

Topic 7

Memory Management ***How to gain 30% performance improvement easily***

資料結構與程式設計
Data Structure and Programming

Sep, 2011

Outline

- ◆ Memory related problems
 - Illegal memory address access
 - Memory leaks
 - Fragmentation
 - Performance issues
- ◆ Memory management
 - Basic concept
 - Categorization
 - How to implement

Memory Related Problems

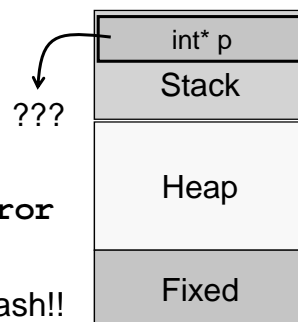
1. Illegal memory address access
2. Memory leaks
3. Fragmentation
4. Performance issues

Illegal Memory Address Access

1. Uninitialized memory read/write
 - Access to the content of a pointer variable that is not yet allocated

```
void f() {  
    MyClass* p;  
    ...  
    int i  
    = p->getData(); // error  
}
```

→ Compilation OK; Runtime crash!!



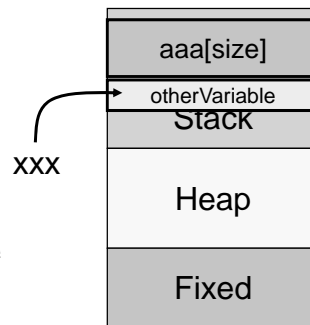
Illegal Memory Address Access

2. Array bound read/write

- Array index is greater than the bound

```
void f() {  
    size_t size;  
    ...  
    size = ...;  
    int aaa[size];  
    ...  
    size_t idx = ...;  
    ...  
    // error if idx >= size  
    aaa[idx] = ...;  
}
```

→ Compilation OK,
but may get strange runtime bug



Illegal Memory Address Access

3. Freed memory read/write

- Access to the just freed memory allocation
- May still get the expected content, but will become garbage when reallocated by others

```
void f() {  
    int* p = new int;  
    cout << p << endl;  
    delete p;  
    // may print out the same address  
    cout << p << endl;  
    *p = 30; // [NOTE] compilation & runtime OK;  
    int j = *p;  
    cout << j << endl;  
    int* q = new int(20);  
    int k = *p;  
    cout << k << endl; // What's the value for k?  
}
```

Illegal Memory Address Access

4. Freeing mismatched memory

- Mixed use of malloc/calloc/free and new/new[]/delete/delete[]

```
▪ int *p = new int(10);  
  int *q = new int[10];  
  int **r = new int*;  
  int **s = new int*[10];  
  delete p or delete []p?  
  delete q or delete []q?  
  delete r or delete []r?  
  delete s or delete []s?
```

Illegal Memory Address Access

5. Doubly freed memory

- Delete the same memory location multiple times

```
▪ int *p = new int(10);  
  int *q = p;  
  delete p;  
  delete q;
```

How to avoid illegal memory access?

1. Allocate and free memory of data members in constructor and destructor

- Use object to wrap the pointer variables

```
class MyClass {
    A* _pp;
public:
    MyClass(int i = 0) { _pp = new A(i); }
    ~MyClass() { delete _pp; }
};

void f() {
    MyClass o; // o._pp is allocated
} // o._pp is deleted automatically
```

- All the operations on _pp should go through class MyClass
 - Can make class A a private class to MyClass (by "friend")
- What about copy constructor or assignment operator?
 - May need "reference count" to avoid double-free error

How to avoid illegal memory access?

2. Paired memory allocation/deletion functions

- Don't allow too many functions to allocate and delete pointers

// [No good] hard to keep track of the memory allocation of _pp

```
class MyClass {
    int* _pp;
public:
    void f1(int i) {
        ...; _pp = new int(i); ... }
    void f2() {
        ...; delete _pp; ... }
    void f3() {
        ...; _pp = new int(j);
        ...; delete _pp; ... }
};
```

How to avoid illegal memory access?

3. Customized array class

- Check index whenever access

```
template <class T>
class MyArray {
    // how many elements in the array
    size_t _size;
    // how much memory is allocated
    size_t _capacity;
    T* _data;
public:
    T& operator [] (size_t i) {
        #ifndef NDEBUG
        if (i >= _size)
            throw ExceptionArraySize(i);
        #endif // NDEBUG
        return _data[i];
    }
};
```

How to avoid illegal memory access?

4. Don't use malloc/calloc/free in C++

- They won't call the constructors/destructors

```
class Temp{
public:
    string c;
};
Temp *test;
int main()
{
    test = (Temp*)malloc(sizeof(Temp));
    cout << test->c << endl; // Garbage...
    ...
}
```

How to avoid illegal memory access?

5. Correctly use of new/new[] and delete/delete[]
6. Memory management

In short, create your own style and strictly abide by your disciplines

What about the overhead generated by the above preventions?

- ◆ Minor overhead is OK; better than debugging tricky memory bugs
 - ◆ Use “#ifndef NDEBUG” to bypass them in optimized mode compilation
 - “Debug build” --- for developer
 - g++ -g xxx.cpp
 - “Optimized build” --- for tool release
 - g++ -O3 -DNDEBUG xxx.cpp
- ```
=====
#ifndef NDEBUG
<codes for debug mode only>
#endif // NDEBUG
=====
```

## Memory Related Problems

1. Illegal memory address access
2. Memory leaks
3. Fragmentation
4. Performance issues

## What is memory leak?

- ◆ Not freeing allocated memory, so as the program runs, the total occupied memory is increasing and cannot be reclaimed
  - ➔ Performance degradation due to thrashing
  - ➔ Program terminated due to memory out



## Why do I have memory leaks?

1. Pointer data members not freed
  - `class A { B *_b; ...; A() { _b = new... };  
A a1; ... A a2 = a1; ...`
2. Local pointers not freed
  - `A *a; a = new ...;  
if (xxx) { ... return; }  
delete a;`
3. Freeing memory mismatch
  - e.g. `p = new MyClass[10]; ...; delete p;`
4. Overwrite on allocated pointer variables

```
void f() {
 int* p = new int;
 ...
 p = g();
 // original p cannot be deleted
 delete p;
}
```

## How do I know if I have memory leak?

- ◆ Well, as the program runs longer, the memory usage is increasing and doesn't seem to saturate.
- ◆ To diagnose
  1. Code review
  2. Using tools
    - Commercial: purify
    - GNU: valgrind (<http://valgrind.org/>)

## How to avoid memory leaks?

- ◆ Good practice makes it all !!
- ◆ Memory management
  - Block allocation and deletion
- ◆ Use “reference count” to keep track whether it is safe to delete an pointer
  - How??
  - $A * p = q;$   
// Who's ref count is incremented by 1?
  - Constructor? Destructor? Object wrapper?

## What is fragmentation?

- ◆ Like the fragmentation in your hard disk, the memory used in your program may have fragmentation too



## What is fragmentation?

- ◆ Like the fragmentation in your hard disk, the memory used in your program may have fragmentation too

- ```
MyClass12Byte* a = new MyClass12Byte ;  
MyClass16Byte* b = new MyClass16Byte;  
MyClassWhatever* c = new MyClassWhatever;  
delete a;  
delete b;  
MyClass16Byte* d = new MyClass16Byte;  
MyClass16Byte* e = new MyClass16Byte;
```

➔ Memory fragmentation of 12 Bytes (where??)

How to avoid memory fragmentation?

- ◆ Memory fragmentation will make your program use more memory than necessary
- ◆ How to fix it?
 - Not easy, unless you use your own memory management and carefully allocate memory pieces with different sizes

Performance Issues

- ◆ Overhead in system calls of memory allocation / deletion
- ◆ What's the runtime difference?
 1.

```
A* a[1 << 20];  
for (int i = 0; i < (1 << 20); i++) {  
    a[i] = new A;  
    *(a[i]) = i;  
}
```
 2.

```
A* a[1 << 20];  
A* b = (A *)calloc(1 << 20, sizeof(int));  
for (int i = 0; i < (1 << 20); i++) {  
    a[i] = b + i;  
    *(a[i]) = i;  
}
```

→ But, will A's constructors be called in the second case?

Basic Concepts of Memory Management

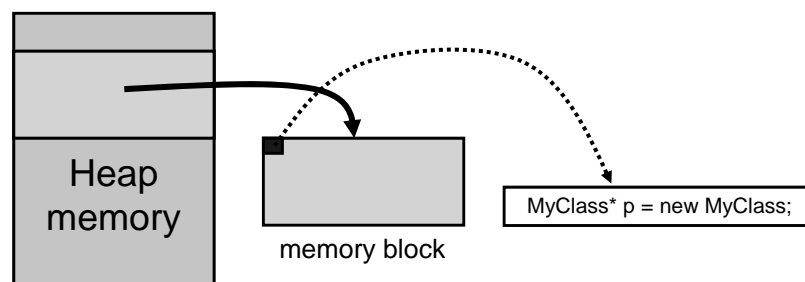
- ◆ Allocate a big chunk of memory from the system at a time
 - Distribute memory to the pointer variables by the memory manager
- ◆ No need to free pointers one by one; free the whole chunk at once
 - Return memory to system when mission is completed
 - Possibly memory leak-free
 - [Optional] Freed pointer memory is recorded in the recycle list (no deletion); can be used for later memory request

Issues about Memory Manager

1. Number of memory blocks
 - Continuous or non-continuous
2. Overload of new/delete operators
 - Use new/delete or customized alloc()/free()
3. Memory manager association (by type or id)
4. Recycle or not
 - Garbage collection?

Memory Blocks in Memory Manager

- ◆ How many memory should we claim from the system each time (i.e. 1 memory block)?
 - Too small: many system calls
 - Too big: waste of memory if not used up
- ➔ Depend on applications, usually 4K – 1MB



1. Continuous Memory Block

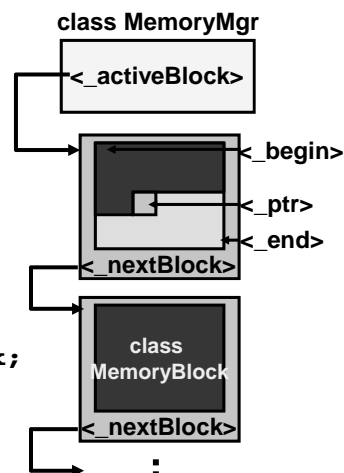
- ◆ Only 1 memory block
 - When `_size >= _capacity`, reallocate a bigger block and copy the original data over
 - Difficult to work with pointer variables (why?)
 - Addresses are continuous; can access by index

```
class MemoryBlock {
    #define S_SIZE_T sizeof(size_t)
public:
    MemoryBlock(size_t B) { // block size = B Bytes
        _begin = _next = (void*)malloc(B);
        _end = _begin + numElm(B);_}
    void* alloc (size_t t) { // t is number of Bytes
        void* tmp = getNext(t);
        if (tmp >= _end) { /* allocate new memory and copy to it */ }
        void* ret = _next; _next = tmp;
        return ret;
    }
private:
    void    *_begin, *_next, *_end;
    void*    getNext(size_t t) const {
        size_t nt = numElm(t); return (size_t*)_next + nt; }
    size_t  numElm(size_t t) const {
        return (t + S_SIZE_T - 1) / S_SIZE_T; }
};
```

2. Non-continuous Memory Block

- ◆ When `size >= capacity`, allocate a new memory block and link it to the previous one

```
class MemoryBlock {
    char*    _begin;
    char*    _ptr;
    char*    _end;
    MemoryBlock* _nextBlock;
};
class MemoryMgr {
    MemoryBlock* _activeBlock;
};
```



Overload of new/delete operators

- ◆ We can overload the new and delete operators of a class
 - `void* operator new (size_t t);`
 - `void* operator new[] (size_t t);`
 - `void operator delete (void* p);`
 - `void operator delete[] (void* p);`(Can also be static functions)

[Note] The parameters 't' and 'p' are passed in by compiler with the "new/delete" calls
- ◆ Advantage
 - Memory manager is transparent to the programmer; can turn on and off easily
- ◆ For more information, please see (for example)
 - <http://www.relisoft.com/book/tech/9new.html>

Example: newOp.cpp

```
class A
{
    int    _a;
    int    _b;
    int    _c;
    short  _d;
    // sizeof(A) = 14 → 16 Bytes

public:
    A() {}
    ~A() {}

    static void* operator new(size_t t) {
        cout << "new (inside A): " << endl;
        cout << ">> size = " << t << endl;
        A* p = (A*)malloc(t);
        cout << ">> ptr = " << p << endl;
        return p;
    }

    static void* operator new[](size_t t) {
        cout << "new[] (inside A): " << endl;
        cout << ">> size = " << t << endl;
        A* p = (A*)malloc(t);
        cout << ">> ptr = " << p << endl;
        return p;
    }

    static void operator delete(void* p) {
        cout << "delete (inside A): " << endl;
        cout << ">> ptr = " << p << endl;
        free(p);
    }

    static void operator delete[](void* p) {
        cout << "delete[] (inside A): " << endl;
        cout << ">> ptr = " << p << endl;
        free(p);
    }
};
```

Example: newOp.cpp

```
int main()
{
    A* a = new A;
    cout << endl;
    cout << "new (in main): " << endl;
    cout << ">> ptr = " << a << endl;
    cout << endl;

    A* b = new A[10];
    cout << endl;
    cout << "new[] (in main): " << endl;
    cout << ">> ptr = " << b << endl;
    cout << endl;

    delete a;
    cout << endl;
    delete []b;
    cout << endl;
}
```

===== Sample output =====

```
new (inside A):
>> size = 16
>> ptr = 0x502010
```

```
new (in main):
>> ptr = 0x502010
```

```
new[] (inside A):
>> size = 168
>> ptr = 0x502030
```

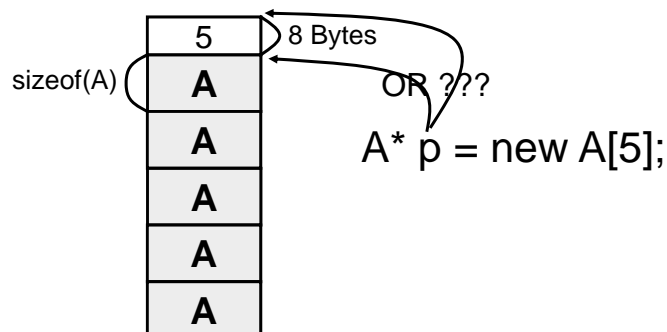
```
new[] (in main):
>> ptr = 0x502038
```

```
delete (inside A):
>> ptr = 0x502010
```

```
delete[] (inside A):
>> ptr = 0x502030
```

What did “newOp.cpp” tell you?

1. Must have destructor... somehow....
 - Try comment out the destructor in newOp.cpp...
2. Size of “new[]” = array size + 8 // for 64-bit platform
 - How do we record the array size for delete?
 - ➔ i.e. delete [] p; // what's the size?
3. Size of class A is promote to 16 Bytes (multiple of SIZE_T)



What did “newOp.cpp” tell you?

1. Must have destructor... somehow....
 - Try comment out the constructor in newOp.cpp...
 2. Size of “new[]” = array size + 8 // for 64-bit platform
 3. Size of class A is promote to 16 Bytes (multiple of SIZE_T)
 4. The ptr in new[] () points to the “-8” address
 5. The ptr in new[] caller points to the array begin
 6. The ptr in delete[] () points to the “-8” address
 7. The ptr in delete[] caller points to the array begin
- In this example, memory is explicitly allocated by “malloc()” (or new char[numBytes])
(Will the constructor and destructor be called?)
- What if we want to use a memory manager (for chunk alloc and delete)?
(Who returns the pointers of new and delete?)
- Can the “new()”, “delete()” be non-static member functions?

Closer look at “new A” and “delete A”

- ◆ A *a = new A;
 1. A::operator new() is called
 2. Constructor of A is called
 3. The return pointer address is copied to ‘a’
- ◆ A *a = new A[10];
 - Similar to “A *a = new A” except that 10 constructors are called
- ◆ delete a;
 1. Destructor of A is called
 2. A::operator delete() is called
- ◆ delete []a;
 - Similar to “delete a” except that several destructors are called

8/4-Byte Aligned

◆ In the previous example, the size of data members in A is $4 + 4 + 4 + 2 = 14$. However...

- `sizeof(A) = 16`
- The parameters to `new()` and `new[10]` are 16 and 168

◆ But, if the class A is changed to:

```
class A {  
    char _data[14];  
};
```

- `sizeof(A) = 14`
 - The parameters to `new()` and `new[10]` are 14 and 148
- ➔ NOT 8-Byte aligned!!

Issues about Memory Manager

1. Number of memory blocks
 - Continuous or non-continuous
2. Overload of `new/delete` operators
 - Use `new/delete` or customized `alloc()/free()`
3. Memory manager association (by type or id)
4. Recycle or not
 - Garbage collection?

Memory Manager Association

- ◆ Which memory manager to call when you allocate a memory in new/delete operator?

(i.e. instead of calling "malloc()" and "free()" directly...)

```
→ void* new(size_t t) {  
    ... memMgr->alloc(t); ...}  
→ void operator delete (void* p) {  
    ... memMgr->free((T*)p);
```

~~Is "memMgr" a data member?~~

Is "memMgr" a global variable?

1. Declared as "static" Data Member

```
class MyClass {  
    static MemoryMgr *const _mem_s;  
public:  
    void* operator new(size_t t) {  
        _mem_s->alloc(t); }  
};
```

- Each class is associated with an unique memory manager
- What if new/delete operators are not overloaded?
- What if we want to associate more than 1 memory managers for a class? (i.e. 1 class → n memMgr)
 - [Reason] Can have options to free portion of the memory
 - Swap with other memory manager (bookkeeping needed)
 - Who control this??

2. Use a Global Map(class id, MemManager);

```
class MyClass {
    static int const _mem_id_s;
public:
    void* operator new(size_t t) {
        ::globalMemMap[_mem_id_s]
        ->alloc(t);
    }
};
```

➔ Memory manager association is controlled by a global function/class

HW#4 Implementation of class MemMgr

```
template <class T>
class MemMgr {
private:
    size_t      _blockSize;
    MemBlock<T>* _activeBlock;
    MemRecycleList<T>
        _recycleList[R_SIZE];
};

template <class T>
class MemBlock {
    friend class MemMgr<T>;
    char*      _begin;
    char*      _ptr;
    char*      _end;
    MemBlock<T>* _nextBlock;
};

template <class T>
class MemRecycleList {
    friend class MemMgr<T>;
    // the array size of the recycled data
    size_t      _arrSize;
    // the first recycled data
    T*          _first;
    // next MemRecycleList
    // with _arrSize + n*R_SIZE
    MemRecycleList<T>* _nextList;
};
```

Using Memory Management

- ◆ Given a class “A” to be managed by class “MemMgr”
 - class A {
 // overload its new/new[]/delete/delete[] operators
 void* operator new(size_t t) {
 return (void*)(_memMgr->alloc(t)); }
 void* operator new[](size_t t);
 void operator delete(void* p);
 void operator delete[](void* p);
 static void memReset(size_t b = 0);
 static void memPrint();
 // Declare _memMgr as a static data member
 static MemMgr* const _memMgr;
};
 - class MemMgr {
 // Implement “alloc”, “allocArr”, “free”, “freeArr”, “print”
 functions
};

Memory Manager Association

- ◆ We know...
 - Static data member must be initialized in .cpp code
 - e.g. MemoryMgr *const A::_mem_s = new ...
- ◆ Can we associate the memory managers of 2 different classes to the same one?
(i.e. n classes → 1 memMgr)
 - i.e. Share the same memory manager
 - Any problem? (Answered later)

Do you remember.... mem manager is

- ◆ Allocate a big chunk of memory from the system at a time
 - Distribute memory to the pointer variables by the memory manager
- ◆ No need to free pointers one by one; free the whole chunk at once
 - Return memory to system when mission is completed
 - ➔ Possibly memory leak-free
 - [Optional] Freed pointer memory is recorded in the recycle list (no deletion); can be used for later memory request

Why do we need to overload
“delete”?

What does it do?

Can we NOT overload “delete”?
(If so, will the destructor be called?)

Recycling Memory

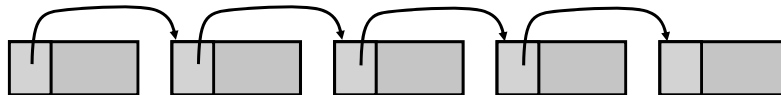
- ◆ Freed pointer memory is recorded in the recycle list (no deletion); can be used for later memory request

How?? Use a “linked list” container class? (No!! extra memory)

1. In the memory manager, keep a recycleList that points to the first recycled memory
2. Reuse the first 4 or 8 Bytes (why?) of each recycled memory, pointing to the address of the next recycled memory

[Restriction]

The size of the managed class should be ≥ 4 (or 8) Bytes



In other words...

```
class MemoryMgr {
    RecycleList _rList;
public:
    void* alloc(size_t t) {
        void* p = _rList.popFront();
        if (P != 0) return p;
        ... // get memory from memory manager
    }
    void delete(void* p) {_rList.pushFront(p); }
};

class MemRecycleList {
    size_t _arrSize; // the array size of the recycle data
    T* _first; // the first recycled data
    MemRecycleList<T>* _nextList;
};

➔ Any problem?
```

Recycling List Implementation

- ◆ [Note] Memory in recycling list is NOT continuous
- ◆ Should the size of the recycled elements in `_rList` be the same? $(3) \rightarrow (1) \rightarrow (5) \rightarrow (1) \dots$
- ◆ If same size \rightarrow Simple implementation
 - `void* _first;` // “`void* _last`” is optional
 - All the elements in the list have the same size (e.g. `= sizeof(A)`)
 - But how do we recycle “`a = new A[n]`”?
 - Just implement `pushFront()` and `popFront()`
 - Don't need to pass in “`size_t`” for `popFront()`
- ◆ If not, how do you find the one you want?
 - For example, “`a = new A[n]`”?

Recycling List with Different Mem Sizes

- ◆ Using linked list?
 - Finding the element of size `S` is $O(n)$
- ◆ Using `map<size, linked list>`?
 - Uh... extra memory
 - $O(\log(m))$ time in “`find()`”
- ◆ Using `array<size, linked list>`?
 - What are the indices? Dynamic or static?
 1. $\{0, 1, 2, 3, 4, 5, 6, \dots, n, \dots\}$
 2. $\{0, 1, 2, 4, 8, 16, 32, 64, \dots, 2^n, \dots\}$
 3. $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 16, 32, \dots, 2^n, \dots\}$
 - Decomposed? (e.g. $13 = 8 + 4 + 1$)
- ◆ Any hybrid idea?

Recycling List Implementation in HW#4

◆ Observation

- Most of the arrays are of small sizes
→ RecycleList[0] ~ RecycleList[255]
for new, new [1], new [2], ..., new [255]

◆ What about new [n], $n \geq 256$?

- Use $m = n \% 256$

Data Structure



MemMgr::_recycleList[R_SIZE]

Source Code

```
class RecycleList {
    // Go through this and _nextList,
    // find out a recycle list whose "_arrSize" == "n"
    MemRecycleList<T>* getList(size_t n) {
        // Find the recycle list whose _arrSize == n
    }
};

class MemMgr {
    size_t getRecycleIdx(size_t t) const{ // t Bytes to recycle
        assert(t >= S); // S: size of the recycled class element
        return (t-SIZE_T)/S; // subtract the size for storing 'n'
    }
    MemRecycleList<T>* getMemRecycleList(size_t n) {
        size_t m = n % R_SIZE;
        return _recycleList[m].getList(n);
    }
};

[e.g.] delete p; // let t = the #Bytes to recycle
→ getMemRecycleList(getMemRecycleList(t)).push_front(p);
```

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Recycling List with Different Classes

- ◆ What if we want to associate the same memory manager to different classes?
(i.e. n classes → 1 memMgr)
(e.g. class Inheritance?)
 - What would be the mem size in the recycling list? GCD of sizes of A & B?
Multiple of sizeof(size_t)?
 - More difficult to manage...!!
 - Suggest to use memory management without recycling

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Memory Management without Recycling

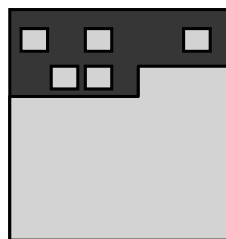
- ◆ Overload "new" to get memory from memMgr
- ◆ Overload "delete" to do nothing but calling destructor
 - ➔ No recycle

```
class MemoryBlock {
    char*      _begin;
    char*      _ptr;
    char*      _end;
    MemoryBlock* _nextBlock;
public:
    MemoryBlock(size_t B) { _begin = (char*)malloc(B); }
};

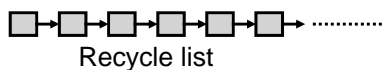
class MemoryMgr {
    MemoryBlock* _activeBlock;
    void* alloc(size_t) {
        // get memory from _activeBlock;
        // If over the limit, new MemoryBlock as _activeBlock }
    };
};
```

Garbage Collection

- ◆ After using memory management for a while, we may have many recycled memory pieces but not much required memory



Memory block



Recycle list

- ◆ Can we rearrange the pointers so that the freed memory can be put together and even returned to system earlier?
 - Pointer value changes? How to keep the associations?
 - Index or pointer?
 - Too many to cover; beyond the scope of this class...

Conclusion

- ◆ Memory related problems are mostly runtime problems
 - You won't see them during compilation
 - Crash during runtime → difficult to debug
 - But please use debugger instead of "cout"
- ◆ Use memory manager to allocate a block of memory instead of piece by piece
 - Don't need to worry about freeing individual memory → no memory leak
 - Still need to properly issue "delete" if the callings of destructors are needed
 - Can achieve better memory locality and thus better performance