

Selection Sort

Implementation:

```

void selection_sort (int arr[], int n)
{
    for (int i = 0; i < (n-1); i++)
    {
        for (int j = i+1; j < n; j++)
        {
            if (arr[i] > arr[j])
                swap (arr[i], arr[j]);
        }
    }
    return;
}

```

Analysis:

	Cost	Time
for (int i = 0; i < (n-1); i++)	c_1	$(n-1)$
{		
for (int j = i+1; j < n; j++)	c_2	$\frac{n(n+1)}{2}$
{		
if (arr[i] > arr[j])	c_3	$\frac{n(n+1)}{2}$
swap (arr[i], arr[j]);	c_4	$\frac{n(n+1)}{2}$
}		
}		

∴ Time function $f(t) = c_1(n-1) + c_2 \frac{n(n+1)}{2} + c_3 \frac{n(n+1)}{2}$

Best case occurs when the array is already sorted

1	2	3	4	5
---	---	---	---	---

value of c_4 will be 0, so time function will be

$$f(t) = \left(\frac{c_2 + c_3}{2} \right) * n^2 + \left(\frac{c_2 + c_3 + 2c_1}{2} \right) * n - c_1$$

This function look like $f(t) = an^2 + bn + c$, which is quadratic function

∴ worst case will be occurred when the array is reversely sorted

5	4	3	2	1
---	---	---	---	---

$$f(t) = \left(\frac{c_1 + c_2 + c_3}{2} \right) * n^2 + \left(\frac{c_1 + c_2 + c_3 + 2c_1}{2} \right) * n - c_1$$

∴ look like $f(t) = (an^2 + bn + c)$; which is quadratic function.

Time complexity:

$$f(t) = an^2 + bn + c \quad [\text{Best case}]$$

$$O(n^2)$$

Worst case: n^2

$$f(n) = an^2 + bn + c$$

$$-2(n^2) \dots$$

$$(i+1) : (i-1) > i : 0 = i \dots$$

$$(i+1) : i > i : i+1 = i \dots$$

$$(i+1) : i < i \dots$$

$$(i+1) : i < i \dots$$

1. n^2

2. $(n+1) : i < i \dots$

$$(i+1) : (i-1) > i : 0 = i \dots$$

$$(i+1) : i > i : i+1 = i \dots$$

Insertion sort

Implementation:

```

void insertion(int arr[], int n)
{
    int i, j, x;
    for(i=0; i<n; i++)
    {
        x = arr[i];
        j = i-1;
        while(j >= 0 && arr[j] > x)
            arr[j+1] = arr[j], j--;
        arr[j+1] = x;
    }
}

```

Analysis:

Cost

time (n)

```

void insertion(int arr[], int n)
{
    int i, j, x;
    for(i=0; i<n; i++)
    {
        x = arr[i];
        j = i-1;
        while(j >= 0 && arr[j] > x)
            arr[j+1] = arr[j], j--;
        arr[j+1] = x;
    }
}

```

C1

n

C2

n-1

C3

n-1

C4

 $\frac{n(n+1)}{2}$

C5

 $\frac{n(n+1)}{2}$


```

    arr[j+1] = x;
  }
}

```

c6

n-1

∴ Time function:

$$f(t) = c_1 n + c_2 (n-1) + c_3 (n-1) + c_4 \frac{n(n+1)}{2} + c_5 \frac{n(n+1)}{2} + c_6 (n-1)$$

Best case occurs when array is already sorted

1	2	3	5	7
---	---	---	---	---

when c_4 will execute for $(n-1)$ time and c_5 will execute for 0 time.

so function will be,

$$f(t) = c_1 n + c_2 (n-1) + c_3 (n-1) + c_4 (n-1) + c_5 \cdot 0 + c_6 (n-1)$$

$$= n(c_1 + c_2 + c_3 + c_4 + c_6) - (c_2 + c_3 + c_4 + c_6)$$

∴ This look like $y = an - b$, That's linear.
 worst case occurs when the array is reverse sorted,

7	5	3	2	1
---	---	---	---	---

∴ Time function will be

$$\begin{aligned} f(t) &= c_1 n + c_2 (n-1) + c_3 (n-1) + c_4 * \frac{n(n+1)}{2} + c_5 * \frac{n(n+1)}{2} \\ &\quad + c_6 (n-1) \\ &= \frac{c_1 + c_2}{2} n^2 + \frac{2(c_1 + c_2 + c_3 + c_6)}{2} n - (c_2 + c_3 + c_6) \end{aligned}$$

∴ look like $y = an^2 + bn + c$; This is a quadratic eqⁿ.

Time complexity:

Best case:

We know, $y = an - b$ [∴ Best case]

So, $\Omega(n)$

Worst case:

We know, $y = an^2 + bn + c$ [∴ worst case]

So, $O(n^2)$