

# Arranging Coins - Detailed Documentation

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### 1. Problem Statement

You are given  $n$  coins, and you want to build a staircase where the  $i$ -th row contains exactly  $i$  coins. The staircase stops when the remaining coins are not enough to complete the next row. Return the number of complete rows of the staircase that can be built.

#### Example 1:

Input:  $n = 5$

Output: 2

Explanation:

```
*  
**  
***
```

Since the third row is incomplete, the output is 2.

**Example 2:**

Input:  $n = 8$

Output: 3

Explanation:

```
*  
**  
***  
****
```

Since the fourth row is incomplete, the output is 3.

**2. Intuition**

The number of coins required to form  $k$  complete rows is the sum of the first  $k$  natural numbers:

$$S_k = k * (k + 1) / 2$$

To find the largest  $k$  such that:

$$k * (k + 1) / 2 \leq n$$

We can solve for  $k$  using the quadratic equation.

**3. Key Observations**

- The sum of the first  $k$  numbers follows the arithmetic sum formula:

$$S_k = k * (k + 1) / 2$$

- Given  $n$ , we need to find the largest integer  $k$  such that:

$$k * (k + 1) \leq 2 * n$$

- This can be solved using the quadratic formula.

#### 4. Approach

- i. Use the quadratic formula: The equation

$$k^2 + k - 2n = 0$$

is a quadratic equation of the form:

$$ax^2 + bx + c = 0$$

Here,

- $a = 1$
- $b = 1$
- $c = -2n$

- ii. Solve for k using the quadratic formula:

$$k = (-1 + \sqrt{1 + 8 * n}) / 2$$

- iii. Take the integer part (floor value) to get the maximum number of complete rows.

#### 5. Edge Cases

Case	Explanation
$n = 1$	Only 1 row can be formed. Output: 1
$n = 0$	No coins available. Output: 0
$n = 2$	Only 1 full row (*) can be formed. Output: 1
Large n (e.g., $2^{31} - 1$ )	Tests efficiency and integer handling.

## 6. Complexity Analysis

### Time Complexity

- The formula uses only one square root calculation:  $O(1)$
- No loops are required.
- Overall Time Complexity:  $O(1)$

### Space Complexity

- No extra space is used.
- Overall Space Complexity:  $O(1)$

## 7. Alternative Approaches

### i. Linear Iteration (Brute Force)

- Start with  $k = 1$  and subtract coins row by row until  $n < k$ .
- Time Complexity:  $O(\sqrt{n})$ .
- Inefficient for large values of  $n$ .

### ii. Binary Search

- Use binary search on the range  $[1, n]$  to find  $k$  such that:

$$k * (k + 1) / 2 \leq n$$

- Time Complexity:  $O(\log n)$ , better than brute force but still slower than  $O(1)$ .

## 8. Test Cases

### Basic Cases

```
assert Solution().arrangeCoins(5) == 2 # Example 1
assert Solution().arrangeCoins(8) == 3 # Example 2
assert Solution().arrangeCoins(1) == 1 # Smallest case
assert Solution().arrangeCoins(0) == 0 # Edge case with no coins
assert Solution().arrangeCoins(2) == 1 # Incomplete second row
```

### Large Input

```
assert Solution().arrangeCoins(2147483647) == 65535 # Large case
```

## 9. Final Thoughts

- The mathematical approach using the quadratic formula is the best solution ( $O(1)$  time complexity).
- Alternative methods (linear, binary search) work but are less efficient.
- The formula-based solution is clean, efficient, and precise for handling large  $n$ .