EECS 390 – Lecture 22

Template Metaprogramming

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Template Metaprogramming

- Uses templates to produce source code at compile time, which is then compiled with the rest of the program's code
- A form of compile-time specialization that takes advantage of the language's rules for template instantiation
- Most common in C++, though it is available in D and a handful of other languages
- Template metaprogramming is Turing complete, with computations expressed recursively

Motivation

- Can be used to compute values at compile time
 - The values can then be used where compile-time constants are required, such as the size of a non-dynamic array
 - However, newer versions of C++ enable compile-time value computation with constexpr rather than template metprogramming
- Must be used to manipulate types
 - Types are compile-time entities, so they can only be manipulated at compile time
 - Example: std::tuple

```
std::tuple<int, double> items = { 3, 4.1 };
int first = std::get<0>(items);
double second = std::get<1>(items);
```

std::get() uses template metaprogramming to determine types of tuple elements

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Template Specialization

Key to template metaprogramming

for specialization

- Allows a specialized definition for instantiating a template with specific arguments
- Example:

```
Generic definition
```

Specialization for int argument

```
template <class T>

struct is_int {
    static const bool value = false;
};

This specialization
    has no template
    struct is_int<int> {
        static const bool value = true;
};

Full argument list
    is_int<double>::value is false
```

is_int<int>::value is true

Reporting a Value

 We can report a value at compile time by arranging for it to be contained in an error message

Compiletime assertion

Pairs

We can represent a pair, whose items are arbitrary types, as:

```
Type aliases
```

```
template <class First, class Second>
struct pair {
    using car = First;
    using cdr = Second;
};
```

We can represent an empty list as:

```
struct nil {
};
```

Alias Templates

We can introduce alias templates to extract the first and second from a pair:

```
template <class Pair>
using car_t = typename Pair::car;

template <class Pair>
using cdr_t = typename Pair::cdr;
```

- The typename keyword is required when we have a nested type whose enclosing type depends on a template parameter
 - Otherwise, the compiler assumes we are referring to a value rather than a type

Type aliases

act as

"variables"

Empty Predicate

Template specialization to determine if a list is empty:

```
template <class List>
struct is_empty {
   static const bool value = false;
};

template <>
struct is_empty<nil> {
   static const bool value = true;
}
```

Compile-time constant can be used as argument to report

```
pair.cpp: In instantiation of 'struct
report<pair<char, pair<int, pair<double,
nil> > >, 0>':
pair.cpp:76:33: required from here
pair.cpp:67:3: error: static assertion
failed: report
pair.cpp: In instantiation of 'struct
report<nil, 1>':
pair.cpp:78:33: required from here
pair.cpp:67:3: error: static assertion
failed: report
```

Variable Templates

C++14 introduced variable templates, which are parameterized variables that hold a value:

```
template <class List>
const bool is_empty_v = is_empty<List>::value;
```

Then is_empty_v<nil> is true, but is_empty_v<pair<int, nil>> is false

```
pair.cpp: In instantiation of 'struct
report<pair<char, pair<int, pair<double,
nil> > >, 0>':
pair.cpp:76:33: required from here
pair.cpp:67:3: error: static assertion
failed: report
pair.cpp: In instantiation of 'struct
report<nil, 1>':
pair.cpp:78:33: required from here
pair.cpp:67:3: error: static assertion
failed: report
```

Pair Length

We can use a recursive template to compute the length of a list:

```
template <class List>
               struct length {
                 static const int value =
                   length<cdr t<List>>::value + 1;
               };
               template <>
   Base
              struct length<nil> {
   case
                 static const int value = 0;
Variable
             template <class List>
template
               const int length v = length<List>::value;
       report<x, length v<x> d;
     pair.cpp: In instantiation of 'struct report<pair<char,</pre>
     pair<int, pair<double, nil> > >, 3>':
     pair.cpp:79:31: required from here
     pair.cpp:67:3: error: static assertion failed: report
```

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Reverse

Reverse defined "tail recursively" as follows:

```
Remaining list
```

Base

case

```
Reversed
template <class List, class SoFar>
struct reverse helper {
                                          so far
  using type =
    typename reverse helper<cdr t<List>,
      pair<car t<List>, SoFar>>::type;
};
template <class SoFar>
struct reverse_helper<nil, SoFar> {
 using type = SoFar;
};
                                       Seed initial
template <class List>
                                         values
using reverse t =
  typename reverse helper<List, nil>::type;
```

Partial Class Template Specialization

 A class template may be partially specialized, accepting a subset of the template parameters

```
template <class SoFar>
struct reverse_helper<nil, SoFar> {
  using type = SoFar;
};
```

using x = pair<char, pair<int, pair<double, nil>>>;
report<reverse_t<x>, 0> e;

```
pair.cpp: In instantiation of 'struct report<pair<double,
pair<int, pair<char, nil> > >, 0>':
pair.cpp:80:32: required from here
pair.cpp:67:3: error: static assertion failed: report
```

Any instantiation

where the first

argument is nil

will use this

Numerical Computations

- We can use C++'s support for integer template arguments to perform numerical computations
- New version of report template:

```
template <long long N> struct report {
   static_assert(N > 0 && N < 0, "report");
};</pre>
```

Ensure that assertion will fail after instantiation, not before

Base

case

Factorial

Recursive computation of factorial:

```
template <int N> struct factorial {
   static const long long value =
     N * factorial<N - 1>::value;
};

template <>
   struct factorial<0> {
   static const long long value = 1;
};

report<factorial<5>::value> a;
```

factorial.cpp: In instantiation of 'struct report<12011>':
factorial.cpp:51:34: required from here
factorial.cpp:47:3: error: static assertion failed: report
 static_assert(N > 0 && N < 0, "report");
 ^</pre>

Command-Line Macros

 We can use a macro to make our computation generic, and then specify the value at the command line

```
#ifndef NUM
#define NUM 5
#endif
```

```
report<factorial<NUM>::value> a;
```

Define a macro from command line

Preventing Negative Input

- Negative input causes unbounded recursion
- We can prevent it as follows:

Helper template does computation

```
template <int N>
struct factorial_helper {
  static const long long value =
    N * factorial helper<N - 1>::value;
};
template <>
struct factorial helper<0> {
  static const long long value = 1;
};
template <int N> struct factorial {
  static assert(N >= 0,
                "argument must be non-negative");
  static const long long value =
    factorial helper<N >= 0 ? N : 0>::value;
```

Prevent instantiation of helper with negative value

Alternative: Default Argument

 Alternatively, we can use a second default argument that prevents unbounded recursion when the first argument is negative:

factorial<5> translates to factorial<5,</pre> true>, which uses the generic version

```
template <int N, bool /*Positive*/ = (N > 0)>
struct factorial {
  static const long long value =
    N * factorial<N - 1>::value;
};
template <int N>
struct factorial<N, false> {
  static const long\long value = 1;
};
```

No name necessary here, since we don't use the parameter for anything else

factorial<0> translates to factorial<0, false>, which matches the specialization

Two base

cases

Fibonacci Numbers

We can compute Fibonacci numbers as follows:

```
template <int N> struct fib {
  static const long long value =
    fib<N - 1>::value + fib<N - 2>::value;
};
template <>
struct fib<1> {
  static const long long value = 1;
template <>
struct fib<0> {
  static const long long value = 0;
};
```

Computation is efficient, since compiler only instantiates a set of arguments once¹

¹This is akin to **memoization** in functional programming.