Tutorial 5

```
Chapter 9: ex. 5, 6, 7, 8, 10, 13
Chapter 11: ex. 25, 26, 29, 30, 31, 34, and 35
Function definitions:
iter :: Int \rightarrow (a \rightarrow a) \rightarrow (a \rightarrow a)
iter n f
  | n > 0 = f \cdot iter (n-1) f
   I otherwise = id
reverse :: [a] \rightarrow [a]
reverse [] = []
reverse (x:xs) = reverse xs ++ [x]
```

Tutorial 5

```
elem :: a \rightarrow [a] \rightarrow a
elem y [] = False
elem y (x:xs) = if y == x then True else elem y xs
length :: [a] \rightarrow Int
length [] = 0
length (x:xs) = 1 + length xs
(++) :: [a] \rightarrow [a] \rightarrow [a]
(++) [] [] = []
(++) [] (y:ys) = y : ((++) [] ys)
(++) (x:xs) ys = x : ((++) xs ys)
```

Tutorial 5

```
map f [] = []
map f (x:xs) = f x : map f xs
abs x
   | x < 0 = abs(-x)
   | otherwise = x
signum x
   | x < 0 = -1
   | x == 0 = 0
   | otherwise = 1
```

to prove for all finite lists xs and ys:

```
sum (xs ++ ys) = sum xs + sum ys
where
sum [] = 0
sum (x:xs) = x+sum xs
[] ++ ys = ys
(x:xs)++ys = x:(xs++ys)
```

```
sum ([] ++ ys)
```

```
sum ([] ++ ys)
= sum (ys)
```

```
sum ([] ++ ys)
= sum (ys)
= 0 + sum ys
```

```
sum ([] ++ ys)
= sum (ys)
= 0 + sum ys
= sum [] + sum ys
```

```
sum ([] ++ ys)
= sum (ys)
= 0 + sum ys
= sum [] + sum ys
Ind:
sum ((x:xs) ++ ys)
```

```
sum ([] ++ ys)
= sum (ys)
= 0 + sum ys
= sum [] + sum ys
Ind:
sum ((x:xs) ++ ys)
= sum (x:(xs+ys))
```

```
sum ([] ++ ys)
= sum (ys)
= 0 + sum ys
= sum [] + sum ys

Ind:
sum ((x:xs) ++ ys)
= sum (x:(xs+ys))
= x + sum(xs++ys)
```

```
sum ([] ++ ys)
= sum (ys)
= 0 + sum ys
= sum [] + sum ys

Ind:
sum ((x:xs) ++ ys)
= sum (x:(xs+ys))
= x + sum(xs++ys)
= x + sum xs + sum ys
```

Base: sum ([] ++ ys) = sum (ys) = 0 + sum ys = sum [] + sum ys Ind: sum ((x:xs) ++ ys) = sum (x:(xs+ys)) = x + sum(xs++ys)

= x + sum xs + sum ys

= sum(x:xs) + sum ys

- 1. to prove for all finite lists xs + + [] = xs
- 2. to prove for all finite lists xs: (xs ++ ys) ++ zs = xs ++ (ys ++ zs)

```
to prove for all finite lists xs : xs ++ [] = xs
```

Base

```
xs = []: [] ++ [] = []
```

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```

Base

```
xs = []: [] ++ [] = []
```

Ind.

```
(x:xs) ++ []
```

```
to prove for all finite lists xs : xs ++ [] = xs
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Base

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xs = []: [] ++ [] = []
Ind.
(x:xs) ++ []
= x:(xs ++ [])
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```
to prove for all finite lists xs: xs ++ [] = xs
```

Base

```
xs = []: [] ++ [] = []
Ind.
(x:xs) ++ []
= x:(xs ++ [])
= x:xs
QED.
```

to prove for all finite lists xs: (xs ++ ys) ++ zs = xs ++ (ys ++ zs)

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```
([] ++ ys) ++ zs
= ys ++ zs
= [] ++ (ys ++ zs)
```

to prove for all finite lists xs: (xs ++ ys) ++ zs = xs ++ (ys ++ zs)

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```
([] ++ ys) ++ zs
= ys ++ zs
= [] ++ (ys ++ zs)
Ind:
((x:xs) ++ ys) ++ zs
= (x :~(xs ++ ys)) ++ zs
```

to prove for all finite lists xs: (xs ++ ys) ++ zs = xs ++ (ys ++ zs)

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([] ++ ys) ++ zs
= ys ++ zs
= [] ++ (ys ++ zs)
Ind:
((x:xs) ++ ys) ++ zs
= (x :~(xs ++ ys)) ++ zs
= x :~((xs ++ ys) ++ zs)
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to prove for all finite lists xs: (xs ++ ys) ++ zs = xs ++ (ys ++ zs)

Base:

([] ++ ys) ++ zs

= ys ++ zs

= [] ++ (ys ++ zs)

Ind:

((x:xs) ++ ys) ++ zs

= (x :~(xs ++ ys)) ++ zs

= x :~((xs ++ ys) ++ zs)
```

 $= x : \sim (xs ++ (ys ++ zs))$

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Base:
([] ++ ys) ++ zs
= ys ++ zs
= [] ++ (ys ++ zs)
Ind:
((x:xs) ++ ys) ++ zs
= (x : \sim (xs ++ ys)) ++ zs
= x : \sim ((xs ++ ys) ++ zs)
= x : \sim (xs ++ (ys ++ zs))
```

= (x:xs) ++ (ys ++ zs)

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= (x : \sim (xs ++ ys)) ++ zs
= x : \sim ((xs ++ ys) ++ zs)
= x : \sim (xs ++ (ys ++ zs))
= (x:xs) ++ (ys ++ zs)
= x : \sim (xs ++ (ys ++ zs))
```

- 1. to prove for all finite lists xs: sum(reverse xs) = sum xs
- 2. to prove for all finite lists xs: length (reverse xs) = length xs

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Ind:
sum (reverse (x:xs))
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1. to prove for all finite lists xs: sum(reverse xs) = sum xs

```
xs = []: sum(reverse []) = sum []
Ind:
sum (reverse (x:xs))
= sum (reverse xs ++ [x]) -- def reverse
```

1. to prove for all finite lists xs: sum(reverse xs) = sum xs

```
xs = []: sum(reverse []) = sum []
Ind:
sum (reverse (x:xs))
= sum (reverse xs ++ [x]) -- def reverse
= sum (reverse xs) + sum [x] -- sum lemma below
```

1. to prove for all finite lists xs: sum(reverse xs) = sum xs

```
xs = []: sum(reverse []) = sum []

Ind:
sum (reverse (x:xs))
= sum (reverse xs ++ [x]) -- def reverse
= sum (reverse xs) + sum [x] -- sum lemma below
= sum (reverse xs) + x -- def sum
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= sum (reverse xs ++ [x]) -- def reverse
= sum (reverse xs) + sum [x] -- sum lemma below
= sum (reverse xs) + x -- def sum
= x + sum (reverse xs) -- commutativity +
```

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xs = []: sum(reverse []) = sum []

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sum (reverse (x:xs))

= sum (reverse xs ++ [x]) -- def reverse

= sum (reverse xs) + sum [x] -- sum lemma below

= sum (reverse xs) + x -- def sum

= x + sum (reverse xs) -- commutativity +

= x + sum xs -- inductive hypothesis
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1. to prove for all finite lists xs: sum(reverse xs) = sum xs

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= sum (reverse xs ++ [x]) -- def reverse
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= sum (reverse xs) + x -- def sum
= x + sum (reverse xs) -- commutativity +
= x + sum xs -- inductive hypothesis
= sum (x:xs) -- definition of sum
```

```
sum ([] ++ ys)
```

```
sum ([] ++ ys)
= sum ys -- def (++)
= 0 + sum ys -- identity of addition
= sum [] ++ sum ys -- def sum
sum ((x:xs) ++ ys)
= sum (x : (xs ++ ys)) -- def (++)
= x + sum (xs ++ ys) -- def sum
= x + (sum xs + sum ys) -- inductive hypothesis
= (x + sum xs) + sum ys -- associativity +
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= sum(x:xs) + sum ys
```

2. to prove for all finite lists xs: length (reverse xs) = length xs

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Base

xs
= []: length(reverse [])

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= []: length(reverse [])
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2. to prove for all finite lists xs: length(reverse xs) = length xs

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= []: length(reverse [])
= length []

Ind:
length (reverse (x:xs))
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2. to prove for all finite lists xs: length (reverse xs) = length xs Base XS = []: length(reverse []) = length [] Ind: length (reverse (x:xs)) = length (reverse xs ++ [x]) -- def reverse = length (reverse xs) + length [x] -- length lemma below = length (reverse xs) + 1 -- def length

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```
length ([] ++ ys)
```

```
length ([] ++ ys)
= length ys
             -- def (++)
= 0 + length ys -- identity of addition
= length [] ++ length ys -- def length
length ((x:xs) ++ ys)
= length (x : (xs ++ ys)) -- def (++)
= 1 + length (xs ++ ys) -- def length
= 1 + (length xs + length ys) -- inductive hypothesis
= (1 + length xs) + length ys -- associativity +
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length ([] ++ ys)
= length ys
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= 1 + length (xs ++ ys) -- def length
= 1 + (length xs + length ys) -- inductive hypothesis
= (1 + length xs) + length ys -- associativity +
= length(x:xs) + length ys
```

```
elem z (xs ++ ys) = elem z xs | | elem z ys
```

```
elem z (xs ++ ys) = elem z xs || elem z ys

Base

xs = []: elem z ([] ++ ys)
```

```
elem z (xs ++ ys) = elem z xs || elem z ys

Base

xs = []: elem z ([] ++ ys)

= elem z ys = False || elem z ys
```

```
elem z (xs ++ ys) = elem z xs || elem z ys

Base

xs = []: elem z ([] ++ ys)

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= False
```

```
elem z (xs ++ ys) = elem z xs || elem z ys

Base

xs = []: elem z ([] ++ ys)

= elem z ys = False || elem z ys

= False

= elem z [] || elem z ys
```

```
elem z (x:xs ++ ys)
```

```
elem z (x:xs ++ ys)
= elem z (x:(xs ++ ys))
```



```
Ind:
elem z (x:xs ++ ys)
= elem z (x:(xs ++ ys))
  case z==x: = True
                   = True || elem z ys
                   = elem z (x:xs) || elem z ys
   case z/=x: = elem z (xs ++ ys)
                   = elem z xs || elem z ys
                   = elem z (x:xs) || elem z ys
```

to prove for all finite lists xs and defined n:

```
take n xs ++ drop n xs = xs
where
```

```
take 0 _ = [], take _ [] = [], take n (x:xs)=x:take (n-1) xs drop 0 xs = xs, drop n [] = [], drop n (x:xs) = drop (n-1) xs
```

to prove for all finite lists xs and defined n:

```
take n xs ++ drop n xs = xs

where
```

```
take 0 _ = [], take _ [] = [], take n (x:xs)=x:take (n-1) xs

drop 0 xs = xs, drop n [] = [], drop n (x:xs) = drop (n-1) xs

Base case for n==0: take 0 xs ++ drop 0 xs = [] ++ xs = xs
```

to prove for all finite lists xs and defined n:

```
take n xs ++ drop n xs = xs

where

take 0 _ = [], take _ [] = [], take n (x:xs)=x:take (n-1) xs

drop 0 xs = xs, drop n [] = [], drop n (x:xs) = drop (n-1) xs

Base case for n==0: take 0 xs ++ drop 0 xs = [] ++ xs = xs
```

Base case for xs==[]: take n [] ++ drop n [] = [] ++ [] = []

```
to prove for all finite lists xs and defined n:
```

```
take n \times s ++ drop \times n \times s = xs
 where
take 0 _{-} = [], take _{-} [] = [], take n (x:xs)=x:take (n-1) xs
drop 0 xs = xs, drop n [] = [], drop n (x:xs) = drop (n-1) xs
Base case for n==0: take 0 xs ++ drop 0 xs = [] ++ xs = xs
Base case for xs==[]: take n[] ++ drop n[] = [] ++ [] = []
Ind.: for n > 0,
take n (x:xs) ++ drop n (x:xs)
```

```
to prove for all finite lists xs and defined n:
take n \times s ++ drop \times n \times s = xs
 where
take 0 = [], take [] = [], take n (x:xs)=x:take (n-1) xs
drop 0 xs = xs, drop n [] = [], drop n (x:xs) = drop (n-1) xs
Base case for n==0: take 0 xs ++ drop 0 xs = [] ++ xs= xs
Base case for xs==[]: take n[] ++ drop n[] = [] ++ [] = []
Ind.: for n > 0,
take n (x:xs) ++ drop n (x:xs)
= (x:take (n-1) xs) ++ drop (n-1) xs
```

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to prove for all finite lists xs and defined n:
take n \times s ++ drop \times n \times s = xs
 where
take 0 = [], take [] = [], take n (x:xs)=x:take (n-1) xs
drop 0 xs = xs, drop n [] = [], drop n (x:xs) = drop (n-1) xs
Base case for n==0: take 0 xs ++ drop 0 xs = [] ++ xs = xs
Base case for xs==[]: take n[] ++ drop n[] = [] ++ [] = []
Ind.: for n > 0,
take n (x:xs) ++ drop n (x:xs)
= (x:take (n-1) xs) ++ drop (n-1) xs
= x: (take (n-1) xs ++ drop (n-1) xs)
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to prove for all finite lists xs and defined n:
take n \times s ++ drop \times n \times s = xs
  where
take 0 = [], take [] = [], take n (x:xs)=x:take (n-1) xs
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Base case for n==0: take 0 xs ++ drop 0 xs = [] ++ xs = xs
Base case for xs==[]: take n[] ++ drop n[] = [] ++ [] = []
Ind.: for n > 0,
take n (x:xs) ++ drop n (x:xs)
= (x:take (n-1) xs) ++ drop (n-1) xs
= x: (take (n-1) xs ++ drop (n-1) xs)
```

= x:xs

```
to prove for all natural numbers n: fac2 n = fac n where fac 0 = 1, fac n = n*fac (n-1)
```

```
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```

A direct infuctive proof gets stuck. We need to generalize this to

$$facAux n p = p*fac n$$

This is easy to prove by induction on n:

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```

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This is easy to prove by induction on n:

Base:

facAux 0 p

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Base:

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```
to prove for all natural numbers n: fac2 n = fac n where fac 0 = 1, fac n = n*fac (n-1)
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A direct infuctive proof gets stuck. We need to generalize this to

$$facAux n p = p*fac n$$

This is easy to prove by induction on n:

Base:

```
facAux 0 p
= p = p*1
= p*fac 0
```

```
facAux (n+1) p
= facAux n ((n+1)*p)
= (n+1) *p* (fac n)
p*fac (n+1) = (n+1)*p*(fac n)
Now, we can conclude:
```

```
fac2 n
= facAux n 1
= 1*fac n
= fac n
```

```
for all x prove: f.(g.h) x == (f.g).h x
```

```
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```

```
f.(g.h) x
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```

```
f.(g.h) x
= f (g.h x)
```

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for all x prove: f.(g.h) x == (f.g).h x
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f.(g.h) x
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= f (g (h x))
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f.(g.h) x
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```
f.(g.h) x
= f (g.h x)
= f (g (h x))
(f.g).h x
= (f.g) (h x)
```

```
for all x prove: f.(g.h) x == (f.g).h x
```

```
f.(g.h) x
= f (g.h x)
= f (g (h x))
(f.g).h x
= (f.g) (h x)
= f (g (h x)) QED.
```

```
for all f prove: id.f = f
```

```
for all f prove: id.f = f
```

```
(id.f) x
```

```
for all f prove: id.f = f
```

```
(id.f) x
= id (f x)
```

```
for all f prove: id.f = f
```

```
(id.f) x
= id (f x)
= f x
QED.
```

to prove for all natural n: iter n id = id

```
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```

Base:

```
iter 0 id = id (by def.)
```

```
to prove for all natural n: iter n id = id
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Base:

```
iter 0 id = id (by def.)
```

Ind.:

```
iter (n+1) id x
```

```
to prove for all natural n: iter n id = id
```

Base:

```
iter 0 id = id (by def.)
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Ind.:

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iter (n+1) id x
= id . iter n id x
```

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iter (n+1) id x = id . iter n id x = id.id x

Ind.:

```
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```

Base: iter 0 id = id (by def.) Ind.: iter (n+1) id x = id . iter n id x = id.id x = id (idx x))

```
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iter 0 id = id (by def.)

Ind.:
iter (n+1) id x
= id . iter n id x
```

= id.id x

= id x

QED.

= id (idx x))

```
to prove: abs.abs x = abs x
```

```
to prove: signum.signum x = signum x
```

to prove: abs.abs x = abs x

```
to prove: abs.abs x = abs x abs.abs x = abs (abs x)
```

```
to prove: abs.abs x = abs x
abs.abs x = abs(abs x)

1) case x < 0: abs(abs x) = abs (-x)</pre>
```

```
to prove: abs.abs x = abs x
abs.abs x = abs(abs x)

1) case x < 0: abs(abs x) = abs (-x)

2) case x >= 0: abs(abs x) = abs x Q.E.D.
```

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to prove: abs.abs x = abs x
abs.abs x = abs(abs x)

1) case x < 0: abs(abs x) = abs (-x)
2) case x >= 0: abs(abs x) = abs x
Q.E.D.

to prove: signum.signum x = signum x
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to prove: abs.abs x = abs x
abs.abs x = abs(abs x)

1) case x < 0: abs(abs x) = abs (-x)
2) case x >= 0: abs(abs x) = abs x Q.E.D.

to prove: signum.signum x = signum x

1) case x == 0: signum.signum 0 = signum (signum 0) = signum0
```

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to prove: abs.abs x = abs x
abs.abs x = abs(abs x)

1) case x < 0: abs(abs x) = abs (-x)
2) case x >= 0: abs(abs x) = abs x Q.E.D.

to prove: signum.signum x = signum x

1) case x == 0: signum.signum 0 = signum (signum 0) = signum 0
2) case x > 0: signum.signum x = signum 1 = 1 = signum x
```

```
to prove: abs.abs x = abs x
abs.abs x = abs(abs x)

1) case x < 0: abs(abs x) = abs (-x)
2) case x >= 0: abs(abs x) = abs x Q.E.D.

to prove: signum.signum x = signum x

1) case x == 0: signum.signum 0 = signum (signum 0) = signum 0
2) case x > 0: signum.signum x = signum 1 = 1 = signum x
3) case x < 0: signum.signum x = signum (-1) = -1 = signum x</pre>
```

```
to prove: map f (ys++zs) = map f ys ++ map f zs
```

```
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Base:
map f ([]++zs)
```

```
to prove: map f (ys++zs) = map f ys ++ map f zs

Base:
map f ([]++zs)
= map f zs
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to prove: map f (ys++zs) = map f ys ++ map f zs

Base:
map f ([]++zs)
= map f zs
= [] ++ map f zs
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Base:
map f ([]++zs)
= map f zs
= [] ++ map f zs
= map f [] ++ map f zs
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```
to prove: map f (ys++zs) = map f ys ++ map f zs

Base:
map f ([]++zs)
= map f zs
= [] ++ map f zs
= map f [] ++ map f zs
Ind.:
map f ((y:ys)++zs)
```

```
to prove: map f (ys++zs) = map f ys ++ map f zs
Base:
map f([]++zs)
= map f zs
= [] ++ map f zs
= map f [] ++ map f zs
Ind.:
map f((y:ys)++zs)
= f (y:(ys++zs))
```

```
to prove: map f (ys++zs) = map f ys ++ map f zs
Base:
map f([]++zs)
= map f zs
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Ind.:
map f((y:ys)++zs)
= f (y:(ys++zs))
= f y: map f (ys++zs)
```

```
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map f([]++zs)
= map f zs
= [] ++ map f zs
= map f [] ++ map f zs
Ind.:
map f((y:ys)++zs)
= f (y:(ys++zs))
= f y: map f (ys++zs)
= f y: (map f ys ++ map f zs)
```

```
map f (y:ys) ++ map f zs
```

```
map f (y:ys) ++ map f zs
= (f y: (map f ys)) ++ map f zs
```

```
map f (y:ys) ++ map f zs
= (f y: (map f ys)) ++ map f zs
= f y : (map f ys ++ map f zs)
QED.
```

```
to prove: concat (map (map f) xs) = map f (concat xs)
```

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Base:
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= concat ((map f x):map (map f) xs)
= (map f x) ++ concat (map (map f) xs)
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```
to prove: (0<) . (+1) = (0<=)
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Proof: ((0<) . (+1) ) x = (0<) . ((+1) x)
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Proof: ((0<) . (+1) ) x = (0<) ((+1) x)
= (0<) ((\arrowvert a -> a+1) x) = (0<) (x+1)
```

```
to prove: (0<) . (+1) = (0<=)

Proof: ((0<) . (+1) ) x = (0<) ((+1) x)
= (0<) ((\setminus a -> a+1) x) = (0<) (x+1)
= (\setminus a -> 0<a) (x+1) = 0<x+1
```

```
to prove: (0<) . (+1) = (0<=)

Proof: ((0<) . (+1) ) x = (0<) . ((+1) . x)

= (0<) . ((-1) . x) . ((-1) . x)

= (0<) . ((-1) . x) . ((-1) . x)

= (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (-1) . (
```

```
to prove: (0<) . (+1) = (0<=)

Proof:

((0<) . (+1)) x = (0<) ((+1) x)

= (0<) ((\a -> a+1) x) = (0<) (x+1)

= (\a -> 0<a) (x+1) = 0<x+1

(0<=) x

= (\a -> 0<=a) x
```

```
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Proof:

((0<) . (+1)) x = (0<) ((+1) x)

= (0<) ((\a -> a+1) x) = (0<) (x+1)

= (\a -> 0<a) (x+1) = 0<x+1

(0<=) x

= (\a -> 0<=a) x

= 0<=x
```

```
to prove: (0<) . (+1) = (0<=)
Proof:
((0<) \cdot (+1)) \times = (0<) ((+1) \times)
= (0<) ((\advarrange a+1) x) = (0<) (x+1)
= (\a -> 0 < a) (x+1) = 0 < x+1
(0 <=) x
= (\a -> 0 <= a) x
= 0 <= x
= 0 < x+1
```