Database systems I

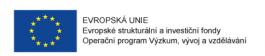
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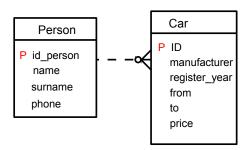
Outline

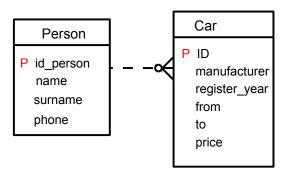
- Introduction
- Conceptual modeling situations
 - Codebooks
 - Historical data
 - Tree and graph data
 - Fact table

Introduction

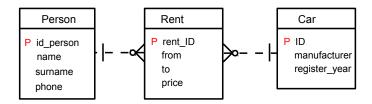
- We described the conceptual modeling tools perspective in the previous lecture
- However, knowledge about the modeling tools is not enough
- We need to
 - learn a typical modeling situations,
 - and we need to to create several database conceptual models by ourself

- Let us have a task: We need to keep track of borrowing information for customers in a car rental company
- What problems can you observe in a following solution?





- What exactly is one entity of the Car entity type?
- The entity is one car or it is a borrow?
- What about duplicity? Can we have a car in many entities?



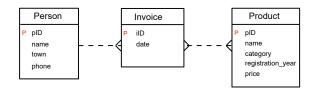
This solution avoid the problems of the previous design

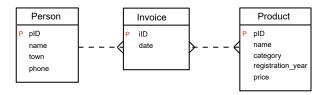
- Entity type Car in the second model is very static and contains limited number of records
- We call such entity type as a Codebook
- Questions that we should answer for each entity type in a database design:
 - What is one record in the relation?
 - Is the name of entity type appropriate?
 - Is there a redundancy?

- Codebook is a entity type that describes some categories
- The most significant feature of a codebook is that we do not expect many DML operations on a result table
- Other examples of codebooks:
 - Car model in a car rent IS
 - Types of payment in a e-shop
 - List of towns or countries
 - List of athletic clubs
 - and so on

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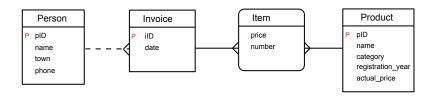
- Let us have a task: We would like to store information about customers, products and purchases. Each purchase belongs to an invoice. The invoices must also be retrospectively accessible (we are interested in a history of purchases).
- What problems can occur?





- There are basically two ways how to solve it:
 - We never update values like price but we insert new records

This solution can be fine if we do not update the price very rarely



- There are basically two ways how to solve it:
 - We never update values like price but we insert new records
 - We can add price into M:N relationship

Much better solution in a case of frequent price updates

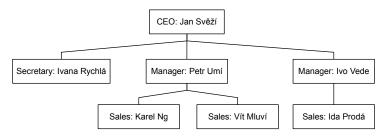
- What about the other attributes?
- We may ask whether any change in the person's phone (or address) should be reflected in all invoices of that person, or whether the old invoices should remain unchanged
- Obviously, the problem is related to whether it is necessary to preserve the history of the entities or whether the current status is sufficient

Historic Data - Other Examples

- We have a car rental company are we interested about car renting history for each car, or is it enough to know who is currently renting a car?
- Let us have an building monitoring application is it sufficient to have just an up-to-date list of people in the building, or do we need these lists retrospectively?
- Let us have an application for a pool lockers is it sufficient to have just an up-to-date list of used lockers, or do we need these lists retrospectively?

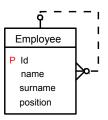
Hierarchical Data

- Let us have a task: We would like to store an information about a fact that an employee can have just one boss
- In other words, we would like to store the following hierarchy

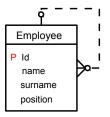


Hierarchical Data

- There is many different approaches to store and query tree data
- In the relational database we can use the following data model:



Hierarchical Data



- Why relationship is obligatory from both sides?
- This representation has its limitations:
 - It can be inefficient to work with large data
 - The number of self-joins during a tree traversal is equal to the depth of traversal

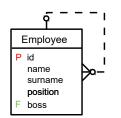
- Find all employees working under 'Peter Pan'
- We may start with the following query

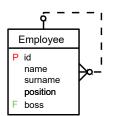
```
SELECT e2.* FROM Employee e1

JOIN Employee e2 ON e1.id = e2.boss

WHERE e1.name = 'Peter'

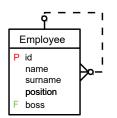
and e2.surname = 'Pan'
```





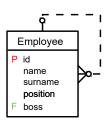
- Find all employees working under 'Peter Pan'
- Then use UNION to add another level

```
SELECT e2.* FROM Employee e1
JOIN Employee e2 ON e1.id = e2.boss
WHERE e1.name = 'Peter'
and e2.surname = 'Pan'
UNION
SELECT e3.* FROM Employee e1
JOIN Employee e2 ON e1.id = e2.boss
JOIN Employee e3 ON e2.id = e3.boss
WHERE e1.name = 'Peter'
and e2.surname = 'Pan'
```



- Find all employees working under 'Peter Pan'
- and another one ...

```
SELECT e2.* FROM Employee e1
JOIN Employee e2 ON e1.id = e2.boss
WHERE e1.name = 'Peter'
and e2.surname = 'Pan'
UNION
SELECT e3.* FROM Employee e1
JOIN Employee e2 ON e1.id = e2.boss
JOIN Employee e3 ON e2.id = e3.boss
WHERE e1.name = 'Peter'
and e2.surname = 'Pan'
UNION
SELECT e4.* FROM Employee e1
JOIN Employee e2 ON e1.id = e2.boss
```



- Find all employees working under 'Peter Pan'
- Generally we need a recursive query

```
WITH rcte AS (
    SELECT e2.* FROM Employee e1
    JOIN Employee e2 ON e1.id = e2.boss
    WHERE e1.name = 'Peter'
    and e2.surname = 'Pan'
    UNION
    SELECT e2.* FROM rcte e1
    JOIN Employee e2 ON e1.id = e2.boss
)

SELECT * FROM rcte
```

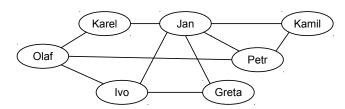
 Which may be quite difficult task (not only for a SQL developer)

Hierarchical data - Other Examples

- Data of a genealogical tree
- Product categories in an internet shop

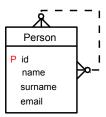
Graph Data

- Let us have a task: We need to store informations about persons and a fact that two people knows each other
- The goal is to store a data having a graph structure



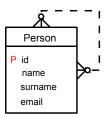
Graph Data

- The solution is very similar to the hierarchical data model
- We need a M:N recursive relationship
- The relationship is obligatory if we allow isolated nodes in a graph



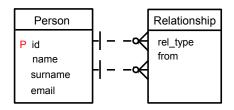
Graph Data with Labeled Edges

- We may want to also store information about the edges of the graph
- Our previous example uses graph with unlabeled graphs
- How to extend this model?



Graph Data with Labeled Edges

 The solution is to do a M:N relationship decomposition and add the necessary information about edges into the binding table



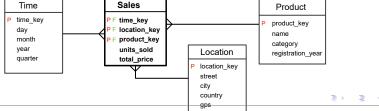
Graph Data

- The data model has its limitations (similarly to hierarchical data model)
- For example let us consider the following queries:
 - Shortest path between two vertexes
 - Centrality of an vertex
 - Find any path between vertexes where the edges has labels from set L
- Most of these queries are extremely difficult to express using SQL (we need recursive queries)
- Moreover their query processing may be slow

Graph Data - Other Examples

- Transportation network (MHD, Trains, ...)
- Mutual matches between sport clubs
- Function calls in a program

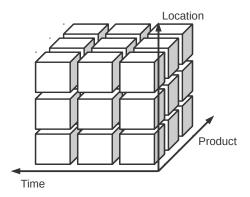
- This schema is often used in a data warehouses/business intelligence
- The main purpose of data warehouse is to provide a tools to analyze data
- The data are organized into two types of tables
 - Fact table Central table that refers to the dimension tables
 - Dimension table 'Codebook' tables surrounding the fact table



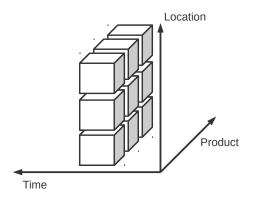
- The data in a data warehouse are usually transfered from a operational database and they are read-only
- The processes responsible for the transfer are called Extract Transform Load (ETL)



• The data in a data warehouse are usually modeled as cubes



- The data in a data warehouse are usually modeled as cubes
- The queries are subsequently modeled as a slices in the cube



Data Warehouse Database Systems

- Columns datastores database systems designed especially for these types of data and queries
- The data are stored according to the columns in a data storage
- It makes aggregate queries very efficient

Reference

• http://dbedu.cs.vsb.cz