

EECS 280 – Lecture 19

Structural Recursion

1

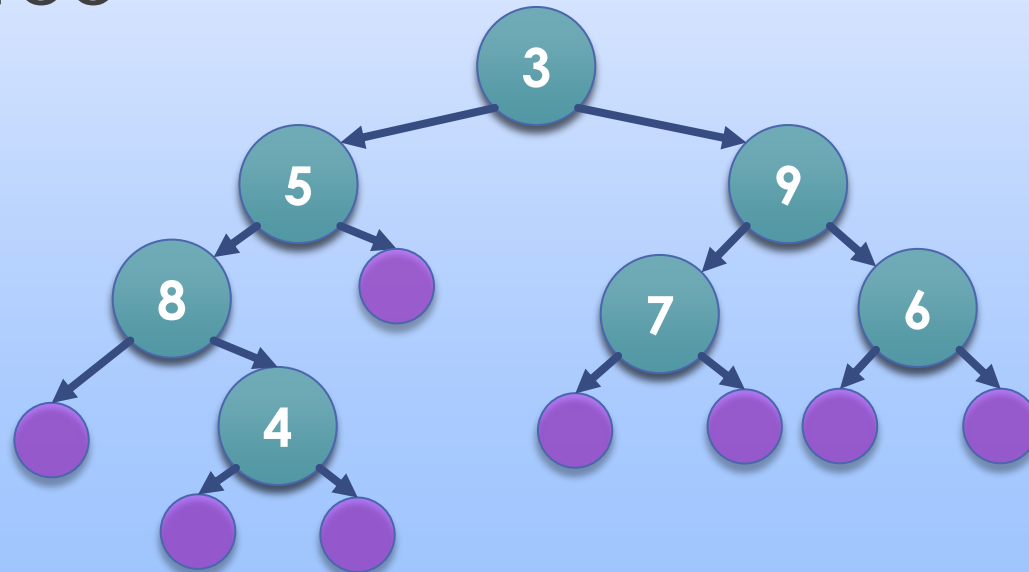
3/23/2022

Recursive Data Structures

➡ List



➡ Tree



Recall: Linked List Data Representation

Assume int elements for now.

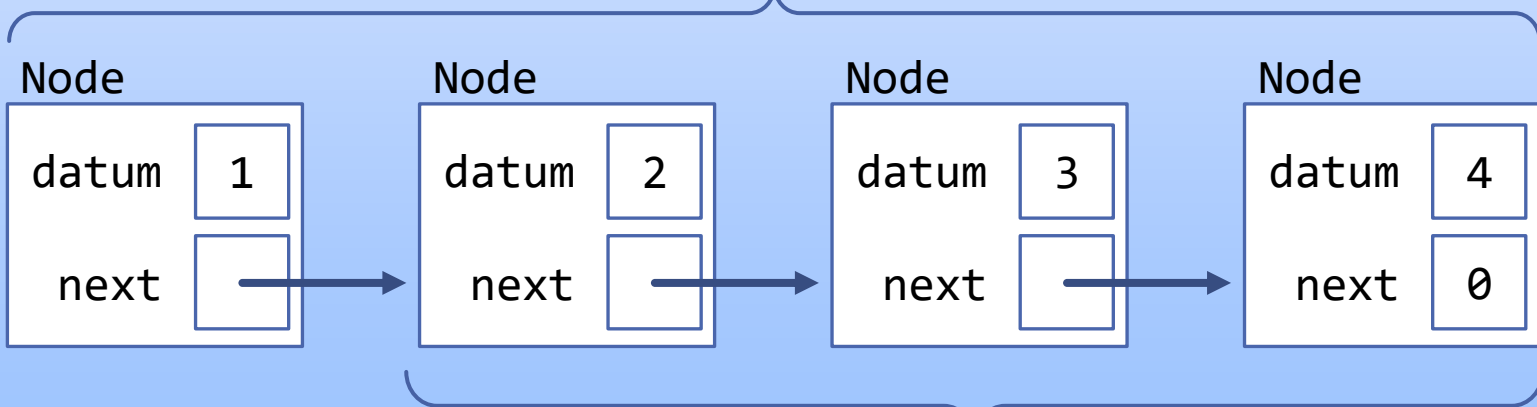
```
struct Node {  
    int datum;  
    Node *next;  
};
```

Used to store an element of the list.

Contains the address of the next node in the list.

- The node structure of a linked list is self-similar!

This is a list...



This piece is also a list!

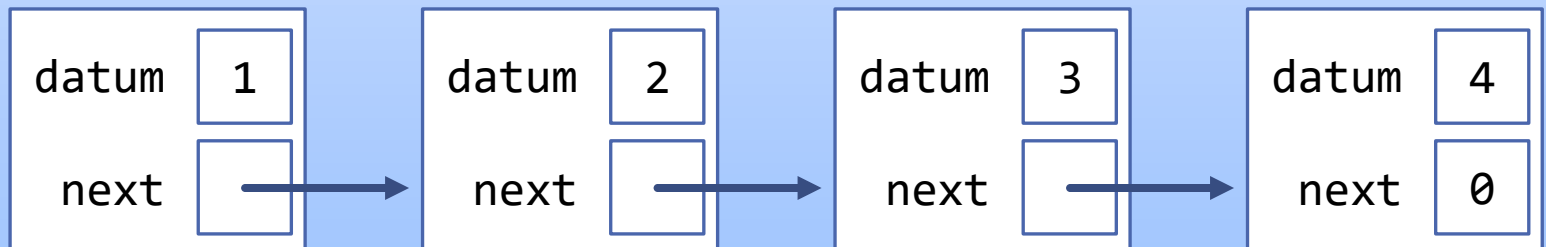
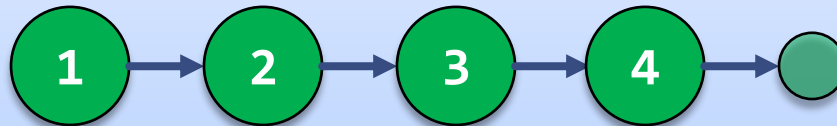
Recursively Defining a List

► Conceptually, any list is either:

1. empty




2. A datum, followed by a sub-list



Processing a List Recursively

- For example, let's **compute the length of a list L**.
- Consider the two cases:

1. empty  easy enough: $\text{length}(L) = 0$

2. A datum, followed by a sub-list



$$\text{length}(L) = 1 + \text{length}(\text{the sub-list})$$

Recursion!

~~$\text{length}(\text{1} \rightarrow \text{2} \rightarrow \text{3} \rightarrow \text{4} \rightarrow \text{ })$~~

~~$= 1 + \text{length}(\text{2} \rightarrow \text{3} \rightarrow \text{4} \rightarrow \text{ })$~~

~~$1 + \text{length}(\text{3} \rightarrow \text{4} \rightarrow \text{ })$~~


~~$1 + \text{length}(\text{4} \rightarrow \text{ })$~~

~~$1 + \text{length}(\text{ })$~~

0

Processing a List Recursively

- ▶ For example, let's compute the length of a list L.
- ▶ Consider the two cases:

1. empty  easy enough: $\text{length}(L) = 0$

2. A datum, followed by a sub-list



$\text{length}(L) = 1 + \text{length}(\text{the sub-list})$

```
int length(Node *node) {  
    if (node == nullptr) { // BASE CASE  
        return 0;  
    }  
    else { // RECURSIVE CASE  
        return 1 + length(node->next);  
    }  
}
```

This could also be written as `!node.`

Exercise: List Recurrence Relations

- ▶ See Exercise 19.1 in the accompanying worksheet:

bit.ly/3fQE21a

Exercise: max

```
struct Node {  
    int datum;  
    Node *next;  
};
```

- ▶ Now, write a function to find the maximum element.
- ▶ Hint: You may need to use a different base case.
- ▶ Hint: Use the recursive leap of faith on the sub-list.



```
// Use this helper function if you want  
int max(int x, int y) {  
    if (x > y) { return x; }  
    else { return y; }  
}
```

```
// REQUIRES: 'node' must not be null (i.e. the list  
//           starting at 'node' may not be empty)  
// EFFECTS: Returns the largest element in the list.  
int max(Node *node) {  
    // YOUR CODE HERE  
}
```


Solution: max

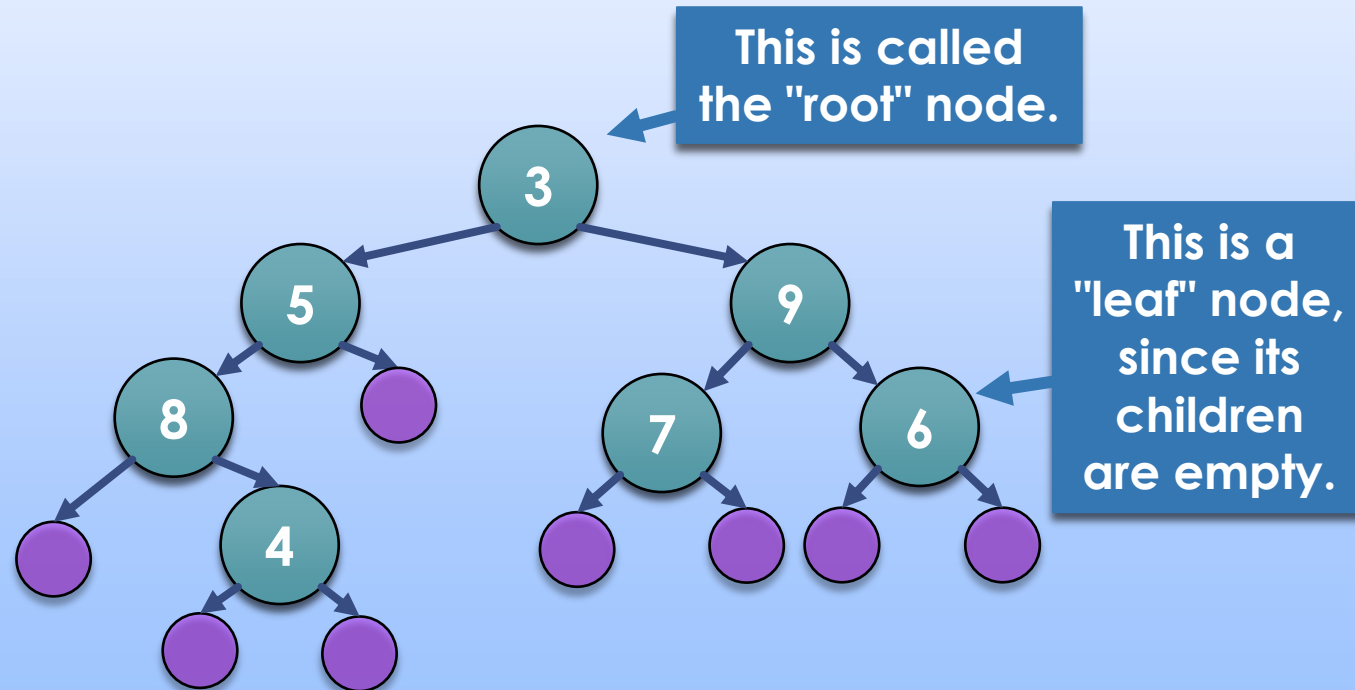
```
struct Node {  
    int datum;  
    Node *next;  
};
```

- Now, write a function to find the maximum element.
- Hint: You may need to use a different base case.
- Hint: Use the recursive leap of faith on the sub-list.



```
// REQUIRES: 'node' must not be null (i.e. the list  
//           starting at 'node' may not be empty)  
// EFFECTS: Returns the largest element in the list.  
int max(Node *node) {  
    // BASE CASE - A single element list  
    if (node->next == nullptr) {  
        return node->datum;  
    }  
    else { // RECURSIVE CASE  
        return std::max(node->datum, max(node->next));  
    }  
}
```

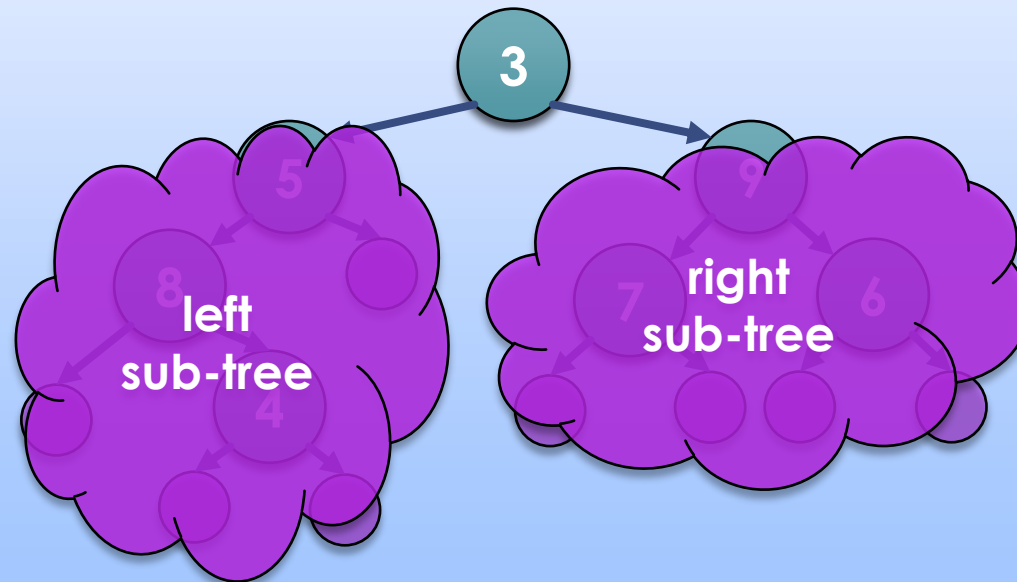
Recursive Data Structures



Recursively Defining a Tree

► Conceptually, a tree is either:

1. empty 
2. A datum, with left and right sub-trees



We will be working on *binary trees*, in which each node has two children.

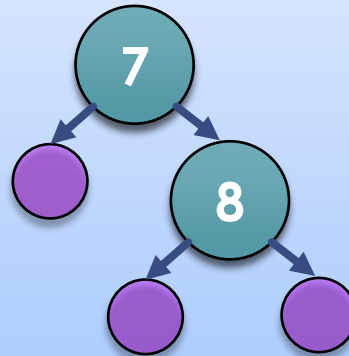
Properties of Trees

- ▶ We can measure a tree in two ways:
 - ▶ Size: The total number of elements
 - ▶ Height¹: The longest chain of nodes from root to leaf.

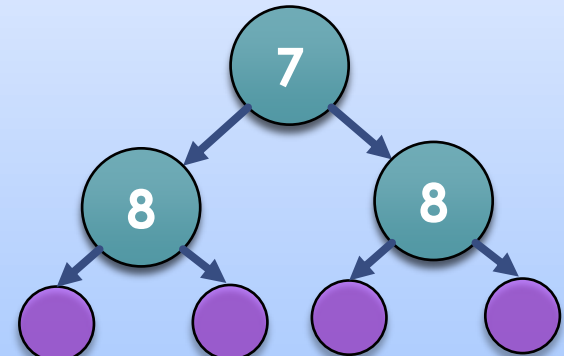
Size: 0
Height: 0



Size: 2
Height: 2



Size: 3
Height: 2



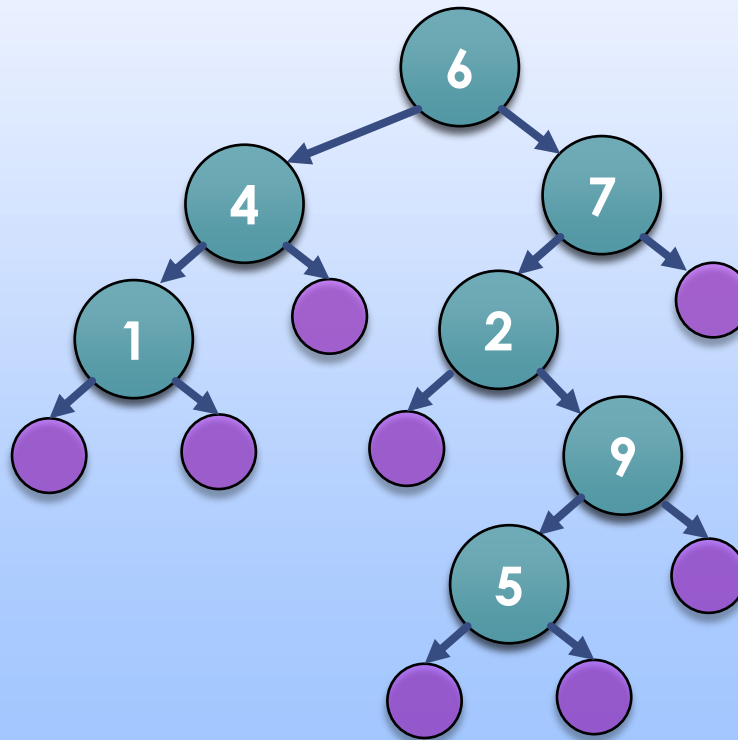
¹ This is sometimes called the “maximum depth” of the tree.

Properties of Trees

- What are the size and height of this tree?

Size: 7

Height: 5

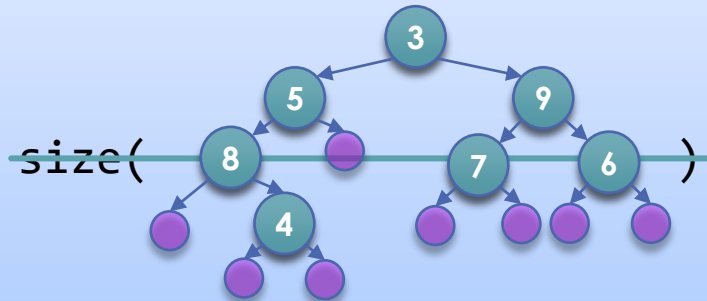
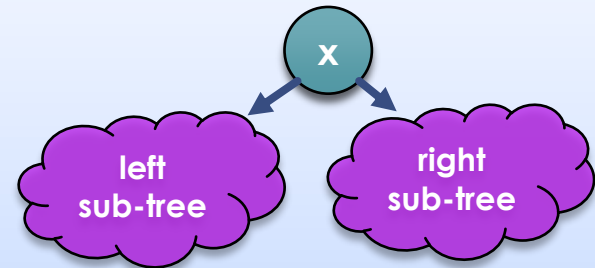


Processing a Tree Recursively

- For example, let's **compute the size of a tree T**:
- Consider the two cases:

1. empty 

2. A datum, with left and right sub-trees



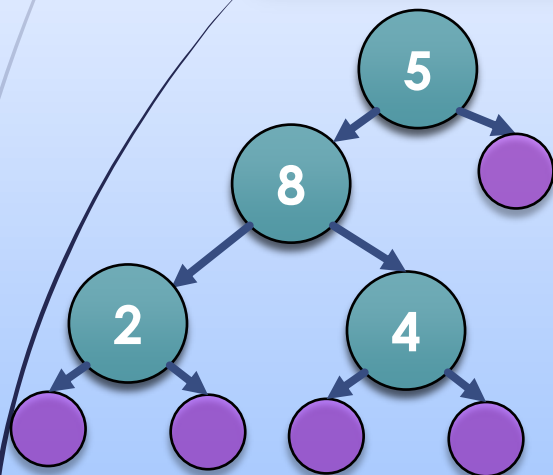
**Need to
count the
root node!**

$$\text{size}(\text{left sub-tree}) + \text{size}(\text{right sub-tree}) + 1$$

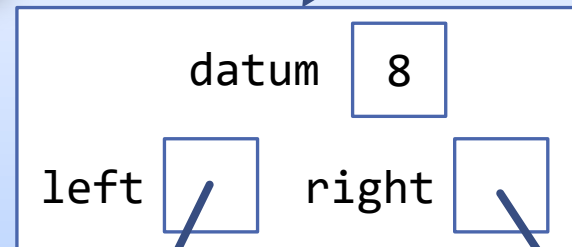
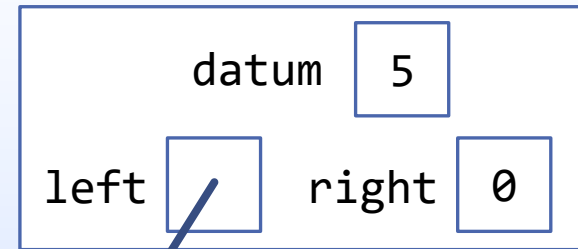
The equation is illustrated with the tree structure from the previous diagram. The left sub-tree is rooted at 8, and the right sub-tree is rooted at 7. The final '+ 1' represents the root node 3.

Tree Data Representation

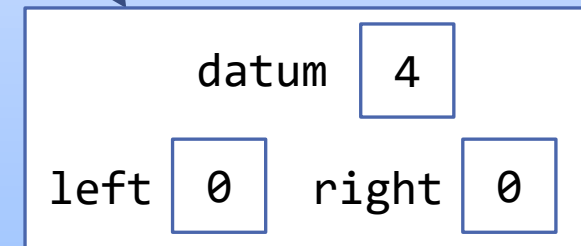
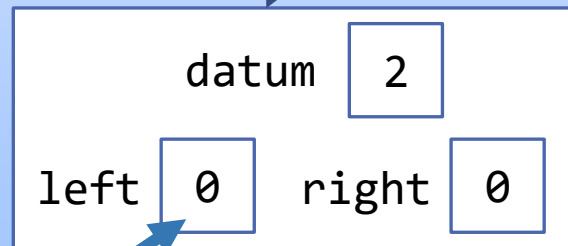
```
struct Node {  
    int datum;  
    Node *left;  
    Node *right;  
};
```



Empty trees are represented by a null pointer.



As with a List, these will be dynamically allocated.




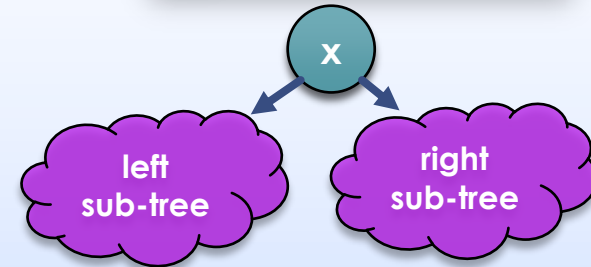
Example: Tree size

```
struct Node {  
    int datum;  
    Node *left;  
    Node *right;  
};
```

➤ Let's **compute the size of a tree T**:

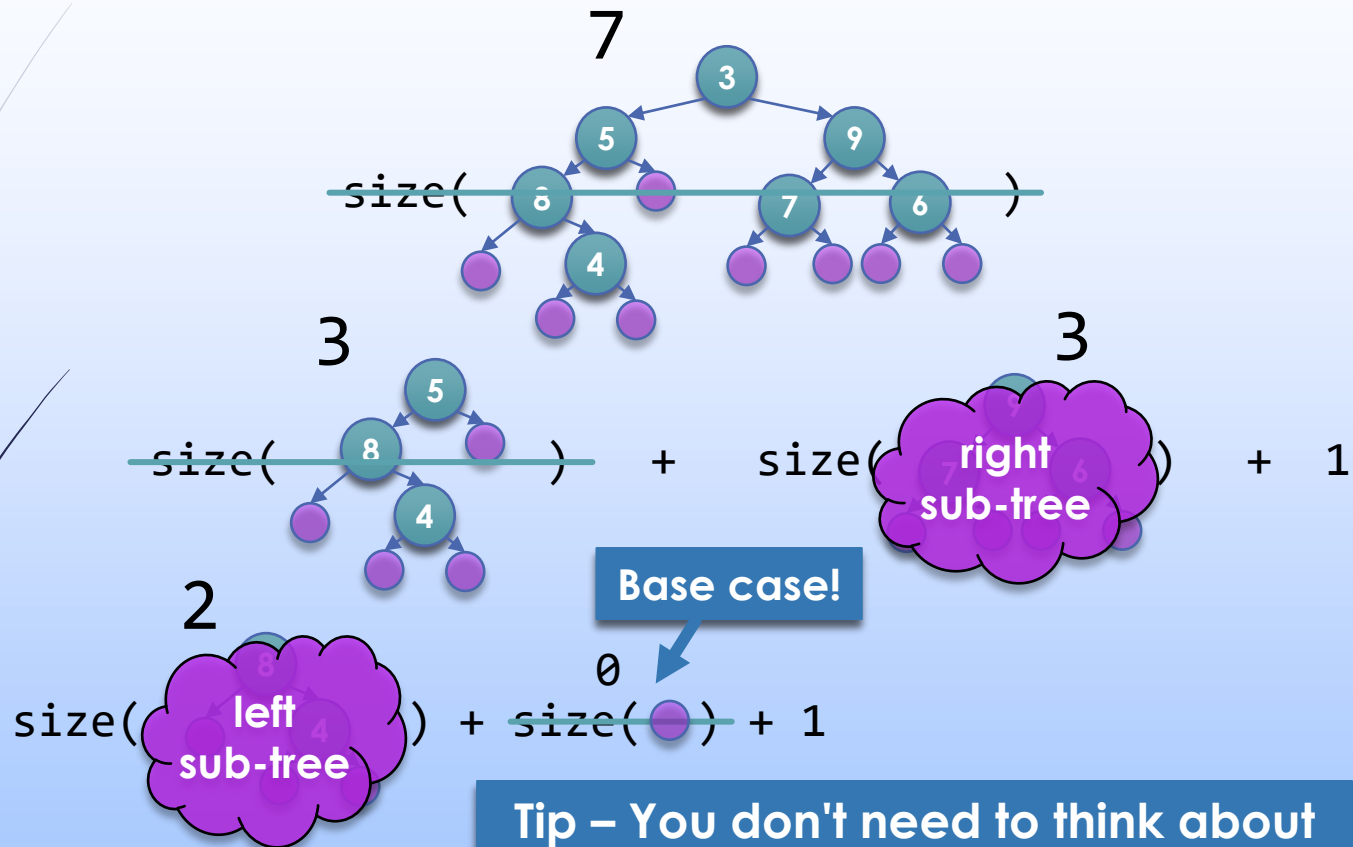
➤ Consider the two cases:

1. empty 
2. A datum, with left and right sub-trees



```
// EFFECTS: Returns the size of the tree rooted at 'node'.  
int size(Node *node) {  
    // BASE CASE - Empty tree has size 0  
    if (!node) {  
        return 0;  
    }  
  
    // RECURSIVE CASE  
    return 1 + size(node->left) + size(node->right);  
}
```


Computing Tree Size



Tip – You don't need to think about the recursion "all the way down". Just take the recursive leap of faith!

Exercise: Tree Recurrence Relations

- ▶ See Exercise 19.2 in the accompanying worksheet:

bit.ly/3fQE21a

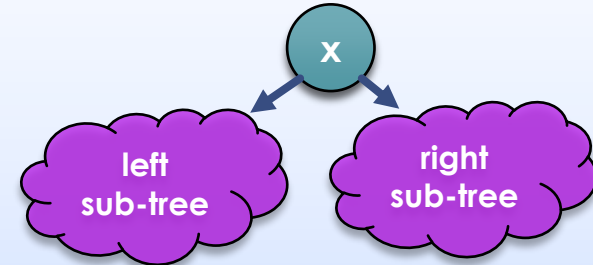
Exercise: Tree height

```
struct Node {  
    int datum;  
    Node *left;  
    Node *right;  
};
```

➤ Write a function to compute the **height** of a tree.

➤ Consider the two cases:

1. empty 
2. A datum, with left and right sub-trees



```
// EFFECTS: Returns the height of the tree rooted at 'node'.  
int height(Node *node) {
```

```
// YOUR CODE HERE
```

```
}
```

Tree Height

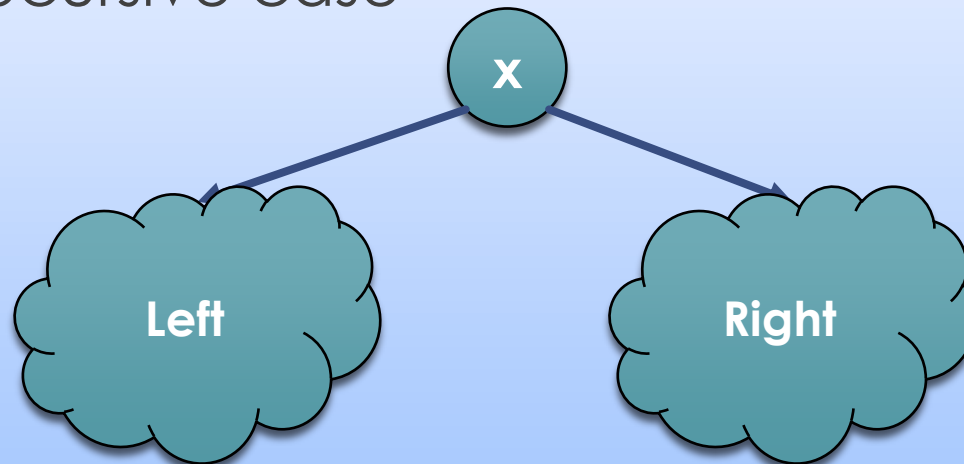
➡ Base case



```
if (!node) {  
    return 0;  
}
```

➡ Recursive case

$\text{height} = 1 + \max(L, R)$



$L = \text{height}(\text{node} \rightarrow \text{left})$

$R = \text{height}(\text{node} \rightarrow \text{right})$

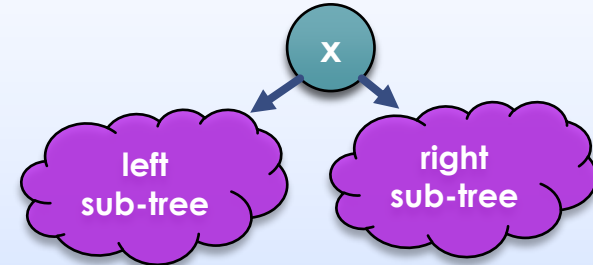
Solution: Tree height

```
struct Node {  
    int datum;  
    Node *left;  
    Node *right;  
};
```

➤ Write a function to compute the **height** of a tree.

➤ Consider the two cases:

1. empty 
2. A datum, with left and right sub-trees




```
// EFFECTS: Returns the height of the tree rooted at 'node'.  
int height(Node *node) {  
    // BASE CASE - Empty tree has size 0  
    if (!node) {  
        return 0;  
    }  
  
    // RECURSIVE CASE  
    return 1 + std::max(height(node->left),  
                        height(node->right));  
}
```

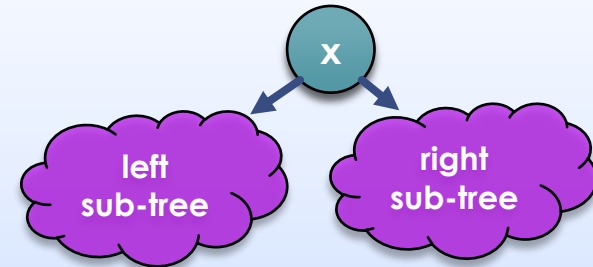
We'll start again in five minutes.



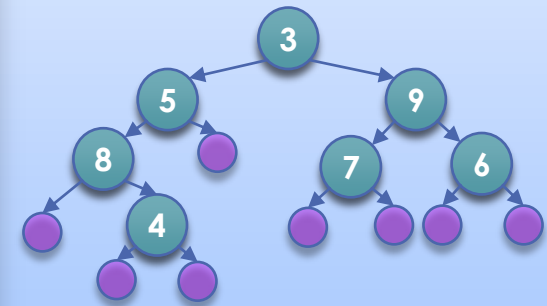
Tree print

```
struct Node {
    int datum;
    Node *left;
    Node *right;
};
```

- Write a function to print the elements of a tree.
- Consider the two cases:
 - empty 
 - A datum, with left and right sub-trees

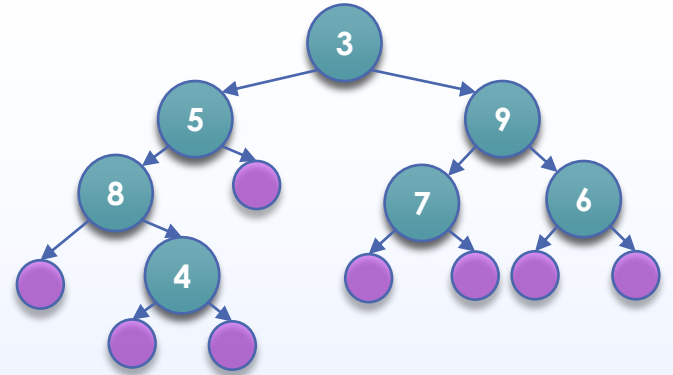


```
// EFFECTS: Prints the elements of
//           the tree rooted at
//           'node', with a space
//           after each element.
void print(Node *node) {
    if (node) { // RECURSIVE CASE
        cout << node->datum << " ";
        print(node->left);
        print(node->right);
    }
}
```



Prints 3 5 8 4 9 7 6

Tree Traversals



- For `print()`, we have a choice of when to process a datum
- A **preorder traversal** processes a datum before the recursive calls.

Prints 3 5 8 4 9 7 6

- An **inorder traversal** processes a datum between the calls.

Prints 8 4 5 3 7 9 6

- A **postorder traversal** processes a datum after the recursive calls.

Prints 4 8 5 7 6 9 3

```
if (node) { // RECURSIVE CASE
    cout << node->datum << " ";
    print(node->left);
    print(node->right);
}
```

```
if (node) { // RECURSIVE CASE
    print(node->left);
    cout << node->datum << " ";
    print(node->right);
}
```

```
if (node) { // RECURSIVE CASE
    print(node->left);
    print(node->right);
    cout << node->datum << " ";
}
```


Tree height

- ▶ Question: Can the height function be implemented tail recursively? If so, how? If not, why not?
 - ▶ Nope! You need to check both sides of the tree, which requires two recursive calls. They can't both be tail calls.

```
// EFFECTS: Returns the height of the tree rooted at 'node'.
int height(Node *node) {
    // BASE CASE - Empty tree has size 0
    if (!node) {
        return 0;
    }

    // RECURSIVE CASE
    return 1 + std::max(height(node->left),
                        height(node->right));
}
```

Types of Recursion

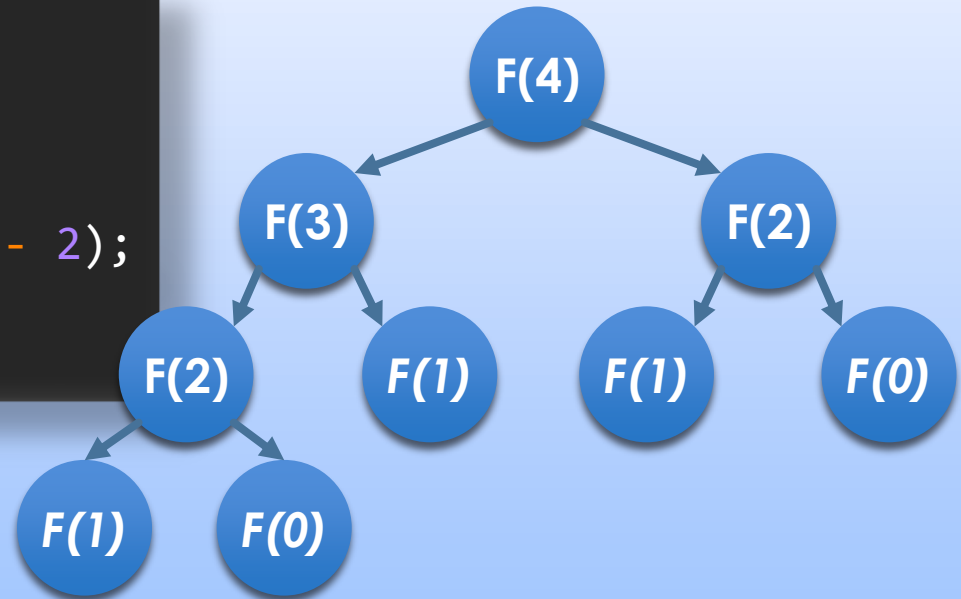
- ▶ A function is **linear recursive** if it makes at most one recursive call each time the function is called.
 - ▶ Example: fact, List max
- ▶ A function is **tail recursive** if it is linear recursive and all recursive calls are tail calls, so that no work is done after a recursive call.
 - ▶ Example: fact_tail, max_tail
- ▶ A function is **tree recursive** if it might make more than one recursive call each time the function is called.
 - ▶ Example: Tree size, height
 - ▶ A function doesn't have to operate on trees to be tree recursive! It is tree recursive if the structure of the recursive calls *branches* and thus looks like a tree.

Subproblem Graph: Fibonacci

$$F(n) = \begin{cases} 0 & n = 0 \\ 1 & n = 1 \\ F(n-1) + F(n-2) & n > 1 \end{cases}$$

```
int fib(int n) {  
    if (n <= 1) {  
        return n;  
    }  
    return fib(n - 1) + fib(n - 2);  
}
```

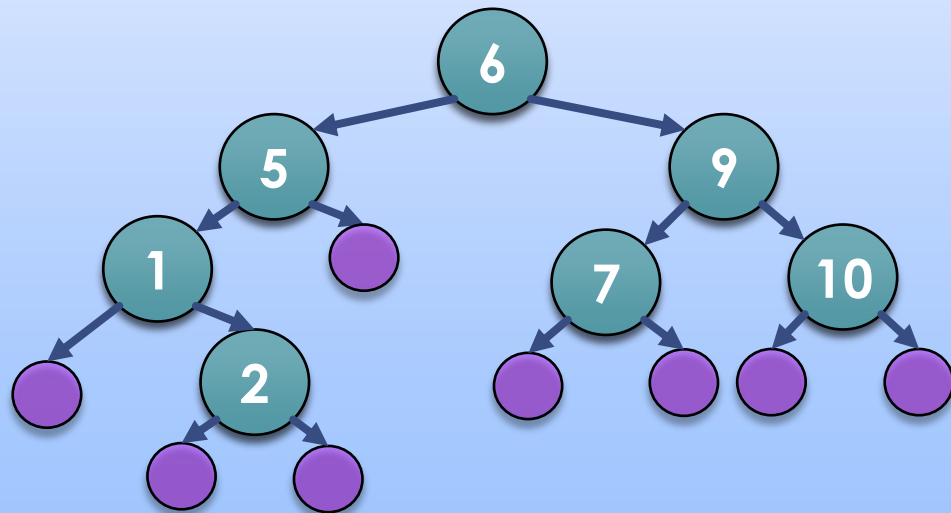
This fib function
is tree recursive.



Binary Search Trees (BSTs)

- ▶ A tree is a binary search tree if...
 - ▶ It is empty
- OR
- ▶ The left and right subtrees are binary search trees.
- ▶ All elements in any **left** subtree are **less** than the root.
- ▶ All elements in any **right** subtree are **greater** than the root.

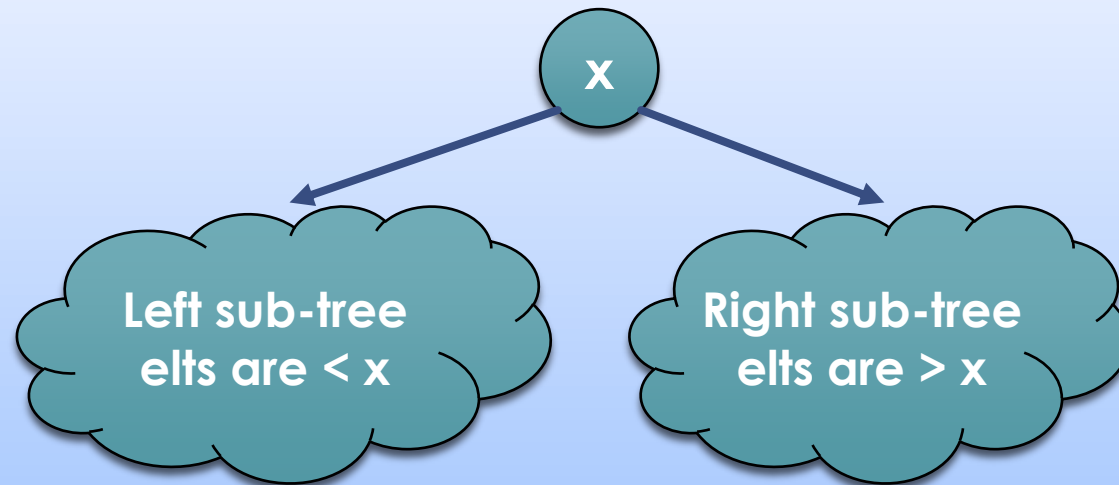
It is so called
because searching
for elements can
be done efficiently.



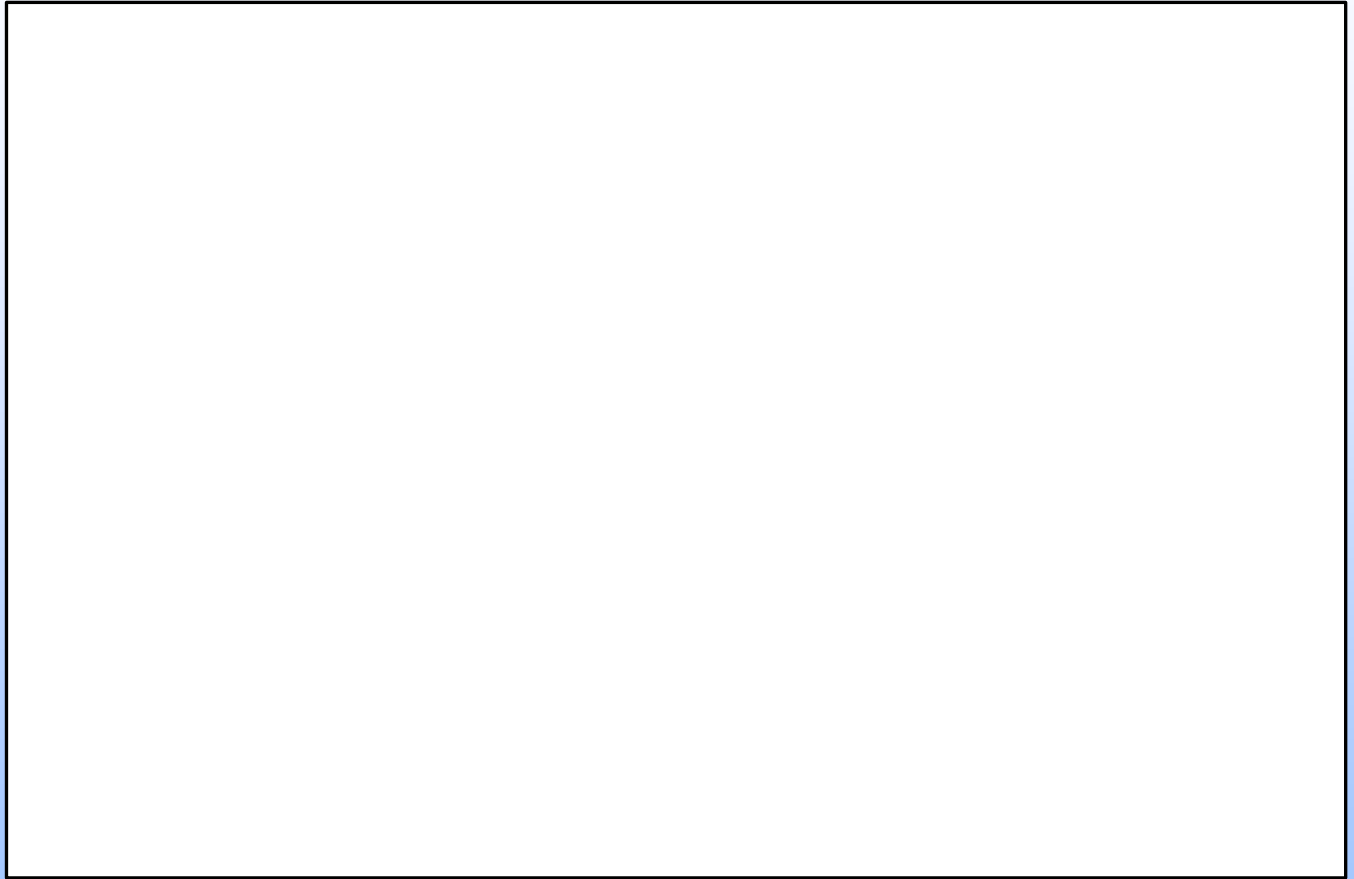
Note: In the slides and on project 5, we make a simplifying assumption that there are no duplicates in our BSTs.

Searching in a Binary Search Tree

- ▶ If we're looking for an element in a BST, comparing with the root tells us where to look.
 - ▶ Thus the name "*Binary Search Tree*".



Building a Binary Search Tree

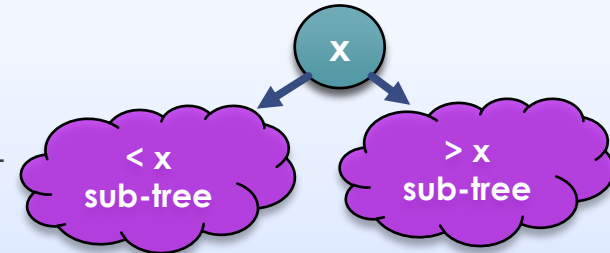


Example: BST max

```
struct Node {  
    int datum;  
    Node *left;  
    Node *right;  
};
```

► The maximum element in a binary search tree is:

1. The datum in the node if the right sub-tree is empty
2. Otherwise, the maximum element in the right sub-tree

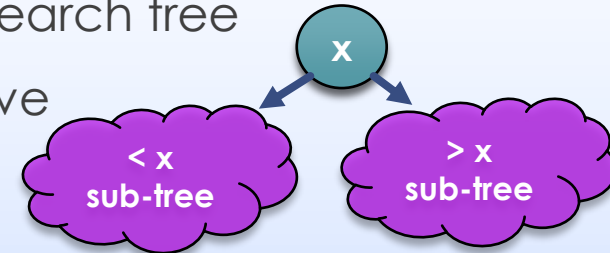


```
// REQUIRES: 'node' must be a binary search tree that  
//           is not empty (i.e. 'node' is not null)  
// EFFECTS: Returns the largest element in the tree.  
int max(Node *node) {  
    // BASE CASE - Right sub-tree is empty  
    if (!node->right) {  
        return node->datum;  
    }  
    else { // RECURSIVE CASE - Right sub-tree not empty  
        return max(node->right);  
    }  
}
```

Exercise: BST contains

```
struct Node {  
    int datum;  
    Node *left;  
    Node *right;  
};
```

- Write a function that determines whether or not an element is contained in a binary search tree
- Your solution should be tail recursive

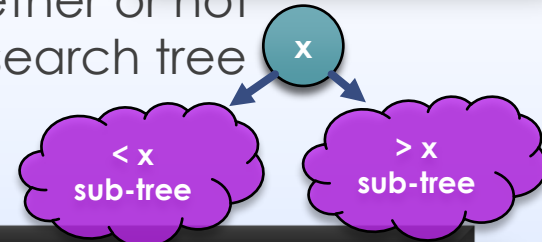


```
// REQUIRES: 'node' must be a binary search tree  
// EFFECTS: Returns whether or not the tree contains  
//           the given item.  
bool contains(Node *node, int item) {  
  
    // YOUR CODE HERE  
  
}
```


Solution: BST contains

```
struct Node {  
    int datum;  
    Node *left;  
    Node *right;  
};
```

- Write a function that determines whether or not an element is contained in a binary search tree
- Your solution should be tail recursive



```
// REQUIRES: 'node' must be a binary search tree  
// EFFECTS: Returns whether or not the tree contains  
//           the given item.  
bool contains(Node *node, int item) {  
    if (!node) {  
        return false;  
    } else if (node->datum == item) {  
        return true;  
    } else if (node->datum > item) {  
        return contains(node->left, item);  
    } else {  
        return contains(node->right, item);  
    }  
}
```

The BinarySearchTree Interface

```
template <typename T>
class BinarySearchTree {
public:
    BinarySearchTree();
    BinarySearchTree(const BinarySearchTree &other);
    BinarySearchTree & operator=(const BinarySearchTree &other);
    ~BinarySearchTree();

    bool empty() const;
    int size() const;
    bool contains(const T &item) const;
    void insert(const T &item);

private:
    struct Node {
        T datum;
        Node *left, *right;
    };

    Node *root;
};
```

BinarySearchTree Implementation

- We can implement the functions that operate on Nodes as private, static member functions

```
template <typename T>
class BinarySearchTree {
public:
    bool contains(const T &item) const {
        return contains_impl(root, item);
    }
private:
    Node *root;
    static bool contains_impl(Node *node, const T &item) {
        if (!node) {
            return false;
        } else if (node->datum == item) {
            return true;
        } else if (node->datum > item) {
            return contains_impl(node->left, item);
        } else {
            return contains_impl(node->right, item);
        }
    }
};
```

static means the function can't access member variables. It shouldn't! The pointer to the root of the node structure is passed in.

Project 5 BinarySearchTree.h

- In project 5, you'll implement a BST.
- The starter files provide a code skeleton and we've implemented several pieces for you:
 - The overall structure
 - The data representation and `Node` struct
 - Iterators
 - The Big Three
- Many of these call static `"_impl"` functions for manipulating the recursive node structure.
 - You write the `"_impl"` functions!