

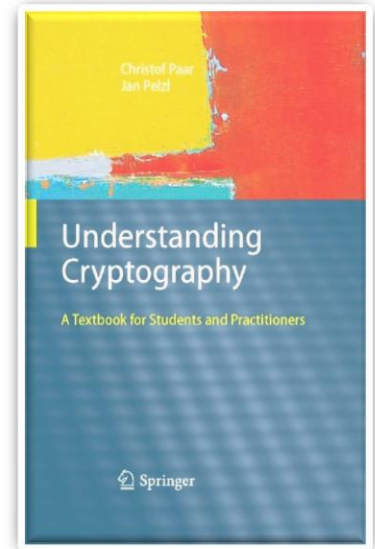
# Introduction to Cryptography

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# Acknowledgement

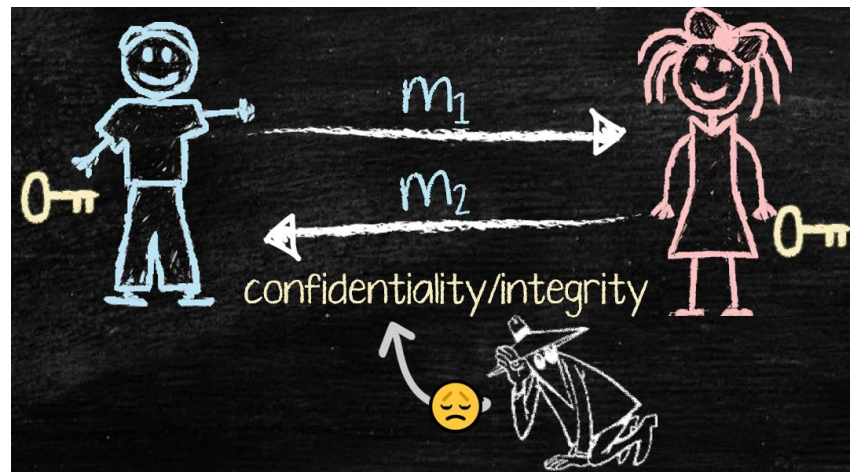
- These slides are based on the following resources (but modified):
  - Slides accompanying the textbook '*Understanding Cryptography*' by Christof Paar and Jan Pelzl  
<http://crypto-textbook.com/>
  - Slides from Dr. Ryan Riley  
<https://vsecurity.info/>



# Outline

- Introduction to Cryptography
- Caesar Cipher
- Substitution Cipher

# Introduction to Cryptography



# Definition

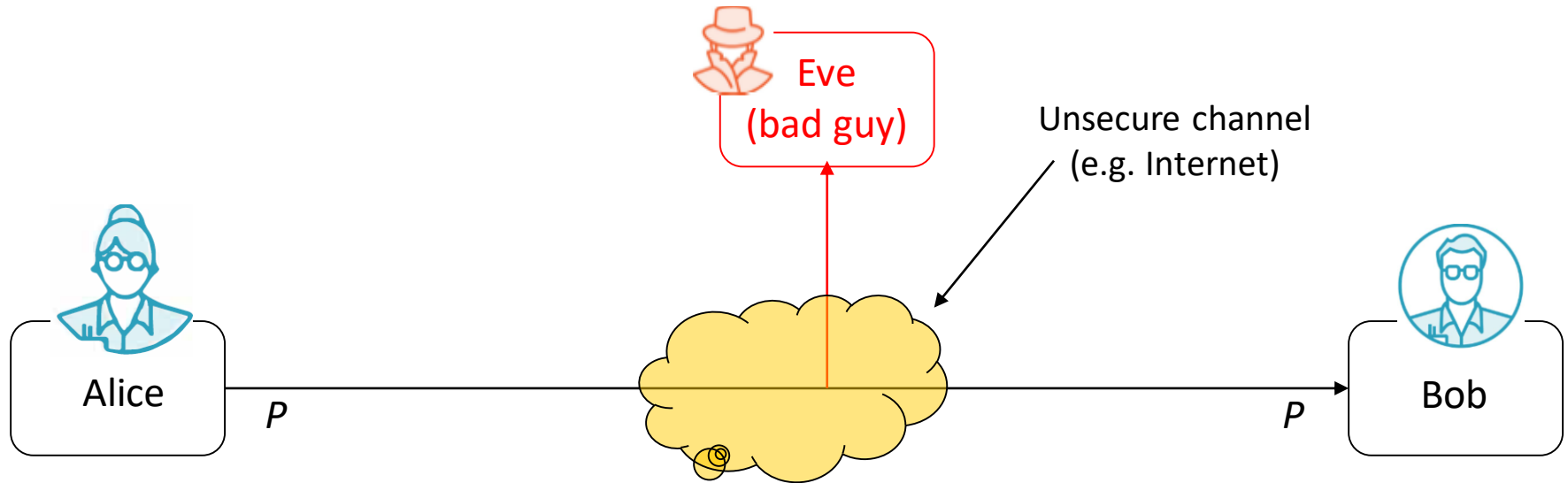
- The word Cryptography is Greek
  - Crypto: Secret + Graphy: Writing
  - Method to send secret messages using a key
- Basic goal is **Secure communication**
  - Send messages that no one but the expected recipient can read
- Many other applications such as:
  - Cryptocurrency, Blockchains
  - Authentication
  - Digital signatures
  - ...

# Terminology

- Plaintext: A message in its original form
- Ciphertext: A message in encrypted form
- Encryption: Transforming PT to CT
- Decryption: Transforming CT to PT
- Encryption Algorithm / Cipher: The method used for encryption

# Symmetric Cryptography

- Alternative names: **private-key**, **single-key** or **secret-key** cryptography.



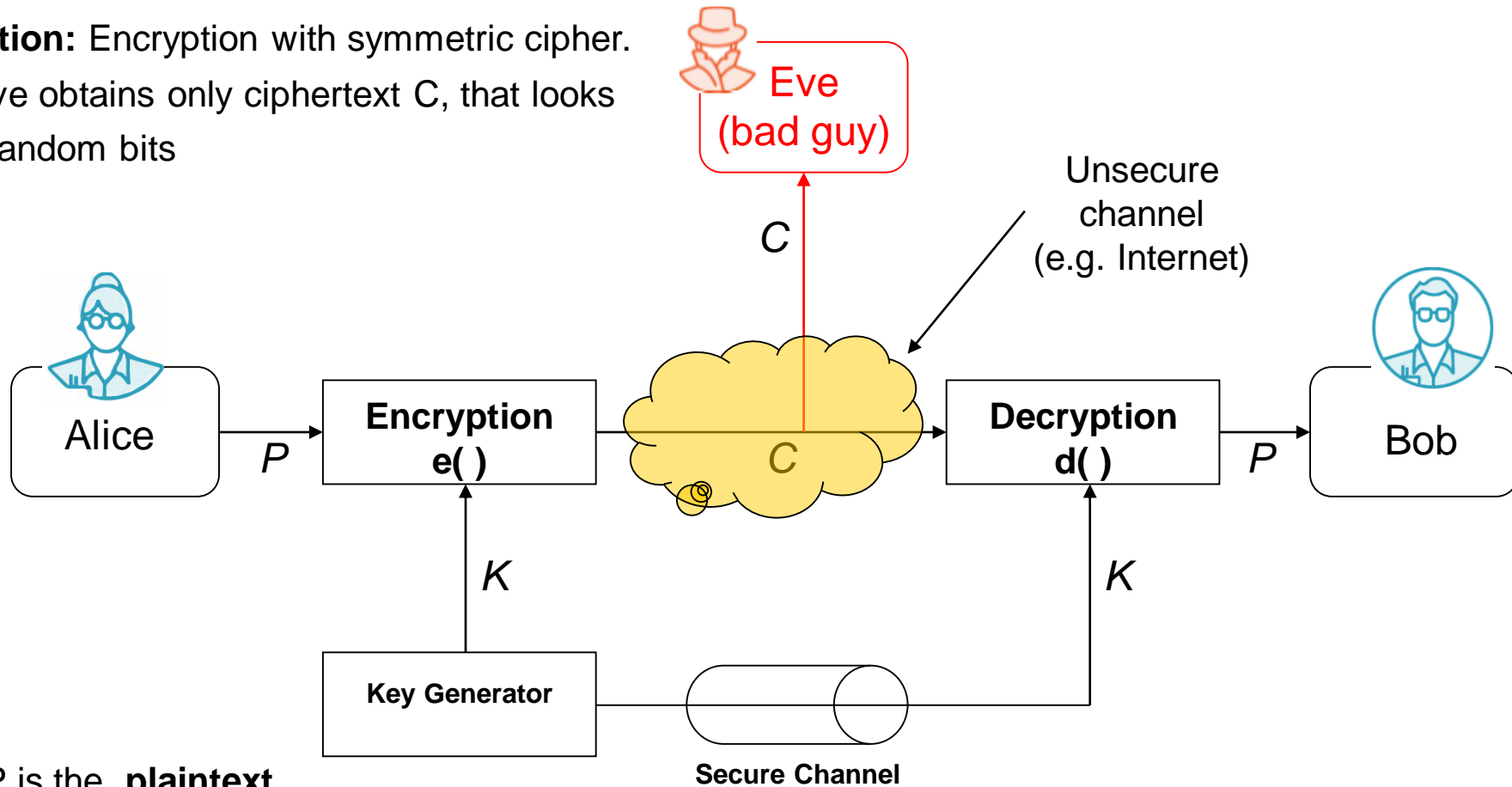
- Problem Statement:**

- 1) Alice and Bob would like to communicate via an unsecure channel (e.g., Internet)
- 2) A malicious third party Eve (the bad guy) has channel access but should not be able to understand the exchanged messages

# Symmetric Cryptography

**Solution:** Encryption with symmetric cipher.

⇒ Eve obtains only ciphertext  $C$ , that looks like random bits



- $P$  is the **plaintext**
- $C$  is the **ciphertext**
- $K$  is the **key**
- Set of all keys  $\{K_1, K_2, \dots, K_n\}$  is the **key space**



# Symmetric Cryptography

- |                       |              |
|-----------------------|--------------|
| • Encryption equation | $C = e_K(P)$ |
| • Decryption equation | $P = d_K(C)$ |

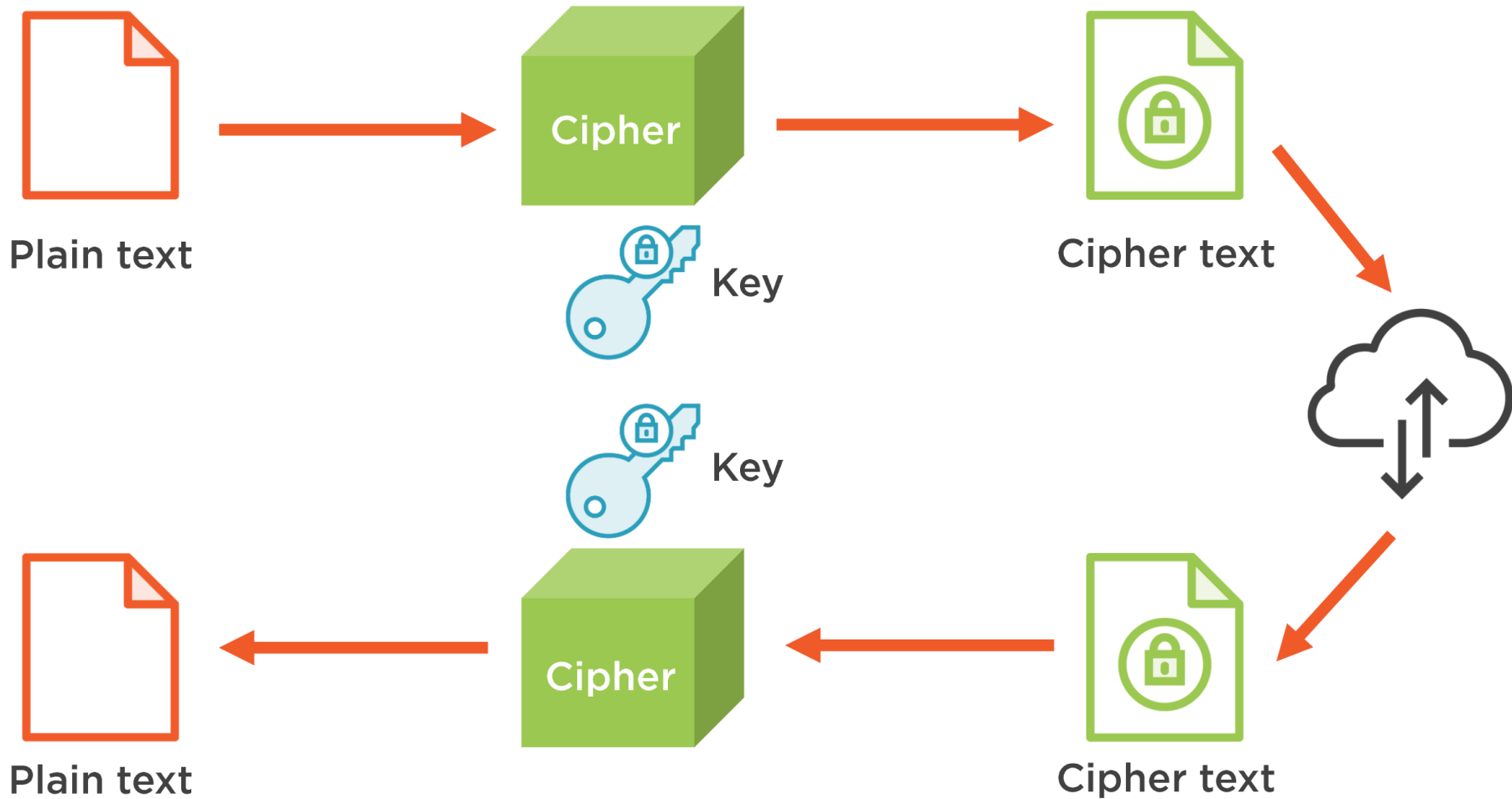
- Encryption and decryption are inverse operations if the same key  $K$  is used on both sides:

$$d_K(C) = d_K(e_K(P)) = P$$

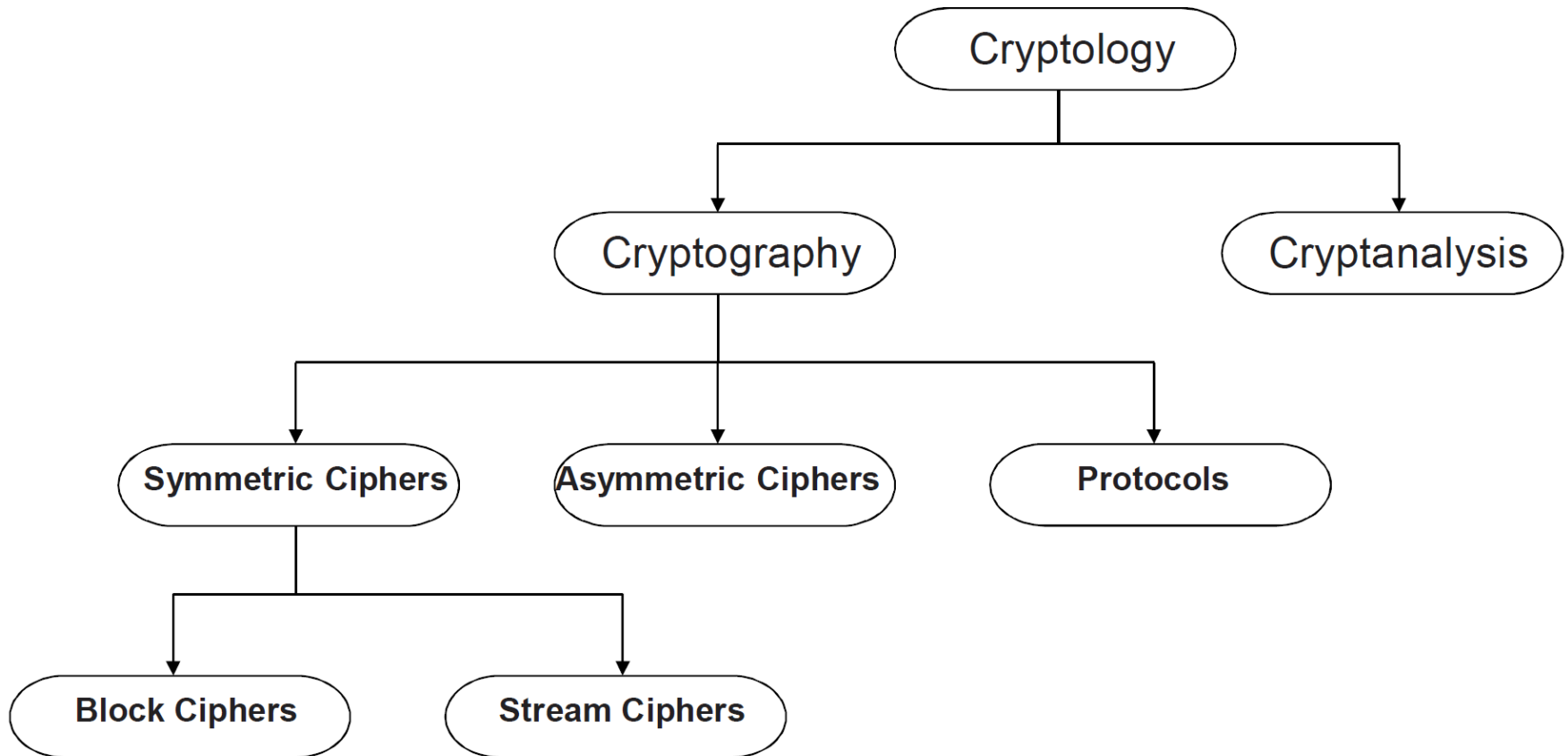
- Important: The key must be transmitted via a **secure channel** between Alice and Bob.
- The secure channel can be realized, e.g., by a human courier or a secure key exchange mechanism (this will be covered later)
- However, the system is only secure if an attacker does not learn the key  $K$ !

⇒ **The problem of secure communication is reduced to secure transmission and storage of the key  $K$**

# Symmetric Cryptography - Summary



# Classification of the Field of Cryptology



**Cryptanalysis** = Trying to break the key and read encrypted messages

# Why do we need Cryptanalysis?

- There is no *mathematical proof of security* for any practical cipher
- The only way to have assurance that a cipher is secure is to try to break it (and fail) !
  - We let lots of really smart people try to break it (cryptanalysis). If they can't, we assume it is secure
  - **But...** We might be wrong

**Kerckhoff Principle** is paramount in modern cryptography:



A cryptosystem should be secure even if the attacker knows all details about the system, **with the exception of the secret key**

- In order to achieve Kerckhoff's Principle in practice:  
**Only use widely known ciphers that have been cryptanalyzed for several years by good cryptographers!**
- **Remark:** It is tempting to assume that a cipher is “more secure” if its details are kept secret. However, history has shown time and again that secret ciphers can almost always be broken once they have been reverse engineered. (Example: Content Scrambling System (CSS) for DVD content protection.)

# Kerckhoff's Principle



Security of a Cryptographic Algorithm should rely  
ONLY on the secrecy of the KEYS, and  
NOT on the secrecy of the METHOD used

“Do not rely on security through obscurity”

# Brute-Force Attack (or Exhaustive Key Search) against Symmetric Ciphers

- Treats the cipher as a black box
- Requires (at least) 1 plaintext-ciphertext pair ( $P_0$ ,  $C_0$ )
- Check all possible keys until condition is fulfilled:

$$d_k(P_0) = C_0$$

Key length in bit	Key space	Security life time
64	$2^{64}$	Short term (few days or less)
128	$2^{128}$	Long-term (several decades in the absence of quantum computers)
256	$2^{256}$	Long-term (also resistant against quantum computers – note that QC do not exist at the moment)

# Caesar Cipher

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# Simple Ciphers

- Originally, cryptography was performed by hand
- Goal was to protect messages sent by couriers
  - From people who might intercept the courier
  - From the courier himself
- War was a popular time to use them

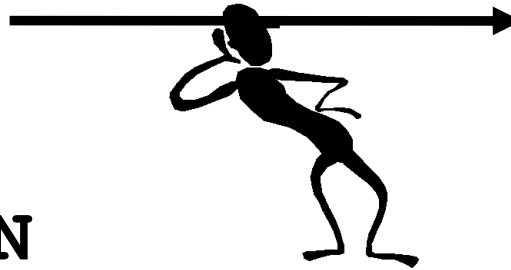


# Caesar Cipher

Julius Ceasar (100-44 BC)



DWWDFN DW GDZQ



**ATTACK AT DAWN**

**Encrypt the message!**



**ATTACK AT DAWN**

**Decrypt the ciphertext!**

- The sender and receiver must know something that the adversary doesn't.
- This is called a **cryptographic key**

# Caesar Cipher

**Secret key:** A random number from  $\{1, \dots, 26\}$ , say **3**



DWWDFN DW GDZQ



## Encryption

Message:      **ATTACK AT DAWN**

**Key: + 3**

↓ ↓ ↓ ↓ ↓ ↓    ↓ ↓    ↓ ↓ ↓ ↓

Ciphertext:    **DWWDFN DW GDZQ**

# Caesar Cipher

**Secret key:** A random number from  $\{1, \dots, 26\}$ , say **3**



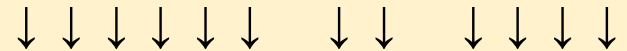
DWWDFN DW GDZQ



Decryption

Ciphertext: DWWDFN DW GDZQ

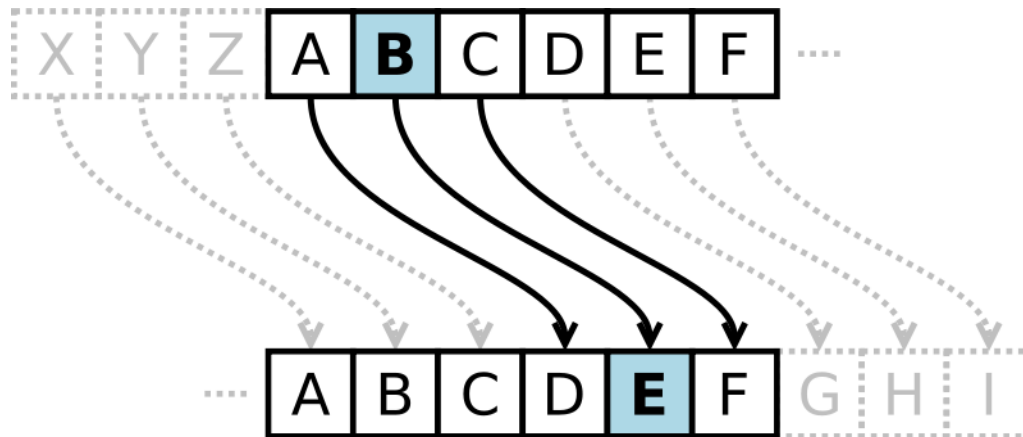
**Key: - 3**



Message: ATTACK AT DAWN

# Caesar Cipher

- Earliest documented cipher was used by Caesar in 50BC !
- Each letter in a message is substituted by another that is 3 letters away
  - **A** becomes **D**, **B** becomes **E**, etc.
  - Note that the letters "wrap around" at the end of the alphabet, which can be mathematically expressed using **mod 26**



# Caesar Cipher Example

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25

ATTACK AT DAWN

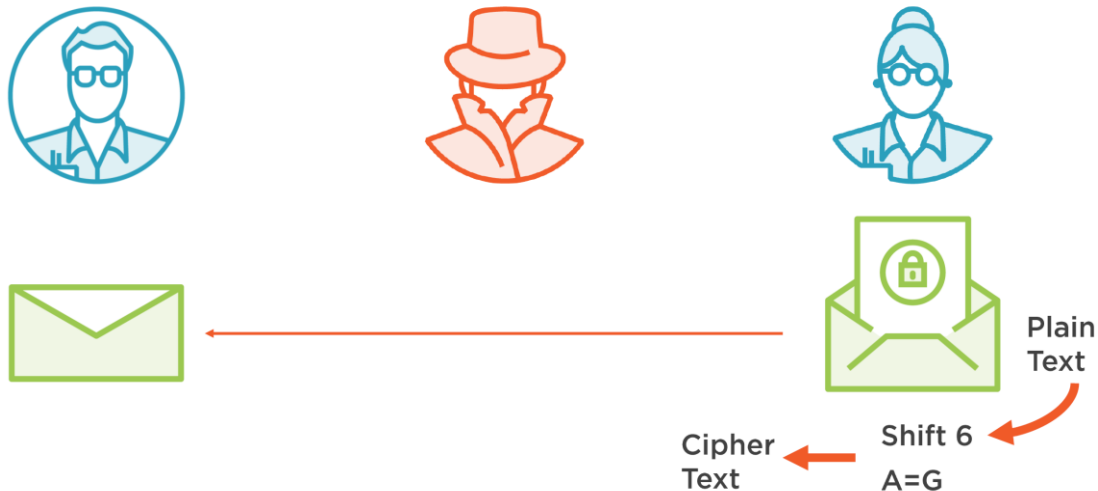
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# Shift Cipher

- Generic version of Caesar cipher
- Each letter is shifted by N. In Caesar, N=3

Let  $k, x, y \in \{0, 1, \dots, 25\}$

- Encryption:  $y = e_k(x) \equiv x + k \text{ mod } 26$
- Decryption:  $x = d_k(x) \equiv y - k \text{ mod } 26$



# Shift Cipher: Example

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25

- Let's do one for  $N=10$ ...

ATTACK AT DAWN

# Shift Cipher Cryptanalysis

- How do we break this?
- Brute-force: Try all possible values for N
  - There are only 26
- Feasibility?
  - Easy by hand
  - Trivial by computer



# Substitution Cipher

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# Substitution Cipher

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
Q	R	A	W	G	N	C	X	M	B	V	L	Z	D	S	J	T	E	K	Y	F	U	I	P	O	H

- Generate a random set of substitutions for each letter
  - Always a 1:1 correspondence

# Substitution Cipher Example

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
Q	R	A	W	G	N	C	X	M	B	V	L	Z	D	S	J	T	E	K	Y	F	U	I	P	O	H

ATTACK AT DAWN

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# Substitution Cipher Cryptanalysis

- Brute-force: Try all possible letter combinations
  - There are  $(26!) = 403291461126605635584000000$
- Exhaustive key search will take a long time ...
- Letter frequency analysis attack can be used against the substitution cipher

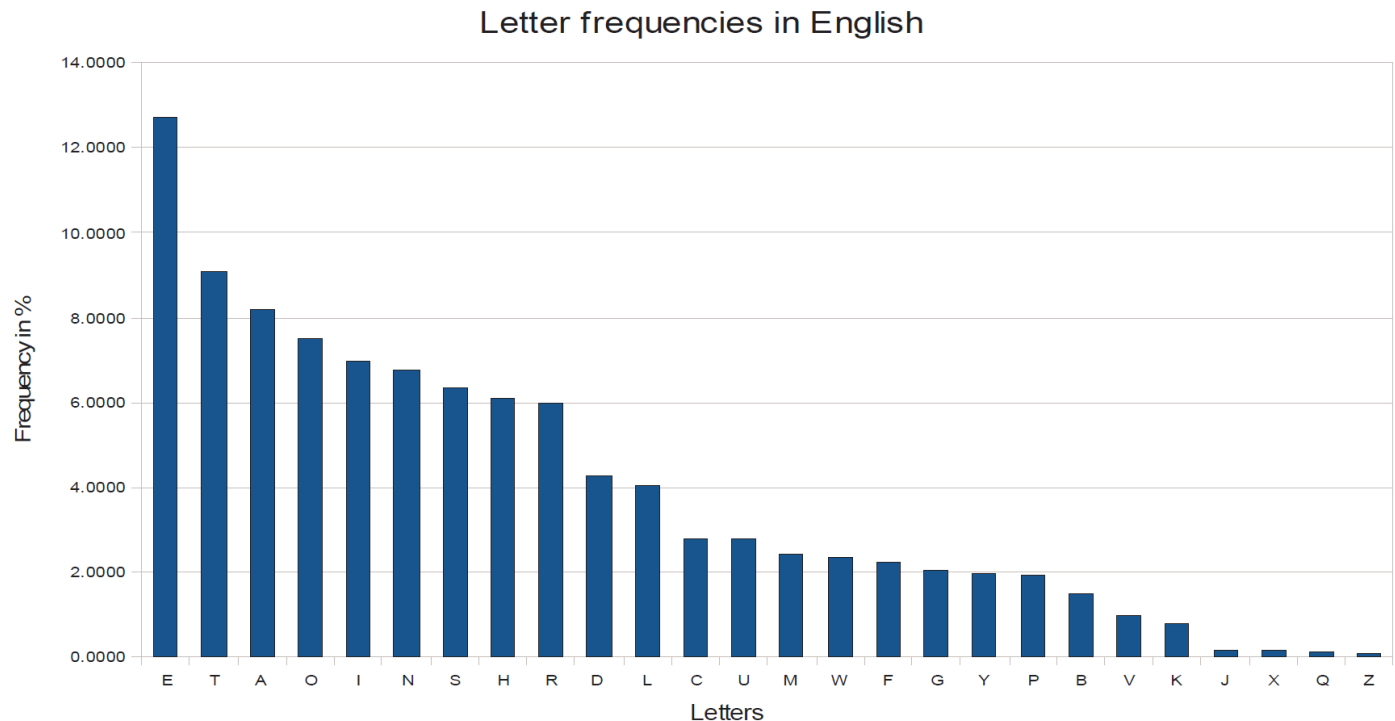
# Substitution Cipher Cryptanalysis

QYYQAV QY WQID

- Key observation: In a substitution cipher, **basic language features** are preserved
  - You can tell how often a letter occurs in the message
  - You can see when letters repeat
  - Etc.
- Use a technique called frequency analysis

# Frequency Analysis

- Not all letters in a language occur with the same frequency. E.g., In English,
  - E is most common
  - Vowels are about 40%
  - Vowels tend to be separated by consonants
  - Q tends to be followed by U
  - Etc.



# Breaking the Substitution Cipher with Letter Frequency Attack

- Let's take an example and identify the most frequent letter:

i<sup>q</sup> ifcc v<sup>qqr</sup> fb rd<sup>q</sup> vfllc<sup>q</sup> na rd<sup>q</sup> cfjwhwz hr bnnb hcc  
hwwhbs<sup>qvqbre</sup> hw<sup>q</sup> vh1<sup>q</sup>

- We replace the ciphertext letter q by E and obtain:

i<sup>E</sup> ifcc v<sup>EER</sup> fb rd<sup>E</sup> vfllc<sup>E</sup> na rd<sup>E</sup> cfjwhwz hr  
bnnb hcc hwwhbs<sup>EvEbre</sup> hw<sup>E</sup> vh1<sup>E</sup>

- By further guessing based on the frequency of the remaining letters we obtain the plaintext:

WE WILL MEET IN THE MIDDLE OF THE LIBRARY AT NOON ALL ARRANGEMENTS  
ARE MADE

- In practice, not only frequencies of individual letters can be used for an attack, but also the frequency of letter pairs (i.e., “th” is very common in English), letter triples, etc.

# Vigenère Cipher

## (1900-1950)

- Poly-alphabetic cipher
  - One plaintext letter can become *different* ciphertext letters
- Uses a text based key and modulo arithmetic to perform the encryption
- Frequency analysis is possible, but much more difficult



# Vigenère Cipher: Example

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25

- Let's choose a key of "MONKEY"

ATTACK AT DAWN

MONKEY MO NKEY

MHGKGI MH QKAL

# One-Time Pad

- Vigenère cipher with a randomly chosen key as long as the message
- Key needs to be shared between parties beforehand
- Key can **never** be re-used
- Provable unbreakable without the key
- This is the only perfect cryptography

# One-Time Pad: Example

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
26	27	28	29	30	31	32	33	34	35	36	37	38	39	40	41	42	43	44	45	46	47	48	49	50	51

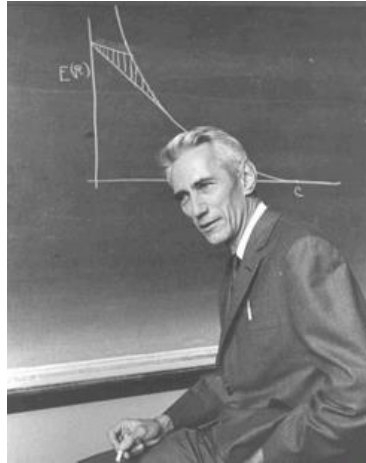
- Our random key is “FOWIFOZMQOAF”

ATTACK AT DAWN

FOWIFO ZM QOAF

# A Mathematical Theory of Secrecy

Claude Shannon (late 1940s)



Perfect Secrecy is guaranteed by the One-time Pad

THE GOOD: Unbreakable regardless of the power of the adversary

THE BAD: Impractical! Needs very, very long shared keys.

THE UGLY: length of key  $\geq$  total length of all messages you ever want to send

# Crypto Components

- All of the previous techniques have two basic components:
  - **Algorithm** (What you do to the message)
  - **Key** (The secret that you need in order to encrypt/decrypt properly)
- When using these algorithms, the **key is secret**
- The algorithm is not

# Summing Up

- We trust a cryptographic algorithm if lots of smart people can't break it
- We looked at three types of simple ciphers:
  - Shift Cipher
  - Substitution Cipher
- They each have an algorithm and a key
- Long key is required for symmetric algorithms in order to prevent exhaustive key-search attacks

# Resources

- Cryptool - Software demonstrating many ancient and modern ciphers

<https://www.cryptool.org/en/>

- An excellent [one-hour video](#) summarizing the last 40 years of modern cryptography by Ron Rivest
- The [International Association of Cryptographic Research](#) is the professional organization of cryptographers.