

### 1. Problem Statement

Given a positive integer  $n$ , find the smallest integer that:

- Has exactly the same digits as  $n$ .
- Is greater than  $n$ .
- Fits in a 32-bit signed integer.

If no such number exists, return -1.

❖ Example 1:

Input:  $n = 12$

Output: 21

❖ Example 2:

Input:  $n = 21$

Output: -1

### 2. Intuition

This is a classic "next permutation" problem:

Find the next number that is greater than the current one by rearranging its digits.

We seek the next lexicographical permutation of the digits in  $n$ .

### 3. Key Observations

- If digits are sorted in descending order (like 321), no greater permutation exists.
- To form the next permutation:
  - Find a pivot: the first digit from the right which is smaller than the digit next to it.

- Find the smallest digit on the right side of the pivot that is greater than the pivot.
- Swap the pivot with this digit.
- Reverse the digits to the right of the pivot.

#### 4. Approach

- Convert the number into a list of digits.
- Traverse from right to left to find the first decreasing digit.
- If not found, return -1.
- Find the smallest digit greater than this pivot on its right.
- Swap and reverse the sublist.
- Convert back to integer and return if within 32-bit limit.

#### 5. Edge Cases

- Input has only one digit  $\rightarrow$  return -1.
- All digits are in descending order  $\rightarrow$  return -1.
- Result exceeds  $2^{31} - 1 \rightarrow$  return -1.
- Duplicates exist  $\rightarrow$  still valid if reordering forms a greater number.

#### 6. Complexity Analysis

□ Time Complexity:

- $O(n)$  where  $n$  is the number of digits.

▣ Space Complexity:

- $O(n)$  for storing digits as a list.

## 7. Alternative Approaches

- Brute Force: Generate all permutations, sort and find next  $\rightarrow$  Inefficient:  $O(n!)$  time.
- Using built-in next\_permutation (C++): Fast but not available in all languages.
- Heap-based methods: Overkill for this problem.

## 8. Algorithm

- Convert number to list of digits.
- Find the first index  $i$  from the end such that  $\text{digits}[i] < \text{digits}[i+1]$ .
- If no such index exists, return -1.
- Find index  $j$  such that  $\text{digits}[j] > \text{digits}[i]$ , starting from the end.
- Swap  $\text{digits}[i]$  and  $\text{digits}[j]$ .
- Reverse the sublist from  $i+1$  to the end.
- Convert the list back to integer and check if it fits in 32-bit.

## 9. Test Cases

Input	Output	Description
12	21	Next permutation exists.
21	-1	No greater permutation.
1234	1243	Smallest next permutation.
4321	-1	Already highest permutation.
1999999999	-1	Result exceeds 32-bit limit.

## 10. Final Thoughts

This problem blends math, logic, and algorithm design. It's a great case of:

- Understanding permutations.
- Applying optimal  $O(n)$  solution.
- Handling constraints like integer bounds.