

EECS 390 – Lecture 2

Basic Elements

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Digression:

Value and Reference Semantics

- **Value semantics:** a variable is nothing more than a name associated with an object
 - The storage for the variable is the same as that of the object itself
 - The association between a variable and an object cannot be broken as long as the variable is in scope
- **Reference semantics:** a variable is an indirect reference to an object
 - A variable has its own storage that is distinct from that of the object it refers to
 - A variable can be modified to be associated with a different object

Digression: Reference Semantics

- In a language with reference semantics, variables behave like C/C++ pointers
 - But can't do arithmetic on them

- Example:

```
>>> x = []  
>>> y = x  
>>> print(id(x))  
4546751752  
>>> print(id(y))  
4546751752
```

Get
unique ID
of object

C++ equivalent:

```
list *x = new list();  
list *y = x;  
cout << x << endl;  
  
cout << y << endl;
```

- Python: everything has reference semantics
- Java: primitives have value semantics, everything else has reference semantics

Programming Paradigms

- Languages can be classified in many ways
- A fundamental classification is by what programming paradigms they support
 - Imperative programming
 - Declarative programming
 - Functional programming
 - Logic programming
 - Object-oriented programming

Imperative Programming

- Program decomposed into explicit computational steps in the form of **statements**
 - A statement executes some operation, generally changing the state of the program
 - Statements (appear to) execute in a well-defined sequence
- Primary paradigm in most commonly used languages (C, C++, Java, Python, etc.)

Declarative Programming

- Expresses computation in terms of what it should compute rather than how
- **Functional programming:** models computation after mathematical functions
 - Generally avoids mutation
 - Primary paradigm in the Lisp family (including Scheme), Haskell, ML
 - Some support in C++, Java, Python
- **Logic programming:** expresses a program in the form of facts and rules
 - Primary paradigm in Prolog, SQL, Make

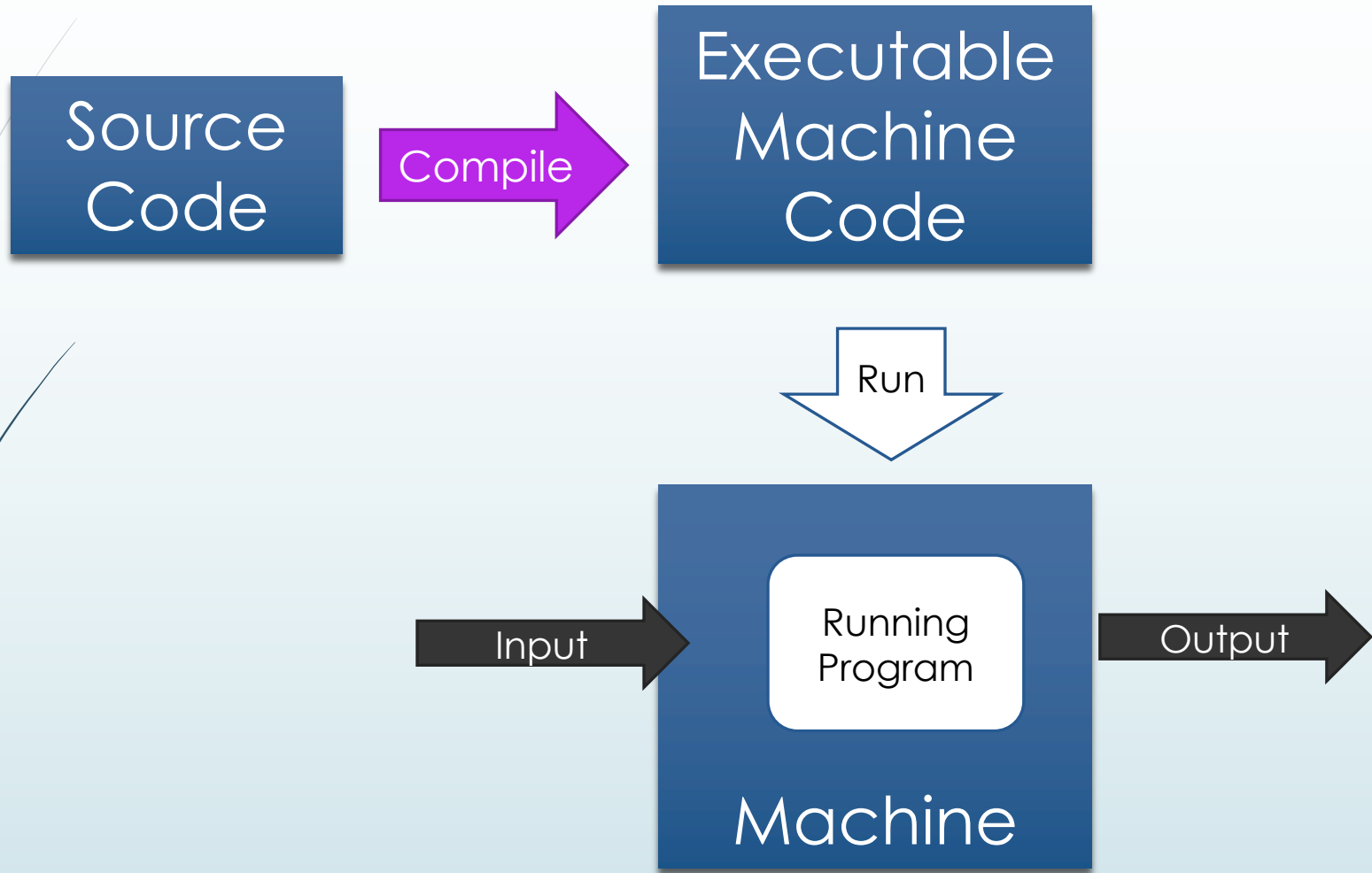
Type Systems

- All data are represented as bits
- Types determine:
 - What data means
 - What operations are valid on that data
 - How to perform those operations
- **Static typing** infers types directly from the source code and checks their use at compile time
- **Dynamic typing** tracks and checks types at runtime

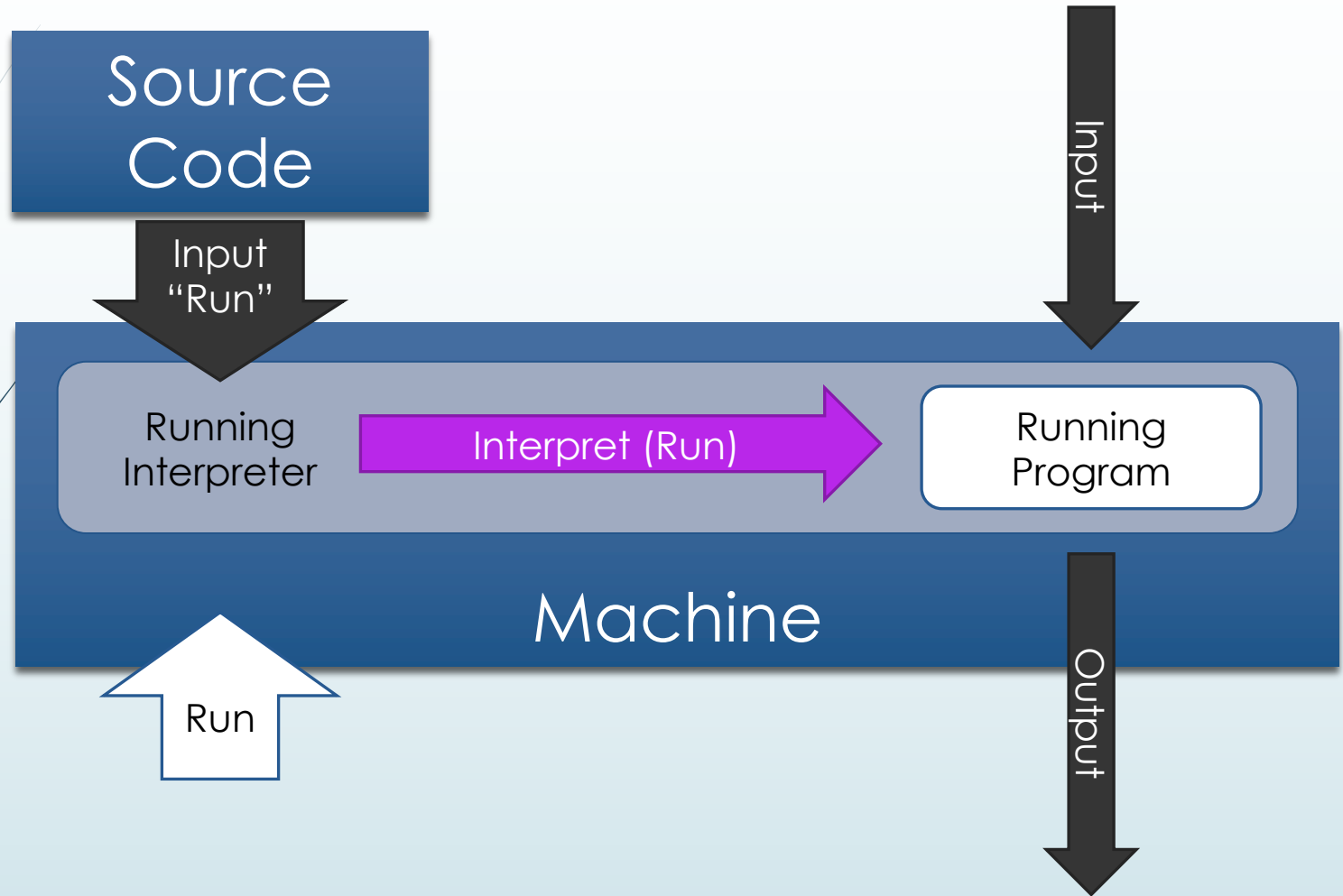
Compilation and Interpretation

- Programs can be compiled, interpreted, or some combination of the two
- Compilation: program translated to a form more suitable for execution on a machine
 - Target is often, but not necessarily, machine code
- Interpretation: program is input to interpreter, which interprets and performs the computation it specifies
 - Generally, code is directly interpreted rather than first translated into a different form

Compilation



Interpretation



Compiled vs. Interpreted

Compiled

- Faster
 - No execution engine
- Less portable
 - Must compile for each machine
- Less flexible
 - Program is fixed at compile-time

Interpreted

- Slower
 - Must go through engine
- More portable
 - As long as each machine has an interpreter
- More flexible
 - Program can change at runtime

Hybrids also exist!

Review: Abstraction

- **Abstraction** is the idea of using something for what it does without needing to know the details of how it does it
- Primary tool for managing complexity
 - Facilitates separation of concerns
 - Results in better modularity, maintainability
- However, there can be performance tradeoffs
 - Higher-level abstractions generally do not provide control over how they are implemented

Levels of Description

- **Grammar:** what phrases are correct
 - **Lexical structure:** what sequences of symbols represent correct words
 - **Syntax:** what sequences of words represent correct phrases
- **Semantics:** what does a correct phrase mean
- **Pragmatics:** how do we use a meaningful phrase
- **Implementation:** how are the actions specified by a meaningful phrase accomplished

Lexical Structure

- A **character set** is the alphabet of a language
 - e.g. ASCII, Unicode, or subsets thereof
- **Tokens** are the “words” in a programming language
 - Smallest element that is meaningful to the compiler or interpreter
 - Lexical analysis is often the first step in interpreting or compiling a program
- A token ends at a character that is invalid for the token, including whitespace
- Types of tokens
 - Literals
 - Identifiers
 - Keywords
 - Operators
 - Separators

Literals and Keywords

- **Literals** represent a particular value directly in source code
 - Examples: 3, 1.4, "hello world"
- Each primitive type often has its own set of literals
- **Keywords** are words that have special meaning in the language
 - Examples: `if`, `while`
- In many languages, keywords are reserved and cannot be used for other purposes (e.g. as identifiers)

Identifiers

- Used to name an entity in a program
- The language specifies what characters can be used in an identifier
 - C++
 - First character: `_`, lowercase and uppercase letters, some escape sequences representing non-ASCII characters
 - Remaining characters: all of the above, plus digits
 - Scheme
 - Allows `! $ % & * + - . / : < = > ? @ ^ _ ~` in identifier!
 - Some implementations are even more permissive

- Something to think about during the break:
Suppose you want to count the number of bits that are 1 in a very long bitstring (e.g. a `vector<int>`)
- Here's a C++ function to do so:

```
std::size_t count(const std::vector<int> &data) {  
    constexpr int num_bits = sizeof(int) * 8;  
    std::size_t count = 0;  
    for (auto item : data) {  
        for (int i = 0; i < num_bits; ++i) {  
            count += (item >> i) & 1;  
        }  
    }  
    return count;  
}
```

- Can you improve this (by a constant factor)?

Syntax

- Concerned with the structure of code fragments
- Specifies what sequences of tokens constitute valid program fragments
 - Example: an expression must have balanced sets of parentheses
- Specified using a formal grammar (future topic!)

Semantics

- Concerned with the meaning of code fragments
 - e.g. what a piece of code defines, what value it computes, or what action it takes
- Further restrict what is valid code
 - Many things are syntactically correct but semantically invalid
- There are formal systems for specifying semantics, but natural language is often used instead

First-Class Entities

- ▶ We use **entity** to denote something that can be named in a program
 - ▶ Other terms also used: **citizen**, **object**
 - ▶ Examples: types, functions, data objects, values
- ▶ A **first-class entity** is an entity that supports all operations generally available to other entities
 - ▶ e.g. can be assigned to a variable, passed to or returned from a function

	C++	Java	Python	Scheme
Functions	sort of	no	yes	yes
Types	no	no	yes	no
Control	no	no	no	yes

Expressions

- An **expression** is a syntactic construct that is *evaluated* to produce a value
 - Examples: `3 + 4`, `foo()`
- Literals are one of the simplest kinds of expressions
 - Evaluate to the value they represent
- An identifier can syntactically be an expression
 - But only semantically valid if it names a first-class entity
 - Evaluates to the entity it names

Data Objects

- Consider the following code. What does the identifier `x` evaluate to when used as an expression?

```
int x = 3;  
... x ...; // x used as an expression
```

Poll: What does `x` evaluate to?

- A) Always the value 3**
- B) It depends on how `x` is used**

L-Values and R-Values

- An object can have two values associated with it
 - Its location in memory, called its ***l-value***
 - The value that it contains, called its ***r-value***
- Some objects, like temporaries, only have r-values
 - They may not actually exist in memory
- When an expression results in an object that has an l-value, it evaluates to the l-value
- The l-value is implicitly converted to an r-value if necessary

Compound Expressions

- **Precedence** and **associativity** rules determine how subexpressions are grouped when multiple operators are involved
- Precedence: divides operators into priority groups
 - Example: $\{*,/, \%\} > \{+,-\}$
- Associativity: how operators in the same precedence group apply
 - Example: $x = y = 3 + 4 - - - 5$
- Order of evaluation is distinct from precedence and associativity
 - Can be specified (mostly left to right in Python, Java), unspecified (function arguments in Scheme), or partially specified (C++)
 - Example: `cout << ++x << x;`

Statements and Side Effects

- Imperative languages have **statements**, which are **executed** to carry out some action
- Generally have **side effects**, which change the state of the machine
- Language syntax determines what constitutes a statement and how it is terminated
 - C family: simple statements terminated by semicolon
 - Python: newline (usually) or semicolon (rare)
 - Scheme?

Declarations and Definitions

- A **declaration** introduces a name into a program, along with properties about what it names
 - Examples

```
extern int x;  
void foo(int, int);  
class SomeClass;
```
- A **definition** additionally specifies the actual data or code that the name refers to
 - C, C++: definitions are declarations, but a declaration need not be a definition
 - Java: no distinction between definitions and declarations
 - Python: no declarations¹, definitions are statements that are executed

¹ Type annotations are not considered declarations. Quoting from PEP 526: "Type annotations should not be confused with variable declarations in statically typed languages."