

Parameters: Our goal is to generate k hidden-bits. $N = \Theta(k^\delta), \delta \in (0, 1)$. The exact weight w is specifically chosen to ensure a proper minimum distance in exact-LPN. Denote Gaussian noise distribution with \mathcal{B}_μ^N and exact-weight noise distribution with X_w^N .

- **Setup**(1^k):

1. $\forall i \in [k], A_i \xleftarrow{\$} \{0, 1\}^{N \times N}, s_i \xleftarrow{\$} \{0, 1\}^N, e_i \leftarrow \mathcal{B}_\mu^N$.
2. Decide w , weight parameter for exact-LPN.
3. Compute $b_i := A_i \cdot s_i + e_i$.
4. Sample $\alpha \xleftarrow{\$} \{0, 1\}^{N \times N}$ (for hiding seed)
5. $\text{crs} := \{(A_i, b_i) \mid i \in [k]\}, w, \alpha, \text{td} := \{s_i \mid i \in [k]\}$.

- **Genbits**($1^k, \text{crs}$):

1. $\text{seed} \xleftarrow{\$} \{0, 1\}^N, \epsilon \leftarrow \mathcal{B}_\mu^N$, hide seed as $\beta := \alpha \cdot \text{seed} + \epsilon$.
2. $\forall i \in [k], e'_i \leftarrow \{0, 1\}^N$, compute $b'_i := A_i \cdot \text{seed} + e'_i$.
3. Compute hidden-bits $r_i := \text{hc}(b_i; \text{seed}), \forall i \in [k]$.
4. Sample $x \xleftarrow{\$} \{0, 1\}^{N-1}, \forall i \in [k], \eta_i \leftarrow X_w^N$ s.t. $|\eta_i| = w$.
5. Compute $B_i := A_i \cdot (x \| r_i) + \eta_i$.
6. $\text{com} := \{x, \beta\}, \pi_i := \{b'_i, B_i\}, \forall i \in [k]$.

- **Verify**($1^k, \text{crs}, \text{com}, i, \pi_i, r, \text{td}_i$):

1. Check $r = \text{hc}(\pi_{i,1}; \text{td}_i)$ i.e. $r = \text{hc}(b'_i; s_i)$.
2. Check $|\pi_{i,2} \oplus A_i \cdot (\text{com}_1 \| r)| = w$ i.e. $|B_i \oplus A_i \cdot (x \| r)| = w$
3. **But how do we validate $\text{com}_2 = \beta$? Or can we prove cheating β will not hurt binding?**
4. Accept if and only if all the above hold.

- **Open**($1^k, \text{crs}, \text{com}$): Inefficiently solve LPN sample β to get seed , then compute all hidden bits from $r_i := \text{hc}(b_i; \text{seed}), \forall i \in [k]$.

There is an error probability because of hardcore computation, which can be reduced to negligible by limiting noise rate.