Documentation on Minimum Genetic Mutation - Complete Documentation

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1. Problem Statement

A gene string is represented as an 8-character long string with characters from {'A', 'C', 'G', 'T'}. Given:

- A startGene (valid mutation starting point).
- An endGene (target mutation).
- A list bank containing valid intermediate gene mutations.

Objective:

Find the **minimum number of mutations** required to convert startGene into endGene, where:

- Each mutation changes exactly **one** character.
- The resulting gene **must be in the bank**.
- If no mutation path exists, return -1.

Example 1:

Input:

```
startGene = "AACCGGTT"
endGene = "AACCGGTA"
bank = ["AACCGGTA"]
```

Output:1

Example 2:

Input:

```
startGene = "AACCGGTT"
endGene = "AAACGGTA"
bank = ["AACCGGTA", "AACCGCTA", "AAACGGTA"]
```

Output: 2

2. Intuition

This problem can be visualized as a graph traversal:

- Each valid gene represents a **node**.
- A mutation that changes one character is an **edge**.
- We need the **shortest path** from startGene to endGene, which suggests using **Breadth-First Search** (**BFS**).

BFS is ideal because:

- It explores mutations **level-by-level**, ensuring the shortest path is found first.
- It avoids unnecessary computations using a visited set.

3. Key Observations

- If endGene is **not in bank**, return -1 immediately.
- Each mutation can only change **one character at a time**.
- There are only 8 positions and 4 possible characters per position → max 24 possible mutations per gene.
- The bank length is **at most 10**, making BFS feasible.

4. Approach

Step 1: Convert bank to a Set for Fast Lookup

• A set provides O(1) lookup time to check valid mutations.

Step 2: Initialize BFS

- Use a queue storing tuples (currentGene, mutationCount).
- Start with (startGene, 0).
- Track visited genes to avoid cycles.

Step 3: Process Queue (BFS Traversal)

- 1. **Extract** a gene from the queue.
- 2. Try mutating each character at every position (A, C, G, T).
- 3. If the new mutation is in bank and not visited:
 - \circ If it matches endGene, return mutations count + 1.
 - Otherwise, add it to the queue for further exploration.
- 4. If the queue is exhausted without finding endGene, return -1.

5. Edge Cases

- ✓ endGene not in bank \rightarrow Return -1.
- \checkmark startGene == endGene → Return 0.

- ✓ No valid mutations possible \rightarrow Return -1.
- ✓ Multiple mutation paths exist \rightarrow BFS ensures the shortest path is chosen.
- ✓ Smallest case: bank is empty \rightarrow Return -1.

6. Complexity Analysis

Time Complexity

- Each gene has **8 positions**.
- Each position has **3 possible mutations** (excluding itself).
- There are at most 10 genes in the bank.
- Worst-case scenario: O(8 * 3 * N) = O(N), which is very efficient.

Space Complexity

- BFS queue stores up to N genes \rightarrow O(N).
- Visited set and bank set \rightarrow O(N).
- Overall Space Complexity: O(N).

7. Alternative Approaches

a. Depth-First Search (DFS)

- Could work but may not find the shortest path first.
- Requires backtracking, leading to **higher time complexity**.

2. Dijkstra's Algorithm

• Overkill for this problem since BFS finds the shortest path efficiently.

3. Bidirectional BFS

• Would optimize performance but is unnecessary given the small constraints.

8. Test Cases

```
# Test Case 1

print(solution.minMutation("AACCGGTT", "AACCGGTA", ["AACCGGTA"])) # Expected Output: 1

# Test Case 2

print(solution.minMutation("AACCGGTT", "AAACGGTA",

["AACCGGTA","AACCGCTA","AAACGGTA"])) # Expected Output: 2

# Test Case 3 (Impossible case)

print(solution.minMutation("AACCGGTT", "AAACGGTA", ["AACCGGTC"])) # Expected Output: -1

# Test Case 4 (EndGene not in bank)

print(solution.minMutation("AACCGGTT", "AACCGGTA", [])) # Expected Output: -1
```

9. Final Thoughts

- **BFS** ensures the shortest mutation path is found efficiently.
- **Set lookups** make operations fast.
- The approach works well given the problem constraints $(N \le 10)$.
- For **larger constraints**, bidirectional BFS could be explored.