

# Cheatsheet - Asymptotic Analysis

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## 1. About

Asymptotic analysis is an alternative way of describing the time or memory requirements of an algorithm.

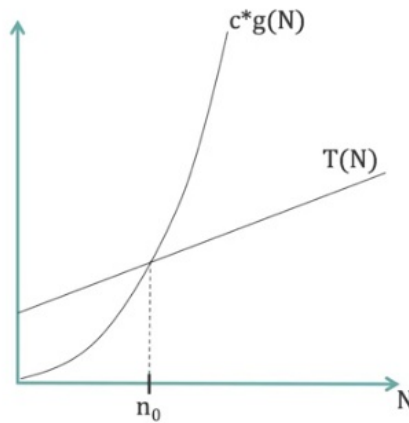
## 2. Big O Notation

Big O notation  $O(x)$  defines a set of functions that act as an **upper bound**  $g(N)$  for  $T(N)$ . Formally defined as:

$T(N)$  is  $O(g(N))$  if there exist positive constants  $c$  and  $n_0$  such that:

$$T(N) \leq c \times g(N) \quad \text{for all } N > n_0$$

Note that there can be **multiple functions**  $g_x(N)$  that act as an **upper bound** for  $T(N)$ . Additionally, do notice that it's **not necessary** that  $c \times g(N)$  is equal to or greater than  $T(N)$  for all values of  $N$ .



For example, consider:

$$T(N) = 10N^2 + 15N + 5$$

$$g(N) = N^2$$

$$c = 1$$

Here,  $c \times g(N)$  is never greater than  $T(N)$ , because there is no solution for:

$$10N^2 + 15N + 5 \leq 1 \times N^2$$

However, consider:

$$c = 25$$

In case of  $N = 1$  we get:

$$\begin{aligned} 10 \times 1^2 + 15 \times 1 + 5 &\leq 25 \times 1^2 \\ &= 10 + 15 + 5 \leq 25 \\ &= 30 \leq 25 \end{aligned}$$

Which is false. However, for  $N = 2$  we get:

$$\begin{aligned} 10 \times 2^2 + 15 \times 2 + 5 &\leq 25 \times 2^2 \\ &= 40 + 30 + 5 \leq 100 \\ &= 75 \leq 100 \end{aligned}$$

Which is true. Therefore:

$$\begin{aligned} T(N) \text{ is } O(N^2) \text{ because} \\ T(N) \leq 25 \times g(N) \text{ for all } N \geq 2 \end{aligned}$$

The choice for  $c$  is **arbitrary**, as long as it satisfies the conditions.

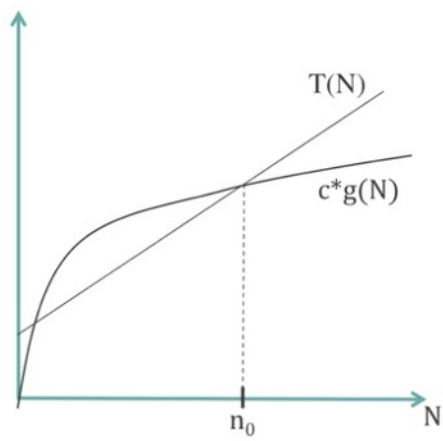
## 3. Omega Notation

The Omega notation  $\Omega(x)$  defines a set of functions that act as a **lower bound**  $g(N)$  for  $T(N)$ . Formally defined as:

$T(N)$  is  $\Omega(g(N))$  if there exist positive constants  $c$  and  $n_0$  such that:

$$T(N) \geq c \times g(N) \text{ for all } N > n_0$$

Similarly to the Big O notation, there can be **multiple functions**  $g_x(N)$  that act as a **lower bound** for  $T(N)$  and it's **not necessary** that  $c \times g(N)$  is equal to or less than  $T(N)$  for all values of  $N$ , but only for the larger values.



4. Theta Notation

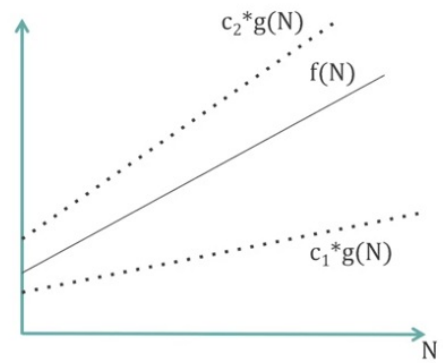
The Theta notation  $\Theta(x)$  defines a **single function** that acts as both an **upper and lower bound** for  $(T(N))$ . Formally defined as:

$T(N)$  is  $\Theta(g(N))$  if there exist positive constants  $c_1, c_2$  and  $n_o$  such that both those conditions hold true:

$$\begin{aligned} T(N) &\geq c_1 \times g(N) \text{ for all } N > n_0 \\ T(N) &\leq c_2 \times g(N) \text{ for all } N > n_0 \end{aligned}$$

Alternatively:

$$c_1 \times g(N) \leq T(N) \leq c_2 \times g(N) \text{ for all } N > n_0$$



As already noted, Theta notation has **only one function**.