CN [Day - 4]

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Question 1: A user on a web browser (Client) makes an HTTPS request to access a webpage from a remote server. The DNS server is queried to resolve the domain name. Once the IP is resolved, a secure connection is established between the client and server using RSA. The following sequence occurs:

- 1. Client sends a DNS query to resolve www.securedata.com.
- 2. DNS resolves the domain to 192.0.2.10.
- 3. The browser initiates an HTTPS GET request using HTTP over TLS.
- 4. During the handshake, the server sends its public key (e, n) = (7, 187).
- 5. The client generates a session key k = 45 and encrypts it using the RSA public key.
- 6. The encrypted session key is sent to the server.
- 7. Communication proceeds using this session key under symmetric encryption.

Questions

- (a) Identify and explain the role of each Application Layer protocol used in this scenario.
- (b)Calculate the encrypted session key using the RSA public key (7, 187).
- (c) Explain why DNS is critical to the Application Layer and how it supports scalability.
- (d) Describe the purpose of the HTTPS handshake and how it ensures secure communication.
- (e) Suppose an attacker intercepts the DNS response and alters the IP address. What type of attack is this, and how can it be mitigated?

Answer:

A) Identify and explain the role of each Application Layer protocol used in this scenario

- 1. DNS (Domain Name System):
 - o Resolves the human-readable domain (www.securedata.com) to an IP address (192.0.2.10).
 - o Operates over UDP (port 53) for queries and is essential for hostname resolution.
- 2. HTTP (HyperText Transfer Protocol):
 - o Application-layer protocol used by the browser to request web content.
 - o In this case, HTTP is wrapped inside TLS, making it HTTPS.
- 3. HTTPS (HTTP Secure):
 - HTTP + TLS/SSL ensures data confidentiality and integrity.
 - Establishes a secure channel using public key cryptography (RSA) and symmetric encryption for the session.

B) Calculate the encrypted session key using the RSA public key (e, n) = (7, 187)

RSA Encryption formula:

cipher = $k^e \setminus mod n$

Given:

Session key k=45

- Public key (e,n)=(7,187)
- cipher=45⁷mod 187

Let's compute step by step:

- $45^2 = 2025$
- 45³=45×2025=91125
- $45^4 = 45 \times 91125 = 4100625$
- Instead, we use modular exponentiation for efficiency:

Fast Modular Exponentiation:

We need to compute:

45⁷ mod 187

Using:

 $45^7 = ((45^2)^3) \times 45$

Step-by-step:

- 45¹mod 187=45
- $45^2 = 2025 \Rightarrow 2025 \mod 187 = 2025 (187 \times 10) = 2025 1870 = 155$
- 45⁴ = (452)2=1552=24025 ⇒ 24025 mod 187=24025 (187×128)=24025 23936=89
- 45⁷=45⁴×45²×45 mod 187=89×155×45 mod 187

Now calculate:

- 89×155=13795
- 13795×45=620775
- 620775 mod 187=620775-(3320×187=620840)=-65+187=122

Encrypted session key = 122

C) Why is DNS critical to the Application Layer, and how does it support scalability?

Role of DNS:

- Translates **human-readable domain names** into **IP addresses** required by the transport and network layers.
- Acts as the "phonebook" of the internet.

Scalability:

- DNS is hierarchical (root → TLD → authoritative DNS).
- Uses caching, replication, and distributed servers to handle millions of queries efficiently.
- Enables the **internet to grow** without being bottlenecked by a central server.

D) Purpose of the HTTPS handshake and how it ensures secure communication

Purpose:

- The TLS/SSL handshake ensures that the client and server:
 - Authenticate each other (server always, client optionally)
 - Establish a shared secret key for symmetric encryption
 - Prevent eavesdropping, tampering, and forgery

Process Overview:

- 1. Server sends its digital certificate containing its RSA public key.
- 2. Client verifies the certificate, then **encrypts a session key** using the public key.
- 3. Only the server can decrypt it using its **private key**.
- 4. All further communication is **encrypted symmetrically** using this shared key.

Ensures confidentiality, integrity, and authentication.

E) If an attacker intercepts the DNS response and alters the IP address...

This is known as a **DNS Spoofing (DNS Cache Poisoning)** attack.

- The attacker sends a **fake IP address** for the domain name.
- The user unknowingly connects to a malicious server instead of the legitimate one.

Type of attack:

- Man-in-the-Middle (MITM)
- Specifically: DNS Spoofing

Mitigation Strategies:

- 1. DNSSEC (Domain Name System Security Extensions):
 - Adds digital signatures to DNS responses to verify authenticity.
- 2. HTTPS with valid certificates:
 - Even if DNS is spoofed, the browser warns if the certificate doesn't match the domain.
- 3. Encrypted DNS (DoH/DoT):
 - o DNS over HTTPS or TLS to prevent tampering by intermediaries.
- 4. Browser-based protection:
 - o Browsers validate certificates using **CA chains** and show warnings.

Question 2: Problem Statement

Consider the following scenario:

A university student is working remotely and performs these tasks:

- 1. Logs into the university server remotely using TELNET to submit an assignment.
- 2. Opens a browser and visits www.university-portal.edu using HTTP.
- 3. The DNS is gueried to resolve the domain name to IP address 203.0.113.8.
- 4. The server responds with a webpage that includes a form submission via POST.
- 5. The server uses RSA encryption for login authentication. The public key is (e = 5, n = 221).
- 6. The user sends an email via SMTP, and then downloads emails via POP3. Questions
- (a) Describe the client-server interaction in each of the above tasks.
- (b) Explain how DNS functions in this scenario and what would happen if DNS fails.
- (c) What port numbers are used by TELNET, HTTP, SMTP, and POP3?
- (d) The password "42" is encrypted using RSA public key (5, 221). What is the ciphertext?
- (e) Explain the security limitations of using TELNET and how SSH addresses them.
- (f) Compare HTTP and E-mail protocols (SMTP, POP3) in terms of their communication pattern.

Answer:

A) Client-Server Interaction in Each Task:

- 1. TELNET Login (Remote Access):
 - The student (client) initiates a TELNET session to log into the university server.
 - TELNET allows terminal-based remote login; the server responds with a login prompt and executes commands sent by the client.
- 2. Accessing Website via HTTP:
 - o The browser (client) sends an HTTP GET request to fetch the webpage.
 - o The web server responds with the requested HTML page.
- 3. DNS Query:
 - The client's machine sends a **DNS query** to a DNS server to resolve www.university-portal.edu to IP 203.0.113.8.
 - o This enables the browser to connect to the correct server.
- 4. Form Submission (HTTP POST):
 - The student fills a form (e.g., assignment upload).
 - o The browser sends an HTTP POST request to the server with form data.
- 5. RSA Login Authentication:
 - \circ The server sends its public RSA key (e = 5, n = 221).

- The client encrypts the password and sends it securely. Only the server (holding the private key) can decrypt it.
- 6. Sending and Receiving Email:
 - o **SMTP** is used to **send** email from client to mail server.
 - o POP3 is used to download or retrieve emails from the mail server to the client.

B) How DNS Functions and What if It Fails:

How DNS works:

- Translates domain names into IP addresses.
- Acts like the internet's phonebook.
- The client sends a query; DNS server responds with the IP.

If DNS fails:

- The client cannot reach the server using the domain name.
- Browser shows "Server not found" or "DNS lookup failed."
- The user would have to manually enter the IP address (if known), which is not practical.

C) Port Numbers Used:

Protocol	Port Number	Description
TELNET	23	Unencrypted terminal access
НТТР	80	Web traffic (unencrypted)
SMTP	25	Sending emails
POP3	110	Retrieving emails

Encrypt password "42" using RSA (e = 5, n = 221)

D) RSA Encryption Formula:

cipher=me mod n

Given:

- Message/password m = 42
- Public key (e = 5, n = 221)

cipher=42⁵mod 221

Step-by-step (modular exponentiation):

RSA Encryption of Plaintext = 42 using Public Key (e = 5, n = 221)

We want to calculate:

Ciphertext = (42^5) mod 221

Step 1: Calculate 42² mod 221

 $42^2 = 1764$

 $1764 \mod 221 = 1764 - (221 \times 7) = 1764 - 1547 = 217$

Step 2: Calculate 42⁴ mod 221

 $42^4 = (42^2)^2 = 217^2 = 47089$

47089 mod 221 = 47089 - (221 × 213) = 47089 - 47073 = 16

Step 3: Calculate 42⁵ mod 221

 $42^5 = 42 \times 42^4 = 42 \times 16 = 672$

 $672 \mod 221 = 672 - (221 \times 3) = 672 - 663 = 9$

Final Answer:	
Encrypted Ciphertext = 9	

(e) Security Limitations of TELNET & How SSH Solves Them

TELNET Limitations:

- Transmits data in plain text including usernames and passwords.
- Vulnerable to eavesdropping, packet sniffing, and MITM attacks.
- No encryption, no authentication of remote server identity.

SSH (Secure Shell) Advantages:

- Encrypts all communication, including credentials.
- Uses **public-key cryptography** to verify the server.
- Provides confidentiality, integrity, and authentication.

SSH is the **secure replacement** for TELNET.

F) Compare HTTP and Email Protocols (SMTP, POP3)

Feature	НТТР	Email (SMTP, POP3)
Direction	Client → Server	SMTP: Client → ServerPOP3: Server → Client
Туре	Pull-based (GET, POST)	SMTP: Push , POP3: Pull
Use Case	Browsing web pages	Sending/receiving emails
Connection	Stateless	Session-based (especially POP3)
Port	80	SMTP: 25, POP3: 110
Protocol Pattern	Request/Response	SMTP: Send → Queue → DeliverPOP3: Fetch