132 Pattern Detection - Detailed Documentation

Table of Contents

- 1. Problem Statement
- 2. Intuition
- 3. Key Observations
- 4. Approach
- 5. Edge Cases
- 6. Complexity Analysis
 - o <u>Time Complexity</u>
 - Space Complexity
- 7. Alternative Approaches
- 8. Test Cases
- 9. Final Thoughts

1. Problem Statement

Given an array of n integers nums, a 132 pattern is a subsequence of three integers nums[i], nums[j], and nums[k] such that:

- i < j < k
- nums[i] < nums[k] < nums[j]

Return True if there is a 132 pattern in nums, otherwise, return False.

Example 1

- Input: nums = [1, 2, 3, 4]
- Output: False
- Explanation: There is no 132 pattern.

Example 2

- Input: nums = [3, 1, 4, 2]
- Output: True
- Explanation: The pattern [1, 4, 2] follows 132 order.

Example 3

- Input: nums = [-1, 3, 2, 0]
- Output: True
- Explanation: Multiple valid patterns: [-1, 3, 2], [-1, 3, 0], [-1, 2, 0].

Constraints

- n == nums.length
- $1 \le n \le 2 * 10^5$
- -10⁹ <= nums[i] <= 10⁹

2. Intuition

The goal is to identify three numbers in the array following the 132 pattern constraint efficiently.

- A brute-force approach checking all triplets (i, j, k) is too slow (O(n³) complexity).
- A more optimal approach is to track potential nums[k] using a monotonic stack and iterate in reverse order.

3. Key Observations

- i. We need three numbers where:
 - a. nums[i] < nums[k] < nums[j]
 - b. i < j < k
- ii. If we iterate backward, we can maintain the largest valid nums[k] efficiently using a monotonic decreasing stack.
- iii. The nums[k] value must be **less than** nums[j] **but greater than** nums[i], which can be efficiently tracked using a **stack**.

4. Approach

We use a monotonic decreasing stack to track elements from right to left:

- 1. Initialize variables:
 - o stack = ☐ (monotonic decreasing stack)
 - o third = $-\infty$ (potential nums[k])
- 2. Iterate from right to left:
 - o If nums[i] < third, return True (we found a valid pattern).
 - While stack is non-empty and stack[-1] < nums[i], update third = stack.pop() (ensuring it remains the largest possible nums[k]).
 - o Push nums[i] onto the stack.
- 3. Return False if no pattern is found.

5. Edge Cases

Case Type	Example	Expected	Reason
		Output	
Minimum	[1]	False	Only one element, cannot form a pattern.
Input			
All Increasing	[1, 2, 3, 4,	False	No valid nums[k] satisfying nums[i] < nums[k]
	5]		< nums[j].
All Decreasing	[5, 4, 3, 2,	False	No valid 132 pattern possible.
	1]		
Negative	[-1, 3, 2, 0]	True	Multiple valid patterns exist.
Values			
Duplicates	[3, 3, 3, 3,	False	All elements are the same, no pattern possible.
	3]		
Large Input	[10 ⁵ ,, 1]	False	Fully decreasing, so no 132 pattern.

6. Complexity Analysis

Time Complexity

• O(n) – Each element is pushed and popped from the stack at most once.

Space Complexity

• O(n) – In the worst case, all elements are stored in the stack.

7. Alternative Approaches

Brute Force (O(n³))

- Iterate through all triplets (i, j, k).
- Compare each nums[i], nums[k]` for the 132 pattern.
- Time Complexity: O(n³) (too slow for large inputs).

Using Two Arrays (O(n²))

- Precompute the minimum nums[i] up to index j.
- Use a nested loop to find nums[k] satisfying the pattern.
- Time Complexity: O(n²), still inefficient.

8. Test Cases

```
def test_find132pattern():
sol = Solution()
assert sol.find132pattern([1, 2, 3, 4]) == False
assert sol.find132pattern([3, 1, 4, 2]) == True
assert sol.find132pattern([-1, 3, 2, 0]) == True
assert sol.find132pattern([3, 3, 3, 3, 3]) == False
assert sol.find132pattern([1, 0, 1, -4, -3]) == False
assert sol.find132pattern([9, 11, 8, 9, 10, 7, 8]) == True
assert sol.find132pattern([10, 20, 30, 40]) == False
assert sol.find132pattern([3, 5, 0, 3, 4]) == True
print("All test cases passed!")
test_find132pattern()
```

9. Final Thoughts

- The monotonic stack approach efficiently detects the 132 pattern in O(n) time.
- The solution works well even for large inputs (n $\approx 2 \times 10^5$).
- Alternative approaches like **brute force** and **nested loops** are too slow.
- This is an excellent example of using stacks to solve sequence-based problems efficiently.