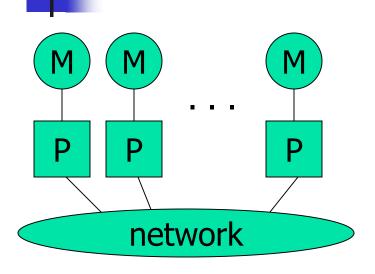


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Message Passing



- Each processor runs a process
- Processes communicate by exchanging messages
- They cannot share memory in the sense that they cannot address the same memory cells
- The above is a programming model and things may look different in the actual implementation (e.g., MPI over Shared Memory)
- Message Passing is popular because it is general:
 - Pretty much any distributed system works by exchanging messages, at some level
 - Distributed- or shared-memory multiprocessors, networks of workstations, uniprocessors
- It is not popular because it is easy (it's not)

MPI Concepts

Fixed number of processors

 When launching the application one must specify the number of processors to use, which remains unchanged throughout execution

Communicator

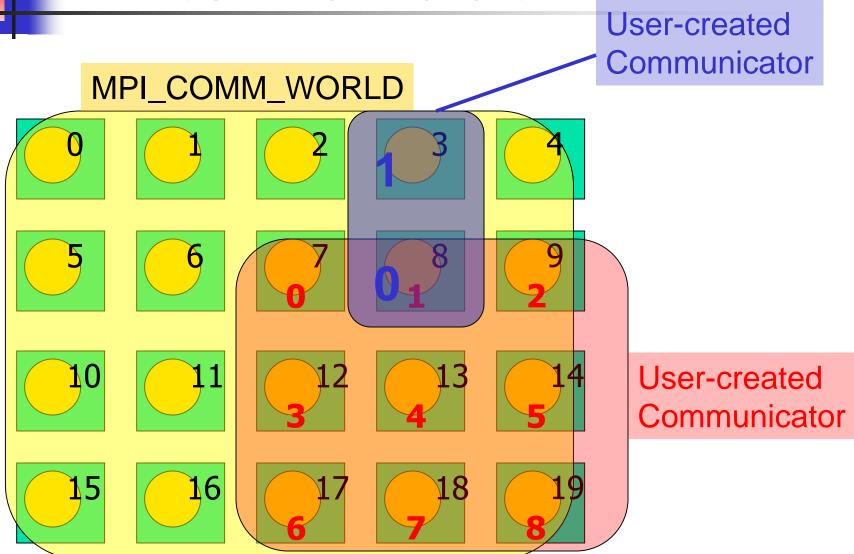
- Abstraction for a group of processes that can communicate
- A process can belong to multiple communicators
- Makes is easy to partition/organize the application in multiple layers of communicating processes
- Default and global communicator: MPI COMM WORLD

Process Rank

- The index of a process within a communicator
- Typically user maps his/her own virtual topology on top of just linear ranks
 - ring, grid, etc.



MPI Communicators



A First MPI Program

```
#include <unistd.h>
#include <mpi.h>
int main(int argc, char **argv) {
  int my rank, n;
                              Has to be called first, and once
  char hostname [128];
 MPI init(&argc, &argv);
 MPI Comm rank (MPI COMM WORLD, &my rank);
 MPI Comm size (MPI COMM WORLD, &n);
  gethostname (hostname, 128);
  if (my rank == 0) { /* master */
    printf("I am the master: %s\n", hostname);
  } else { /* worker */
    printf("I am a worker: %s (rank=%d/%d)\n",
           hostname, my rank, n-1);
 MPI Finalize();
                              Has to be called last, and once
  exit(0); }
```

Compiling/Running it

- Compile with mpicc
- Run with mpirun

```
% mpirun -np 4 my_program <args>
```

- requests 4 processors for running my_program with command-line arguments
- see the mpirun man page for more information
- in particular the -machinefile option that is used to run on a network of workstations
- Some systems just run all programs as MPI programs and no explicit call to mpirun is actually needed
- Previous example program:

```
% mpirun -np 3 -machinefile hosts my_program
I am the master: somehost1
I am a worker: somehost2 (rank=2/2)
I am a worker: somehost3 (rank=1/2)
    (stdout/stderr redirected to the process calling mpirun)
```

Point-to-Point Communication



- Data to be communicated is described by three things:
 - address
 - data type of the message
 - length of the message
- Involved processes are described by two things
 - communicator
 - rank
- Message is identified by a "tag" (integer) that can be chosen by the user

Point-to-Point Communication

- Two modes of communication:
 - Synchronous: Communication does not complete until the message has been received
 - Asynchronous: Completes as soon as the message is "on its way", and hopefully it gets to destination
- MPI provides four versions
 - synchronous, buffered, standard, ready

Synchronous/Buffered sending in MPI

- Synchronous with MPI_Ssend
 - The send completes only once the receive has succeeded
 - copy data to the network, wait for an ack
 - The sender has to wait for a receive to be posted
 - No buffering of data
- Buffered with MPI_Bsend
 - The send completes once the message has been buffered internally by MPI
 - Buffering incurs an extra memory copy
 - Does not require a matching receive to be posted
 - May cause buffer overflow if many bsends and no matching receives have been posted yet

Standard/Ready Send

- Standard with MPI_Send
 - Up to MPI to decide whether to do synchronous or buffered, for performance reasons
 - The rationale is that a correct MPI program should not rely on buffering to ensure correct semantics
- Ready with MPI_Rsend
 - May be started only if the matching receive has been posted
 - Can be done efficiently on some systems as no handshaking is required

Example: Sending and Receiving

```
#include <unistd.h>
#include <mpi.h>
int main(int argc, char **argv) {
  int i, my rank, nprocs, x[4];
 MPI Init(&argc, &argv);
 MPI Comm rank (MPI COMM WORLD, &my rank);
                                                          destination
  if (my rank == 0) { /* master */
                                                            and
    x[0]=42; x[1]=43; x[2]=44; x[3]=45;
                                                           source
    MPI Comm size (MPI COMM WORLD, &nprocs);
    for (i=1;i<nprocs;i++)</pre>
      MPI_Send(x, 4, MPI INT, i, 0, MPI_cOMM WORLD);
                                                         user-defined
  } else { /* worker */
                                                             tag
    MPI Status status;
    MPI Recv(x, 4, MPI INT 0, 0, MPI COMM WORLD, &status);
  MPI Finalize();
                         Max number of
                                              Can be examined via calls
  exit(0);
                        elements to receive
                                              like MPI_Get_count(), etc.
```

Non-blocking Communication

MPI_Issend, MPI_Ibsend, MPI_Isend, MPI_Irsend, MPI_Irecv

```
MPI_Request request;
MPI_Isend(&x,1,MPI_INT,dest,tag,communicator,&request);
MPI_Irecv(&x,1,MPI_INT,src,tag,communicator,&request);
```

 Functions to check on completion: MPI_Wait, MPI_Test, MPI_Waitany, MPI_Testany, MPI_Waitall, MPI_Testall, MPI_Waitsome, MPI_Testsome.

```
MPI_Status status;
MPI_Wait(&request, &status) /* block */
MPI_Test(&request, &status) /* doesn't block */
```

Collective Communication

- Operations that allow more than 2 processes to communicate simultaneously
 - barrier
 - broadcast
 - reduce
- All these can be built using point-to-point communications, but typical MPI implementations have optimized them, and it's a good idea to use them
- In all of these, all processes place the same call (in good SPMD fashion), although depending on the process, some arguments may not be used

Barrier

- Synchronization of the calling processes
 - the call blocks until all of the processes have placed the call
- No data is exchanged

```
MPI_Barrier (MPI_COMM_WORLD)
```

Broadcast

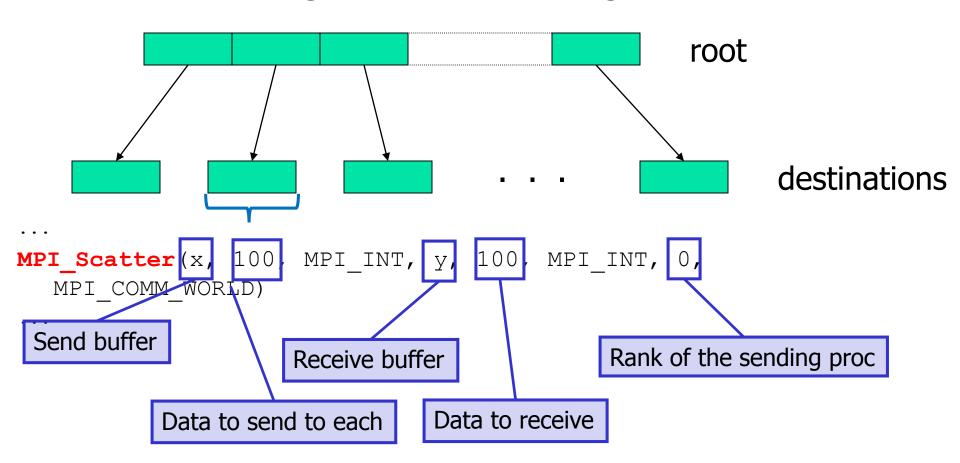
- One-to-many communication
- Note that multicast can be implemented via the use of communicators (i.e., to create processor groups)

```
MPI_Bcast(x, 4, MPI_INT, 0
MPI_COMM_WORLD)
....

Rank of the root
```

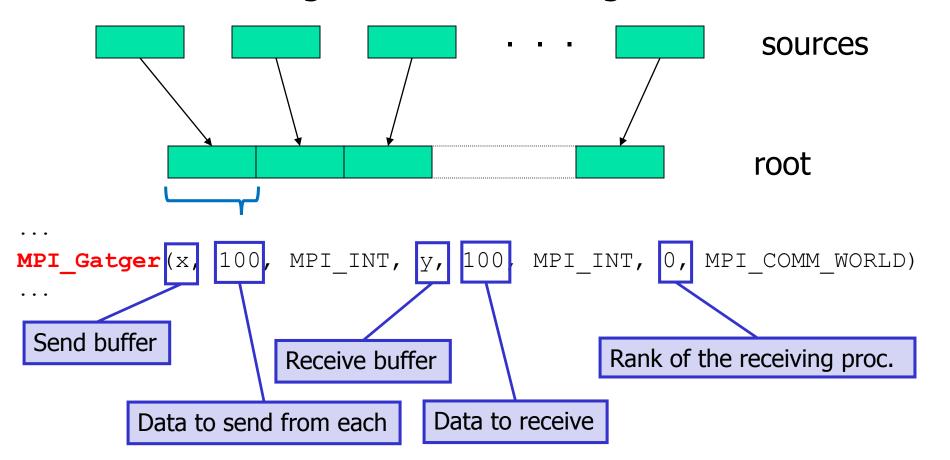
Scatter

- One-to-many communication
- Not sending the same message to all



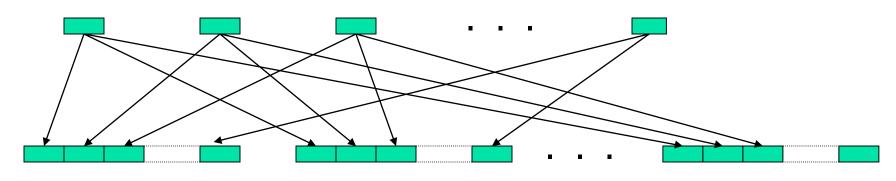
Gather

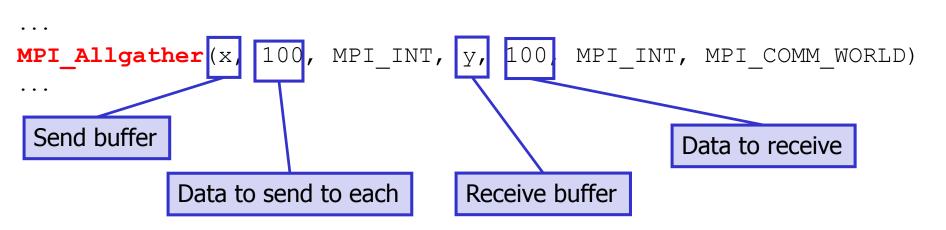
- Many-to-one communication
- Not sending the same message to the root



Gather-to-all

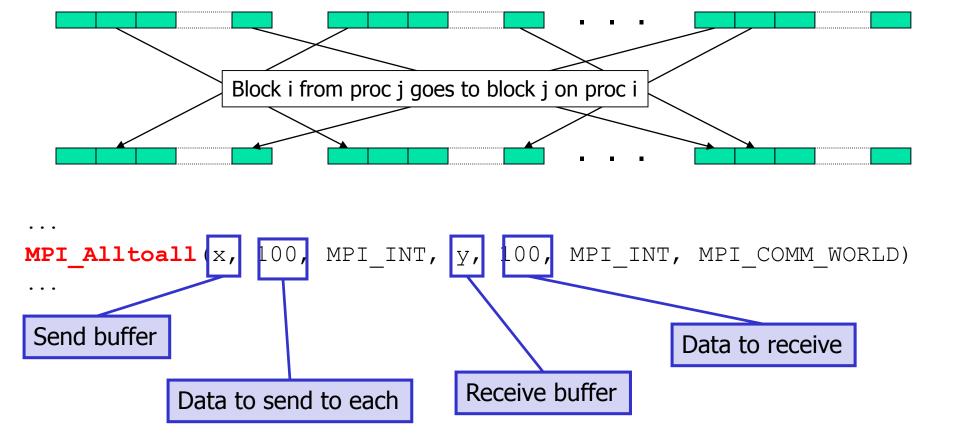
- Many-to-many communication
- Each process sends the same message to all
- Different Processes send different messages





All-to-all

- Many-to-many communication
- Each process sends a different message to each other process



All-to-all example

task1

task2

task3

← sendbuf (before)

recvbuf (after)

Reduction Operations

- Used to compute a result from data that is distributed among processors
 - often what a user wants to do anyway
 - so why not provide the functionality as a single API call rather than having people keep re-implementing the same things
- Predefined operations:
 - MPI_MAX, MPI_MIN, MPI_SUM, etc.
- Possibility to have user-defined operations



MPI_Reduce, MPI_Allreduce

- MPI_Reduce: result is sent out to the root
 - the operation is applied element-wise for each element of the input arrays on each processor
- MPI_Allreduce: result is sent out to everyone

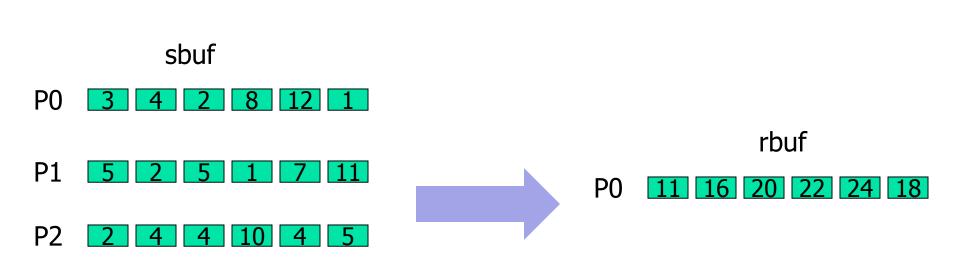
```
MPI_Reduce(x, r, 10, MPI_INT, MPI_MAX, 0 MPI_COMM_WORLD)
...

input array output array array size root

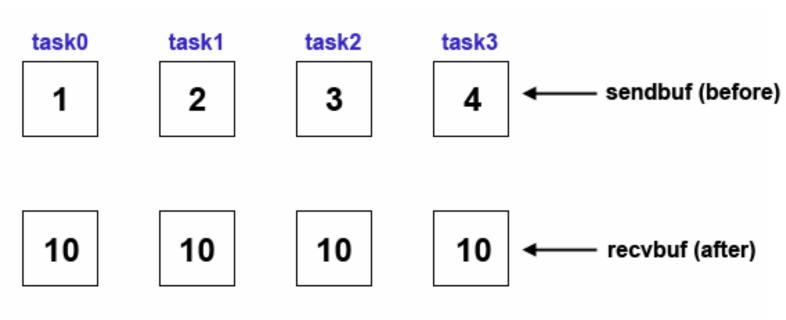
MPI_Allreduce(x, r, 10 MPI_INT, MPI_MAX, MPI_COMM_WORLD)
```

MPI_Reduce example

MPI_Reduce(sbuf,rbuf,6,MPI_INT,MPI_SUM,0,MPI_COMM_WORLD)



MPI_Allreduce example



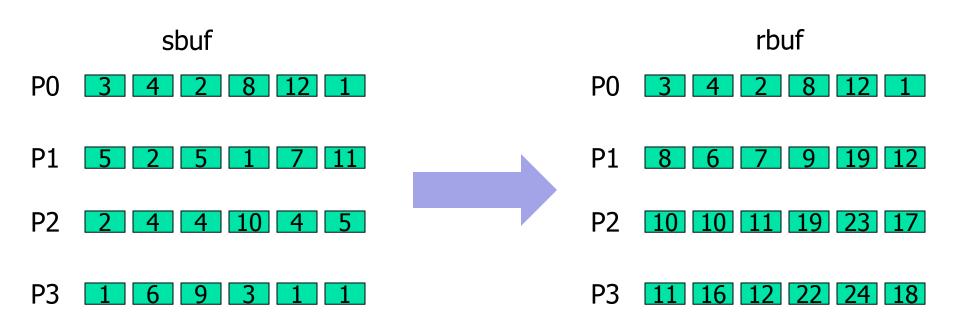
MPI_Reduce_scatter example

	task3	task2	task1	sk0	task0
←—— sendbuf (before)	1	1	1	1	1
	2	2	2	2	2
	3	3	3	3	3
	4	4	4	4	4
← recvbuf (after)	16	12	8	4	4



MPI_Scan: Prefix reduction

process i receives data reduced on process 0 to i.



MPI_Scan(sbuf,rbuf,6,MPI_INT,MPI_SUM,MPI_COMM_WORLD)

User-defined reduce operations

- pointer to a function with a specific prototype
- commute (0 or 1) allows for optimization if true



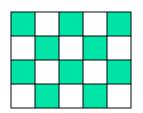
MPI_Op_create example

```
void myfunc(void *a, void *b, int *len, MPI Datatype *dtype) {
  int i;
  for (i = 0; i < *len; ++i)
    ((int*)b)[i] = ((int*)b)[i] + ((int*)a)[i];
int main(int argc, char *argv[]) {
  int myrank, nprocs, sendb, recvb;
 MPI Op myop;
 MPI Init(&argc, &argv);
 MPI Comm size (MPI COMM WORLD, &nprocs);
  MPI Comm rank (MPI COMM WORLD, &myrank);
  MPI Op create (myfunc, 1, &myop);
  sendbuf = 2*myrank+1;
                                               // odd numbers
  MPI Reduce (&sendb, &recvb, 1, MPI INT, myop, 0, MPI COMM WORLD);
  if (myrank == 0) printf("%d^2 = %d\n", nprocs, recvb);
  MPI Finalize();
  return 0;
```



More Advanced Messages

Regularly strided data



Blocks/Elements of a matrix

Data structure

```
struct {
    int a;
    double b;
}
```

A set of variables

```
int a; double b; int x[12];
```

Derived Data Types

- A data type is defined by a "type map"
 - set of <type, displacement> pairs
- Created at runtime in two phases
 - Construct the data type from existing types
 - Commit the data type before it can be used
- Simplest constructor: contiguous type

MPI_Type_contiguous example

```
int buffer[100];
MPI_Datatype chvec;
MPI_Type_contiguous(20, MPI_CHAR, &chvec);
MPI_Type_commit(&chvec);
...
MPI_Send(buffer,1,chvec,1,44,MPI_COMM_WORLD);
MPI_Type_free(&chvec);
```

MPI_Type_indexed()

```
int MPI_Type_indexed(int count,
    int *array_of_blocklengths,
    int *array_of_displacements,
    MPI_Datatype oldtype,
    MPI_Datatype *newtype)
```

len[0]

Len[1]

MPI_Type_indexed example

```
int rank, int vec[3][2] = \{ 11, 12, 21, 22, 31, 32 \};
int ivec[3][2] = \{ 0 \};
int blocklen[2] = \{2, 1\}; int displ[2] = \{4, 0\};
MPI Datatype mytype;
MPI Status mystat;
MPI Init ....
MPI Type indexed(2, blocklen, displ, MPI INT, &mytype);
MPI Type commit(&mytype);
if (rank == 0)
  MPI Send(vec, 1, mytype, 1, 0, MPI COMM WORLD);
else {
  printm(vec);
  MPI Recv(ivec, 1, mytype, 0, 0, MPI COMM WORLD, &mystat);
  printm(ivec);
```

MPI_Type_indexed example cont.

```
void printm(int v[3][2]) {
  int i, j;
  for (i = 0; i < 3; i++) {
    printf("\n");
    for (j=0; j < 2; j++) printf("%02d ", v[i][j]);
  } printf("\n");
}
mpirun -np 2 indexed:
/* int blocklen[2] = {2, 1}; int displ[2] = {4, 0}; */
11 12
21 22
31 32
11 00
00 00
31 32
```

MPI_Type_struct()

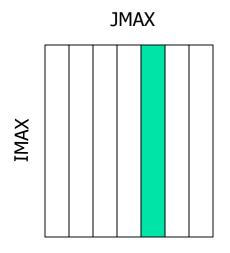
```
int MPI_Type_struct(int count,
   int *array_of_blocklengths,
   MPI_Aint *array_of_displacements,
   MPI_Datatype *array_of_types,
   MPI_Datatype *newtype)
```

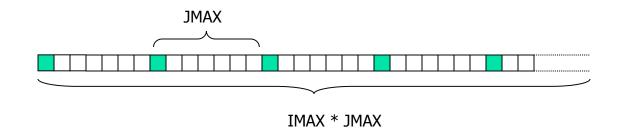
MPI_INT MPI_DOUBLE My_weird_type

MPI_Type_vector example

Sending the 5th column of a 2-D matrix:

```
double results[IMAX][JMAX];
MPI_Datatype newtype;
MPI_Type_vector (IMAX, 1, JMAX, MPI_DOUBLE, &newtype);
MPI_Type_Commit (&newtype);
MPI_Send(&(results[0][5]), 1, newtype, dest, tag, comm);
```





MPI_Group_...

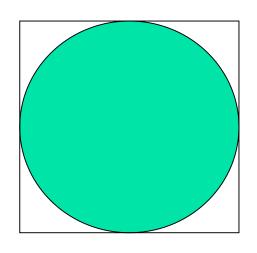
```
int new rank, ranks1[4]={0,1,2,3};
MPI Group orig_grp, new_grp;
MPI Comm new_comm;
MPI Comm group (MPI COMM WORLD, &orig grp);
MPI Group incl(orig grp, 4, ranks1, &new grp);
MPI Comm create (MPI COMM WORLD, new_grp, &new_comm);
MPI Group rank (new grp, &new rank);
```

PI example

$$A_c = \pi r^2$$

$$A_s = (2r)^2$$

$$\pi = 4 \frac{A_c}{A_s}$$



```
int isInside() {
   double x = ((double) rand()) / RAND_MAX;
   double y = ((double) rand()) / RAND_MAX;
   if ((x*x + y*y) <= 1.0) return(1);
   else return(0);
}</pre>
```

PI example cont.

MPI Init(&argc, &argv); // MPI Comm size(MPI COMM WORLD, &size); MPI Comm rank(MPI COMM WORLD, &rank); my trials = (trials + size - 1)/size; trials = my trials * size; srand(rank*17); for (i=0; i < my_trials; i++) inside += isInside();</pre> MPI Reduce (&inside, &sum, 1, MPI LONG, MPI SUM, 0, MPI COMM WORLD) if (rank == 0)printf("PI: %f (%ld)\n", 4.0*sum/trials);