Answers Exam Functional Programming – December 3rd 2014

- 1. $(5 \times 2 = 10 \text{ points})$
 - (a) What is the type of the following Haskell function wtel?

```
wtel [] = []
wtel (x:xs) = if x == [] then wxs else x:wxs
where wxs = wtel xs
```

```
Answer: wtel :: Eq a => [[a]] -> [[a]]
```

(b) What is the type of the following Haskell function cl?

```
cl ps = ps ++ [(p,s) | (p,q) <- ps, (r,s) <- ps, q==r]

Answer: cl :: Eq a => [(a, a)] -> [(a, a)]
```

(c) What is the type of the standard Haskell indexing operator !! (as an example [0..10]!!3 = 3)?

```
Answer: (!!) :: [a] -> Int -> a
```

(d) What is the type of the following Haskell function map 2?

```
map2 f [] [] = []
map2 f (x:xs) (y:ys) = (f x y) : map2 f xs ys

Answer: map2 :: (a -> b -> c) -> [a] -> [b] -> [c]
```

(e) What is the type of the following Haskell function tw?

```
tw = (\f -> (\x -> (f.f) x))

Answer: tw :: (a -> a) -> a -> a
```

2. (15 points) Consider a positive integer N. We denote its decimal digits by X_0 , X_1 , ..., X_k . The number N is called a *funny number* if you can select at most three (but at least one) of its digits such that N is a divisor of the number $(X_0 + X_1 + ... + X_k - S)^S$, where S is the sum of the selected digits. As an example, 1458 is a funny number since $((1+4+5+8)-(1+5))^{1+5}=12^6=2985984$ is divisible by 1458. Note that we selected the two digits 1 and 5.

Write a Haskell function isFunny (including its type) that takes an integer number as its argument, and returns True if and only if this argument is a funny number.

Solution:

3. (3+3+4=10 points)

• Use a list comprehension to define a function inverse:: [(a, b)]->[(b, a)] such that elem (x,y) ps if and only if elem (y,x) (inverse ps).

```
Answer: inverse ps = [(b,a) | (a,b) < -ps]
```

• Use a list comprehension to make your own implementation of the standard Haskell function replicate. The call replicate n x yields a list of length n with x being the value of every element. So, replicate 5 'a' returns "aaaaa".

```
Answer: replicate n \times = [x \mid k \leftarrow [1..n]]
```

• Define a function doubleReverse which takes a list of strings as its argument and reverses each element of the list and then reverses the resulting list. The implementation of doubleReverse must use a list comprehension. As an example, doubleReverse ["hello", "world"] = ["dlrow", "olleh"].

```
Answer: doubleReverse xss = reverse [reverse xs | xs <- xss]
```

4. (3+3+4=10 points)

• The function powers n returns the infinite list $[n^0, n^1, n^2, n^3, ...]$. Give a *recursive* Haskell implementation (including its type) of the function powers.

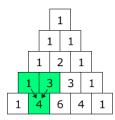
Answer:

```
powers :: Integer -> [Integer]
powers n = 1 : map (n *) (powers n)
```

• The sequence a_k is defined as follows: $a_0 = 1$, $a_1 = 2$, $a_k = 3a_{k-1} + 2a_{k-2}$ for integer k > 1. Define the infinite list seqa, which is the list $[a_0, a_1, a_2, a_3, a_4, ...]$, so seqa!! k should yield a_k .

```
Answer: seqa = 1 : 2 : zipWith (+) (map (*2) seqa) ((map (*3)) (tail seqa))
```

• In the following figure you see the first 5 rows of Pascal's triangle:



To build the triangle, start with the row [1] at the top (we call this row 0), then continue placing numbers below it in a triangular pattern. Each row consists of elements that are the sum of the two numbers above it (except for the boundaries, which are all 1). In the figure, it is highlighted that the 4 in row 4 is obtained by adding the numbers 1 and 3 from row 3.

Give a definition of the infinite list pascalTriangle :: [[Integer]], such that pascalTriangle!!n yields the nth row of Pascal's triangle (i.e. pascalTriangle!!4 = [1, 4, 6, 4, 1]).

Answer:

```
pascalTriangle = iterate nextRow [1]
where nextRow row = zipWith (+) ([0] ++ row) (row ++ [0])
```

5. (15 points) The abstract data type (ADT) Set tp implements a data type for the storage of *sets* of the type tp, where tp is of the class Ord (i.e. the elements are ordered).

Implement a module Set that exports the ADT Set. You can choose a concrete implementation yourself, however this implementation must be hidden from the user of this module.

The following operations on the data type Set must be implemented:

- empty returns an empty set.
- isEmpty returns True for an empty set, otherwise False.
- insert: returns the set after insertion of an element.
- delete: returns the set after removal of an element.
- union: returns the union of two sets.
- intersection: returns the intersection of two sets.

Answer:

```
module Set (Set, empty, isEmpty, insert, delete, union, intersection) where
data Set a = S [a]
empty = S[]
isEmpty (S xs) = null xs
insert x (S xs) = S (ins x xs)
delete x (S xs) = S (del x xs)
  where
    del x [] = []
    del x (y:ys)
      | x < y
      | x < y = y:
| x == y = ys
                  = y: (del x ys)
      | otherwise = y:ys
union (S xs) (S [])
                    = (S xs)
union (S \times S) (S (y:yS)) = union (S (ins y xS)) (S yS)
intersection (S xs) (S [])
                                    = S []
intersection (S []) (S ys)
                                    = S []
intersection (S (x:xs)) (S (y:ys))
                                    = intersection (S xs) (S (y:ys))
  | x < y
  | x > y
                                    = intersection (S (x:xs)) (S ys)
  | otherwise
                                    = insert x (intersection (S xs) (S ys))
-- Note: ins is not exported
ins x [] = [x]
ins x (y:ys)
 | x < y
            = x:y:ys
           = y:ys
  | x == y
  | otherwise = y:(ins x ys)
```

6. (15 points) Given are the following Haskell definitions of the functions f and g:

```
f:: Integer -> Integer
f 0 = 0
f 1 = 1
f n = 5*(f (n-1)) - 6*(f (n-2))

g:: Integer -> Integer -> Integer
g n 0 = 1
g n e = n*(g n (e - 1))
```

Prove for all natural numbers n: f n = g 3 n - g 2 n

Answer: It is easy to see that $g n e = n^e$. We start by proving this first:

```
Base case (e=0): g n 0 = 1 = n^0
Inductive step: g n (e+1) = n*(g n e) = n*n^e = n^{e+1} QED.
```

So, we need to prove: $f n = 3^n - 2^n$.

OED.

7. (15 points) Given are the definitions of the Haskell functions sum, and reverse:

```
sum :: [Integer] -> Integer
sum [] = 0
sum (x:xs) = (sum xs) + x

reverse :: [a] -> [a]
reverse [] = []
reverse (x:xs) = reverse xs ++ [x]
```

Prove that sum (reverse xs) = sum xs for all finite lists xs.

Answer:

```
-- def reverse
sum (reverse [])
                 = sum []
sum (reverse (x:xs)) = sum (reverse xs ++ [x]) -- def reverse
                   = sum (reverse xs) + sum [x] -- Lemma below
                   = sum (reverse xs) + x -- def sum
                                              -- inductive hypothesis
                   = x + sum xs
                                              -- definition of sum
                   = sum (x:xs)
Lemma: sum (xs ++ ys) == sum xs + sum ys
sum ([] ++ ys)
                 = sum ys
                                   -- def (++)
                 = 0 + sum ys -- identity of addition
                 = sum [] ++ sum ys -- def sum
sum ((x:xs) ++ ys) = sum (x : (xs ++ ys)) -- def (++)
                 = x + sum (xs ++ ys) -- def sum
                 = x + sum xs + sum ys -- inductive hypothesis
                 = sum (x:xs) + sum ys -- def sum
```

OED.