

# Extended Transition Function

## ✓ Extended transition function

An extended transition function  $\hat{\delta}$  traces the path of an automaton and determines the final state when an initial state  $q$  and an input string  $x$  are passed through it.

The difference between a simple transition function and the extended transition function is that the former performs a transition of a single character/instance. In contrast, the latter performs the transitions on a complete string.

**A recursive algorithm is used to reach the final state, which is as follows:**

**Base condition:**

$$\hat{\delta}(q, \epsilon) \rightarrow q$$

**Recursion rule:**

$$\hat{\delta}(q, xa) \rightarrow \delta(\hat{\delta}(q, x), a)$$

Here,  $x \in \Sigma^*$  and  $a \in \Sigma$ . Also,  $x$  is a string of characters belonging to the set of the input symbols and  $aa$  is a single character.

The input string is reduced from the right side, character by character, until the base condition—when all the strings are reduced and we are left with a null character (epsilon)—is reached. Then simple transitions are applied to the broken-down string.

A transition table is formed to show each transition in the DFA. If the output is a final state, the given string will be accepted by the DFA.

## Outlines:

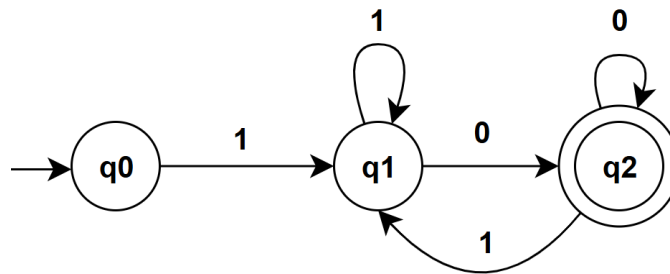
Extended transition Function for DFA: example 1,2

Extended transition Function for NFA: example 1,2

# Extended Transition Function

## Extended Transition function for DFA:

**Example - 1: Construct a DFA to accept the string that start with 1 and end with 0, over  $\Sigma = \{0,1\}$  and check the string 1010 is accepted by DFA using extended transition function.**



## Transition function for DFA:

$\delta$	0	1
$\rightarrow q0$	-	q1
q1	q2	q1
*q2	q2	q1

## Extended Transition function for DFA:

$$\hat{\delta}(q, \epsilon) = q0$$

$$\hat{\delta}(q0, 1) = \delta(\hat{\delta}(q0, \epsilon), 1) = \delta(q0, 1) = q1$$

$$\hat{\delta}(q0, 10) = \delta(\hat{\delta}(q0, 1), 0) = \delta(q1, 0) = q2$$

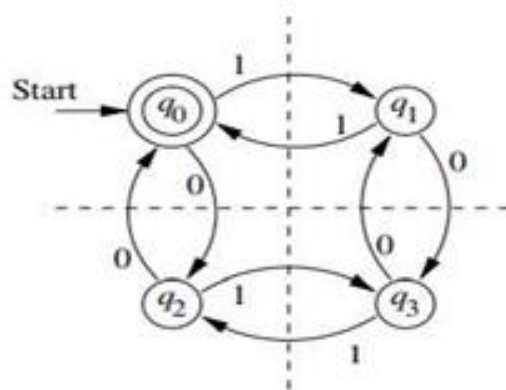
$$\hat{\delta}(q0, 101) = \delta(\hat{\delta}(q0, 10), 1) = \delta(q2, 1) = q1$$

$$\hat{\delta}(q0, 1010) = \delta(\hat{\delta}(q0, 101), 0) = \delta(q1, 0) = q2$$

Since q2 is in Final state, 1010 is accepted by DFA

## Extended Transition Function

**Example - 2: check the string 110101 is accepted by DFA or not, using extended transition function.**



Transition function for DFA:

$\delta$	0	1
$\rightarrow^*q_0$	q2	q1
q1	q3	q0
q2	q0	q3
q3	q1	q2

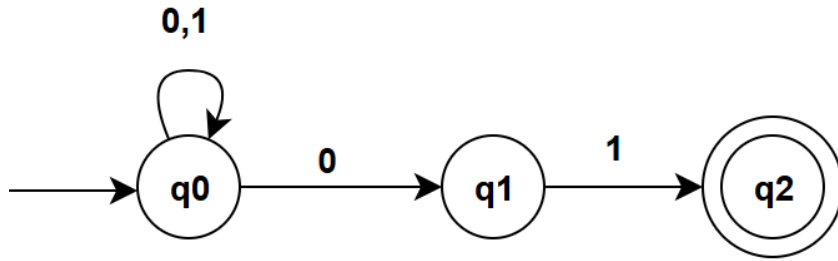
The check involves computing  $\hat{\delta}(q_0, w)$  for each prefix  $w$  of 110101, starting at  $\epsilon$  and going in increasing size. The summary of this calculation is:

- $\hat{\delta}(q_0, \epsilon) = q_0$ .
- $\hat{\delta}(q_0, 1) = \delta(\hat{\delta}(q_0, \epsilon), 1) = \delta(q_0, 1) = q_1$ .
- $\hat{\delta}(q_0, 11) = \delta(\hat{\delta}(q_0, 1), 1) = \delta(q_1, 1) = q_0$ .
- $\hat{\delta}(q_0, 110) = \delta(\hat{\delta}(q_0, 11), 0) = \delta(q_0, 0) = q_2$ .
- $\hat{\delta}(q_0, 1101) = \delta(\hat{\delta}(q_0, 110), 1) = \delta(q_2, 1) = q_3$ .
- $\hat{\delta}(q_0, 11010) = \delta(\hat{\delta}(q_0, 1101), 0) = \delta(q_3, 0) = q_1$ .
- $\hat{\delta}(q_0, 110101) = \delta(\hat{\delta}(q_0, 11010), 1) = \delta(q_1, 1) = q_0$ .

## Extended Transition Function

### Extended Transition function for NFA:

**Example-1: check the NFA is accepted or not for the input 00101.**



Transition function for NFA:

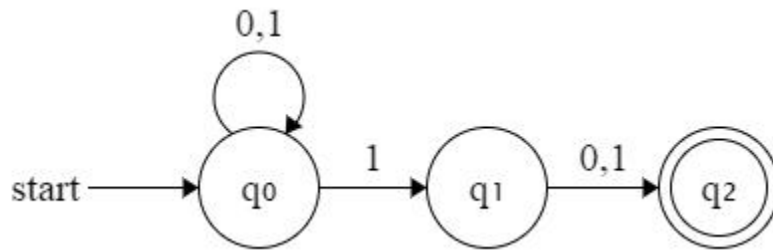
$\delta$	0	1
$\rightarrow q0$	$\{q0, q1\}$	$q0$
$q1$	-	$q2$
$*q2$	-	-

1.  $\hat{\delta}(q_0, \epsilon) = \{q_0\}$ .
2.  $\hat{\delta}(q_0, 0) = \delta(q_0, 0) = \{q_0, q_1\}$ .
3.  $\hat{\delta}(q_0, 00) = \delta(q_0, 0) \cup \delta(q_1, 0) = \{q_0, q_1\} \cup \emptyset = \{q_0, q_1\}$ .
4.  $\hat{\delta}(q_0, 001) = \delta(q_0, 1) \cup \delta(q_1, 1) = \{q_0\} \cup \{q_2\} = \{q_0, q_2\}$ .
5.  $\hat{\delta}(q_0, 0010) = \delta(q_0, 0) \cup \delta(q_2, 0) = \{q_0, q_1\} \cup \emptyset = \{q_0, q_1\}$ .
6.  $\hat{\delta}(q_0, 00101) = \delta(q_0, 1) \cup \delta(q_1, 1) = \{q_0\} \cup \{q_2\} = \{q_0, q_2\}$ .

Line (1) is the basis rule. We obtain line (2) by applying  $\delta$  to the lone state,  $q_0$ , that is in the previous set, and get  $\{q_0, q_1\}$  as a result. Line (3) is obtained by taking the union over the two states in the previous set of what we get when we apply  $\delta$  to them with input 0. That is,  $\delta(q_0, 0) = \{q_0, q_1\}$ , while  $\delta(q_1, 0) = \emptyset$ . For line (4), we take the union of  $\delta(q_0, 1) = \{q_0\}$  and  $\delta(q_1, 1) = \{q_2\}$ . Lines (5) and (6) are similar to lines (3) and (4).  $\square$

## Extended Transition Function

**Example-2: Check  $w_1 = 001$  and  $w_2 = 01010$  is accepted by the NFA using extended transition function or not.**



Transition Table:

$\delta$	0	1
$\rightarrow q_0$	$\{q_0\}$	$\{q_0, q_1\}$
$q_1$	$\{q_2\}$	$\{q_2\}$
* $q_2$	$\emptyset$	$\emptyset$

Extended transition Function Input Processing:

For input,  $w = 011$

$$\hat{\delta}(q_0, \epsilon) = \{q_0\}$$

$$\hat{\delta}(q_0, 0) = \{q_0\}$$

$$\begin{aligned} \hat{\delta}(q_0, 01) &= \delta(\hat{\delta}(q_0, 0), 1) \\ &= \delta(q_0, 1) \\ &= \{q_0, q_1\} \end{aligned}$$

$$\begin{aligned} \hat{\delta}(q_0, 011) &= \delta(\hat{\delta}(q_0, 01), 1) \\ &= \delta(\{q_0, q_1\}, 1) \\ &= \delta(q_0, 1) \cup \delta(q_1, 1) \\ &= \{q_0, q_1\} \cup \{q_2\} \\ &= \{q_0, q_1, q_2\} \end{aligned}$$

For input,  $w = 01010$ . Since sub-string 01 is already calculated I will reuse it.

$$\begin{aligned} \hat{\delta}(q_0, 010) &= \delta(\hat{\delta}(q_0, 01), 0) \\ &= \delta(\{q_0, q_1\}, 0) \\ &= \delta(q_0, 0) \cup \delta(q_1, 0) \\ &= \{q_0\} \cup \{q_2\} \\ &= \{q_0, q_2\} \end{aligned}$$

$$\begin{aligned} \hat{\delta}(q_0, 0101) &= \delta(\hat{\delta}(q_0, 010), 1) \\ &= \delta(\{q_0, q_2\}, 1) \\ &= \delta(q_0, 1) \cup \delta(q_2, 1) \\ &= \{q_0, q_1\} \cup \emptyset \\ &= \{q_0, q_1\} \end{aligned}$$

$$\begin{aligned} \hat{\delta}(q_0, 01010) &= \delta(\hat{\delta}(q_0, 0101), 0) \\ &= \delta(\{q_0, q_1\}, 0) \\ &= \delta(q_0, 0) \cup \delta(q_1, 0) \\ &= \{q_0\} \cup \{q_2\} \\ &= \{q_0, q_2\} \end{aligned}$$