

📖 Target Sum – Complete Documentation

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1. ★ Problem Statement

You are given an array of integers `nums` and an integer `target`. Your task is to build expressions by adding either a '+' or '-' sign before each number in the array, and calculate how many different expressions evaluate to the target.

📌 Example:

- Input: `nums = [1,1,1,1,1]`, `target = 3`
- Output: 5

Constraints:

- $1 \leq \text{nums.length} \leq 20$
- $0 \leq \text{nums}[i] \leq 1000$
- $-1000 \leq \text{target} \leq 1000$

2. 💡 Intuition

Instead of brute-forcing all 2^n combinations of '+' and '-', we can reframe the problem using a mathematical transformation that connects it to the Subset Sum problem—a classic dynamic programming question.

3. □ Key Observations

- Let the total sum of all numbers be S .
- If we split numbers into two groups:
 - $P \rightarrow$ numbers with +
 - $N \rightarrow$ numbers with -
- Then:
 - $P - N = \text{target}$
 - $P + N = S$
- Adding these two:
 - $2P = \text{target} + S$
 - $P = (\text{target} + S) / 2$

So, the problem reduces to finding how many subsets of nums sum to P .

4. ✂ Approach

- Compute the total sum S of the array.
- Check if $(\text{target} + S)$ is even and $P = (\text{target} + S) / 2$ is valid.
- Use 1D Dynamic Programming to count the number of subsets that sum to P .

□ DP Array:

- $\text{dp}[i]$ will store the number of ways to get a sum of i .
- Start with $\text{dp}[0] = 1$ (1 way to make sum 0: use nothing).
- For each number in `nums`, iterate backwards to update `dp`.

5. Edge Cases

- If $(\text{target} + \text{total_sum})$ is odd, return 0 \rightarrow subset sum is not an integer.
- If $\text{abs}(\text{target}) > \text{total_sum}$, return 0 \rightarrow target not achievable.
- If nums contains zeros, multiple combinations may contribute to same sum.

6. Complexity Analysis

□ Time Complexity:

- $O(n * \text{subset_sum})$ where n is the length of the array, and $\text{subset_sum} = (\text{target} + \text{total_sum}) / 2$

☞ Space Complexity:

- $O(\text{subset_sum})$ due to 1D DP array

7. Alternative Approaches

- Recursive with Memoization:
 - Recursively add and subtract current number
 - Use a dictionary to memoize overlapping states
- Brute Force (Not Recommended):
 - Try all 2^n combinations of '+' and '-'
 - Inefficient for $n > 15$

8. □ Test Cases

✓ Example 1:

Input: $\text{nums} = [1, 1, 1, 1, 1]$, $\text{target} = 3$

Output: 5

✓ Example 2:

Input: $\text{nums} = [1]$, $\text{target} = 1$

Output: 1

✓ Edge Case:

Input: $\text{nums} = [0,0,0,0,0,0,0,1]$, $\text{target} = 1$

Output: 256

9. □ Final Thoughts

- Transforming the problem into a Subset Sum question makes it significantly more efficient.
- This technique is commonly useful in problems involving binary choices (like '+' or '-').
- Knowing how to reframe problems mathematically can unlock easier, optimized solutions.