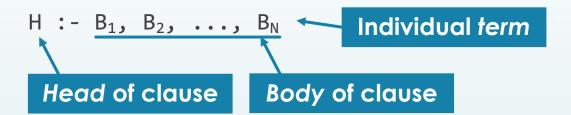
EECS 390 – Lecture 18

Logic Programming II

Review: Horn Clauses

- A logic program is expressed as a set of axioms that are assumed to be true
- An axiom takes the form of a Horn clause, which specifies a reverse implication:



This is equivalent to

$$(B_1 \wedge B_2 \wedge \dots \wedge B_N) \Rightarrow H$$

with implicit quantifiers.

Review: Queries

A goal is a query that the system attempts to prove from the axioms

Possible reasoning:

```
Goal
```

S = bill is also a valid solution given the axioms.

3/27/24

Prolog

- Prolog is the foundational language of logic programming and is the most widely used
- A Prolog program consists of a set of Horn clauses, using the syntax on the preceding slides
- A Horn clause has a head term and optional body terms
- A term may be atomic, compound, or a variable
 - **► Atomic**: atoms and numbers
 - **► Atom**: Scheme-like symbol or quoted string
 - If an atom starts with a letter, it must be lowercase
 - hello =< + 'logic programming'
 - ► Variables: symbols that start with an uppercase letter

Hello X

Compound Terms

 A compound term consists of a *functor*, which is an atom, followed by a list of one or more argument terms

```
pair(1, 2) wizard(harry) writeln(hello(world))
```

- A compound term is interpreted as a predicate, with a truth value, if it is a head term, a body term, or the goal
- Otherwise, the compound term is interpreted as data
 - e.g. hello(world) in writeln(hello(world))

Facts and Rules

A Horn clause with no body is a fact, since it is always true

```
mother(molly, bill).
mother(molly, charlie).
end of clause
```

A Horn clause with a body is a rule

```
parent(P, C) :- mother(P, C).
sibling(A, B) :- parent(P, A), parent(P, B).
```

- Meaning:
 - If mother(P, C) is true, then so is parent(P, C)
 - If parent(P, A) and parent(P, B) are true, then so is sibling(A, B)
- A program is a set of Horn clauses

Goals and Queries

- A goal is a predicate that the interpreter attempts to prove
- Loading the program from the previous slide and entering the goal sibling(bill, S) produces:

```
?- sibling(bill, S).
S = bill;
S = charlie.
Ask for more solutions
```

- A semicolon asks for more solutions
- A period ends a query
 - Can be entered by the user
 - Can be produced by the interpreter, in which case it is certain no more solutions exist

Implementing Lists

- Compound terms can represent data structures
- Example: use pair(A, B) to represent a pair
 - This won't be a head or body term, so it will be treated as data
- Relations on pairs:

cons(A, B, pair(A, B)).
cdr(pair(_, B), B).
car(pair(A, _), A).
is_null(nil).

Relates a first and second item to a pair

Anonymous variable

```
?- cons(1, nil, X).
X = pair(1, nil).
?- car(pair(1, pair(2, nil)), X).
X = 1.
?- cdr(pair(1, pair(2, nil)), X).
X = pair(2, nil).
?- cdr(pair(1, pair(2, nil)), X),
   car(X, Y), cdr(X, Z).
X = pair(2, nil), Y = 2, Z = nil.
?- is_null(nil).
true.
?- is_null(pair(1, nil)).
false.
```

Singleton Variables

- A singleton variable is a variable that only appears once in an axiom
- Singleton variables can occur inadvertently as a result of a typo:

 Oops

```
cons(First, Second, pair(Frist, Second)).
```

- To address this, the Prolog interpreter warns about the occurrence of a singleton variable
- We can inform the interpreter about an intentional singleton by using a name that begins with an underscore
 Named, intentional singleton variable

```
cdr(pair(_First, Second), Second).
car(pair(First, _), First).
```

Anonymous variable – does not match any other occurrence of _

Prolog Lists

 Prolog also provides built-in linked lists, specified as elements between square brackets

```
[] [1, a] [b, 3, foo(bar)]
```

The pipe symbol acts like a dot in Scheme, separating some elements from the rest of the list

```
?- writeln([1, 2 | [3, 4]]).
[1,2,3,4]
true.
```

This allows us to write predicates like the following:

Numbers and Comparisons

- Prolog includes integer and floating-point numbers
- Comparison predicates can be written in infix order

The = operator specifies explicit unification, not equality

Arithmetic

 Arithmetic operators represent compound terms and are not evaluated

Comparisons perform evaluation on both operands

```
?- 7 =:= 3 + 4. 7 is equal to the result of evaluating +(3, 4)
```

The is operator unifies its first argument with the arithmetic result of its second argument

List Length

We can now define a predicate for length on our list representation:
 Unify Length with the

```
len(nil, 0).
len(pair(_, Second), Length) :-
   len(Second, SLength), Length is SLength + 1.
```

```
?- len(nil, X).
X = 0.
?- len(pair(1, pair(b, nil)), X).
X = 2.
```

Must be second body term so that SLength is sufficiently instantiated for arithmetic

■ Built-in lists have a built-in length predicate

```
?- length([1, a, 3], X).
X = 3.
```

Side Effects

- Prolog provides I/O predicates, including reading from standard input and writing to standard output
- We will only use write and writeln:

```
?- X = 3, write('The value of X is: '),
   writeln(X).
The value of X is: 3
X = 3.
```

Unification and Search

- A logic solver is built around the processes of unification and search
- Search in Prolog uses backward chaining
 - Start with a set of goal terms
 - Look for a clause whose head can unify with a goal term
 - If unification succeeds, replace the old goal term with the body terms of the clause
 - Search succeeds when no more goal terms remain
- Unification attempts to unify two terms, which may require recursively unifying subterms
 - May require instantiating variables to values

Unification

- An atomic term only unifies with itself (or an uninstantiated variable)
- An uninstantiated variable unifies with any term
 - If the other term is not a variable, then the variable is *instantiated* with the value of the other term, i.e. all occurrences of the variable are replaced with the value
 - If the other term is a variable, the two variables are bound together such that later instantiating one with a value also instantiates the other with the same value
- A compound term unifies with another compound term if the functors and number of arguments are the same, and the arguments recursively unify

$$X = 3$$

 $Y = foo(1, Z)$
 $foo(1, A) = foo(B, 3)$ % unifies $B = 1$, $A = 3$

Search Order

- In pure logic programming, search order is irrelevant as long as the search terminates
- In Prolog, clauses are applied in program order, and terms within a body are resolved in left-to-right order
- Example:

```
sibling(A, B) :-
  mother(P, A), mother(P, B).

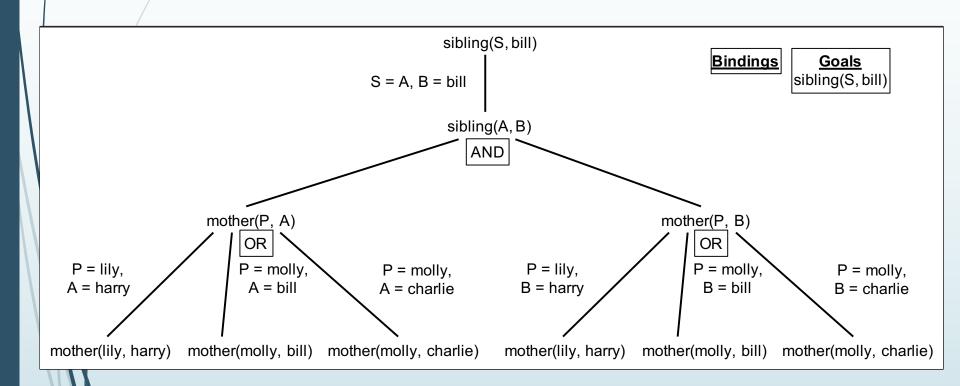
mother(lily, harry).
mother(molly, bill).
mother(molly, charlie).

?- sibling(S, bill)
S = bill
```

Search Tree

```
sibling(A, B) :-
  mother(P, A), mother(P, B).
mother(lily, harry).
mother(molly, bill).
mother(molly, charlie).
```

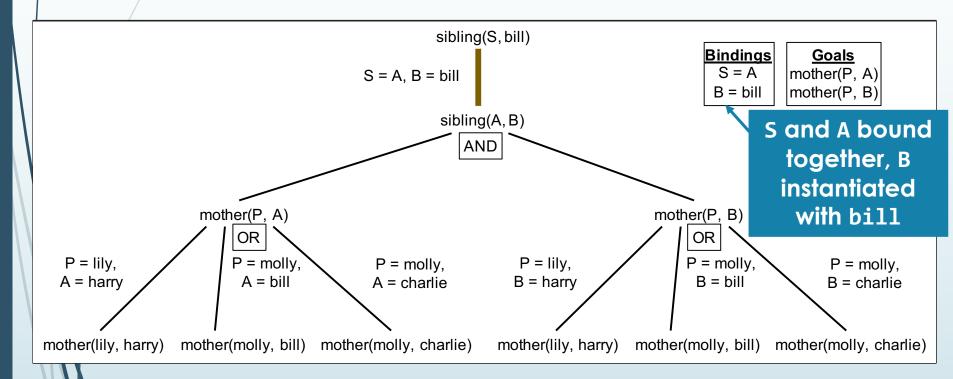
 Search encounters choice points, and backtracking is required on failure or if the user asks for more solutions



Search Tree

```
sibling(A, B) :-
  mother(P, A), mother(P, B).
mother(lily, harry).
mother(molly, bill).
mother(molly, charlie).
```

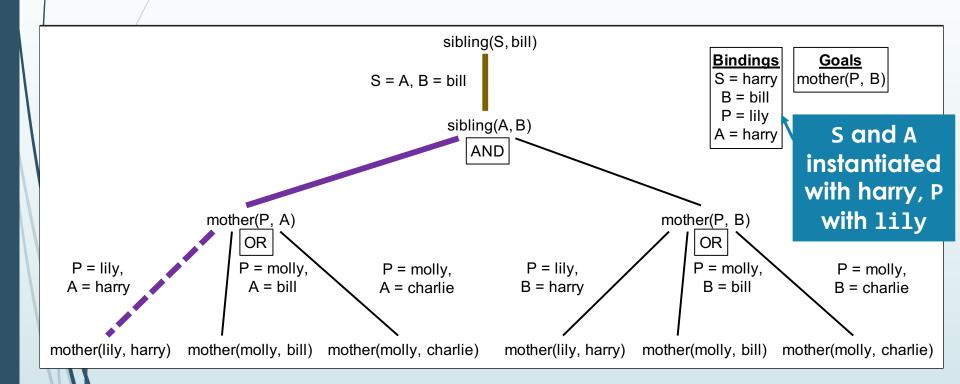
First, sibling(S, bill) is unified with the head term sibling(A, B), and the body terms of the clause are added to the goals



Search Tree

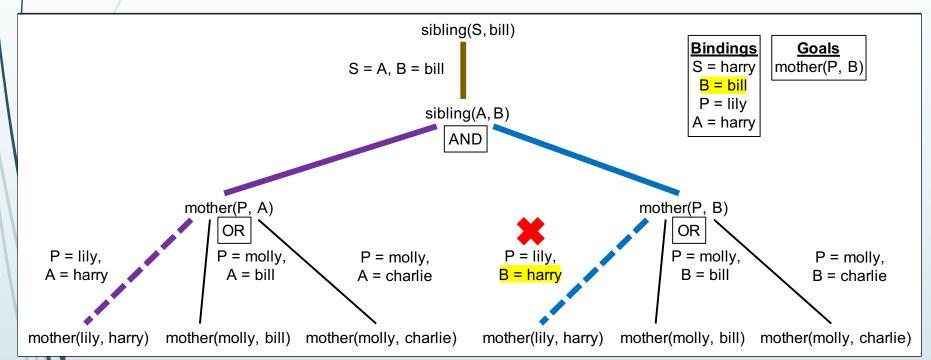
```
sibling(A, B) :-
  mother(P, A), mother(P, B).
mother(lily, harry).
mother(molly, bill).
mother(molly, charlie).
```

■ The goal mother(P, A) is solved first, with an initial choice of applying the fact mother(lily, harry)



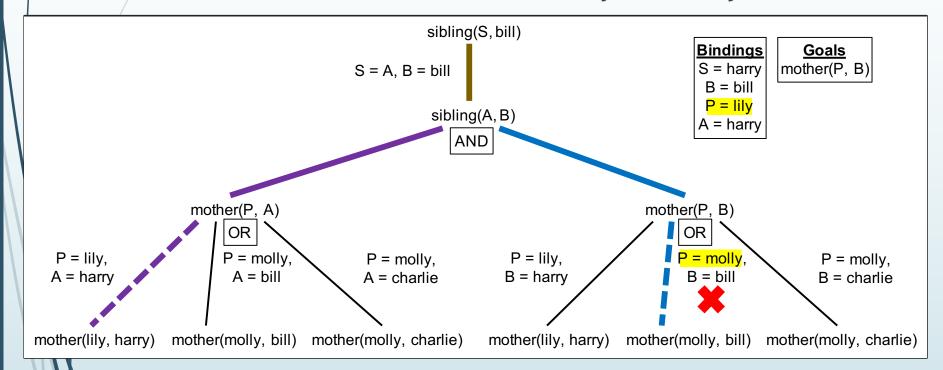
```
sibling(A, B) :-
  mother(P, A), mother(P, B).
mother(lily, harry).
mother(molly, bill).
mother(molly, charlie).
```

- Then the goal mother(P, B) is solved, with an initial choice of applying the fact mother(lily, harry)
- However, unification of B = bill with harry fails



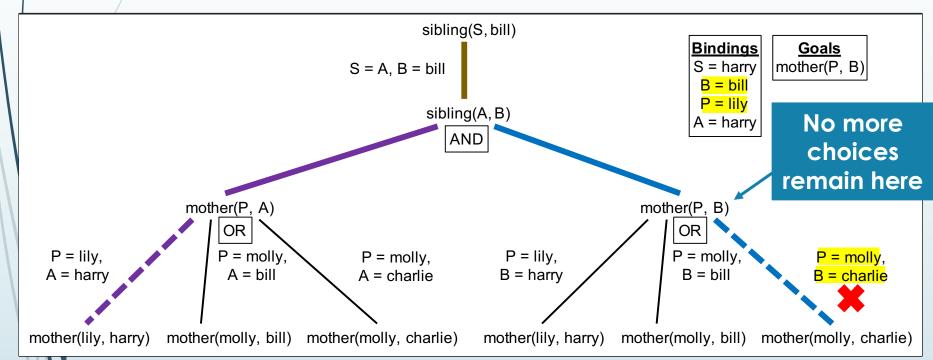
```
sibling(A, B) :-
  mother(P, A), mother(P, B).
mother(lily, harry).
mother(molly, bill).
mother(molly, charlie).
```

- The search backtracks to the previous choice point, attempting to apply the fact mother(molly, bill)
- However, unification of P = lily with molly fails



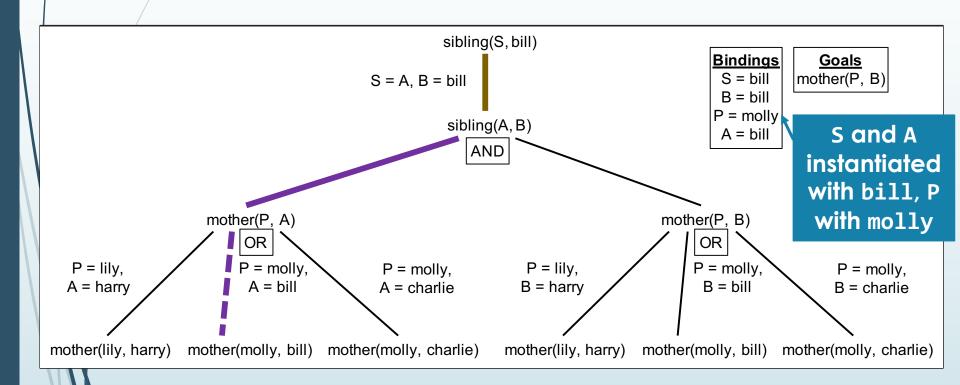
```
sibling(A, B) :-
  mother(P, A), mother(P, B).
mother(lily, harry).
mother(molly, bill).
mother(molly, charlie).
```

- The search backtracks once again, attempting to apply the fact mother(molly, charlie)
- However, unification of P = lily with molly fails



```
sibling(A, B) :-
  mother(P, A), mother(P, B).
mother(lily, harry).
mother(molly, bill).
mother(molly, charlie).
```

The search backtracks to the preceding choice point, unifying mother(P, A) with mother(molly, bill)

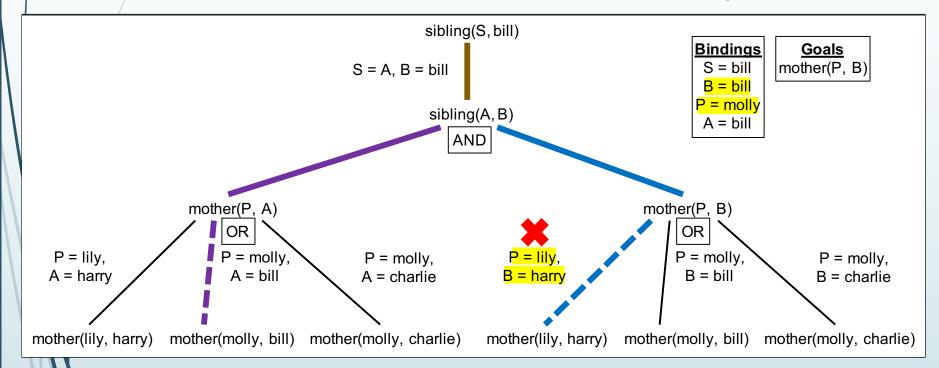


mother(molly, bill).

mother(molly, charlie).

Search Tree

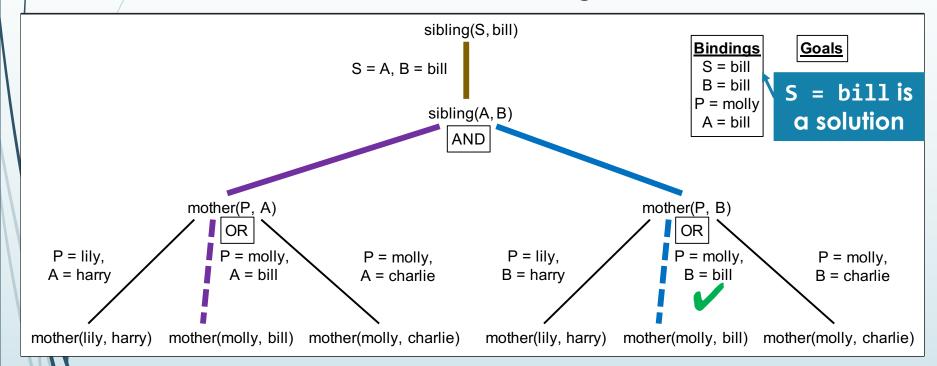
- Then the goal mother(P, B) is solved, with an initial choice of applying the fact mother(lily, harry)
- However, unification of B = bill with harry fails



First Solution

```
sibling(A, B) :-
  mother(P, A), mother(P, B).
mother(lily, harry).
mother(molly, bill).
mother(molly, charlie).
```

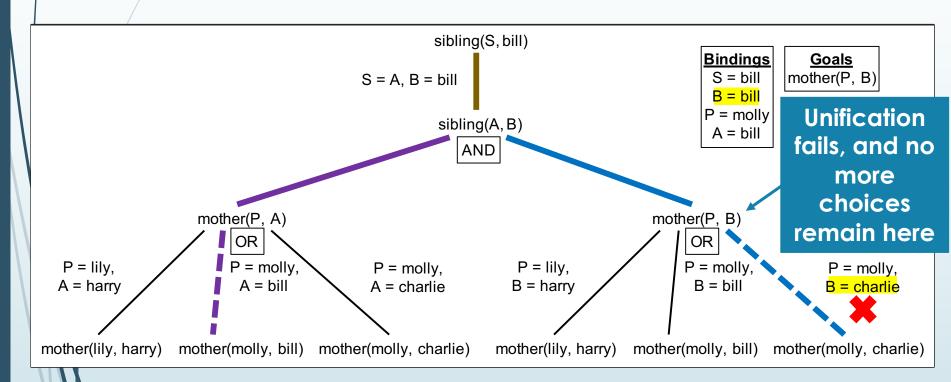
- The search backtracks to the previous choice point, attempting to apply the fact mother(molly, bill)
- Unification succeeds, and no goal terms remain



Continuing the Search

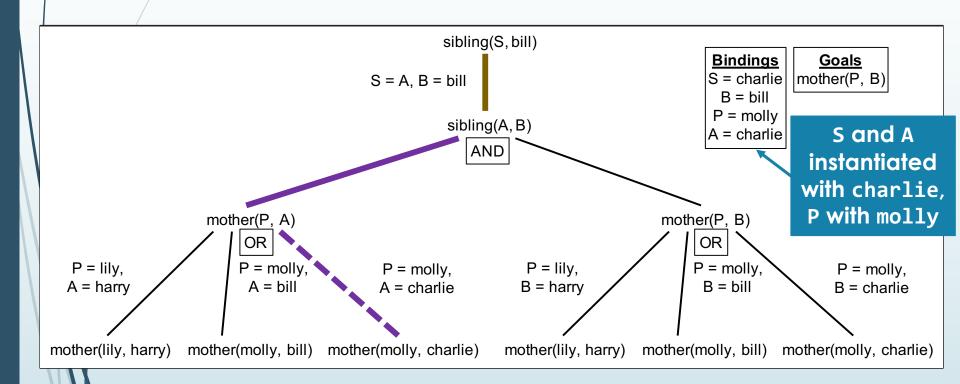
```
sibling(A, B) :-
  mother(P, A), mother(P, B).
mother(lily, harry).
mother(molly, bill).
mother(molly, charlie).
```

■ If we ask the interpreter for another solution, it backtracks to the previous choice point, attempting to apply the fact mother(molly, charlie)



```
sibling(A, B) :-
  mother(P, A), mother(P, B).
mother(lily, harry).
mother(molly, bill).
mother(molly, charlie).
```

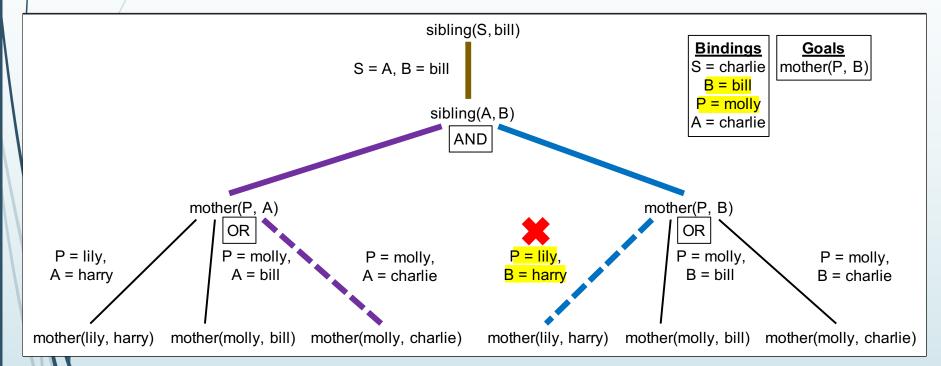
The search backtracks to the preceding choice point, unifying mother(P, A) with mother(molly, charlie)



mother(molly, charlie).

Search Tree

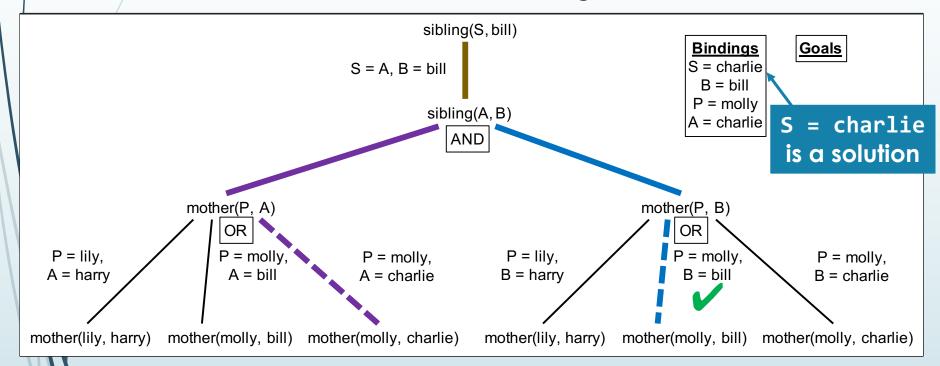
- Then the goal mother(P, B) is solved, with an initial choice of applying the fact mother(lily, harry)
- However, unification of B = bill with harry fails



Second Solution

```
sibling(A, B) :-
  mother(P, A), mother(P, B).
mother(lily, harry).
mother(molly, bill).
mother(molly, charlie).
```

- The search backtracks to the previous choice point, attempting to apply the fact mother(molly, bill)
- Unification succeeds, and no goal terms remain



No More Solutions

```
sibling(A, B) :-
  mother(P, A), mother(P, B).
mother(lily, harry).
mother(molly, bill).
mother(molly, charlie).
```

If we ask the interpreter for another solution, it backtracks to the previous choice point, attempting to apply the fact mother(molly, charlie)

