Game Architecture Scripting

Today's Agenda

- What is scripting?
- Wait... why are we doing this?
- Some things to consider
- Parser & generator concepts
 - Big topic: just enough to whet your appetite
- Virtual machines
- Integrating scripting with an engine
- Visual script
- Case study: Lua

What is scripting?

- Way of expressing (a portion of) game logic:
 - o Al
 - Level mechanics
 - UI behavior
 - Weapons
 - o Ftc.

Often:

- Higher level than engine language
- Allows faster iteration than core engine code
- Not performance critical
- More than half of 'game code' is script

Implementation:

- 'Type-y' language (e.g. Lua)
- Visual programming (e.g. Unreal Blueprint)

Wait... why are we doing this?

Efficiency?

- Drops barriers for less technical developers
- Updates on-the-fly without reloading the game
- High level / domain specific script can do more with less
- Memory; script can be easily loaded and unloaded

Flexibility?

- Player generated content
- Post-ship functional updates without 1st party approval

Robustness?

Constrained script environment avoids OOM errors, NULL pointers, etc.

So a good engine must have a scripting system?

Absolutely not.

A scripting system is complexity. Avoid it unless you really need it.

Small teams and small games often don't need it.

- Coroutines & cooperative multithreading
- Events vs polling
- Safety
- Role of script in the engine

- Coroutines & cooperative multithreading
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For example, some AI script:

```
// Initially idle
self.behavior = idle

// Wait until hurt
yield_until(self.health < 1)

// Go on the attack
self.behavior = attack</pre>
```

- Coroutines & cooperative multithreading
- Events vs polling
- Safety
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```
With events:
on(health message msg):
     if (self.health > 0.5 && msg.health <= 0.5)</pre>
          self.behavior = attack2
on(player near msg):
     if (self.behavior == idle)
          self.behavior = attack1
With polling:
update():
     if (self.behavior == idle && player near())
          self.behavior = attack1
     if (self.health <= 0.5 && last health > 0.5)
          self.behavior = attack2
     last health = self.health
```

- Coroutines & cooperative multithreading
- Events vs polling
- Safety
- Role of script in the engine

- Permit memory allocation?
- Permit loops?
- Invalidating references to objects?

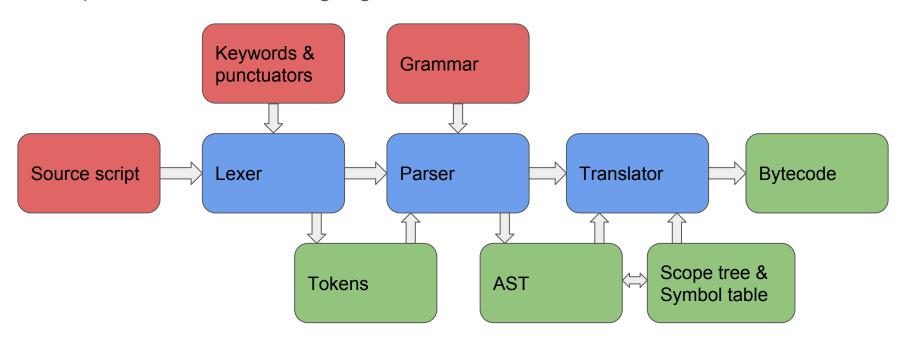
- Coroutines & cooperative multithreading
- Events vs polling
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- Role of script in the engine

Range of responsibility:

- Scripts are components on entities -- extend interfaces
- ...
- Scripts implement game state machine -- own main loop

Parsers and generators

A compiler for text based language...



Lexer

Our mission:

- Convert stream of characters into stream of tokens.
 - One token per language keyword. E.g.; if, else, while, goto, int, etc.
 - One token per language punctuator. E.g.; ++, ==, (, <, >>, etc.
 - One token per constant type. E.g.; strings, floats, integers
 - One token for some 'identifier'

How we do it:

- LL(1) Recursive-Descent Lexer
 - LL(1) = read from Left to right, visit children Left to right, 1 character lookahead
 - Recursive = functions can call themselves
 - Descent = top-down

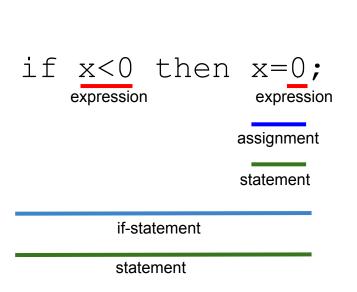
LL(1) Recursive-Descent Lexer

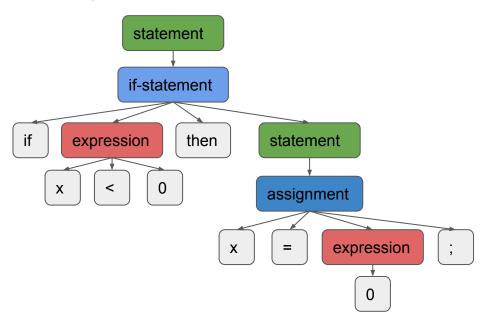
```
char* lex get next(
     const char* stream,
     token t* out token,
     const char** out text,
     int* out len)
     char* p = skip whitespace(stream);
     if (!p) { *out text = p; *out token = token eof; *out len = 0; }
     else if (strcmp(p, "if") == 0) { *out text = p; *out token = token if; *out len = 2; }
     else if (strcmp(p, "do") == 0) { *out text = p; *out token = token do; *out len = 2; }
     else if (is digit(p)) { *out text = p; *out token = token integer; *out len = xxx; }
     else if (is nondigit(p)) { *out text = p; *out token = token ident; *out len = xxx; }
     else if (*p == '+') { *out text = p; *out token = token plus; *out len = 1; }
    return p + *out len;
```

Parser

Our mission:

- Recognize language phrases
- Convert a list of tokens into an abstract syntax tree (AST)





LL(k) Recursive-Descent Parser

```
void statement()
     ast emit(node statement);
     if parse match(token_if) if_statement();
     else assignment();
void if statement()
     ast emit(node if statement);
     expression();
     parse ensure(token then);
     statement();
```

```
void expression() { /* stuff */ }

void assignment()
{
    ast_emit(node_assignment);
    parse_ensure(token_identifier);
    parse_ensure(token_equal);
    expression();
}
```

Types of ASTs

- Abstract syntax trees should not rigidly follow language structure
- Abstract syntax trees should be:
 - Dense; no unnecessary nodes
 - Convenient; easy to walk
 - Meaningful; emphasize operators and operations
- Homogenous AST; one node type with child list
 - Easy to walk, hard to work with
- Normalized heterogeneous AST; N node types with child list
 - Easy to walk, easier to work with
- Irregular Heterogenous AST; N node types
 - Hard to walk, easier to work with

Translation

Our mission:

Convert AST into bytecode for virtual machine

Implementation:

- Repeated AST walks
- Build scope tree
- Gather symbols and resolve
- Emit bytecode

Walking the AST

Usually depth first traversal

- Sometimes looking for node of a particular type (e.g. all assignment nodes)
- Sometimes looking at structure of subtrees

As each node is visited, can perform action:

- On the way down
- On the way up

Symbol tables

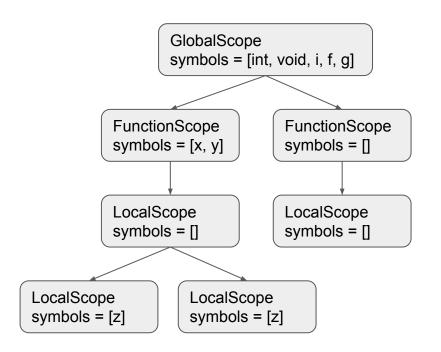
- Walk AST looking for identifier declarations
- Create entry in a table for each unique declaration
 - Variable symbols
 - Function symbols
 - Struct/class symbols
 - Enum symbols
 - Built-in symbols

Scope trees

- Walk AST looking for constructs that introduce scope
 - Global scope
 - Namespace scope
 - Class scope
 - Function scope
 - Local scope
- Create a tree of scopes
- Resolve symbol by searching up in the scope tree

Scope Tree Example

```
int i = 42;
int f(int x, int y)
    int i;
    { int z = x+y; i = z; }
    { int z = i+1; i = z; }
    return i;
void g()
    f(i, 2);
```



Generating bytecode

- Walk the AST and generate a stream of virtual machine instructions
- Use scope tree and symbol table to reserve virtual stack/registers
- Heavily dependent on virtual machine architecture...

Virtual Machines

A virtual CPU with a made up instruction set that runs our script.

- Easy to target (vs real CPU).
- Fast enough (pack instructions tight).

Two flavors:

- Stack based
- Register based

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```
iconst 1; Push integer 1 on stack
iconst 2; Push integer 2 on stack
iadd ; Pop & add from stack
; Push result on stack
print ; Pop & print top of stack
```

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A virtual CPU with a made up instruction set that runs our script.

- Easy to target (vs real CPU).
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Two flavors:

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The Big Switch of Death

```
switch (*bytecode++)
case opcode iload:
     stack.push int from mem(bytecode);
    bytecode += sizeof(int);
    break;
case opcode iadd:
     stack.push int(stack.pop int() + stack.pop int());
     break;
case opcode print:
     switch (stack.top().type())
     case type int:
          vm printf("%d\n", stack.pop int());
         break;
```

Our engine/script binding must allow:

- Engine to invoke script functions
- Script to invoke engine functions
- Script to use engine defined structures

- Engine able to work with script structures
- Script is able to be updated on-the-fly

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```
In script:
script add(x, y):
     return x + y;
In engine:
int engine add(int x, int y)
     vm func t func = vm getfunc("script add");
     vm call t call = vm new call(func);
     call.push arg int(x);
     call.push arg int(y);
     vm run(call);
     return call.pop return int();
```

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```
int engine add(int x, int y)
     return x + y;
void engine add wrapper(vm call t* call)
     int result = engine add(
          call->pop arg int(),
          call->pop arg int());
     call->push return int(result);
void vm init()
     vm register func(engine add wrapper);
```

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```
struct foo
{
    int x;
    float y;
};

void vm_init()
{
    vm_struct_t s = vm_new_struct("foo");
    s.add_field("x", type_int);
    s.add_field("y", type_float);
}
```

Our engine/script binding must allow:

A dynamic type system, engine-side? Ouch.

- Engine to invoke script functions
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Our engine/script binding must allow:

Left as an exercise to the listener.

- Engine to invoke script functions
- Script to invoke engine functions
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- Engine able to work with script structures
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Final thoughts on script/engine integration

- Binding script to engine is lots of boilerplate code
- Some engines have custom header file preprocessors
 - Generate script binding boilerplate automagically
- Some engines use a mix of C preprocessor macros and typing
- Some engines just do it all by hand
- Depends on size of engine, extent of script usage, etc.

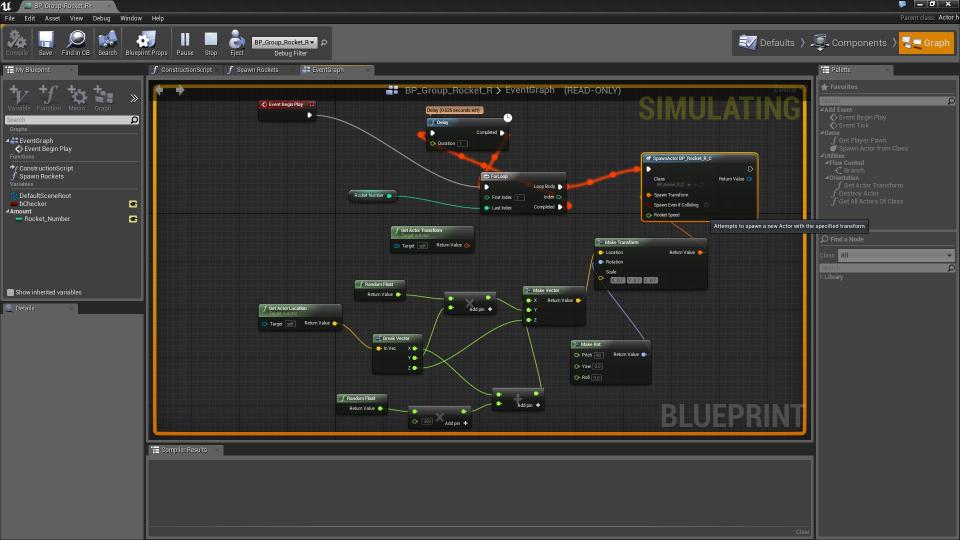
Visual Script

For some:

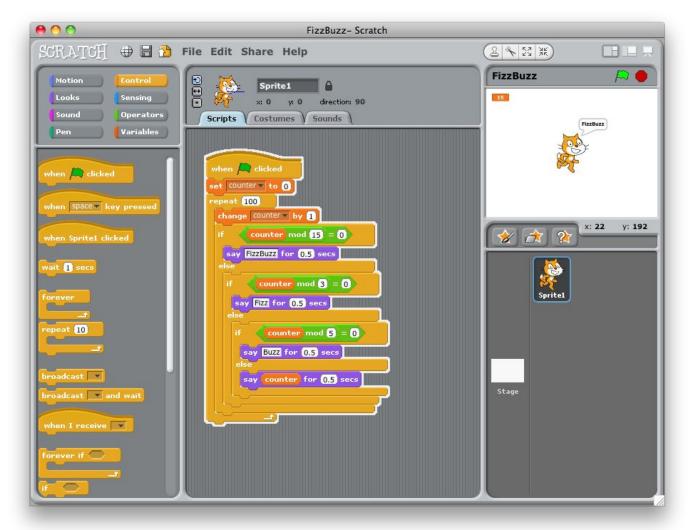
- Typing can be intimidating
- Syntax is a big challenge
- Knowing all possible valid actions at any point is helpful
- Exploration is preferred (over understanding)

For those folks, we have visual languages. Visual languages take many forms:

- Graphs (nodes and edges)
- Fill-the-blank statement composition
- Tiles (e.g. Scratch)



Let's run and	d jump!		
			\rightarrow
are	you drinking?		
Some of	friends are in th	he park.	
Let's play ba	sketball! Are you _	?	
1000			
1	outside	my	
=	outside What	my	
	What	robot	



Visual script backend

Assume the frontend exists. Options for backend:

- Visuals map to script code in a text file. Just compile and run the text.
- Visuals map to objects in a high level state machine. Make those objects.

That's all we will say about that.

Case Study: Lua

- Handwritten recursive descent parser
 - No AST, just emit bytecode during parse
- Register based VM
- Instruction width is 32 bits

Lua 5.1 Instruction Set

MOVE Copy a value between registers

LOADK Load a constant into a register

LOADBOOL Load a boolean into a register

LOADNIL Load nil values into a range of registers

GETUPVAL Read an upvalue into a register

GETGLOBAL Read a global variable into a register

GETTABLE Read a table element into a register

SETGLOBAL Write a register value into a global variable

SETUPVAL Write a register value into an upvalue

SETTABLE Write a register value into a table element

NEWTABLE Create a new table

SELF Prepare an object method for calling

ADD Addition operator

SUB Subtraction operator

MUL Multiplication operator

DIV Division operator

MOD Modulus (remainder) operator

POW Exponentiation operator

UNM Unary minus operator

NOT Logical NOT operator

LEN Length operator

CONCAT Concatenate a range of registers

JMP Unconditional jump

EQ Equality test

LT Less than test

LE Less than or equal to test

TEST Boolean test, with conditional jump

TESTSET Boolean test, with conditional jump and assignment

CALL Call a closure

TAILCALL Perform a tail call

RETURN Return from function call

FORLOOP Iterate a numeric for loop

FORPREP Initialization for a numeric for loop

TFORLOOP Iterate a generic for loop

SETLIST Set a range of array elements for a table

CLOSE Close a range of locals being used as upvalues

CLOSURE Create a closure of a function prototype

VARARG Assign vararg function arguments to registers

Lua likes Tables

In Lua, lots of things are tables (associative arrays):

- Vectors
- Maps
- Objects
- Modules...

For example:

```
Point = { x=5, y=10 }
print(Point["x"], Point["y"])
```

Let's look at our homework assignment...

Summary

- Features of scripting languages in games
- Parser and generator concepts
- Virtual machines
- Integrating scripting into an engine
- Visual script
- Lua

End lecture

A lot of the content in this lecture was adapted from:

Language Implementation Patterns: Create Your Own Domain-Specific and General Programming Language. Terence Parr.

Start writing your own language, compiler and VM:

- <u>http://llvm.org/docs/tutorial/index.html</u>
- http://www.iro.umontreal.ca/~felipe/IFT2030-Automne2002/Complements/tinyc.c