

Disjoint Sets: Naive Implementations

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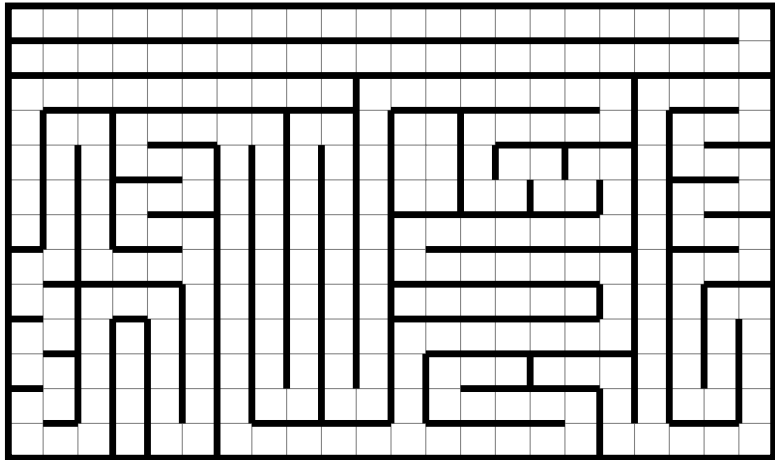
Data Structures
Data Structures and Algorithms

Outline

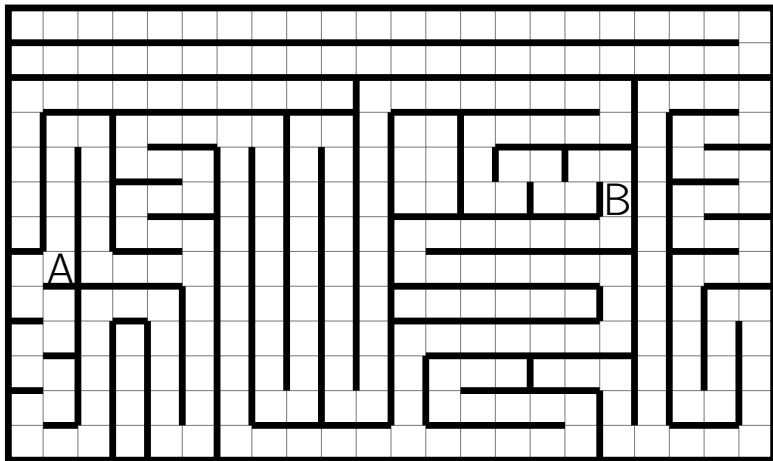
1 Overview

2 Naive Implementations

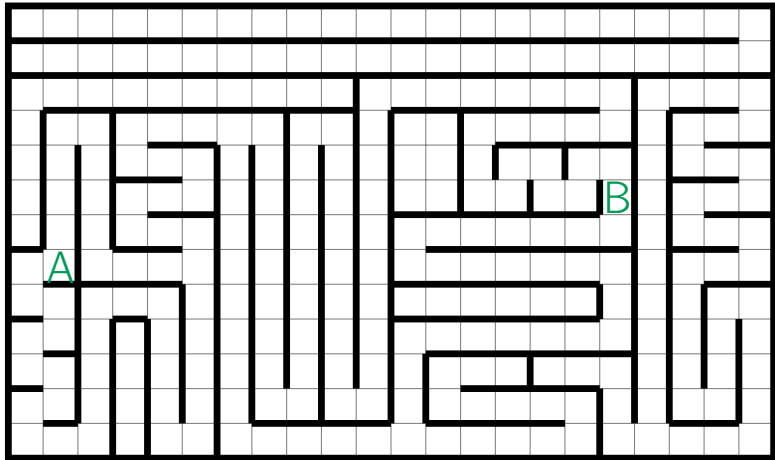
Maze: Is B Reachable from A?



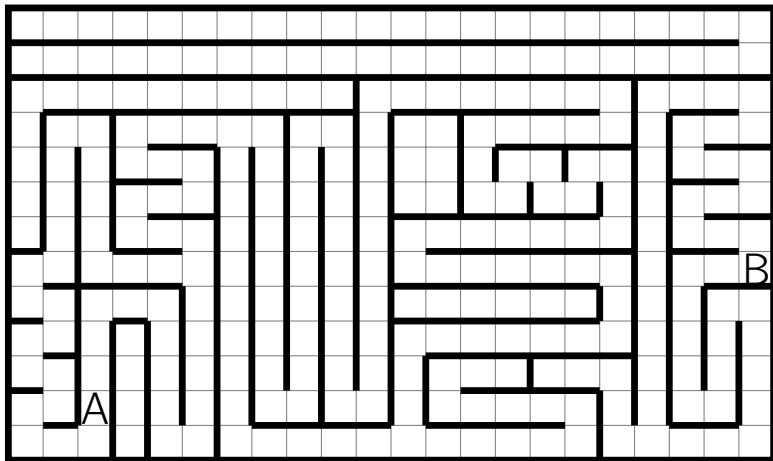
Maze: Is B Reachable from A?



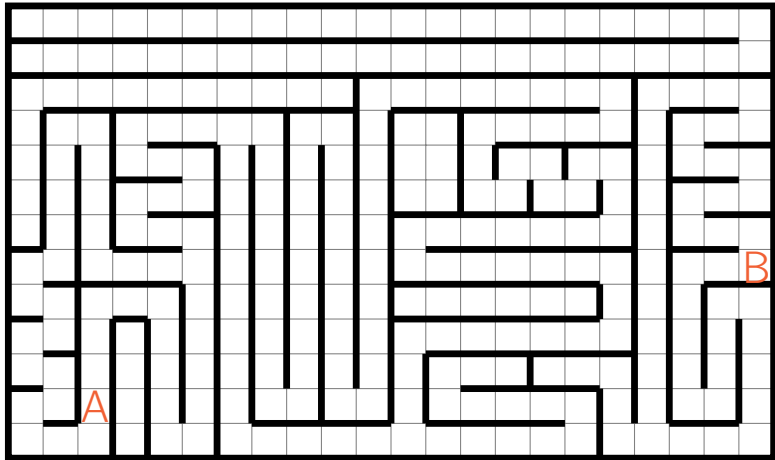
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 - if x and y lie in the same set, then $\text{Find}(x) = \text{Find}(y)$
 - otherwise, $\text{Find}(x) \neq \text{Find}(y)$
- $\text{Union}(x, y)$ merges two sets containing x and y

Preprocess(*maze*)

for each cell c in *maze*:

 MakeSet(c)

for each cell c in *maze*:

 for each neighbor n of c :

 Union(c, n)

Preprocess(*maze*)

```
for each cell  $c$  in  $maze$ :  
    MakeSet( $c$ )  
for each cell  $c$  in  $maze$ :  
    for each neighbor  $n$  of  $c$ :  
        Union( $c, n$ )
```

IsReachable(A, B)

```
return Find( $A$ ) = Find( $B$ )
```

Building a Network

Building a Network



MakeSet(1)

Building a Network



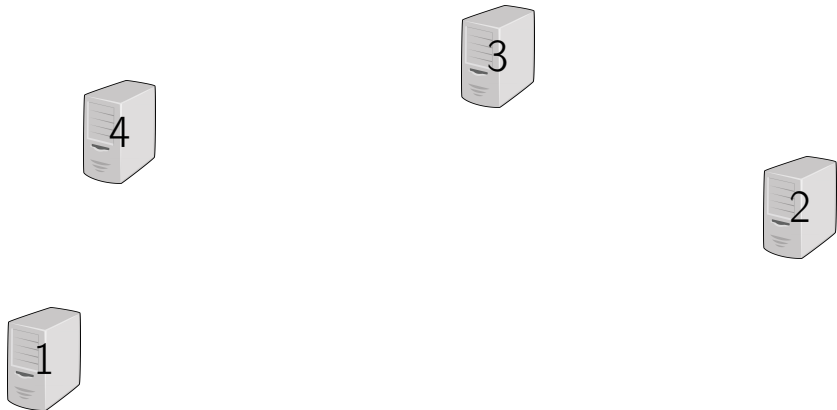
MakeSet(2)

Building a Network



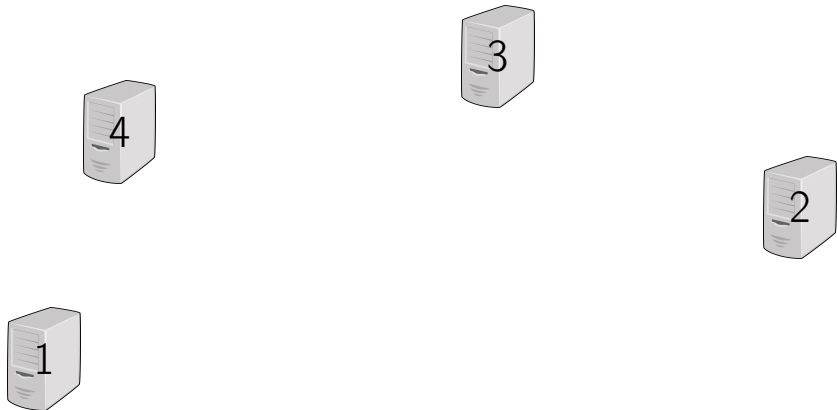
MakeSet(3)

Building a Network



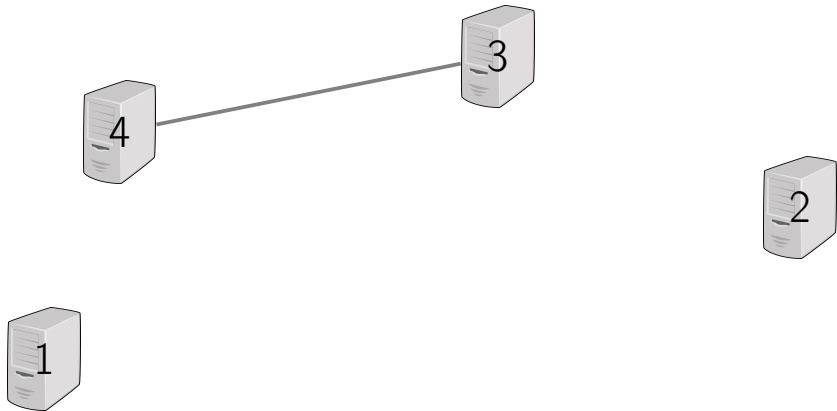
MakeSet(4)

Building a Network

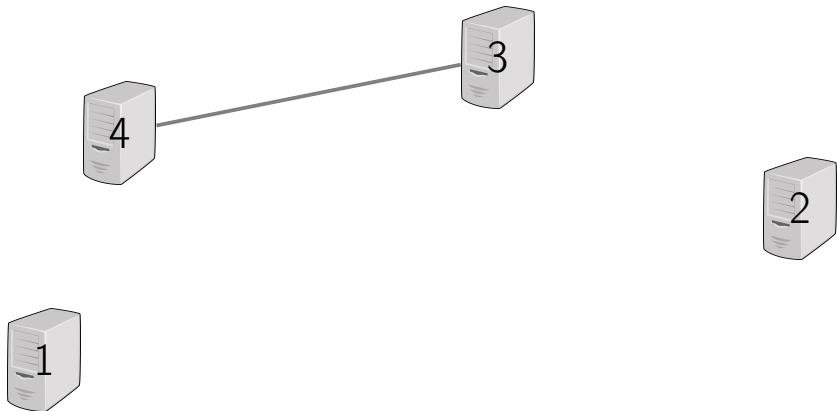


$\text{Find}(1) = \text{Find}(2) \rightarrow \text{False}$

Building a Network

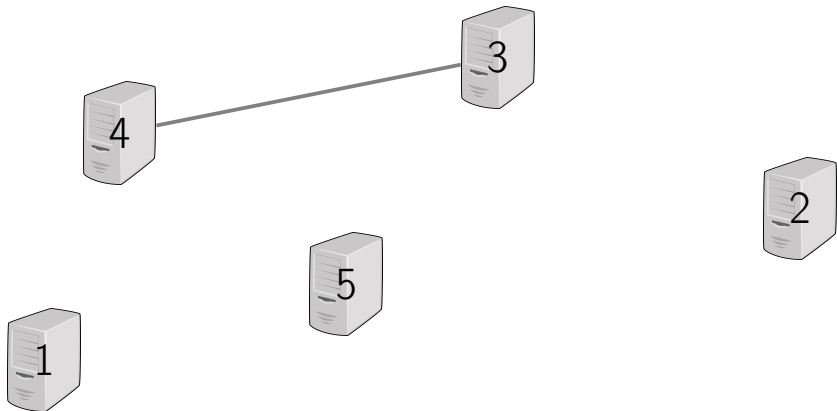


Building a Network



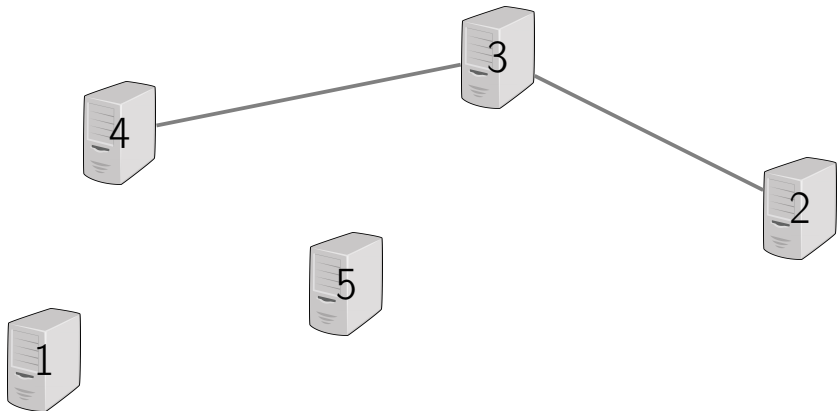
`Union(3, 4)`

Building a Network

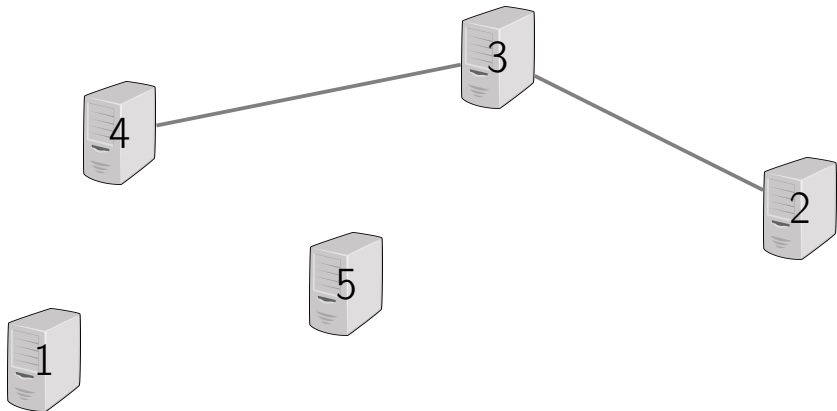


MakeSet(5)

Building a Network

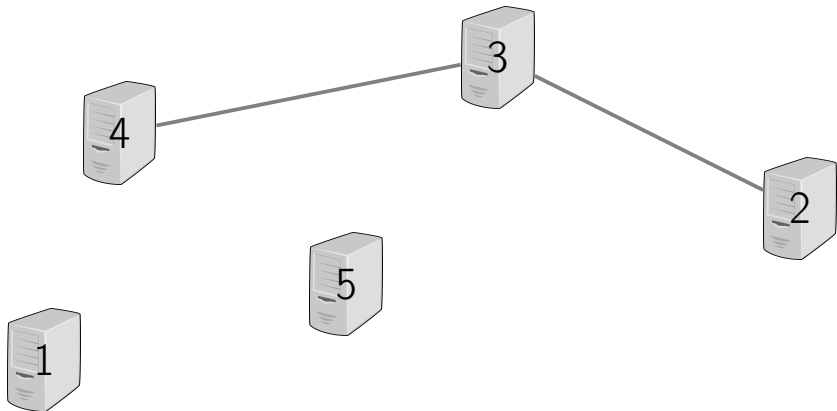


Building a Network



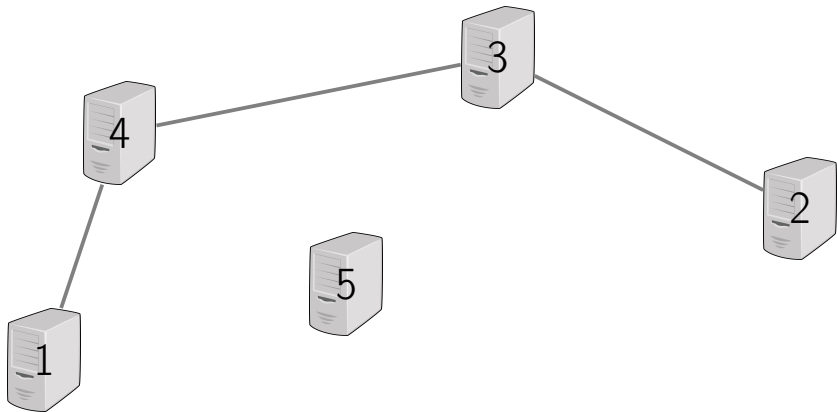
Union(3, 2)

Building a Network

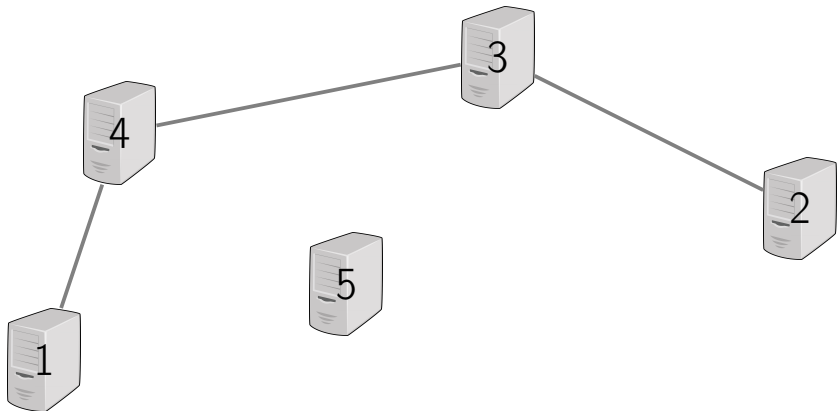


$\text{Find}(1) = \text{Find}(2) \rightarrow \text{False}$

Building a Network

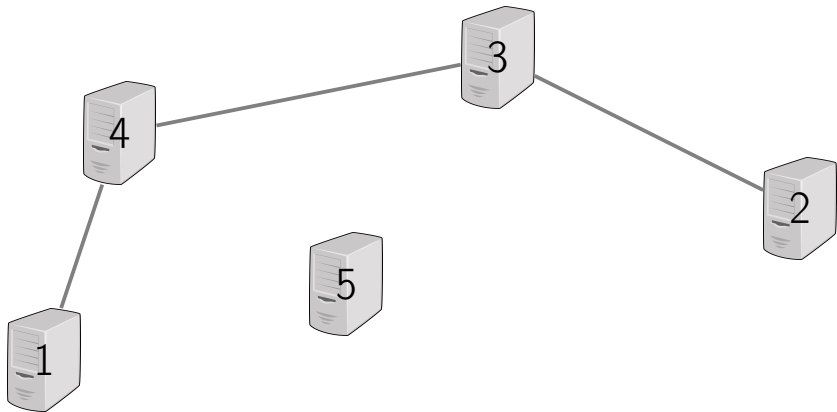


Building a Network



Union(1, 4)

Building a Network



$\text{Find}(1) = \text{Find}(2) \rightarrow \text{True}$

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For simplicity, we assume that our n objects are just integers $1, 2, \dots, n$.

Using the Smallest Element as ID

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- Use array `smallest[1...n]`:
`smallest[i]` stores the smallest element in the set i belongs to

Example

$\{9, 3, 2, 4, 7\}$ $\{5\}$ $\{6, 1, 8\}$

	1	2	3	4	5	6	7	8	9
smallest	1	2	2	2	5	1	2	1	2

MakeSet(i)

smallest[i] $\leftarrow i$

Find(i)

return smallest[i]

MakeSet(i)

$\text{smallest}[i] \leftarrow i$

Find(i)

return $\text{smallest}[i]$

Running time: $O(1)$

Union(i, j)

$i_id \leftarrow \text{Find}(i)$

$j_id \leftarrow \text{Find}(j)$

if $i_id = j_id$:

 return

$m \leftarrow \min(i_id, j_id)$

for k from 1 to n :

 if $\text{smallest}[k] \in \{i_id, j_id\}$:

$\text{smallest}[k] \leftarrow m$

Union(i, j)

```
 $i\_id \leftarrow \text{Find}(i)$   
 $j\_id \leftarrow \text{Find}(j)$   
if  $i\_id = j\_id$ :  
    return  
 $m \leftarrow \min(i\_id, j\_id)$   
for  $k$  from 1 to  $n$ :  
    if  $\text{smallest}[k]$  in  $\{i\_id, j\_id\}$ :  
         $\text{smallest}[k] \leftarrow m$ 
```

Running time: $O(n)$

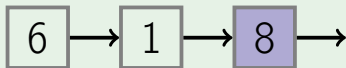
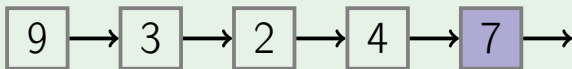
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- What basic data structure allows for efficient merging?

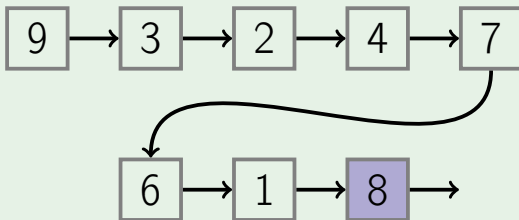
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- Linked list!

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- What basic data structure allows for efficient merging?
- Linked list!
- Idea: represent a set as a linked list, use the list tail as ID of the set

Example: merging two lists



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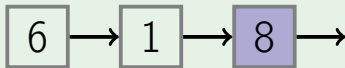
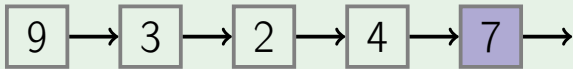
- Pros:

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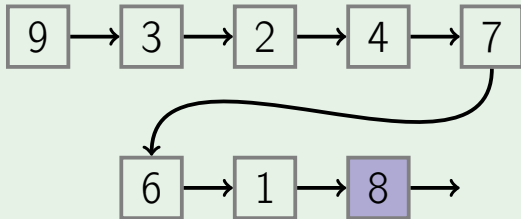
- Cons:

- Running time of Find is $O(n)$ as we need to traverse the list to find its tail
- Union(x, y) works in time $O(1)$ **only** if we can get the tail of the list of x and the head of the list of y in constant time!

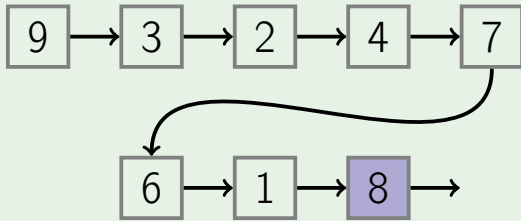
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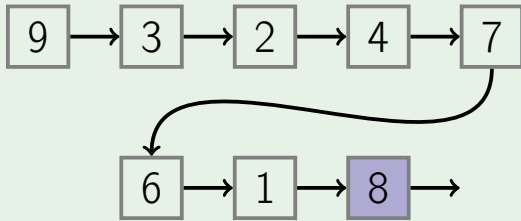


Example: merging two lists



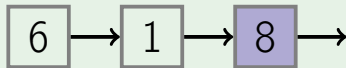
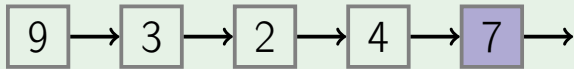
`Find(9)` goes through all elements

Example: merging two lists

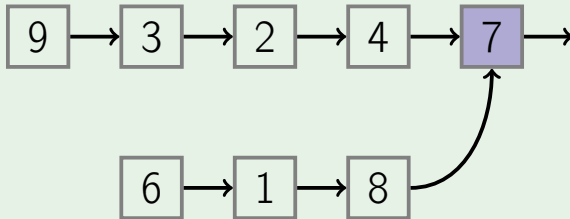


can we merge in a different way?

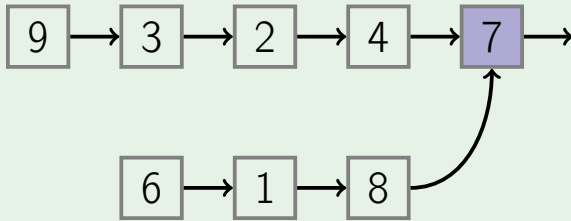
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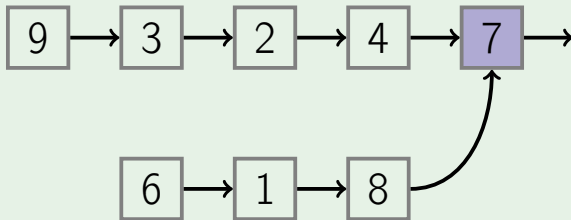


Example: merging two lists



instead of a list we get a tree

Example: merging two lists



we'll see that representing sets as trees gives a very efficient implementation: nearly constant amortized time for all operations