

■ Beautiful Arrangement - Documentation

1. Problem Statement

Given an integer n , return the number of beautiful arrangements that can be constructed using the integers from 1 to n .

A permutation perm of length n is called a beautiful arrangement if for every index i (1-indexed), either:

- $\text{perm}[i] \% i == 0$ or
- $i \% \text{perm}[i] == 0$

2. Intuition

We are generating permutations of numbers from 1 to n under specific divisibility rules. Instead of generating all $n!$ permutations, we can use backtracking to only explore valid candidates at each position based on the condition, which greatly reduces computation.

3. Key Observations

- The constraint $\text{perm}[i] \% i == 0$ or $i \% \text{perm}[i] == 0$ acts as a filter during permutation generation.
- We can prune invalid paths early during recursion.
- Using a boolean visited array avoids placing the same number more than once.

4. Approach

- Start from position 1 to n .
- For each position i , try all numbers 1 to n .
- Only place a number at index i if:
 - num hasn't been used ($\text{not visited}[\text{num}]$) and
 - $\text{num} \% i == 0$ or $i \% \text{num} == 0$.

- Recurse to the next position.
- When $\text{pos} > n$, increment the count of valid arrangements.

5. Edge Cases

- $n = 1$: Only one number \rightarrow valid by default.
- $n = 15$: Maximum value, performance still acceptable due to pruning.

6. Complexity Analysis

□ Time Complexity

- Worst case: $O(n!)$ without pruning.
- With pruning: Much less than $n!$, exact depends on divisibility conditions.

□ Space Complexity

- $O(n)$ for visited array and recursion stack.

7. Alternative Approaches

- Dynamic Programming with Bitmasking: Store state as a bitmask and use memoization to avoid recomputation.
 - More efficient for large n .
 - More complex to implement.

8. Test Cases

Input	Output	Explanation
$n = 1$	1	Only one permutation: $[1]$ ✓
$n = 2$	2	$[1,2]$ and $[2,1]$ both valid ✓
$n = 3$	3	Valid permutations: $[1,2,3]$, $[3,2,1]$, $[2,1,3]$ ✓

9. Final Thoughts

This problem is a classic example of using backtracking with constraints to avoid unnecessary computation. The core idea is not generating permutations blindly, but making informed choices using divisibility rules. While the brute-force approach is infeasible for $n = 15$, the optimized backtracking handles it efficiently.