

# **Advanced Data Structures and Algorithms**

*Comprehensive Assignment Solutions*

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## 1 Complexity Classes

**Question 1.** Define the class Co-NP. Explain the type of problems that belong to this complexity class.

**Detailed Solution:**

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### Definition

**Co-NP** is a complexity class containing decision problems for which the "No" instances can be efficiently verified. Formally, a language  $L$  is in Co-NP if its complement,  $\bar{L}$  (the set of strings not in  $L$ ), is in the class **NP**.

In simpler terms:

- **NP Problem:** "If the answer is YES, there exists a short proof (certificate) that I can check in polynomial time."
- **Co-NP Problem:** "If the answer is NO, there exists a short proof (counter-example) that I can check in polynomial time."

### Visual Relationship of Complexity Classes

The relationship between P, NP, and Co-NP is often visualized using a Venn diagram. P is contained in the intersection of NP and Co-NP.

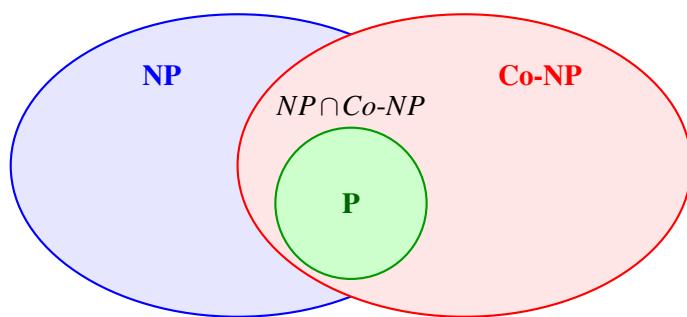


Diagram assumes  $P \neq NP$  and  $NP \neq Co-NP$

### Type of Problems

Co-NP problems typically involve properties that must hold for **all** possible structures or assignments (Universality).

**1. TAUTOLOGY (Logic):**

- *Problem:* Given a boolean formula  $\phi$ , is it true for **every** possible assignment of truth values to variables?
- *Why Co-NP:* If the answer is NO, there exists at least one assignment where  $\phi$  is False. This single assignment serves as a succinct "counter-example" (certificate) that can be verified easily.

**2. UNSAT (Logic):**

- *Problem:* Given a boolean formula  $\phi$ , is it impossible to satisfy? (i.e., Is it false for all assignments?)
- *Why Co-NP:* This is the direct complement of SAT. If the answer is NO (meaning it IS satisfiable), we can present a satisfying assignment as proof.

**3. GRAPH NON-ISOMORPHISM:**

- *Problem:* Given two graphs  $G_1$  and  $G_2$ , are they **not** isomorphic?
- *Why Co-NP:* The complement problem is Graph Isomorphism (in NP). If the answer to Non-Isomorphism is NO (meaning they ARE isomorphic), the certificate is simply the mapping (bijection) between the vertices.

**4. NON-HAMILTONIAN GRAPH:**

- *Problem:* Given a graph  $G$ , is it true that there is **no** Hamiltonian cycle?
- *Why Co-NP:* If the answer is NO (meaning there IS a cycle), the certificate is simply the sequence of vertices forming the cycle.

**Conclusion**

The class Co-NP complements NP. While NP characterizes problems where "Yes" answers are easy to verify (existence proofs), Co-NP characterizes problems where "No" answers are easy to verify (universal truths). Most computer scientists believe  $NP \neq Co-NP$ , meaning that just because we can verify a solution quickly doesn't mean we can verify the \*absence\* of a solution quickly.

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