On old laidge on a busy lighway is too so nashow to parmit two way traffic, so one way traffic is to.

```
living senaphores:
                                        (no of coss that can be )
    Simplifier space (INT-MAX)
   Semaphore mutex(1)
    enum · Direction (LEFT, RIGHT3;
    Direction currely = LEFT;
                                                    (initialized to zero)
     Semaphore direction [Direction]; =
    int active cals = 0, waiting cals = 0;
   Direction senesse (Direction A)

[ mutaxsacktupe of wait (mutex);

if (curdir!= A & & active cass > 0)
             waitingcals ++; signal(); direction [A] -> Press usboupt)
                                                     ( reed to wait)
         3
         else
                                                 11 can travel
            culldir = A;
             activecals ++; signal mutex -> oleapichuspend ();
           space -> suspend();
  void exit (Direction A)
       space > supporting signal (space);
       muters tomap(); wait (mutex);
                                                       (If we are changing the directions)
        if (-- active cas==0)
        ( while (waitingcorss)
            [ waitingcols --;
                active case ++;
          direction Everesse(A)? -> mait()
        mutex's signal (mutex);
```

- 0.2. Compare the deadlock prevention also by preventing circular wait with the deadlock avoidance algorithm:
- A deadlock avaidance algorithm tends to increase the suntime anaheads due to the cost of mentoring the current resource allocation.

  A deadlock avaidance algorithm over sequires additional information about the process and these algorithms are usually differing in amount of info sequired. The Resource allocation state is determines the execution. Info sequired the Resource allocation state is determined the execution.

  A circular wait orderne prevents the formation of deadlock by preventing A circular wait orderne prevents the formation of deadlock by preventing or what is less concurrent than deadlock algorithm.

  By this, we find that a deadlock avaidance ochome increases system throughputs
- 23. Identify whether the following statement is TRUE or FALSE. If the statement is
- A: B this statement is FALSE. Deadlock prevention obsertions limit the no of sequests that can be made by a process. These limits ensure that attenst one of the necessary conditions for deallocks cannot occur. But by this one of the necessary conditions for deallocks cannot occur. But by this one of the necessary conditions for deallocks cannot occur. But by this one is lot of ownhead and results in overall low device (CPV utilization and reduced system throughput since resource may be allocated but and reduced system throughput since resource may be allocated but unused for a long time. Also starvation is possible since a process may need to wait indefinitely.

Thus, the correct statement is: Deadlock prevention pretocal reduce CPU utilization.