

The code is printing the count of the substrings, Guess the output.*

```
public class Main {
    public static void main(String[] args) {
        String text = "abababab";
        String substring = "aba";
        int count = 0;

        for (int i = 0; i <= text.length() - substring.length(); i++) {
            if (text.substring(i, i + substring.length()).equals(substring)) {
                count++;
            }
        }

        System.out.println(count);
    }
}
```

2 text: a b a b a b a b
 3

4 substring: a b a
 5

Bruteforce Algorithms

Recap:

- Array indexing
- Substrings
- Subarrays
- Subsequences.

Bruteforce
(Search) algorithm

"Try all possible
solutions (systematically)."



10,000 possible
solutions!

0	0	0	0
0	0	0	1
0	0	0	2
:			
0	0	0	9
0	0	1	0
:			
9	9	9	9

Systematically
try all
possible
solutions

Features of Brute force algorithms

- Blind search : No strategy involved.
(naive)
- Guaranteed to find solution (if it exists)
- Often easiest to understand & implement
- Mostly very slow. (take a lot of time to run)

$$4 \text{ digits} = 10^4 \text{ possible soln}$$

$$5 \text{ digits} = 10^5 \text{ " "$$

Pattern of a Brute force algorithm :

- Generate all possible candidate solutions
- For each candidate solution,
 - check if it satisfies constraints of the problem
- (opt) - pick the best solution.

Ex-1: Find all the divisors of a number n (factors)

$$n = 16$$

$$n < 10,000$$

Generating
all possible
solutions

```

 $\frac{2}{1}$   $\frac{n \% i}{16 \% 1 = 0} \checkmark$  1, 2, 4, 8, 16 .
2  $16 \% 2 = 0 \checkmark$ 
3  $16 \% 3 = 1$ 
4  $16 \% 4 = 0 \checkmark$  for  $i = 1$  to  $n \{$ 
 $\vdots$  if ( $n \% i == 0 \{$ 
16  $16 \% 16 = 0 \checkmark$  print ( $i \}$ )
 $\}$ 

```

Ex 2: Find the lowest common multiple (l.c.m) of two numbers n & m .

e.g.: $n = 6$

$$m = 8$$

Multiples of 6 : 6, 12, 18, 24, 30, 36,
Multiples of 8 : 8, 16, 24, 32, 40, 48...

i
max(n, m) check if i is a multiple
of both n & m .

$$6\pi^{\frac{1}{2}} \approx 48 \quad n \neq m$$

1. Generate all possible solutions
 2. check if it satisfies constraints
(if it is a multiple of both n & m)

$n = 6, m = 8$

```

for (int i = Math.max(n, m); i <= n*m; i++) {
    if ((i % n == 0) && (i % m == 0)) {
        S.O.P(i);
        break;
    }
}

```

<u>i</u>	<u>$i \leq 6*8$</u>	<u>$(i \% n == 0) \& (i \% m == 0)$</u>
8	$8 \leq 48 \checkmark$	$8 \% 6 \times$
9	$9 \leq 48 \checkmark$	$9 \% 6 == 0 \times \& (9 \% 8 == 0)$
<u>i</u>	<u>$i \leq 48$</u>	<u>$i \% 6 \& i \% 8$</u>

e.g.: $n = 2, m = 3$.

$i = \max(n, m), i \leq n*m$

<u>i</u>	<u>$2 * 3$</u>
2	6
3	
4	
5	
6	

Count Such Pairs

A	3	0	6	2	7
---	---	---	---	---	---

A.length = 5

$$k = 9 \quad \left| \begin{array}{l} \{3, 6\}, 3+6=9 \\ \{2, 7\}, 2+7=9 \\ \text{Output: } 2 \end{array} \right.$$

Brute force:

1. Generate all possible pairs
2. For each pair,
check if its sum == k

int count = 0;

```
for(int i=0; i < A.length-1; i++) {
    for(int j = 4i+1; j < A.length; j++){
        // {A[i], A[j]} is the current pair
        if(A[i] + A[j] == k)
            count++;
}
```

<u>i</u>	<u>j</u>	<u>Pairs</u>	<u>Sum</u>
0	1	3, 0	3
0	2	3, 6	9 ✓
0	3	3, 2	5
0	4	3, 7	10
1	2	0, 6	6
1	3	0, 2	2
1	4	0, 7	7
2	3	6, 2	8
2	4	6, 7	13
3	4	2, 7	9 ✓

Intersection of Arrays

<u>arr1</u>	<u>i</u>	<u>3 5 2</u>	<u>j</u>
<u>arr2</u>		<u>5 -10 8</u>	

" For every element in arr1
 - check if it appears
 in arr2.

```
for (i = 0; i < n; i++) {
```

x = arr1[i]

// Check if x appears in arr2.

```
for (j = 0; j < n; j++) {
```

if (x == arr2[j])

S.O.P(x);
break;

3

7

5

3

if (arr1[i] == arr2[j]) {
S.O.P(arr1[i]);
break}

3

