

Numerical Methods in Engineering (Solution of Linear Equations-2)

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CH-2-16(MO) Numerical Methods, Lecture 5 Set of Linear Equations-II

- ❑ Began the Solution of Linear equations
 - ❑ The motivation was from several engineering applications
 - ❑ Understood the definitions of diagonal, tri-diagonal, upper triangular and lower triangular matrices
 - ❑ Understood the logic for solution of equations when the coefficient matrix is either diagonal, upper triangular or lower triangular

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Agenda for Today

- ❑ Understand Direct Solvers
- ❑ Non-Iterative
- ❑ Subject to Round-off errors
- ❑ Used for a smaller system ($N < 10$), unless the Matrix is conditioned
- ❑ Involves reduction of the coefficient matrix to a diagonal matrix or an upper triangular matrix, or a lower triangular matrix.
- ❑ This followed by the sweeps discussed earlier

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Cramer's Rule-I

$$\begin{bmatrix} 3 & -1 & 2 \\ 1 & 2 & 3 \\ 2 & -2 & -1 \end{bmatrix} \begin{Bmatrix} x_1 \\ x_2 \\ x_3 \end{Bmatrix} = \begin{Bmatrix} 12 \\ 11 \\ 2 \end{Bmatrix} \quad \text{Sol} = \begin{Bmatrix} 3 \\ 1 \\ 2 \end{Bmatrix}$$

Solution

$$x_1 = \frac{\begin{vmatrix} 12 & -1 & 2 \\ 11 & 2 & 3 \\ 2 & -2 & -1 \end{vmatrix}}{\begin{vmatrix} 3 & -1 & 2 \\ 1 & 2 & 3 \\ 2 & -2 & -1 \end{vmatrix}} \quad x_2 = \frac{\begin{vmatrix} 3 & 12 & 2 \\ 1 & 11 & 3 \\ 2 & 2 & -1 \end{vmatrix}}{\begin{vmatrix} 3 & -1 & 2 \\ 1 & 2 & 3 \\ 2 & -2 & -1 \end{vmatrix}}$$

CH-2-16(MO) Numerical Methods, Lecture 5 Set of Linear Equations-II

Cramer's Rule-II

$$x_3 = \frac{\begin{vmatrix} 3 & -1 & 12 \\ 1 & 2 & 11 \\ 2 & -2 & 2 \end{vmatrix}}{\begin{vmatrix} 3 & -1 & 2 \\ 1 & 2 & 3 \\ 2 & -2 & -1 \end{vmatrix}}$$

- ❑ Number of multiplication operations $\sim (n-1)(n+1)!$
- ❑ For $n = 10$, $M \sim 3.6 \times 10^8$
- ❑ For $n = 100$, $M \sim 10^{157}$

Severe Round
Off Error

Gauss Elimination-I

$$\begin{bmatrix} 3 & -1 & 2 \\ 1 & 2 & 3 \\ 2 & -2 & -1 \end{bmatrix} \begin{Bmatrix} x_1 \\ x_2 \\ x_3 \end{Bmatrix} = \begin{Bmatrix} 12 \\ 11 \\ 2 \end{Bmatrix} \quad Sol = \begin{Bmatrix} 3 \\ 1 \\ 2 \end{Bmatrix}$$

Augmented Matrix

$$\begin{bmatrix} 3 & -1 & 2 & 12 \\ 1 & 2 & 3 & 11 \\ 2 & -2 & -1 & 2 \end{bmatrix}$$

Gauss Operations

Augmented Matrix

$$\begin{bmatrix} 3 & -1 & 2 & 12 \\ 1 & 2 & 3 & 11 \\ 2 & -2 & -1 & 2 \end{bmatrix} \Rightarrow \begin{bmatrix} 3 & -1 & 2 & 12 \\ 0 & 7 & 7 & 21 \\ 0 & -4 & -7 & -18 \end{bmatrix}$$

$3R_2 - R_1$
 $3R_3 - 2R_1$

$$\Rightarrow \begin{bmatrix} 3 & -1 & 2 & 12 \\ 0 & 7 & 7 & 21 \\ 0 & 0 & -21 & -42 \end{bmatrix} \Rightarrow Sol = \begin{Bmatrix} 3 \\ 1 \\ 2 \end{Bmatrix}$$

$7R_3 + 4R_2$

Comments on Gauss Elimination

- ❑ Involves $\frac{1}{3} N^3 - N^2 - \frac{1}{3} N$ (approx) Multiplicative operations
- ❑ For $N=10$, This is eq ~ 430
- ❑ Errors propagate and hence, not used for $N > 10$
- ❑ One should always check for Residues
- ❑ Pivoting and Scaling reduces error propagation

Pivoting and Scaling

- ❑ Diagonal dominance reduces error propagation
- ❑ Rearranging the coefficient to get diagonal dominance is called pivoting
- ❑ Switching of only rows to improve diagonal dominance is called partial pivoting
- ❑ Dividing all the coefficients by the maximum value in a row is called scaling
- ❑ These can easily be incorporated and is always recommended

Gauss Jordan Elimination

Augmented Matrix

$$\begin{bmatrix} 3 & -1 & 2 & 12 \\ 1 & 2 & 3 & 11 \\ 2 & -2 & -1 & 2 \end{bmatrix} \Rightarrow \begin{array}{l} R_1^* = R_1/3 \\ R_2 - R_1^* \\ R_3 - 2R_1^* \end{array} \begin{bmatrix} 1 & -1/3 & 2/3 & 4 \\ 0 & 7/3 & 7/3 & 7 \\ 0 & -4/3 & -7/3 & -6 \end{bmatrix}$$

$$R_2^{**} = 3R_2/7$$

$$\begin{bmatrix} 1 & -1/3 & 2/3 & 4 \\ 0 & 1 & 1 & 3 \\ 0 & -4/3 & -7/3 & -6 \end{bmatrix} \begin{array}{l} R_1 - R_2^{**}/(-3) \\ R_3 - 4R_2^{**}/(-3) \end{array} \begin{bmatrix} 1 & 0 & 1 & 5 \\ 0 & 1 & 1 & 3 \\ 0 & 0 & -1 & -2 \end{bmatrix}$$

Comments on Gauss Jordan Elimination

$$\begin{array}{l} R_1 - R_3^* \\ R_2 = R_3^* \\ R_3^* = R_3/-1 \end{array} \begin{bmatrix} 1 & 0 & 0 & 3 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 2 \end{bmatrix}$$

- ❑ $M_{GJ} = 0.5N^3 + N^2 - 0.5N$, For $N=10$, $M_{GJ}=595$
- ❑ Requires 50% more effort than that for Gauss procedure
- ❑ Procedure is useful for getting inverse

Towards Solution of Linear Equations-III

$$\begin{bmatrix} 3 & -1 & 2 & 1 & 0 & 0 \\ 1 & 2 & 3 & 0 & 1 & 0 \\ 2 & -2 & -1 & 0 & 0 & 1 \end{bmatrix}$$

$$\Rightarrow \begin{bmatrix} 1 & 0 & 0 & -0.5714 & 0.7143 & 1 \\ 0 & 1 & 0 & -1 & 1 & 1 \\ 0 & 0 & 1 & 0.8571 & -0.5714 & -1 \end{bmatrix}$$

L-U Decomposition-I

- If the problem has to be repeated with several source vectors for the same coefficient vector, L-U decomposition is recommended

$$[A] = [L][U]$$

- Such a decomposition speeds up calculation
- In general two methods are available
 - Crout's Decomposition
 - Dolittle's Decomposition

Crout's Decomposition

$$\begin{bmatrix} l_{11} & 0 & 0 \\ l_{21} & l_{22} & 0 \\ l_{31} & l_{32} & l_{33} \end{bmatrix} \begin{bmatrix} 1 & u_{12} & u_{13} \\ 0 & 1 & u_{23} \\ 0 & 0 & 1 \end{bmatrix} = \begin{matrix} \textcircled{1} & \textcircled{3} & \textcircled{5} \\ \textcircled{2} & & \\ \textcircled{4} & & \end{matrix} \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{bmatrix}$$

- $a_{11} = l_{11}$, $a_{21} = l_{21}$, $a_{31} = l_{31}$
- $a_{12} = l_{11} u_{12}$, $a_{13} = l_{11} u_{13}$
- $a_{22} = l_{21} u_{12} + l_{22}$, $a_{32} = l_{31} u_{12} + l_{32}$
- $a_{23} = l_{21} u_{13} + l_{22} u_{23}$
- $a_{33} = l_{31} u_{13} + l_{32} u_{23} + l_{33}$

Logic for Crout's Method

$$\begin{aligned} l_{i1} &= a_{i1}, \text{ for } i = 1, n \\ u_{1j} &= a_{1j} / l_{11}, \text{ for } j = 2, n \\ \text{for } j &= 2, n-1 \\ l_{ij} &= a_{ij} - \sum_{k=1}^{j-1} l_{ik} u_{kj} \quad (\text{for } i = j, n) \\ u_{ji} &= \left(a_{ji} - \sum_{k=1}^{j-1} l_{jk} u_{ki} \right) / l_{jj} \quad (\text{for } i = j+1, n) \\ l_{nn} &= a_{nn} - \sum_{k=1}^n l_{nk} u_{kn} \end{aligned}$$

Solution for Crout's Method

$$[A]\{x\} = \{b\} \Rightarrow [L][U]\{x\} = \{b\}$$

- Introducing $[U]\{x\} = \{d\}$ 1
- This implies $[L]\{d\} = \{b\}$ 2
- Since [L] and {b} are known, {d} can be found from Eq.(2) by forward sweep
- Once {d} is found out, {x} can be found from Eq. (1) by backward sweep

Comments on Crout's Method-I

- ☐ $M_{\text{Crout}} = M_{\text{Gauss}}$
- ☐ But, back substitution $M = n^2 - n$
- ☐ Therefore for a large set one may save substantial effort (n^3 vs n^2)
- ☐ It is possible to store the coefficients of [L] and [U] in [A] itself as [A] is no longer required. This saves memory
- ☐ Other decompositions are similar

Comments on Crout's Method-II

- ☐ It is possible to store L and U in A itself and conserve memory and logic written accordingly

$$\begin{bmatrix} l_{11} & u_{12} & u_{13} \\ l_{21} & l_{22} & u_{23} \\ l_{31} & l_{32} & u_{33} \end{bmatrix}$$

- ☐ Indices have to be carefully addressed
- ☐ Since memory is cheap, this no longer may be required