamma, !A} \text{prom}$$

In the above inference rule, $?\Gamma := \{ ?B : B \in \Gamma \}$. We can only promote a formula A to $!A$ when every other formula on the same derivation line has ? as their main connective.

We collect all our rules below for ease of access.

$$\begin{array}{ccccccc} \frac{}{A, A^\perp} \text{id} & \frac{\Gamma, A, B, \Delta}{\Gamma, B, A, \Delta} \text{ex} & \frac{\Gamma, A \quad \Delta, A^\perp}{\Gamma, \Delta} \text{cut} & & & & \\ \frac{\Gamma, A \quad \Gamma, B}{\Gamma, A \& B} \& & \frac{\Gamma, A}{\Gamma, A \oplus B} \oplus_0 & \frac{\Gamma, A}{\Gamma, A \oplus B} \oplus_1 & \frac{}{\Gamma, \top} \top & & \\ \frac{\Gamma, A \quad \Delta, B}{\Gamma, \Delta, A \otimes B} \otimes & \frac{\Gamma, A, B}{\Gamma, A \wp B} \wp & & & \frac{}{1} 1 & \frac{\Gamma}{\Gamma, \perp} \perp & \\ \frac{\Gamma}{\Gamma, ?A} \text{wk} & \frac{\Gamma, ?A, ?A}{\Gamma, ?A} \text{ctr} & \frac{\Gamma, A}{\Gamma, ?A} \text{der} & \frac{?\Gamma, A}{?\Gamma, !A} \text{prom} & & & \end{array}$$

1.3 Example derivations

We should see a few examples of derivations using these rules.

Example 1.3.1

Let us prove the modus ponens rule: $A \otimes (A \multimap B) \vdash B$.

To show this, first recall that $A \multimap B$ means $A^\perp \wp B$. So we have to show that $A \otimes (A^\perp \wp B) \vdash B$. Recalling that we only have a one-sided sequent calculus and that we interpret $\Gamma \vdash C$ as $\vdash \Gamma^\perp, C$, we have to show that $\vdash (A \otimes (A^\perp \wp B))^\perp, B$. Applying the $(-)^\perp$ operation repeatedly, this is asking us to show that $\vdash A^\perp \wp (A \otimes B^\perp), B$. Now that everything is in the syntactic language, we proceed with the proof. We will show the uses of the exchange rules for this example and this example only.

$$\frac{\frac{\frac{\overline{A, A^\perp} \text{ id}}{A^\perp, A} \text{ ex} \quad \frac{\overline{B, B^\perp} \text{ id}}{B, B^\perp} \text{ id}}{A^\perp, B, A \otimes B^\perp} \otimes}{\frac{A^\perp, B, A \otimes B^\perp}{B, A^\perp, A \otimes B^\perp} \text{ ex}} \wp$$

$$\frac{B, A^\perp \wp (A \otimes B^\perp)}{A^\perp \wp (A \otimes B^\perp), B} \text{ ex}$$

This completes the proof. □

Exercise 1.3.2

Show that $\vdash A \multimap A$.

Example 1.3.3

Let us prove that $(A \& B)^\perp \vdash A^\perp \oplus B^\perp$. This is one of de Morgan's laws.

We need to show that $\vdash A \& B, A^\perp, B^\perp$.

$$\frac{\frac{\overline{A, A^\perp} \text{ id}}{A, A^\perp \oplus B^\perp} \oplus_1 \quad \frac{\overline{B, B^\perp} \text{ id}}{B, A^\perp \oplus B^\perp} \oplus_0}{A \& B, A^\perp \oplus B^\perp} \&$$

In the proof above, we distinguished between the \oplus_0 rule and the \oplus_1 rule. We will not do this anymore and simply write \oplus for either rule from now onwards.

Let us now prove the converse direction to the de Morgan's law above: $A^\perp \oplus B^\perp \vdash (A \& B)^\perp$. This time, we need to show that $\vdash (A^\perp \oplus B^\perp)^\perp, (A \& B)^\perp$, or equivalently, $\vdash A \& B, A^\perp \oplus B^\perp$. But this is precisely the de Morgan's law we have just established; we get this direction for free. □

If A and B are formulas $A \vdash B$ and $B \vdash A$, then we write $A \dashv\vdash B$ and say that A and B are linearly equivalent. The de Morgan's law established in [Example 1.3.3](#) can thus be succinctly written as the linear equivalence $(A \& B)^\perp \dashv\vdash A^\perp \oplus B^\perp$.

Exercise 1.3.4

Show the remaining de Morgan's laws:

1. $(A \oplus B)^\perp \dashv\vdash A^\perp \& B^\perp$,
2. $(A \otimes B)^\perp \dashv\vdash A^\perp \wp B^\perp$, and
3. $(A \wp B)^\perp \dashv\vdash A^\perp \otimes B^\perp$.

Example 1.3.5

Let us prove that $!(A \& B) \vdash !A \otimes !B$, i.e. the exponential $!$ connective turns additive conjunction to multiplicative conjunction.

We need to show that $\vdash (!(A \& B))^\perp, !A \otimes !B$. Fleshing this out, our goal is $\vdash ?(A^\perp \oplus B^\perp), !A \otimes !B$.

$$\begin{array}{c}
 \frac{\frac{\frac{}{A, A^\perp} \text{id}}{A, A^\perp \oplus B^\perp} \oplus}{A, ?(A^\perp \oplus B^\perp)} \text{der} \quad \frac{\frac{\frac{}{B, B^\perp} \text{id}}{B, A^\perp \oplus B^\perp} \oplus}{B, ?(A^\perp \oplus B^\perp)} \text{der} \\
 \frac{}{!A, ?(A^\perp \oplus B^\perp)} \text{prom} \quad \frac{}{!B, ?(A^\perp \oplus B^\perp)} \text{prom} \\
 \hline
 \frac{}{?(A^\perp \oplus B^\perp), ?(A^\perp \oplus B^\perp), !A \otimes !B} \otimes \\
 \hline
 \frac{}{?(A^\perp \oplus B^\perp), !A \otimes !B} \text{ctr}
 \end{array}$$

This completes the proof. □

Exercise 1.3.6

Show that $!(A \& B) \dashv\vdash !A \otimes !B$. Also show that $?(A \oplus B) \dashv\vdash ?A \wp ?B$.

Exercise 1.3.7

Show that $1 \dashv\vdash !\top$ and that $\perp \dashv\vdash ?0$.

Exercise 1.3.8

Show that $A^\perp \dashv\vdash A \multimap \perp$.

Exercise 1.3.9

Show that \otimes is associative, i.e. $(A \otimes B) \otimes C \dashv\vdash A \otimes (B \otimes C)$.

2 Categorical semantics

Everyone loves a bit of category theory.

2.1 Category-theoretic preliminaries: symmetric monoidal closed categories

A symmetric monoidal category is precisely what its name suggests: we equip a category with a bifunctor $(-) \otimes (-)$ which gives it the structure of a symmetric monoid. This will be the natural categorical generalisation of equipping a set with a binary operation which turns it into a symmetric monoid. We will, in a categorical sense, require that \otimes is associative, symmetric, and has a unit.

Definition 2.1.1

Equip a category \mathcal{C} with

1. a functor $(-) \otimes (-): \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$, called the tensor product,
2. an object $\mathbf{1} \in \text{ob } \mathcal{C}$, called the unit object,
3. a natural isomorphism $a: ((-) \otimes (-)) \otimes (-) \rightarrow (-) \otimes ((-) \otimes (-))$, called the associator,
4. a natural isomorphism $\lambda: \mathbf{1} \otimes (-) \rightarrow \text{id}_{\mathcal{C}}$, called the left unitor,
5. a natural isomorphism $\rho: (-) \otimes \mathbf{1} \rightarrow \text{id}_{\mathcal{C}}$, called the right unitor,
6. and a natural isomorphism $\sigma: (-) \otimes (-) \rightarrow (-) \otimes (-)$, called the braiding/symmetry of the tensor product.

This category \mathcal{C} is said to be a symmetric monoidal category if, for all $x, y, z, w \in \text{ob } \mathcal{C}$, the following four diagrams in \mathcal{C} commute.

1.

$$\begin{array}{ccc}
 & ((x \otimes y) \otimes z) \otimes w & \\
 a_{x,y,z} \otimes \text{id}_w \swarrow \cong & & \searrow \cong a_{x \otimes y, z, w} \\
 (x \otimes (y \otimes z)) \otimes w & & (x \otimes y) \otimes (z \otimes w) \\
 \downarrow \cong a_{x,y \otimes z, w} & & \downarrow \cong a_{x,y,z \otimes w} \\
 x \otimes ((y \otimes z) \otimes w) & \xrightarrow{\text{id}_x \otimes a_{y,z,w} \cong} & x \otimes (y \otimes (z \otimes w))
 \end{array}$$

2.

$$\begin{array}{ccc}
 (x \otimes \mathbf{1}) \otimes y & \xrightarrow[\cong]{a_{x,\mathbf{1},y}} & x \otimes (\mathbf{1} \otimes y) \\
 \downarrow \cong \rho_x \otimes \text{id}_y & & \downarrow \cong \text{id}_x \otimes \lambda_y \\
 & x \otimes y &
 \end{array}$$

3.

$$\begin{array}{ccc}
& (x \otimes y) \otimes z & \\
\sigma_{x,y} \otimes \text{id}_z \swarrow \cong & & \searrow \cong a_{x,y,z} \\
(y \otimes x) \otimes z & & x \otimes (y \otimes z) \\
a_{y,x,z} \downarrow \cong & & \downarrow \cong \sigma_{x,y \otimes z} \\
y \otimes (x \otimes z) & & (y \otimes z) \otimes x \\
& \searrow \cong \text{id}_y \otimes \sigma_{x,z} & \swarrow \cong a_{y,z,x} \\
& y \otimes (z \otimes x) &
\end{array}$$

4.

$$\begin{array}{ccc}
x \otimes y & \xrightarrow[\cong]{\sigma_{x,y}} & y \otimes x \\
& \searrow \cong \text{id}_{x \otimes y} & \downarrow \cong \sigma_{y,x} \\
& & x \otimes y
\end{array}$$

The assertions that the first three diagrams commute for all $x, y, z, w \in \text{ob } \mathcal{C}$ are respectively called the pentagon identity, the triangle identity, and the hexagon identity.

In practice, one simply specifies the triple $(\mathcal{C}, \otimes, \mathbf{1})$ when equipping a category \mathcal{C} with a symmetric monoidal structure.

We will soon become very familiar with the category **Rel** of relations. The objects of **Rel** are sets, and morphisms $A \xrightarrow{R} B$ in **Rel** are precisely subsets $R \subseteq A \times B$. It will turn out that this category **Rel** can be equipped with a lot more than just a symmetric monoidal structure, but let us establish this fact as a first stepping stone.

Example 2.1.2

Let **Rel** be the category of sets and relations between sets. Let us show that $(\mathbf{Rel}, \times, \{*\})$ is a symmetric monoidal category, where \times denotes the usual cartesian product and $\{*\}$ denotes a singleton set.

We will only define the associator $a: ((-) \times (-)) \times (-) \rightarrow (-) \times ((-) \times (-))$ and the left unitor $\lambda: \{*\} \times (-) \rightarrow \text{id}_{\mathbf{Rel}}$, as the rest are similar.

For $A, B, C \in \text{ob } \mathbf{Rel}$, define the morphism $(A \times B) \times C \xrightarrow{a_{A,B,C}} A \times (B \times C)$ to be the relation

$$a_{A,B,C} := \left\{ ((x, y), z), (x, (y, z)) : x \in A, y \in B, z \in C \right\}.$$

For $A \in \text{ob } \mathbf{Rel}$, define the morphism $\{*\} \times A \xrightarrow{\lambda_A} A$ to be the relation

$$\lambda_A := \left\{ ((*, x), x) : x \in A \right\}.$$

The verification that these, along with the appropriate definitions for the right unitor and the symmetry, are natural isomorphisms satisfying the commuting diagrams in Definition 2.1.1 is straightforward. \square

Exercise 2.1.3

Show that the category **Set** with the usual cartesian product and any singleton set $\{*\}$ is a symmetric monoidal category.

Exercise 2.1.4

Fix any field K and consider the category \mathbf{Vect}_K of vector spaces over K . Let $(-) \otimes_K (-): \mathbf{Vect}_K \times \mathbf{Vect}_K \rightarrow \mathbf{Vect}_K$ denote the usual tensor product over K of vector spaces over K . Also recall that we can view the field K as a vector space over itself. Show that $(\mathbf{Vect}_K, \otimes_K, K)$ is a symmetric monoidal category.

Just as we can define a homomorphism between monoids, there is a canonical notion of a morphism between symmetric monoidal categories: a symmetric monoidal functor. This is simply a functor between symmetric monoidal categories which preserves the symmetric monoidal structure of the domain category.

Definition 2.1.5

Let $(\mathcal{C}, \otimes, \mathbf{1})$ and $(\mathcal{D}, \oplus, \mathbf{0})$ be symmetric monoidal categories. A symmetric monoidal functor from $(\mathcal{C}, \otimes, \mathbf{1})$ to $(\mathcal{D}, \oplus, \mathbf{0})$ consists of a functor $F: \mathcal{C} \rightarrow \mathcal{D}$, a natural transformation $\mu: F(-) \oplus F(-) \rightarrow F((-) \otimes (-))$, and a morphism $\mathbf{0} \xrightarrow{\mu} F\mathbf{1}$ making the following four diagrams in \mathcal{C} commute for all $x, y, z \in \text{ob } \mathcal{C}$.

1.

$$\begin{array}{ccccc}
 & & (Fx \oplus Fy) \oplus Fz & & \\
 & \swarrow \mu_{x,y} \oplus \text{id}_{Fz} & & \searrow a_{Fx, Fy, Fz}^{\mathcal{D}} & \\
 F(x \otimes y) \oplus Fz & & & & Fx \oplus (Fy \oplus Fz) \\
 \downarrow \mu_{x \otimes y, z} & & & & \downarrow \text{id}_{Fx} \oplus \mu_{y,z} \\
 F((x \otimes y) \otimes z) & & & & Fx \oplus F(y \otimes z) \\
 & \swarrow \text{id} & & \searrow \mu_{x, y \otimes z} & \\
 & Fa_{x,y,z}^{\mathcal{C}} & F(x \otimes (y \otimes z)) & &
 \end{array}$$

where $a^{\mathcal{C}}$ and $a^{\mathcal{D}}$ are the associators for $(\mathcal{C}, \otimes, \mathbf{1})$ and $(\mathcal{D}, \oplus, \mathbf{0})$ respectively.

2.

$$\begin{array}{ccc}
 \mathbf{0} \oplus Fx & \xrightarrow[\cong]{\lambda_{Fx}^{\mathcal{D}}} & Fx \\
 \downarrow \mu \oplus \text{id}_{Fx} & & \uparrow \cong F\lambda_x^{\mathcal{C}} \\
 F\mathbf{1} \oplus Fx & \xrightarrow{\mu_{\mathbf{1},x}} & F(\mathbf{1} \otimes x)
 \end{array}$$

where $\lambda^{\mathcal{C}}$ and $\lambda^{\mathcal{D}}$ are the left unitors for $(\mathcal{C}, \otimes, \mathbf{1})$ and $(\mathcal{D}, \oplus, \mathbf{0})$ respectively.

3.

$$\begin{array}{ccc}
 Fx \oplus \mathbf{0} & \xrightarrow[\cong]{\rho_{Fx}^{\mathcal{D}}} & Fx \\
 \downarrow \text{id}_{Fx} \oplus \mu & & \uparrow \cong F\rho_x^{\mathcal{C}} \\
 Fx \oplus F\mathbf{1} & \xrightarrow{\mu_{x,\mathbf{1}}} & F(x \otimes \mathbf{1})
 \end{array}$$

where $\rho^{\mathcal{C}}$ and $\rho^{\mathcal{D}}$ are the right unitors for $(\mathcal{C}, \otimes, \mathbf{1})$ and $(\mathcal{D}, \oplus, \mathbf{0})$ respectively.

4.

$$\begin{array}{ccc}
Fx \oplus Fy & \xrightarrow[\cong]{\sigma_{Fx, Fy}^{\mathcal{D}}} & Fy \oplus Fx \\
\downarrow \mu_{x, y} & & \downarrow \mu_{y, x} \\
F(x \otimes y) & \xrightarrow[\cong]{F\sigma_{x, y}^{\mathcal{C}}} & F(y \otimes x)
\end{array}$$

where $\sigma^{\mathcal{C}}$ and $\sigma^{\mathcal{D}}$ are the symmetries for $(\mathcal{C}, \otimes, \mathbf{1})$ and $(\mathcal{D}, \oplus, \mathbf{0})$ respectively.

Note the abuse of notation by using the same letter μ for both the morphism $\mathbf{0} \xrightarrow{\mu} \mathbf{1}$ and the natural transformation $\mu: F(-) \oplus F(-) \rightarrow F((-) \otimes (-))$.

Recall that a cartesian closed category is a category \mathcal{C} with all finite products such that the functor $(-) \times y: \mathcal{C} \rightarrow \mathcal{C}$ has a right adjoint $(-)^y: \mathcal{C} \rightarrow \mathcal{C}$. The category **Set** is cartesian closed, as we can perform currying to identify functions $f: A \times B \rightarrow C$ with functions $\bar{f}: A \rightarrow B^C$, where B^C denotes the set of all functions from C to B . We remark that this object B^C is precisely the set $\text{hom}_{\mathbf{Set}}(C, B)$; for any two $B, C \in \text{ob } \mathbf{Set}$, the category **Set** has another object B^C which acts as the collection $\text{hom}_{\mathbf{Set}}(C, B)$ of morphisms from C to B .

The notion of cartesian closedness is now adapted for a symmetric monoidal category $(\mathcal{C}, \otimes, \mathbf{1})$, with \otimes in place of \times .

Definition 2.1.6

A symmetric monoidal closed category is a symmetric monoidal category $(\mathcal{C}, \otimes, \mathbf{1})$ equipped with a functor $(-) \multimap (-): \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{C}$, called the internal-hom functor, such that, for each $y \in \text{ob } \mathcal{C}$, the functor $(-) \otimes y: \mathcal{C} \rightarrow \mathcal{C}$ is left adjoint to the functor $y \multimap (-): \mathcal{C} \rightarrow \mathcal{C}$.

Example 2.1.7

We know (from [Example 2.1.2](#)) that **Rel** can be given the structure of a symmetric monoidal category. Let us now equip it with an internal-hom functor to make it a symmetric monoidal closed category.

For $A, B, C \in \text{ob } \mathbf{Rel}$, recall that a morphism $A \times B \xrightarrow{R} C$ in **Rel** is simply a subset $R \subseteq (A \times B) \times C$. Hence we have an isomorphism

$$\begin{array}{ccc}
\text{hom}_{\mathbf{Rel}}(A \times B, C) & \xrightarrow{\cong} & \text{hom}_{\mathbf{Rel}}(A, B \times C) \\
\downarrow \Psi & & \downarrow \Psi \\
R & \longmapsto & \left\{ (a, (b, c)) : ((a, b), c) \in R \right\}
\end{array}$$

So we can take $B \multimap C$ to be $B \times C$. The verification that this isomorphism is natural in A and C is routine work. Following on from [Example 2.1.2](#), this equips **Rel** with the structure of a symmetric monoidal closed category. \square

Exercise 2.1.8

Let \mathcal{C} be a cartesian closed category and let $\mathbf{1}$ denote the terminal object of \mathcal{C} . Show that $(\mathcal{C}, \times, \mathbf{1}, (-)^{(-)})$ is a symmetric monoidal closed category.

Exercise 2.1.9

Fix any field K and consider the category \mathbf{fdVect}_K of finite-dimensional vector spaces over K . Recall also that the field K itself can be viewed as a one-dimensional vector space over K . Let \otimes_K denote the usual tensor product over K of vector spaces over K . Let $[-, -]: \mathbf{fdVect}_K^{\text{op}} \times \mathbf{fdVect}_K \rightarrow \mathbf{fdVect}_K$ be

the functor such that $[V, W]$ is the K -vector space of all linear maps from V to W , whenever $V, W \in \text{ob } \mathbf{fdVect}_K$.

Show that \mathbf{fdVect}_K , together with \otimes_K , K , and $[-, -]$, is a symmetric monoidal closed category, even though \mathbf{fdVect}_K is not cartesian closed.

Definition 2.1.10

A $*$ -autonomous category is a symmetric monoidal closed category $(\mathcal{C}, \otimes, \mathbf{1}, \multimap)$ equipped with a dualising object $\perp \in \text{ob } \mathcal{C}$ such that, for all $x \in \text{ob } \mathcal{C}$, the canonical morphism $x \xrightarrow{\partial_x} (x \multimap \perp) \multimap \perp$ below is an isomorphism,

$$\begin{array}{ccc}
 \text{id}_{x \multimap \perp} \in \text{hom}_{\mathcal{C}}(x \multimap \perp, x \multimap \perp) & \xrightarrow[\cong]{(-) \otimes x \dashv x \multimap (-)} & \text{hom}_{\mathcal{C}}((x \multimap \perp) \otimes x, \perp) \\
 & & \downarrow \cong (-) \circ \sigma_{x, x \multimap \perp} \\
 & & \text{hom}_{\mathcal{C}}(x \otimes (x \multimap \perp), \perp) \\
 & & \downarrow \cong (-) \otimes x \dashv x \multimap (-) \\
 & \searrow & \downarrow \\
 & \partial_x \in \text{hom}_{\mathcal{C}}(x, (x \multimap \perp) \multimap \perp) &
 \end{array}$$

where $\sigma : (-) \otimes (-) \rightarrow (-) \otimes (-)$ is the symmetry of $(\mathcal{C}, \otimes, \mathbf{1})$. In this case, we write $x^* := x \multimap \perp$, giving rise to the dualisation functor $(-)^* : \mathcal{C}^{\text{op}} \rightarrow \mathcal{C}$.

Exercise 2.1.11

Let $(\mathcal{C}, \otimes, \mathbf{1}, \multimap, \perp)$ be a $*$ -autonomous category. For $x \in \text{ob } \mathcal{C}$, let $x \xrightarrow{\partial_x} x^{**}$ denote the canonical morphism as defined in [Definition 2.1.10](#). Show that the dualisation functor $(-)^* : \mathcal{C}^{\text{op}} \rightarrow \mathcal{C}$ is an equivalence of categories. Furthermore, show that $\partial : \text{id}_{\mathcal{C}} \rightarrow (-)^{**}$ is a natural isomorphism.

Example 2.1.12

We know (from [Example 2.1.7](#)) that **Rel** can be given the structure of a symmetric monoidal closed category. Let us now show that it can be given the structure of a $*$ -autonomous category.

Following on from [Example 2.1.7](#), we take any singleton set $\{*\}$ as the dualising object. The diagram chase in [Definition 2.1.10](#) then yields that the canonical morphism $A \xrightarrow{\partial_A} (A \times \{*\}) \times \{*\}$ is simply the relation

$$\partial_A = \left\{ (x, ((x, *), *)) : x \in A \right\},$$

which is clearly an isomorphism for any $A \in \text{ob } \mathbf{Rel}$. □

Exercise 2.1.13

Fix a field K . Following on [Exercise 2.1.9](#), show that the field K viewed as a one-dimensional vector space over itself is a dualising object in \mathbf{fdVect}_K , making \mathbf{fdVect}_K a $*$ -autonomous category.

2.2 Category-theoretic preliminaries: comonads and coKleisli categories

Definition 2.2.1

Fix a category \mathcal{C} . A comonad on \mathcal{C} consists of a functor $T : \mathcal{C} \rightarrow \mathcal{C}$ and natural transformations

$\varepsilon: T \rightarrow \text{id}_{\mathcal{C}}$ and $\delta: T \rightarrow TT$ making the following diagrams

$$\begin{array}{ccc}
 T & \xleftarrow{\varepsilon_T} & TT & \xrightarrow{T\varepsilon} & T \\
 & \searrow \text{id}_{\mathcal{C}} & \uparrow \delta & \nearrow \text{id}_{\mathcal{C}} & \\
 & & T & &
 \end{array}
 \qquad
 \begin{array}{ccc}
 T & \xrightarrow{\delta} & TT \\
 \delta \downarrow & & \downarrow \delta_T \\
 TT & \xrightarrow{T\delta} & TTT
 \end{array}$$

in the functor category $[\mathcal{C}, \mathcal{C}]$ commute.

The following [Exercise 2.2.2](#) shows that any adjoint pair of functors give rise to a comonad.

Exercise 2.2.2

Given an adjunction $\mathcal{C} \xrightleftharpoons[\text{G}]{\text{F}} \mathcal{D}$ with unit $\eta: \text{id}_{\mathcal{C}} \rightarrow GF$ and counit $\varepsilon: FG \rightarrow \text{id}_{\mathcal{D}}$, show that the triple $(FG, \varepsilon, F\eta_G)$ constitute a comonad on \mathcal{D} .

Definition 2.2.3

The coKleisli category for a comonad (T, ε, δ) on a category \mathcal{C} is the category \mathcal{C}_T consisting of the following.

1. We declare $\text{ob } \mathcal{C}_T := \text{ob } \mathcal{C}$.
2. For each $A, B \in \text{ob } \mathcal{C}_T$, we declare $\text{hom}_{\mathcal{C}_T}(A, B) := \text{hom}_{\mathcal{C}}(TA, B)$.
3. Given two morphisms $A \xrightarrow{f} B \xrightarrow{g} C$ in \mathcal{C}_T , the composite morphism $A \xrightarrow{gf} C$ in \mathcal{C}_T is defined to be the composite morphism

$$TA \xrightarrow{\delta_A} TTA \xrightarrow{Tf} TB \xrightarrow{g} C$$

in \mathcal{C} .

4. For each $A \in \text{ob } \mathcal{C}_T$, the identity morphism $A \overset{\text{f}}{\rightsquigarrow} \text{id}_A$ in \mathcal{C}_T is the morphism $TA \xrightarrow{\varepsilon_A} A$ in \mathcal{C} .

The use of squiggly arrows $A \overset{f}{\rightsquigarrow} B$ for morphisms in the coKleisli category are for cosmetic reasons; they are there to remind us that when we draw an arrow $A \overset{f}{\rightsquigarrow} B$ in the coKleisli category, we really mean an arrow $A \xrightarrow{f} TB$ in the base category.

If we were really strict about notation, we should write $\mathcal{C}_{\mathbb{T}}$ for the coKleisli category for a comonad $\mathbb{T} = (T, \varepsilon, \delta)$ on a category \mathcal{C} ; it could be the case that two different comonads have the same underlying endofunctor.

The coKleisli category is one of those rare instances where the fact that it is a category is not actually that obvious.

Exercise 2.2.4

Let (T, ε, δ) be a comonad on a category \mathcal{C} , and let \mathcal{C}_T denote the coKleisli category for this comonad. Show that \mathcal{C}_T is indeed a category, i.e. show that composition in \mathcal{C}_T (as defined in [Definition 2.2.3](#)) is associative and that composing any morphism $A \overset{f}{\rightsquigarrow} B$ in \mathcal{C}_T with either of the identity morphisms

$$A \overset{\text{f}}{\rightsquigarrow} \text{id}_A \text{ or } B \overset{\text{f}}{\rightsquigarrow} \text{id}_B \text{ in } \mathcal{C}_T \text{ yield } f.$$

Recall (from [Exercise 2.2.2](#)) that any adjunction gives rise to a comonad. The following [Exercise 2.2.5](#) shows that any comonad stems from an adjunction in sense of [Exercise 2.2.2](#).

Exercise 2.2.5

Let (T, ε, δ) be a comonad on a category \mathcal{C} , and let \mathcal{C}_T denote the coKleisli category for this comonad.

Find an adjoint pair of functors $\mathcal{C}_T \xrightleftharpoons[\mathcal{G}]{\mathcal{F}} \mathcal{C}$ with unit $\tilde{\eta}$ and counit $\tilde{\varepsilon}$ such that $(FG, \tilde{\varepsilon}, F\tilde{\eta}_G) = (T, \varepsilon, \delta)$.

2.3 Seely categories

Definition 2.3.1

A Seely category is a $*$ -autonomous category $(\mathbb{L}, \otimes, \mathbf{1}, \multimap, \perp)$ which has all finite products with terminal object $\top \in \text{ob } \mathbb{L}$, equipped with a comonad $(!, \varepsilon, \delta)$ on \mathbb{L} , an isomorphism $\mathbf{1} \xrightarrow[\cong]{\mu} !\top$, and a natural isomorphism $\mu: !(-) \otimes !(-) \rightarrow !((-) \times (-))$ such that all of the following hold.

1. The triple $(\mathbb{L}, \times, \top)$ is a symmetric monoidal category, where \times denotes the categorical product in \mathbb{L} .
2. The pair $(!, \mu)$ is a symmetric monoidal functor from $(\mathbb{L}, \times, \top)$ to $(\mathbb{L}, \otimes, \mathbf{1})$.
3. For each $x, y \in \text{ob } \mathbb{L}$, the following diagram

$$\begin{array}{ccc}
 !x \otimes !y & \xrightarrow[\cong]{\mu_{x,y}} & !(x \times y) \\
 \downarrow \delta_x \otimes \delta_y & & \downarrow \delta_{x \times y} \\
 & & !! (x \times y) \\
 & & \downarrow !(\pi_1, \pi_2) \\
 !x \otimes !y & \xrightarrow[\mu_{!x, !y}]{\cong} & !(x \times y)
 \end{array}$$

in \mathbb{L} commutes, where $A \xleftarrow{\pi_1} A \times B \xrightarrow{\pi_2} B$ are the projections associated with the product $A \times B$.

Note the abuse of notation by using the same letter μ to denote the isomorphism $\mathbf{1} \xrightarrow[\cong]{\mu} !\top$ as well as the natural isomorphism $\mu: !(-) \otimes !(-) \rightarrow !((-) \times (-))$. These isomorphisms denoted by μ are known as the Seely isomorphisms.

We should see an example of a Seely category. We have been working with the category **Rel** of sets and relations so far, making it into a $*$ -autonomous category (cf. [Example 2.1.12](#)). We will now upgrade **Rel** to a Seely category. The endofunctor of the comonad will be the finite multiset functor, which we define below.

A finite multiset on a set S is a function $m: S \rightarrow \mathbb{N}$ such that $m(m^{-1}(\mathbb{N} \setminus \{0\}))$ is a finite subset of \mathbb{N} . We write

$$m = \left[\underbrace{x_1, \dots, x_1}_{n_1 \text{ copies of } x_1}, \dots, \underbrace{x_k, \dots, x_k}_{n_k \text{ copies of } x_k} \right]$$

for the finite multiset m on a set $S \supseteq \{x_1, \dots, x_k\}$ with $m(x_i) = n_i$ for all $1 \leq i \leq k$. In this case, we abuse notation and say that $x_1, \dots, x_k \in m$ and that $|m| = n_1 + \dots + n_k$.

For instance, if $S = \{x, y, z\}$, we write $m = [x, x, x, y]$ to denote the finite multiset $m: S \rightarrow Y$ where $m(x) = 3$, $m(y) = 1$, and $m(z) = 0$, and we say that $x, y \in m$, $z \notin m$, and $|m| = 4$. We do not distinguish between permutations of elements in a finite multiset m , so $m = [x, x, x, y] = [x, x, y, x] = [x, y, x, x] = [y, x, x, x]$. The notation $[]$ refers to the empty finite multiset, i.e. the multiset which sends all elements of S to 0.

Given finite multisets m and m' on a set S , we define $m + m'$ to be the finite multiset on S given by $(m + m')(x) := m(x) + m'(x)$ for all $x \in S$. For example, if $m = [x, x, x, y]$ and $m' = [x, y, z]$, then $m + m' = [x, x, x, x, y, y, z]$.

We use $\mathcal{M}_{\text{fin}}(S)$ to denote the set of all finite multisets on a set S . Defining

$$\begin{aligned} \mathcal{M}_{\text{fin}}(A \xrightarrow{R} B) \\ := \left\{ (m, m') \in \mathcal{M}_{\text{fin}}(A) \times \mathcal{M}_{\text{fin}}(B) : |m| = |m'| \text{ and there exist permutations } m = [x_1, \dots, x_k] \right. \\ \left. \text{and } m' = [x'_1, \dots, x'_k] \text{ such that } x_i R x'_i \text{ for all } 1 \leq i \leq k \right\} \end{aligned}$$

for relations $R \subseteq A \times B$ makes \mathcal{M}_{fin} into an endofunctor on **Rel**.

Example 2.3.2

We know (from [Example 2.1.12](#)) that **Rel** can be given the structure of a $*$ -autonomous category. Let us now show that **Rel** can be given the structure of a Seely category.

We will take $! := \mathcal{M}_{\text{fin}}$ for the endofunctor of our comonad. For a set A define the morphism $\mathcal{M}_{\text{fin}}(A) \xrightarrow{\varepsilon_A} A$ in **Rel** to be the relation

$$\varepsilon_A := \{ ([x], x) : x \in A \}.$$

Furthermore, we define the morphism $\mathcal{M}_{\text{fin}}(A) \xrightarrow{\delta_A} \mathcal{M}_{\text{fin}}(\mathcal{M}_{\text{fin}}(A))$ in **Rel** to be the relation

$$\delta_A := \left\{ (m, [m_1, \dots, m_k]) \in \mathcal{M}_{\text{fin}}(A) \times \mathcal{M}_{\text{fin}}(\mathcal{M}_{\text{fin}}(A)) : m = m_1 + \dots + m_k \right\}.$$

These make $\varepsilon: \mathcal{M}_{\text{fin}} \rightarrow \text{id}_{\mathbf{Rel}}$ and $\delta: \mathcal{M}_{\text{fin}} \rightarrow \mathcal{M}_{\text{fin}}\mathcal{M}_{\text{fin}}$ into natural transformations, and it is elementary to check that the diagrams

$$\begin{array}{ccc} [x_1, \dots, x_k] & \xleftarrow{\quad} & [[x_1, \dots, x_k]] \\ & \varepsilon \swarrow & \nearrow \wr \\ & \mathcal{M}_{\text{fin}}(A) & \xleftarrow{\varepsilon_{\mathcal{M}_{\text{fin}}(A)}} \mathcal{M}_{\text{fin}}\mathcal{M}_{\text{fin}}(A) \\ & \swarrow \text{id}_{\mathcal{M}_{\text{fin}}(A)} & \uparrow \delta_A \\ & & \mathcal{M}_{\text{fin}}(A) \\ & & \searrow \wr \\ & & [x_1, \dots, x_k] \end{array}$$

$$\begin{array}{ccc}
[[x_1], \dots, [x_k]] & \xrightarrow{\quad} & [x_1, \dots, x_k] \\
\uparrow \scriptstyle \in & & \scriptstyle \ni \\
& \mathcal{M}_{\text{fin}} \mathcal{M}_{\text{fin}}(A) \xrightarrow{\mathcal{M}_{\text{fin}} \varepsilon_A} \mathcal{M}_{\text{fin}}(A) & \\
& \uparrow \scriptstyle \delta_A \quad \nearrow \scriptstyle \text{id}_{\mathcal{M}_{\text{fin}}(A)} & \\
& \mathcal{M}_{\text{fin}}(A) & \\
\downarrow \scriptstyle \in & & \\
[x_1, \dots, x_k] & &
\end{array}$$

$$\begin{array}{ccc}
\sum_{i=1}^k (\sum_{j=1}^{n_i} m_{i,j}) & \xrightarrow{\quad} & [m_{1,1}, \dots, m_{1,n_1}, \dots, m_{k,1}, \dots, m_{k,n_k}] \\
\downarrow \scriptstyle \in & & \scriptstyle \ni \\
& \mathcal{M}_{\text{fin}}(A) \xrightarrow{\delta_A} \mathcal{M}_{\text{fin}} \mathcal{M}_{\text{fin}}(A) & \\
& \downarrow \scriptstyle \delta_A \quad \downarrow \scriptstyle \delta_{\mathcal{M}_{\text{fin}}(A)} & \\
& \mathcal{M}_{\text{fin}} \mathcal{M}_{\text{fin}}(A) \xrightarrow{\mathcal{M}_{\text{fin}} \delta_A} \mathcal{M}_{\text{fin}} \mathcal{M}_{\text{fin}} \mathcal{M}_{\text{fin}}(A) & \\
\downarrow \scriptstyle \in & & \scriptstyle \ni \\
[\sum_{j=1}^{n_1} m_{1,j}, \dots, \sum_{j=1}^{n_k} m_{k,j}] & \xrightarrow{\quad} & [[m_{1,1}, \dots, m_{1,n_1}], \dots, [m_{k,1}, \dots, m_{k,n_k}]]
\end{array}$$

in **Rel** commute for all sets A , making $(\mathcal{M}_{\text{fin}}, \varepsilon, \delta)$ a comonad on **Rel**.

Recall that the categorical product in **Rel** is the set-theoretic disjoint union \sqcup , so the terminal object in **Rel** is the empty set \emptyset . This makes $(\mathbf{Rel}, \sqcup, \emptyset)$ a symmetric monoidal category. Then there is an obvious isomorphism $\{*\} \xrightarrow[\cong]{\mu} \mathcal{M}_{\text{fin}}(\emptyset)$ in **Rel**, namely $\mu := \{(*, [])\}$. Also, for sets A and B which we may assume without loss of generality to be disjoint, we define the morphism $\mathcal{M}_{\text{fin}}(A) \times \mathcal{M}_{\text{fin}}(B) \xrightarrow{\mu_{A,B}} \mathcal{M}_{\text{fin}}(A \sqcup B)$ in **Rel** to be the relation

$$\mu_{A,B} := \left\{ ((m', m''), m) \in (\mathcal{M}_{\text{fin}}(A) \times \mathcal{M}_{\text{fin}}(B)) \times \mathcal{M}_{\text{fin}}(A \sqcup B) : m' + m'' = m \right\}.$$

Observe that this $\mu: \mathcal{M}_{\text{fin}}(-) \times \mathcal{M}_{\text{fin}}(-) \rightarrow \mathcal{M}_{\text{fin}}((-) \sqcup (-))$ is a natural isomorphism since A and B are disjoint: any finite multiset on $A \sqcup B$ can be viewed as a finite multiset on A together with a finite multiset on B .

It is then easy to check that $(\mathcal{M}_{\text{fin}}, \mu)$ is a symmetric monoidal functor from $(\mathbf{Rel}, \times, \{*\})$ to $(\mathbf{Rel}, \sqcup, \emptyset)$. Finally, for all sets A and B , which we may again assume without loss of generality

to be disjoint, the diagram

$$\begin{array}{ccc}
\mathcal{M}_{\text{fin}}(A) \times \mathcal{M}_{\text{fin}}(B) & \xrightarrow{\mu_{A,B}} & \mathcal{M}_{\text{fin}}(A \sqcup B) \\
\downarrow \delta_A \times \delta_B & & \downarrow \delta_{A \sqcup B} \\
& & \mathcal{M}_{\text{fin}} \mathcal{M}_{\text{fin}}(A \sqcup B) \\
& & \downarrow \mathcal{M}_{\text{fin}}((\mathcal{M}_{\text{fin}} \pi_A) \cup (\mathcal{M}_{\text{fin}} \pi_B)) \\
\mathcal{M}_{\text{fin}}(\mathcal{M}_{\text{fin}}(A)) \times \mathcal{M}_{\text{fin}}(\mathcal{M}_{\text{fin}}(B)) & \xrightarrow{\mu_{\mathcal{M}_{\text{fin}}(A), \mathcal{M}_{\text{fin}}(B)}} & \mathcal{M}_{\text{fin}}(\mathcal{M}_{\text{fin}}(A) \sqcup \mathcal{M}_{\text{fin}}(B))
\end{array}$$

in **Rel** commutes, where $A \xleftarrow{\pi_A} A \sqcup B \xrightarrow{\pi_B} B$ are the projections associated with the categorical product $A \sqcup B$ in **Rel**, i.e. $\pi_A := \{ (x, x) : x \in A \} \subseteq (A \sqcup B) \times A$ and $\pi_B := \{ (y, y) : y \in B \} \subseteq (A \sqcup B) \times B$. The diagonal of this commuting square relates a pair $(m', m'') \in \mathcal{M}_{\text{fin}}(A) \times \mathcal{M}_{\text{fin}}(B)$ to (some permutation of) a multiset $[m_1, \dots, m_k] \in \mathcal{M}_{\text{fin}}(\mathcal{M}_{\text{fin}}(A) \sqcup \mathcal{M}_{\text{fin}}(B))$ if and only if there is some $i \in \{1, \dots, k\}$ such that $m' = m_1 + \dots + m_i$ and $m'' = m_{i+1} + \dots + m_k$.

All of the above makes **Rel** into a Seely category. \square

2.4 Modelling linear logic with Seely categories

We now start to model linear logic with Seely categories. We observe that the structure of a Seely category already gives us canonical ways to define the interpretations of many connectives. The remaining connectives will be defined via their duality properties. For instance, we will interpret $A \oplus B$ as $(A^\perp \& B^\perp)^\perp$.

Definition 2.4.1

Fix a Seely category \mathbb{L} . We define the interpretation $\llbracket A \rrbracket$ of a formula A inductively as follows.

1. For propositional constants p , we are free to choose an assignment $\llbracket p \rrbracket \in \text{ob } \mathbb{L}$.
2. If $A \equiv 1$, then $\llbracket 1 \rrbracket$ is the unit of the tensor product in \mathbb{L} from the $*$ -autonomous structure of \mathbb{L} .
3. If $A \equiv \perp$, then $\llbracket \perp \rrbracket$ is the dualising object in \mathbb{L} from the $*$ -autonomous structure of \mathbb{L} .
4. If $A \equiv \top$, then $\llbracket \top \rrbracket$ is the terminal object in \mathbb{L} .
5. If $A \equiv B \multimap C$, then $\llbracket B \multimap C \rrbracket := \llbracket B \rrbracket \multimap \llbracket C \rrbracket$, where the \multimap on the right-hand side is the internal-hom functor from the $*$ -autonomous structure of \mathbb{L} .
6. If $A \equiv 0$, then $\llbracket 0 \rrbracket := \llbracket \top \multimap \perp \rrbracket$.
7. If $A \equiv p^\perp$ for some propositional constant p , then $\llbracket p^\perp \rrbracket := \llbracket p \multimap \perp \rrbracket$.
8. If $A \equiv B \& C$, then $\llbracket B \& C \rrbracket := \llbracket B \rrbracket \times \llbracket C \rrbracket$, where \times is the categorical product in \mathbb{L} .
9. If $A \equiv B \oplus C$, then $\llbracket B \oplus C \rrbracket := \llbracket ((B \multimap \perp) \& (C \multimap \perp)) \multimap \perp \rrbracket$.
10. If $A \equiv B \otimes C$, then $\llbracket B \otimes C \rrbracket := \llbracket B \rrbracket \otimes \llbracket C \rrbracket$, where the \otimes on the right-hand side is the tensor product from the $*$ -autonomous structure of \mathbb{L} .
11. If $A \equiv B \wp C$, then $\llbracket B \wp C \rrbracket := \llbracket ((B \multimap \perp) \otimes (C \multimap \perp)) \multimap \perp \rrbracket$.

12. If $A \equiv !B$, then $\llbracket !B \rrbracket := !\llbracket B \rrbracket$, where the $!$ on the right-hand side is the endofunctor in the comonad on \mathbb{L} .

13. If $A \equiv ?B$, then $\llbracket ?A \rrbracket := \llbracket !(B \multimap \perp) \rrbracket \multimap \perp$.

We chose to define dual connectives via the $(-) \multimap \perp$ operator. We should check that this agrees with the $(-)^\perp$ operator.

Lemma 2.4.2

Fix a Seely category \mathbb{L} , fix an interpretation of propositional constants in \mathbb{L} , and fix a formula A . Then $\llbracket A^\perp \rrbracket \cong \llbracket A \multimap \perp \rrbracket$ in \mathbb{L} .

Proof. We induct on the complexity of A .

If $A \equiv p$ for some propositional constant p , then we have $\llbracket p^\perp \rrbracket = \llbracket p \multimap \perp \rrbracket$ by definition.

If $A \equiv 1$, then for any $x \in \text{ob } \mathbb{L}$ we have

$$\text{hom}_{\mathbb{L}}(x, \llbracket 1^\perp \rrbracket) = \text{hom}_{\mathbb{L}}(x, \llbracket \perp \rrbracket) \cong \text{hom}_{\mathbb{L}}(x \otimes 1, \llbracket \perp \rrbracket) \cong \text{hom}_{\mathbb{L}}(x, \llbracket 1 \multimap \perp \rrbracket),$$

and so $\llbracket 1^\perp \rrbracket = \llbracket \perp \rrbracket \cong \llbracket 1 \multimap \perp \rrbracket$.

If $A \equiv \perp$, then for any $x \in \text{ob } \mathbb{L}$ we have

$$\text{hom}_{\mathbb{L}}(x, \llbracket \perp^\perp \rrbracket) = \text{hom}_{\mathbb{L}}(x, \llbracket 1 \rrbracket) \cong \text{hom}_{\mathbb{L}}(x, \llbracket (1 \multimap \perp) \multimap \perp \rrbracket) \cong \text{hom}_{\mathbb{L}}(x, \llbracket \perp \multimap \perp \rrbracket)$$

since $\llbracket 1 \rrbracket \cong \llbracket (1 \multimap \perp) \multimap \perp \rrbracket$ and $\llbracket (1 \multimap \perp) \rrbracket \cong \llbracket \perp \rrbracket$. Thus $\llbracket 1^\perp \rrbracket = \llbracket \perp \rrbracket \cong \llbracket 1 \multimap \perp \rrbracket$.

If $A \equiv \top$, then $\llbracket \top^\perp \rrbracket = \llbracket 0 \rrbracket = \llbracket \top \multimap \perp \rrbracket$ by definition.

If $A \equiv 0$, then $\llbracket 0^\perp \rrbracket = \llbracket \top \rrbracket \cong \llbracket (\top \multimap \perp) \multimap \perp \rrbracket = \llbracket 0 \multimap \perp \rrbracket$ since $\top \cong (\top \multimap \perp) \multimap \perp$.

If $A \equiv B \& C$, then

$$\begin{aligned} \llbracket (B \& C)^\perp \rrbracket &= \llbracket B^\perp \oplus C^\perp \rrbracket \\ &= \llbracket ((B^\perp) \multimap \perp) \& (C^\perp \multimap \perp) \rrbracket \multimap \perp \\ &\cong \llbracket (B \& C) \multimap \perp \rrbracket \end{aligned}$$

where, in the last line, we used the inductive hypothesis which asserts that $B^\perp \cong B \multimap \perp$ and $C^\perp \cong C \multimap \perp$.

If $A \equiv B \oplus C$, then

$$\begin{aligned} \llbracket (B \oplus C)^\perp \rrbracket &= \llbracket B^\perp \& C^\perp \rrbracket \\ &\cong \llbracket ((B^\perp \& C^\perp) \multimap \perp) \multimap \perp \rrbracket, & \text{since } x \cong (x \multimap \perp) \multimap \perp \text{ for all } x \in \text{ob } \mathbb{L}, \\ &\cong \llbracket (((B \multimap \perp) \& (C \multimap \perp)) \multimap \perp) \multimap \perp \rrbracket, & \text{by the inductive hypothesis,} \\ &= \llbracket (B \oplus C) \multimap \perp \rrbracket. \end{aligned}$$

If $A \equiv B \otimes C$, then

$$\begin{aligned} \llbracket (B \otimes C)^\perp \rrbracket &= \llbracket B^\perp \wp C^\perp \rrbracket \\ &= \llbracket ((B^\perp \multimap \perp) \otimes (C^\perp \multimap \perp)) \multimap \perp \rrbracket \\ &\cong \llbracket (B \otimes C) \multimap \perp \rrbracket, \end{aligned}$$

where we used the inductive hypothesis in the last line.

If $A \equiv B \wp C$, then

$$\begin{aligned}
\llbracket (B \wp C)^\perp \rrbracket &= \llbracket B^\perp \otimes C^\perp \rrbracket \\
&\cong \llbracket ((B^\perp \otimes C^\perp) \multimap \perp) \multimap \perp \rrbracket, & \text{since } x \cong (x \multimap \perp) \multimap \perp \text{ for all } x \in \text{ob } \mathbb{L}, \\
&\cong \llbracket (((B \multimap \perp) \otimes (C \multimap \perp)) \multimap \perp) \multimap \perp \rrbracket, & \text{by the inductive hypothesis,} \\
&= \llbracket (B \wp C) \multimap \perp \rrbracket.
\end{aligned}$$

If $A \equiv !B$, then

$$\begin{aligned}
\llbracket (!B)^\perp \rrbracket &= \llbracket ?B^\perp \rrbracket \\
&= \llbracket !(B^\perp \multimap \perp) \multimap \perp \rrbracket \\
&\cong \llbracket !B \multimap \perp \rrbracket
\end{aligned}$$

where we used the inductive hypothesis in the last line.

If $A \equiv ?B$, then

$$\begin{aligned}
\llbracket (?B)^\perp \rrbracket &= \llbracket !B^\perp \rrbracket \\
&\cong \llbracket !(B^\perp \multimap \perp) \multimap \perp \rrbracket \\
&\cong \llbracket !(B \multimap \perp) \multimap \perp \multimap \perp \rrbracket \\
&= \llbracket ?B \multimap \perp \rrbracket,
\end{aligned}$$

again using that $x \cong (x \multimap \perp) \multimap \perp$ for all $x \in \text{ob } \mathbb{L}$ as well as the inductive hypothesis. \square

Exercise 2.4.3

Fix a Seely category \mathbb{L} , fix an interpretation of propositional constants in \mathbb{L} , and fix formulas A and B . Show that $\llbracket A \multimap B \rrbracket \cong \llbracket A^\perp \wp B \rrbracket$ in \mathbb{L} .

For brevity and the ease of reading, we will introduce an abuse of notation. For a Seely category \mathbb{L} and a formula A , we shall drop the double square brackets and simply write A in place of $\llbracket A \rrbracket$ for the interpretation of A in \mathbb{L} , when the context is clear. So the statement of [Lemma 2.4.2](#) will be expressed as follows: “Fix a Seely category \mathbb{L} and a formula A . Then $A^\perp \cong A \multimap \perp$ in \mathbb{L} .”

If Seely categories are to be any good as a model of linear logic, we should at least have a soundness theorem¹ for it.

Exercise 2.4.4

For formulas $A_1, \dots, A_k, A'_1, \dots, A'_\ell, B_1, \dots, B_m, B'_1, \dots, B'_n$, show that the following are equivalent:

- (i) $A_1, \dots, A_k, A'_1, \dots, A'_\ell \vdash B_1, \dots, B_m, B'_1, \dots, B'_n$
- (ii) $A_1 \otimes \dots \otimes A_k \otimes A'_1 \otimes \dots \otimes A'_\ell \vdash B_1 \wp \dots \wp B_m \wp B'_1 \wp \dots \wp B'_n$
- (iii) $A_1 \otimes \dots \otimes A_k \otimes B'_1 \otimes \dots \otimes B'_n \vdash B_1 \wp \dots \wp B_m \wp A'_1 \wp \dots \wp A'_\ell$
- (iv) $1 \vdash A_1^\perp \wp \dots \wp A_k^\perp \wp (A'_1)^\perp \wp \dots \wp (A'_\ell)^\perp \wp B_1 \wp \dots \wp B_m \wp B'_1 \wp \dots \wp B'_n$.

We will interpret a proof of $A \vdash B$ as a morphism from A to B in a Seely category. A proof of $\vdash A$ will be interpreted as a morphism from 1 to A . The equivalences seen in the above [Exercise 2.4.4](#) will also be present in any Seely category \mathbb{L} . Indeed,

$$\text{hom}_{\mathbb{L}}(A, B) \cong \text{hom}_{\mathbb{L}}(1 \otimes A, B) \cong \text{hom}_{\mathbb{L}}(1, A \multimap B) \cong \text{hom}_{\mathbb{L}}(1, A^\perp \wp B).$$

¹The reader should be aware that “soundness” here does not have the same meaning as is often seen in the literature.

A particular instantiation of the above is

$$\text{hom}_{\mathbb{L}}(A^\perp, B) \cong \text{hom}_{\mathbb{L}}(1, A^{\perp\perp} \wp B) \cong \text{hom}_{\mathbb{L}}(1, A \wp B),$$

since $A^{\perp\perp} \cong A$ in \mathbb{L} . More generally,

$$\text{hom}_{\mathbb{L}}(A_1^\perp \otimes \cdots \otimes A_k^\perp, B_1 \wp \cdots \wp B_n) \cong \text{hom}_{\mathbb{L}}(1, A_1 \wp \cdots \wp A_k \wp B_1 \wp \cdots \wp B_n).$$

Theorem 2.4.5 (Soundness of Seely categories)

Fix a Seely category \mathbb{L} and fix an interpretation of propositional constants in \mathbb{L} . For a finite set of formulas $\{A_1, \dots, A_n\}$, if $\vdash A_1, \dots, A_n$ then there exists a morphism $1 \rightarrow A_1 \wp \cdots \wp A_n$ in \mathbb{L} .

Proof. We induct on the complexity of proofs. Throughout, let $\Gamma = \{G_1, \dots, G_k\}$ and $\Delta = \{D_1, \dots, D_n\}$, for formulas $G_1, \dots, G_k, D_1, \dots, D_n$.

Our first base case is the identity rule $\frac{}{A, A^\perp} \text{id}$, giving a proof of $A \vdash A$. The identity morphism $A \xrightarrow{\text{id}_A} A$ is a morphism from A to A .

Another base case is the additive truth rule $\frac{}{\Gamma, \top} \top$. A morphism $1 \rightarrow G_1 \wp \cdots \wp G_k \wp \top$ can be identified with a morphism $G_1^\perp \otimes \cdots \otimes G_k^\perp \rightarrow \top$, and there is a unique such morphism since \top is the terminal object in \mathbb{L} .

The final base case is the multiplicative truth rule $\frac{}{\top} 1$, giving a proof of $\vdash 1$. For this, we use the morphism $1 \xrightarrow{\text{id}_1} 1$.

Suppose we have a proof ending with the exchange rule $\frac{\Gamma, A, B, \Delta}{\Gamma, B, A, \Delta} \text{ex}$. Inductively, we have a morphism $A^\perp \otimes B^\perp \rightarrow G_1 \wp \cdots \wp G_k \wp D_1 \cdots D_n$. We can then precompose this with a symmetry map from the symmetric monoidal structure of \mathbb{L} to obtain a morphism $A^\perp \otimes B^\perp \rightarrow G_1 \wp \cdots \wp G_k \wp D_1 \cdots D_n$.

Suppose we have a proof ending with the cut rule $\frac{\Gamma, A \quad \Delta, A^\perp}{\Gamma, \Delta} \text{cut}$. Then, inductively, we have morphisms $G_1^\perp \otimes \cdots \otimes G_k^\perp \xrightarrow{f} A$ and $A \xrightarrow{g} D_1 \wp \cdots \wp D_n$. So the composition $G_1^\perp \otimes \cdots \otimes G_k^\perp \xrightarrow{gf} D_1 \wp \cdots \wp D_n$ is the desired morphism.

Suppose we have a proof ending with the additive conjunction rule $\frac{\Gamma, A \quad \Gamma, B}{\Gamma, A \& B} \&$. This rule turns the proofs of $\Gamma^\perp \vdash A$ and $\Gamma^\perp \vdash B$ to a proof of $\Gamma^\perp \vdash A \& B$. Inductively, we have morphisms $A \xleftarrow{f} G_1^\perp \otimes \cdots \otimes G_k^\perp \xrightarrow{g} B$. The universal property of the product then gives us a morphism $G_1^\perp \otimes \cdots \otimes G_k^\perp \xrightarrow{(f,g)} A \& B$.

Suppose we have a proof ending with the one of the additive disjunction rules $\frac{\Gamma, A}{\Gamma, A \oplus B} \oplus_0$. Inductively, we have a morphism $A^\perp \xrightarrow{f} G_1 \wp \cdots \wp G_k$. Letting $A^\perp \& B^\perp \xrightarrow{\pi_1} A^\perp$ denote the product projection to A^\perp , we have a morphism $A^\perp \& B^\perp \xrightarrow{f\pi_1} G_1 \wp \cdots \wp G_k$.

The argument for a proof ending with the other additive disjunction rule $\frac{\Gamma, B}{\Gamma, A \oplus B} \oplus_1$ is similar to the argument above.

Suppose we have a proof ending with the multiplicative conjunction rule $\frac{\Gamma, A \quad \Delta, B}{\Gamma, \Delta, A \otimes B} \otimes$. We inductively have morphisms $G_1^\perp \otimes \cdots \otimes G_k^\perp \xrightarrow{f} A$ and $D_1^\perp \otimes \cdots \otimes D_n^\perp \xrightarrow{g} B$. We can then form the morphism $G_1^\perp \otimes \cdots \otimes G_k^\perp \otimes D_1^\perp \otimes \cdots \otimes D_n^\perp \xrightarrow{f \otimes g} A \otimes B$.

Suppose we have a proof ending with the multiplicative disjunction rule $\frac{\Gamma, A, B}{\Gamma, A \wp B} \wp$. Inductively, we have a morphism $1 \rightarrow G_1 \wp \cdots \wp G_k \wp A \wp B$. This gives rise to a morphism $1 \rightarrow G_1 \wp \cdots \wp G_k \wp (A \wp B)$ simply by using the associator maps from the symmetric monoidal structure of \mathbb{L} .

Suppose we have a proof ending with the multiplicative falsity rule $\frac{\Gamma}{\Gamma, \perp} \perp$. The inductive hypothesis gives a morphism $1 \rightarrow G_1 \wp \dots \wp G_k$. We know that $1 \otimes 1 \cong 1$, so we have a morphism $1 \otimes 1 \rightarrow G_1 \wp \dots \wp G_k$, which can be identified with a morphism $1 \rightarrow G_1 \wp \dots \wp G_k \wp 1^\perp$, that is, a morphism $1 \rightarrow G_1 \wp \dots \wp G_k \wp \perp$.

Suppose we have a proof ending with the weakening rule $\frac{\Gamma}{\Gamma, ?A} \text{wk}$. We inductively have a morphism $1 \xrightarrow{f} G_1 \wp \dots \wp G_k$. Then we can get a morphism from $(?A)^\perp = !A^\perp$ to $G_1 \wp \dots \wp G_k$ via the following composite morphism:

$$!A^\perp \xrightarrow{\delta_{A^\perp}} !!A^\perp \xrightarrow{!k} !\top \xrightarrow[\cong]{\mu^{-1}} 1 \xrightarrow{f} G_1 \wp \dots \wp G_k$$

where the natural transformation $\delta: ! \rightarrow !!$ comes from the comonad $(!, \varepsilon, \delta)$ on \mathbb{L} , the morphism $!A^\perp \xrightarrow{k} \top$ is the unique morphism from $!A^\perp$ to the terminal object \top , and the isomorphism $1 \xrightarrow[\cong]{\mu} !\top$ comes from the Seely structure of \mathbb{L} .

Suppose we have a proof ending with the contraction rule $\frac{\Gamma, ?A, ?A}{\Gamma, ?A} \text{ctr}$. Let $!A^\perp \otimes !A^\perp \xrightarrow{f} G_1 \wp \dots \wp G_k$ be a morphism obtained from the inductive hypothesis. Then we can get a morphism from $(?A)^\perp = !A^\perp$ to $G_1 \wp \dots \wp G_k$ via the following composite morphism:

$$!A^\perp \xrightarrow{\delta_{A^\perp}} !!A^\perp \xrightarrow{!(\varepsilon_{A^\perp}, \varepsilon_{A^\perp})} !(A^\perp \& A^\perp) \xrightarrow[\cong]{\mu_{!A^\perp, !A^\perp}^{-1}} !A^\perp \otimes !A^\perp \xrightarrow{f} G_1 \wp \dots \wp G_k$$

where the natural transformations $\varepsilon: ! \rightarrow \text{id}_{\mathbb{L}}$ and $\delta: ! \rightarrow !!$ come from the comonad $(!, \varepsilon, \delta)$ on \mathbb{L} , and the natural isomorphism $\mu: !(-) \otimes !(-) \rightarrow !((-) \& (-))$ comes from the Seely structure on \mathbb{L} .

Suppose we have a proof ending with the dereliction rule $\frac{\Gamma, A}{\Gamma, ?A} \text{der}$. Then, inductively, we have a morphism $A^\perp \xrightarrow{f} G_1 \wp \dots \wp G_k$. The composite morphism

$$!A^\perp \xrightarrow{\varepsilon_{A^\perp}} A^\perp \xrightarrow{f} G_1 \wp \dots \wp G_k$$

is then the desired morphism, where the natural transformation $\varepsilon: ! \rightarrow \text{id}_{\mathbb{L}}$ comes from the comonad $(!, \varepsilon, \delta)$ on \mathbb{L} .

Finally, suppose we have a proof ending with the promotion rule $\frac{? \Gamma, A}{? \Gamma, !A} \text{prom}$. The inductive hypothesis yields a morphism $!G_1^\perp \otimes \dots \otimes !G_k^\perp \xrightarrow[\cong]{\mu_{G_1^\perp, \dots, G_k^\perp}} !(G_1^\perp \& \dots \& G_k^\perp) \xrightarrow{f} A$, where we have abused notation for the natural isomorphism $\mu: !(-) \otimes !(-) \rightarrow !((-) \& (-))$ from the Seely structure of \mathbb{L} . Then the morphism

$$!(G_1^\perp \& \dots \& G_k^\perp) \xrightarrow{\delta_{G_1^\perp \& \dots \& G_k^\perp}} !!(G_1^\perp \& \dots \& G_k^\perp) \xrightarrow{!f} !A$$

is the desired morphism, where the natural transformation $\delta: ! \rightarrow !!$ comes from the comonad $(!, \varepsilon, \delta)$ on \mathbb{L} . \square

The category **Rel** is degenerate with regards to the soundness theorem ([Theorem 2.4.5](#)); any two objects in **Rel** have at least one morphism between them. To show that certain formulas are not provable, we need a Seely category which does not have this degeneracy.

Example 2.4.6

Let us show that $\not\vdash p$ and $\not\vdash p^\perp$ for any propositional constant p .

Consider a category \mathbb{L} with two objects, a and 1 , and exactly one morphism $a \xrightarrow{f} 1$ from a to 1 . We will equip this category \mathbb{L} with the appropriate structure to turn it into a Seely category.

This category \mathbb{L} already has all finite products: the object 1 is the terminal object; the products $a \times a$, $a \times 1$, and $1 \times a$, are equal to a , while the product 1×1 is equal to 1 . We define the tensor product $(-) \otimes (-)$ on \mathbb{L} to coincide with the categorical product $(-) \times (-)$ in \mathbb{L} . We also define the functor $(-) \multimap (-) : \mathbb{L}^{\text{op}} \times \mathbb{L} \rightarrow \mathbb{L}$ on objects as follows:

$$\begin{aligned} a \multimap a &:= 1, & a \multimap 1 &:= 1, \\ 1 \multimap a &:= a, & 1 \multimap 1 &:= 1. \end{aligned}$$

A brute case check shows that we have a natural bijection $\text{hom}_{\mathbb{L}}(x \otimes y, z) \cong \text{hom}_{\mathbb{L}}(x, y \multimap z)$ for all $x, y, z \in \text{ob } \mathbb{L}$. Furthermore, the object a acts as the dualising object in \mathbb{L} : we define $a^* := a \multimap a = 1$ and $1^* := 1 \multimap a = a$ and observe that

$$a^{**} = (a \multimap a) \multimap a = a \quad \text{and} \quad 1^{**} = (1 \multimap a) \multimap a = 1.$$

So far, we have made \mathbb{L} into a $*$ -autonomous category with finite products.

The (underlying functor of the) comonad we will prescribe is the identity functor $\text{id}_{\mathbb{L}}$. The Seely isomorphisms are also given by the appropriate identity morphisms. All of the above turn \mathbb{L} into a Seely category.

Now interpret the propositional constant p in \mathbb{L} as $\llbracket p \rrbracket := a \in \text{ob } \mathbb{L}$. As there is no morphism from 1 to $\llbracket p \rrbracket$ in the Seely category \mathbb{L} above, the soundness theorem ([Theorem 2.4.5](#)) tells us that there is no proof of $\vdash p$.

On the other hand, if we interpret the propositional constant p as $\llbracket p \rrbracket := 1 \in \text{ob } \mathbb{L}$. Then $\llbracket p^\perp \rrbracket = a$, and there is no morphism from 1 to $\llbracket p^\perp \rrbracket$ in \mathbb{L} . Hence, by the soundness theorem ([Theorem 2.4.5](#)), there is no proof of $\vdash p^\perp$. \square

Exercise 2.4.7

Let p and q be two distinct propositional constants. Show that $p \& (p \multimap q) \not\vdash q$.

3 Phase semantics

3.1 Algebraic Preliminaries: Monoids

Definition 3.1.1

A monoid is a comonad (T, ε, δ) on \mathbf{Set}^{op} for which the underlying endofunctor $T: \mathbf{Set}^{\text{op}} \rightarrow \mathbf{Set}^{\text{op}}$ has a left adjoint.

I'm kidding.

3.2 Phase Models

Definition 3.2.1

Let P be a countable set of propositional constants. An (enriched) phase model \mathcal{M} consists of:

1. a commutative monoid M , whose elements are called phases, where we use the period $(-).(-)$ to denote the multiplication of the commutative monoid and we use the notation 1_M for the identity element of the monoid;
2. a subset $\perp \subseteq M$, whose elements are called antiphases, from which we define the dual operation $(-)^{\perp}$ on the collection of subsets of M by

$$X^{\perp} := \{ y \in M : x.y \in \perp \text{ for all } x \in X \}$$

for all subsets $X \subseteq M$;

3. a valuation function $v: P \rightarrow \{ X \subseteq M : X = X^{\perp\perp} \}$; and
4. a submonoid J of M satisfying $\{x\}^{\perp\perp} \subseteq \{x.x\}^{\perp\perp}$ for all $x \in J$,

For a phase model (M, \perp, v, J) , we define $X.Y := \{ x.y : x \in X \text{ and } y \in Y \}$ for all subsets $X, Y \subseteq M$. With this, we have $X^{\perp} = \{ y \in M : X.\{y\} \subseteq \perp \}$ for all subsets $X \subseteq M$.

Exercise 3.2.2

Let (M, \perp, v, J) be a phase model. Show that if we have subsets $X \subseteq Y \subseteq M$, then $X^{\perp} \supseteq Y^{\perp}$.

Exercise 3.2.3

Let (M, \perp, v, J) be a phase model. Show that, for any subset $X \subseteq M$, we have $X \subseteq X^{\perp\perp}$, with equality if and only if $X = Y^{\perp}$ for some subset $Y \subseteq M$.

Exercise 3.2.4

Let (M, \perp, v, J) be a phase model and let \mathcal{X} be a family of subsets of M . Show that

$$\bigcap_{X \in \mathcal{X}} X^{\perp} = \left(\bigcup_{X \in \mathcal{X}} X \right)^{\perp}$$

but, in general, we only have

$$\bigcup_{X \in \mathcal{X}} X^{\perp} \subseteq \left(\bigcap_{X \in \mathcal{X}} X \right)^{\perp}.$$

We should see an example of a phase model.

Example 3.2.5

##?

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