#### ■ Most Frequent Subtree Sum - Documentation

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#### 1. **♦** Problem Statement

Given the root of a binary tree, return the most frequent subtree sum(s).

- A subtree sum is the sum of all values in a subtree rooted at a node (including the node itself).
- If there are multiple sums with the same maximum frequency, return them in any order.

## **♦** Constraints:

- $1 \le \text{Number of nodes} \le 10^4$
- $-10^5 \le Node.val \le 10^5$

# 2. Intuition

Every node defines a subtree, and each subtree has a sum. To solve the problem:

- Traverse the tree
- Calculate the sum of each subtree
- Track how many times each sum appears
- Return the sum(s) that appear most frequently

# 3. Q Key Observations

- A post-order traversal (left-right-root) is ideal since we need to calculate the subtree sum after computing the sums of left and right children.
- We can use a hash map to count how often each subtree sum occurs.

#### 4. Approach

- Use a recursive post-order traversal (dfs) to:
  - o Compute the left and right subtree sums
  - o Add them with the current node's value
  - O Update a Counter to store frequency of each subtree sum
- After the traversal:
  - Identify the maximum frequency
  - Return all subtree sums that match that frequency

#### 5. ▲ Edge Cases

- Tree has only one node → the subtree sum is the node value itself.
- All nodes have value  $0 \rightarrow$  subtree sum of any node is also 0.
- Negative values → must handle them without assumptions about positivity.

# **6.** □ Complexity Analysis $\square$ Time Complexity: • O(n) where n is the number of nodes Each node is visited exactly once ☐ Space Complexity: O(n) for: Recursion stack (in worst case: skewed tree) HashMap storing frequencies of subtree sums 7. Alternative Approaches Iterative DFS or BFS: Possible but complex due to the need to compute child sums first. Using a global variable instead of returning sums: Less clean, more side-effects. Serialization-based approach: Inefficient for large trees, not recommended. 8. Test Cases Test Case 1: Input: root = [5, 2, -3]Output: [2, -3, 4] Test Case 2: Input: root = [5, 2, -5]Output: [2] Test Case 3:

Input: root = [1]

Output: [1]

## Test Case 4:

Input: root = [0, 0, 0]Output: [0]

# 9. ☐ Final Thoughts

- This problem highlights the power of post-order traversal for bottom-up computations in trees.
- It's a great example of combining DFS with hash maps to track frequencies.
- Easily extendable to related problems like "Most Frequent Path Sum", "Max Subtree Product", etc.