

UNIT 3 (I)

- POSETS AND LATTICES

- BOOLEAN ALGEBRA

↳ truth table, k-Map, circuit diagrams, circuit minimization

Relations : Partial Order Sets

Set $A \times B$

$$R = \{ (-, -), (-, -), (-, -), (-, -) \}$$

reflexive, antisymmetric, transitive

$$\text{Poset} = \{ (-, -), (-, -), (-, -), (-, -) \}$$

collection of sets satisfying relation R

poset $\rightarrow (Z^+, |)$ divisibility/divides

$$Z^+ \times Z^+$$

$$\begin{array}{|l} \therefore \text{it's a poset} \\ a \in Z^+ \\ a|a \\ a|b \wedge b|c \Rightarrow a|c \text{ transitive} \\ \forall a \forall b ((aRb) \wedge (bRa) \rightarrow a=b) \text{ antisymmetric} \\ \text{or} \\ \forall a \forall b (((a,b) \in R) \wedge ((b,a) \in R) \rightarrow a=b) \end{array}$$

A relation R on a set A is called partial ordering if it is reflexive, antisymmetric and transitive.

$$1. (Z, =) \rightarrow P$$

$$2. (Z, \neq) \rightarrow \text{not a P. } \frac{(a \neq b) \wedge (b \neq a) \rightarrow a=b}{T} \quad \frac{\therefore \text{not AS}}{F}$$

$$3. (Z, \gg) \rightarrow P$$

Comparable Elements

$$\{ (), (), (), () \}$$

$$a, b \in A$$

$$a \leq b \quad \text{or} \quad b \leq a \quad \rightarrow b \text{ is partially related to } a$$

a and b are comparable to each other

$$a \not\leq b \quad b \not\leq a$$

a and b are incomparable

The elements a and b of a poset (S, \leq) are called comparable if either

\downarrow
set
 \downarrow
partial order

$$a \leq b \quad \text{or} \quad b \leq a \quad \text{when} \quad a \leq b$$

and $b \not\leq a$ then a and b are called incomparable

$$\text{Eg: } (\mathbb{Z}^+, |) \quad \{ (2, 4), (5, 7), \dots \}$$

$$2|4 \quad \text{but} \quad 4 \nmid 2$$

$$5 \nmid 7, \quad 7 \nmid 5$$

Toset : Totally ordered set [every element is comparable]
either $a \leq b$ or $b \leq a$

$$(\mathbb{Z}, \leq)$$

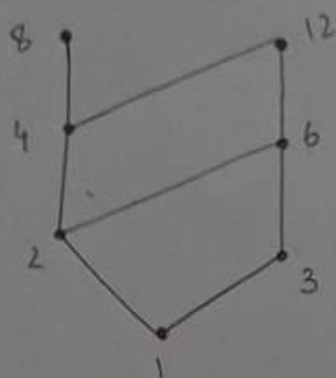
$$(\mathbb{Z}, \leq)$$

Hasse Diagram

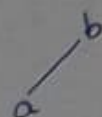
Diagram to represent PosET and TosET

$$(S, \leq) \quad (\mathbb{Z}^+, |)$$

1. $(\{1, 2, 3, 4, 6, 8, 12\}, |)$



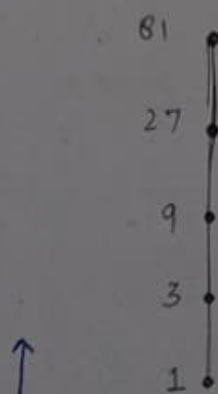
$$a \leq b$$



$$a \not\leq b$$

no downward travel
in Hasse Diagram

2. $(\{1, 3, 9, 27, 81\}, |) \rightarrow \text{TosET}$

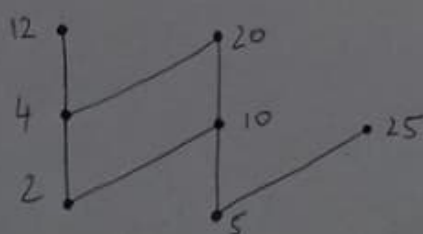


toset: straight line (chain)

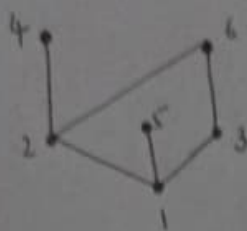
$$(3|27 \wedge 27 \nmid 3) \rightarrow 3 = 27?$$

F

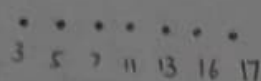
3. $(\{2, 4, 5, 10, 12, 20, 25\}, |)$



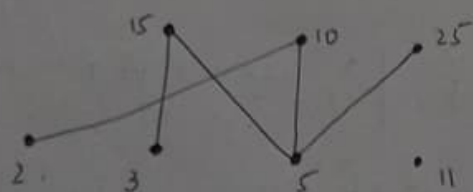
Q1. $(\{1, 2, 3, 4, 5, 6\}, 1)$



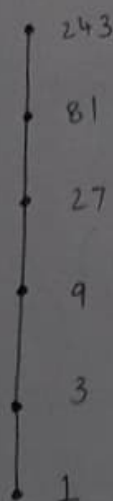
Q2. $(\{3, 5, 7, 11, 13, 16, 17\}, 1)$



Q3. $(\{2, 3, 5, 10, 11, 15, 25\}, 1)$



Q4. $(\{1, 3, 9, 27, 81, 243\}, 1)$

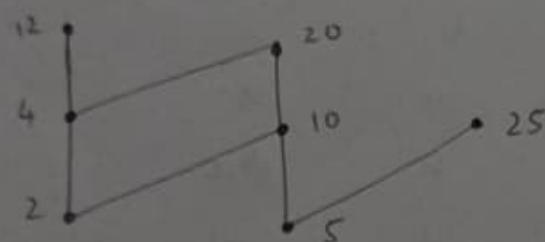


Special Elements of Hasse Diagram

1. Minimal and Maximum Elements

'a' is minimal if there is no element $b \in S$ such that $b \neq a$ for $a, b \in S$

'a' is maximal if there is no element $b \in S$ such that $a \neq b$ for $a, b \in S$



minimal : 2, 5
maximal : 12, 20, 25
greatest \rightarrow x
least \rightarrow x

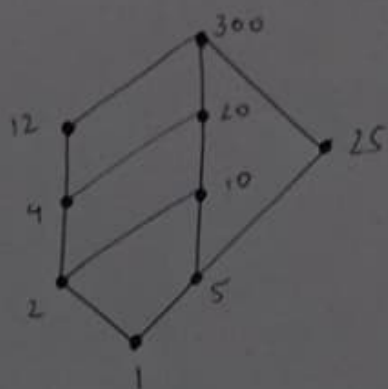
2. Greatest and Least Elements

'a' is the greatest element of poset (S, \leq) if $b \leq a \quad \forall \quad b \in S$.

greatest element is unique if it exists.

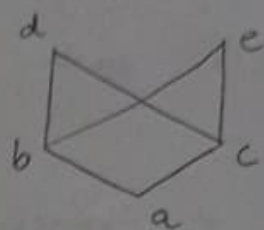
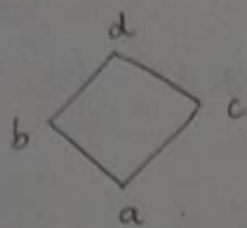
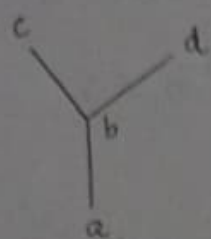
'a' is the least element of poset (S, \leq) if $a \leq b \quad \forall \quad a \in S$

least element is unique if it exists



greatest \rightarrow 300
least \rightarrow 1

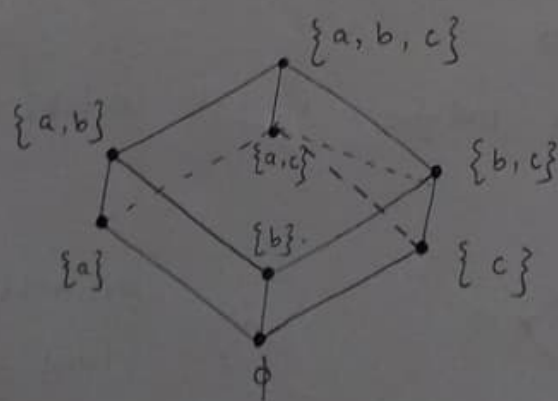
* Trick : If your figure is closed or straight line, then there will be greatest and least element possible



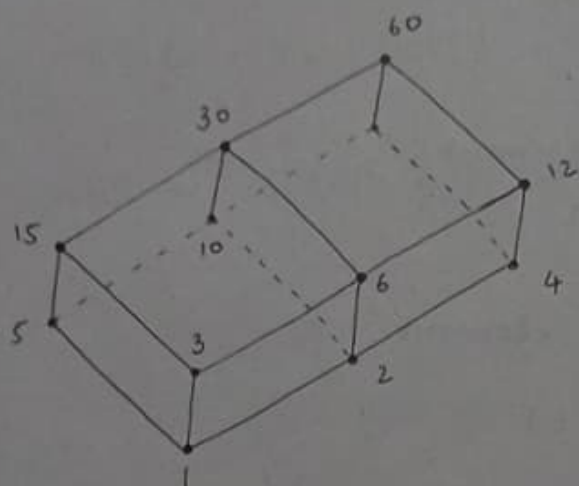
greatest \rightarrow	x	d	b	x
least \rightarrow	a	a	a	a
maximal \rightarrow	c, d	d	b	d, c
minimal \rightarrow	a	a	a	a

Q1. $(P(S), \subseteq)$, $S = \{a, b, c\}$

$$P(S) = \{ \phi, \{a\}, \{b\}, \{c\}, \{a, b\}, \{b, c\}, \{c, a\}, \{a, b, c\} \}$$



$D_{60} \rightarrow$ factors of $\cancel{D} 60$



Q. Show that there is exactly one greatest element of a poset, if such an element exists.

THEOREM
Let g_1 and g_2 be two greatest elements of poset (S, \leq)

where $g_1, g_2 \in S$

If g_1 is the greatest element

$$g_2 \leq g_1 \quad \text{--- (1)}$$

If g_2 is the greatest element

$$g_1 \leq g_2 \quad \text{--- (2)}$$

$$\left| \begin{array}{l} a|b \text{ and } b|a \\ a=b \end{array} \right.$$

$$a \leq b, b \leq a$$

$$a=b$$

$$\boxed{g_1 = g_2}$$

Q. Show that there is exactly 1 least element of a poset, if such an element exists.

THEOREM

Let l_1 and l_2 be two least element of poset (S, \leq)

where $l_1, l_2 \in S$

If l_1 is the least element;

$$l_1 \leq l_2 \quad \text{--- (1)}$$

If l_2 is the least element;

$$l_2 \leq l_1 \quad \text{--- (2)}$$

$$\Rightarrow \boxed{l_1 = l_2}$$

Q. Give a poset that has:

- (i) a minimal element but not a maximal element.
- (ii) a maximal element but not a minimum element.
- (iii) neither minimal nor maximal element.

$$(i) \quad (\mathbb{N}, \overset{\geq}{\cancel{\leq}}), (\mathbb{Z}^+, \overset{\geq}{\cancel{\leq}}), (\mathbb{R}^+, \overset{\geq}{\cancel{\leq}})$$

$$(ii) \quad (\mathbb{Z}^-, \overset{\geq}{\cancel{\leq}}), (\mathbb{R}^-, \overset{\geq}{\cancel{\leq}})$$

$$(iii) \quad (\mathbb{Z}, \overset{\geq}{\cancel{\leq}}), (\mathbb{R}, \overset{\geq}{\cancel{\leq}})$$

Product Order

(A, \leq_1) and (B, \leq_2) are posets.

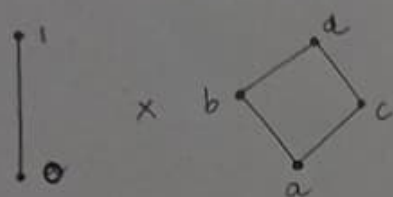
Then $(A \times B, \leq)$ is also a poset.

If $a \leq_1 a'$ in A

and $b \leq_2 b'$ in B

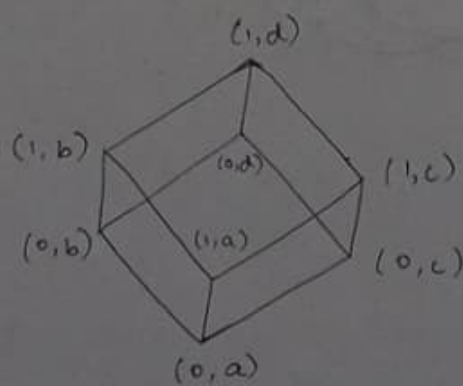
then $(a, b) \leq (a', b')$ in $A \times B$

Q.



$$\{0, 1\} \times \{a, b, c, d\}$$

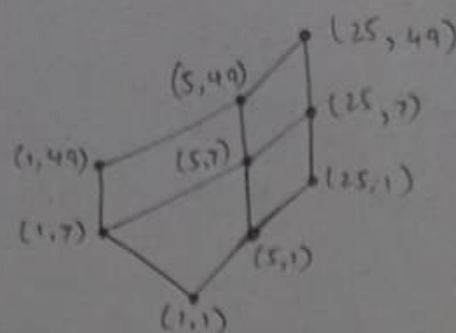
$$\{(0, a), (0, b), (0, c), (0, d), (1, a), (1, b), (1, c), (1, d)\}$$



Q. $D_{25} \times D_{49}$

$$(\{1, 5, 25\}, 1) \times (\{1, 7, 49\}, 1)$$

$$\{(1, 1), (1, 7), (1, 49), (5, 1), (5, 7), (5, 49), (25, 1), (25, 7), (25, 49)\}$$



THEOREM If (A, \leq_1) and (B, \leq_2) be posets. then $(A \times B, \leq)$ is also a poset.

$$\begin{array}{ll} (A, \leq_1) & (B, \leq_2) \\ a, a' \in A & b, b' \in B \\ a'' & b'' \end{array}$$

$$(A \times B, \leq)$$

$$(a, b) \leq (a, b)$$

$$a \leq_1 a$$

$$b \leq_2 b$$

Hence, $(A \times B, \leq)$ is reflexive

$$(a, b), (a', b'), (a'', b'')$$

$$\text{If } (a, b) \leq (a', b')$$

$$\text{and } (a', b') \leq (a'', b'')$$

$$\text{then } (a, b) \leq (a'', b'')$$

$$a, a', a'' \in A$$

$$b, b', b'' \in B$$

$$a \leq_1 a \text{ and}$$

$$a' \leq_1 a''$$

$$\Rightarrow a \leq_1 a''$$

$$b \leq_2 b' \text{ and}$$

$$b' \leq_2 b''$$

$$\Rightarrow b \leq_2 b''$$

Hence, $(A \times B, \leq)$ is transitive

$$\text{If } (a, b) \leq (a', b')$$

$$\text{and } (a', b') \leq (a, b)$$

$$\text{then } (a, b) = (a', b')$$

$$a \leq_1 a' \text{ and } a' \leq_1 a$$

$$\Rightarrow a = a'$$

$$b \leq_2 b' \text{ and } b' \leq_2 b$$

$$\Rightarrow b = b'$$

Hence, $(A \times B, \leq)$ is antisymmetric

$\therefore (A \times B, \leq)$ is a poset.

Upper Bounds and Lower Bounds

Let B be a subset of a poset (A, \leq) and element u belong to A , this is known as an upper bound of B if u succeeds every element of B .

$$(A, \leq)$$

$$(B \subseteq A)$$

$$U \in A$$

$$x \leq U$$

$$\forall x \in B$$

U is upper bound

Let B be a subset of a poset (A, \leq) and element l belongs to A , this is known as lower bound of B if l precedes every element of B .

$$(A, \leq)$$

$$(B \subseteq A)$$

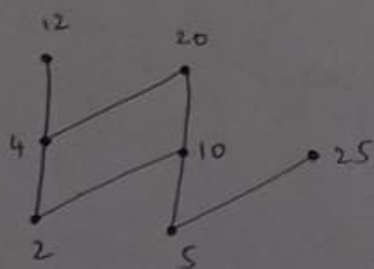
$$L \in A$$

$$L \leq x$$

$$\forall x \in B$$

L is lower bound

Eg:



$$\{2, 4, 5, 10, 12, 20, 25\}$$

ϕ is the lower bound of $\{5, 12\}$

ϕ is the upper bound of $\{5, 12\}$

• upper bound $\rightarrow \{5, 12\}$

kisse relation
rakh raha he?

lower bound $\rightarrow \{5, 12\}$

Se kon
relation
rakh raha
he?

$\{2\}$ is the lower bound of $\{2, 4\}$

$\{4, 12, 20\}$ is the upper bound of $\{2, 4\}$

Least Upper Bound (LUB) : Supremum
Greatest Lower Bound (GLB) : Infimum

$$\{u_1, u_2, u_3, u_4\}$$

$$\{l_1, l_2, l_3, l_4\}$$

$$u_3 \leq u_1$$

$$u_3 \leq u_2$$

$$u_3 \leq u_4$$

$$l_2 \leq l_1$$

$$l_3 \leq l_1$$

$$l_4 \leq l_1$$

LUB: Suppose we are having upper bound
 $\{u_1, u_2, u_3, u_4\}$ and suppose u_3 is
the LUB, then:

$$u_3 \leq u_1$$

$$u_3 \leq u_2$$

$$u_3 \leq u_4$$

if any of the conditions
fail, it will not be
the LUB

GLB: Suppose we are having lower bound
 $\{l_1, l_2, l_3, l_4\}$ and suppose l_3 is
the GLB, then:

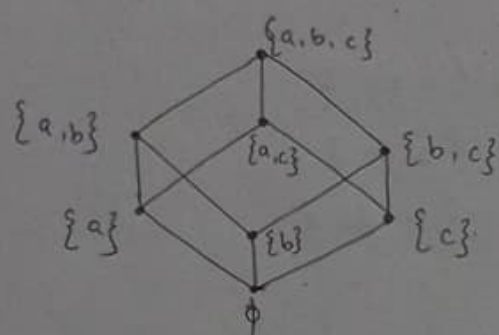
$$l_1 \leq l_3$$

$$l_2 \leq l_3$$

$$l_4 \leq l_3$$

LATTICES

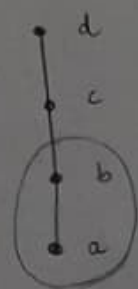
A poset in which every pair of element has both LUB and GLB is known as lattice.



$$\begin{aligned} \phi &\leq \{ \phi, \{a\} \} \leq \{a\} \\ &\{ \phi, \{b\} \} \\ &\{ \phi, \{c\} \} \\ &\{ \phi, \{a, b\} \} \\ &\{ \phi, \{b, c\} \} \\ &\{ \phi, \{a, c\} \} \\ &\{ \phi, \{a, b, c\} \} \end{aligned}$$

NOTE: $a \vee b = a$ join b
(LUB of a and b)
 $a \wedge b = a$ meet b
(GLB of a and b)

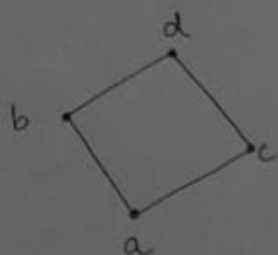
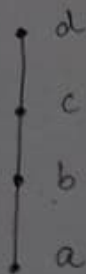
Suppose we are having any toset,



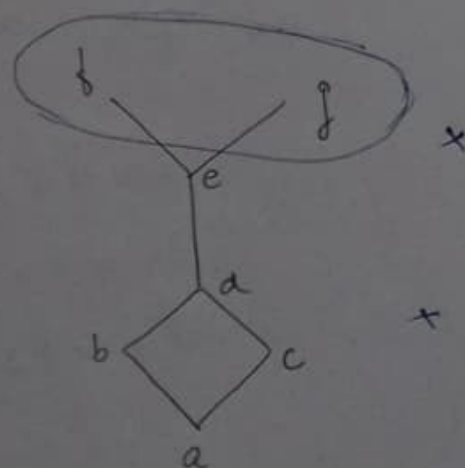
If we take $\{a, b\} \Rightarrow$ LUB = $\{b\}$
GLB = $\{a\}$

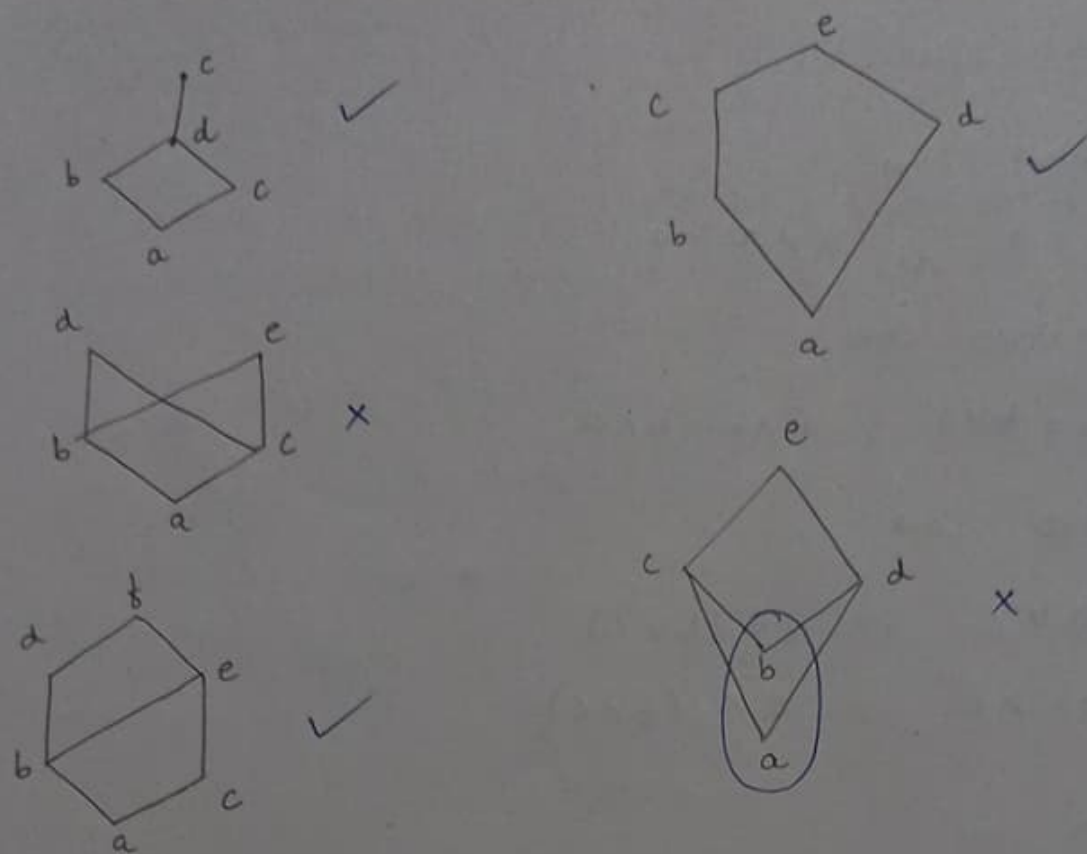
$\{b, c\} \Rightarrow$ LUB = $\{c\}$
GLB = $\{b\}$

Eg:



\downarrow ✓
LUB = d
GLB = a





Sublattices

A non empty subset 's' of lattice 'L' is called a sub-lattice of lattice of \forall 'L' if (L, \leq)
 $S \subseteq L$ (S, \leq)

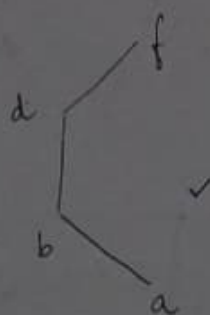
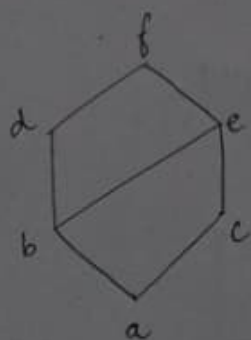
$$a \vee b \in S$$

and

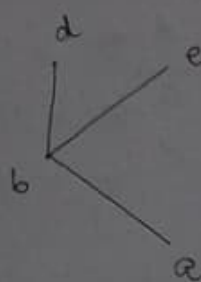
$$a \wedge b \in S$$

whenever $a, b \in S$

Eg:



is a sub-lattice



not a sub-lattice
 as d and e have
 no LUB

Properties of Lattices :

1. Idempotent Law

$$a \vee a = a, \quad a \wedge a = a$$

2. Commutative Law

$$a \vee b = b \vee a, \quad a \wedge b = b \wedge a$$

3. Associative Law

$$(a \vee b) \vee c = a \vee (b \vee c)$$

$$(a \wedge b) \wedge c = a \wedge (b \wedge c)$$

4. Absorption Law

$$a \vee (a \wedge b) = a$$

$$a \wedge (a \vee b) = a$$

THEOREM : For any a, b, c, d in a lattice (L, \leq)

(i) $a \leq a \vee b$ and $a \wedge b \leq a$

(ii) if $a \leq b$ and $c \leq d$ then $a \vee c \leq b \vee d$
and
 $a \wedge c \leq d \wedge d$

(ii) $a \leq b$

$$b \leq b \vee d$$

$$d \leq b \vee d$$

$$a \leq b \vee d$$

$$c \leq b \vee d$$

$$c \leq d$$

$$d \leq b \vee d$$

$$a \vee c \leq b \vee d$$

Types of Lattices :

1. Bounded Lattice :

A lattice is said to be bounded if it has a greatest element and a least element.

$$0, a, 1 \in L$$

$$0 \leq a \leq 1$$

* identity property

$$a \vee 1 = 1, \quad a \vee 0 = a$$

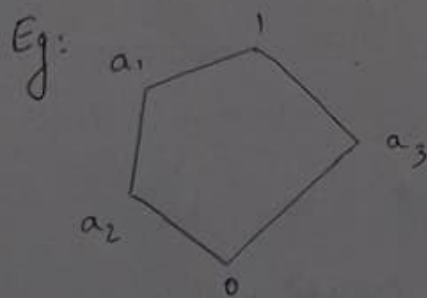
$$a \wedge 1 = a, \quad a \wedge 0 = 0$$

2. Distributive Lattice :

A Lattice (L, \leq) is said to be distributive if $(a \vee b) \wedge c = (a \wedge c) \vee (b \wedge c)$

for ~~$a \vee (b \wedge c)$~~

$$\begin{aligned} & a \vee (b \wedge c) \\ & a \wedge (b \vee c) \\ & (a \wedge b) \vee c \end{aligned}$$



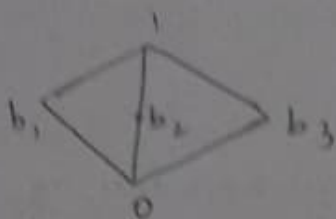
$$a_1 \wedge (a_2 \vee a_3) = (a_1 \wedge a_2) \vee (a_1 \wedge a_3)$$

$$a_1 \wedge 1 = a_2 \vee 0$$

$$a_1 \neq a_2$$

not a distributive lattice.

Eg.



THEOREM : In a distributive lattice (L, \leq) ,
if $a \vee b = a \vee c$ and $b \wedge b = a \wedge c$
imply that $b = c$

$$\begin{aligned}
 b &= b \wedge (b \vee a) && \text{absorption} \\
 &= b \wedge (a \vee c) && \text{given} \\
 &= (b \wedge a) \vee (b \wedge c) && \text{distributive} \\
 &= (a \wedge c) \vee (b \wedge c) && \text{given} \\
 &= (a \vee b) \wedge c && \text{distributive} \\
 &= (a \vee c) \wedge c && \text{given} \\
 &= c && \text{absorption}
 \end{aligned}$$

Hence, proved

3. Modular Lattice:

a lattice (L, \leq) is said to be
modular if $a \vee (b \wedge c) = (a \vee b) \wedge c$
whenever $a \leq c$

~~THEOREM - In a distributive lattice (L, \leq) if $a \vee b = a \vee c$ and $a \wedge b = a \wedge c$, imply that $b = c$~~

THEOREM : Every distributive lattice is modular

$$\begin{array}{l} c = a \vee c \\ | \\ a = a \wedge c \end{array}$$

$$\begin{aligned} a \vee (b \wedge c) &= (a \vee b) \wedge (a \vee c) \\ &= (a \vee b) \wedge c \end{aligned}$$

4. Complemented Lattice

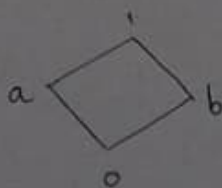
It is a lattice in which every element has one complement.

$$(L, \wedge, \vee, 0, 1)$$

$$b, a \in L$$

$$1. \quad 0 \leq a \leq 1$$

$$\begin{aligned} 2. \quad a \vee b &= 1 \\ a \wedge b &= 0 \end{aligned}$$



$(a)' = b$	$(0)' = 1$
$(b)' = a$	$(1)' = 0$

$$0 \vee 1 = 1$$

$$0 \wedge 1 = 0$$

$$a \vee b = 1$$

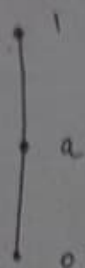
$$a \wedge b = 0$$

$$a \vee b = b \quad \times$$

$$0 \wedge b = 0$$

$$a \vee 0 = a \quad \times$$

$$a \wedge 0 = 0$$



$$1 \vee 0 = 1$$

$$1 \wedge 0 = 0$$

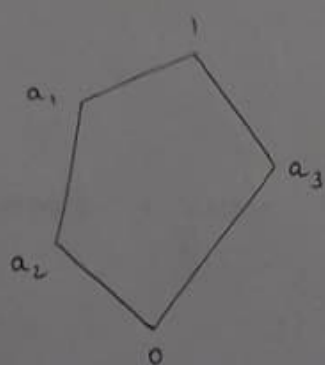
$$0 \vee a = a \quad \times$$

$$0 \wedge a = 0$$

$$a \vee 1 = 1$$

$$a \wedge 1 = a \quad \times$$

not a complementary lattice because there is no complement of a .



$$a_1 \vee a_3 = 1$$

$$a_1 \wedge a_3 = 0$$

$$a_2 \vee a_3 = 1$$

$$a_2 \wedge a_3 = 0$$

$$a_2 \vee 0 = a_2 \quad \times$$

$$a_1 \vee 0 = a_1 \quad \times$$

$$a_3 \vee 0 = a_3 \quad \times$$

$$0 \vee 1 = 1$$

$$0 \wedge 1 = 0$$

$$(a_3)' = a_2, a_1$$

$$(a_1)' = a_3$$

$$(a_2)' = a_3$$

$$(0)' = 1$$

$$(1)' = 0$$

THEOREM : In a distributive lattice if an element has a complement, then this complement is unique.

$$a, a_1, a_2 \in L$$

$$a \vee a_1 = 1$$

$$a \vee a_2 = 1$$

$$a \wedge a_1 = 0$$

$$a \wedge a_2 = 0$$

$$\begin{aligned} a_1 &= a_1 \wedge 1 && \text{identity} \\ &= a_1 (a \vee a_2) && \text{given} \end{aligned}$$

$$= (a_1 \wedge a) \vee (a_1 \wedge a_2)$$

$$= 0 \vee (a_1 \wedge a_2)$$

$$= (a \wedge a_2) \vee (a_1 \wedge a_2)$$

$$= (a \vee a_1) \wedge a_2$$

$$= 1 \wedge a_2$$

$$= a_2$$

THEOREM : In a complemented and distributive lattice

$$(i) \quad a \leq b$$

$$(ii) \quad a \wedge b' = 0$$

$$(iii) \quad a' \vee b = 1$$

$$(iv) \quad b' \leq a'$$

* from (i) prove (ii)

(ii) (iii)

(iii) (iv)

(iv) (i)

$$(ii) \quad b = a \vee b$$

$$a = a \wedge b$$

$$a \vee b = b$$

$$(a \vee b) \wedge b' = b \wedge b'$$

$$(a \wedge b') \vee (b \wedge b') = b \wedge b'$$

$$(a \wedge b') \vee 0 = 0$$

$$(a \wedge b') = 0$$

(iii)

$$(a \wedge b') = 0$$

$$(a \wedge b') = (0)'$$

$$a' \vee b = 1$$

$$\star \quad a \vee (b \wedge c) \leq (a \vee b) \wedge (a \vee c)$$

$$\begin{array}{l} a \leq a \vee b \\ b \wedge c \leq b \leq a \vee b \end{array}$$

$$a \leq a \vee c$$

$$b \wedge c \leq c \leq a \vee c$$

$$a \leq (a \vee b) \wedge (a \vee c)$$

$$b \wedge c \leq (a \vee b) \wedge (a \vee c)$$

$$\Rightarrow a \vee (b \wedge c) \leq (a \vee b) \wedge (a \vee c)$$

$$a \vee (b \wedge c)$$

$$\leq (a \vee b) \wedge c$$

$$\text{if } a \leq c$$

$$a \vee c = a$$

UNIT III (Ind)

BOOLEAN ALGEBRA

Duality Law

U_4, U_3 , boolean

lattice

P, Q	a, b	x, y
\vee	\vee join	$+$ sum
\wedge	\wedge meet	\cdot product
T	1 greatest	1
F	0 least	0
\neg	$'$	$'$, $-$

$$\neg P \wedge (Q \rightarrow R) \equiv \neg P \wedge (\neg Q \vee R) \equiv \neg P \vee (\neg Q \wedge R)$$

$$x \cdot y \bar{z} + \bar{y} z = (x + y + \bar{z}) \cdot (\bar{y} + z)$$

Truth Table

	x	y	z	\bar{z}	\bar{y}	$x y \bar{z}$	$\bar{y} z$	$x y \bar{z} + \bar{y} z$
0	0	0	0	1	1	0	0	0
1	0	0	1	0	1	0	1	1
2	0	1	0	1	0	0	0	0
3	0	1	1	0	0	0	0	0
4	1	0	0	1	1	0	0	0
5	1	0	1	0	1	0	1	1
6	1	1	0	1	0	1	0	1
7	1	1	1	0	0	0	0	0

$$x + y + \bar{z}$$

1
0
1
1
1
1
1
1

$$\bar{y} + z$$

1
1
0
1
1
1
0
1

$$(x + y + \bar{z}) \cdot (\bar{y} + z)$$

1
0
0
1
1
1
0
1

Laws :

$$1. \quad \begin{aligned} x + x &= x \\ x \cdot x &= x \end{aligned}$$

idempotent

$$2. \quad \begin{aligned} x + y &= y + x \\ x \cdot y &= y \cdot x \end{aligned}$$

commutative

$$3. \quad \begin{aligned} x + (y + z) &= (x + y) + z \\ x \cdot (y \cdot z) &= (x \cdot y) \cdot z \end{aligned}$$

associative

$$4. \quad \begin{aligned} x \cdot (y + x) &= x \\ x + (y \cdot x) &= x \end{aligned}$$

absorption

$$5. \quad \begin{aligned} \overline{x + y} &= \bar{x} \cdot \bar{y} \\ \overline{x \cdot y} &= \bar{x} + \bar{y} \end{aligned}$$

de morgan's

$$6. \quad \overline{\bar{x}} = x$$

double complement

$$7. \quad \begin{aligned} x + \bar{x} &= 1 \\ x \cdot \bar{x} &= 0 \end{aligned}$$

identity

$$8. \quad \begin{array}{l} x + 1 = 1 \\ x \cdot 0 = 0 \end{array} \quad \left| \begin{array}{l} \text{unity property} \\ \text{zero property} \end{array} \right.$$

$$9. \quad \begin{array}{l} x + 0 = x \\ x \cdot 1 = x \end{array} \quad \left| \begin{array}{l} \text{domination law} \end{array} \right.$$

$$10. \quad \begin{array}{l} x \cdot (y + z) = x \cdot y + x \cdot z \\ x + (y \cdot z) = (x + y) \cdot (x + z) \end{array} \quad \left| \begin{array}{l} \text{distributive} \end{array} \right.$$

Canonical SOP and POS :

$$\text{min} = x \cdot y \cdot \bar{z}, \quad x \bar{y} z \rightarrow \text{sum} \rightarrow \text{SOP}$$

$$\text{max} = x + y + \bar{z}, \quad x + \bar{y} + z \rightarrow \text{product} \rightarrow \text{POS}$$

$$\text{Eg: } f(x, y, z) = \bar{x}$$

$$= \bar{x} + 0 \quad \text{domination}$$

$$= \bar{x} + (y \cdot \bar{y}) \quad \text{identity}$$

$$= (\bar{x} + y) \cdot (\bar{x} + \bar{y}) \quad \text{distributive}$$

$$= (\bar{x} + y + 0) \cdot (\bar{x} + \bar{y} + 0)$$

$$= (\bar{x} + y + z \bar{z}) \cdot (\bar{x} + \bar{y} + z \bar{z})$$

$$= (\bar{x} + y + z) \cdot (\bar{x} + y + \bar{z}) \cdot (\bar{x} + \bar{y} + z)$$

$$(\bar{x} + \bar{y} + \bar{z})$$

POS form

$$= \bar{x} \cdot 1$$

$$= \bar{x} \cdot (y + \bar{y})$$

$$= \bar{x} \cdot y + \bar{x} \cdot \bar{y}$$

$$= \bar{x} \cdot y \cdot 1 + \bar{x} \cdot \bar{y} \cdot 1$$

$$= \bar{x} \cdot y \cdot (z + \bar{z}) + \bar{x} \cdot \bar{y} \cdot (z + \bar{z})$$

$$= \bar{x} \cdot y \cdot z + \bar{x} \cdot y \cdot \bar{z} + \bar{x} \cdot \bar{y} \cdot z + \bar{x} \cdot \bar{y} \cdot \bar{z}$$

↓
SOP form

Eg: $(x+y) \cdot z$

$$\text{SOP} \rightarrow = (x \cdot 1 + y \cdot 1) z$$

$$= (x \cdot (y + \bar{y}) + y \cdot (x + \bar{x})) z$$

$$= (x \cdot y + x \cdot \bar{y} + y \cdot x + y \cdot \bar{x}) z$$

$$= (xy + x\bar{y} + yx + \bar{x}y) z$$

$$= xyz + x\bar{y}z + \bar{x}yz$$

$$\text{POS} \rightarrow = (x + y + 0) \cdot (z + 0 + 0)$$

$$= (x + y + z\bar{z}) \cdot (z + \cancel{x\bar{x}} + y\bar{y})$$

$$= (x + y + z) (x + y + \bar{z}) \cdot (z + x + y)$$

$$(z + \bar{x} + y)$$

$$= (z + x + \bar{y}) \cdot (x + y + y)$$

Functionally Complete set of connective

$\{+, '\}$

↓
nor

$\{., '\}$

↑
nand

↓ nor

↑ nand

Q. Show that

$$(a) \quad \bar{x} = x \downarrow x = x \uparrow x$$

$$(b) \quad xy = (x \downarrow x) \downarrow (y \downarrow y) = (x \uparrow x) \uparrow (y \uparrow y)$$

$$(c) \quad x+y = (x \downarrow y) \downarrow (x \downarrow y) = (x \uparrow x) \uparrow (y \uparrow y)$$

$$(a) \quad \bar{x} = \overline{x+x} \quad [x = x+x, \text{ idempotent}]$$

$$\bar{x} = x \downarrow x$$

$$(b) \quad xy = \overline{\overline{xy}}$$

$$= \overline{\overline{x+y}}$$

$$= \overline{x \downarrow y}$$

$$= \overline{x+x} \downarrow \overline{y+y}$$

$$= (x \downarrow x) \downarrow (y \downarrow y)$$

$$xy = \overline{\overline{xy}}$$

$$= \overline{(xy) + (xy)}$$

$$= \overline{(\bar{x}\bar{y}) + (\bar{x}\bar{y})}$$

$$= \bar{x}\bar{y} \uparrow \bar{x}\bar{y}$$

$$= (x \uparrow x) \uparrow (y \uparrow y)$$

$$\begin{aligned}
 (c) \quad x+y &= \overline{\overline{x+y}} \\
 &= \overline{\overline{x} \cdot \overline{y}} \\
 &= \overline{x} \uparrow \overline{y} \\
 &= (\overline{x \cdot x}) \uparrow (\overline{y \cdot y}) \\
 &= (x \uparrow x) \uparrow (y \uparrow y)
 \end{aligned}$$

$$\begin{aligned}
 x+y &= \overline{\overline{x+y}} \\
 &= \overline{(\overline{x+y}) \cdot (\overline{x+y})} \\
 &= \overline{(\overline{x+y}) + (\overline{x+y})} \\
 &= (\overline{x+y}) \downarrow (\overline{x+y}) \\
 &= (x \downarrow y) \downarrow (x \downarrow y)
 \end{aligned}$$

KMAP \rightarrow

$f(x, y) =$

	\overline{y}	y
\overline{x}	0 0 0	0 1 1
x	1 0 2	1 1 3

$f(x, y, z) =$

	$\overline{y}\overline{z}$	$\overline{y}z$	$y\overline{z}$	yz
\overline{x}	0 0 0 0	0 0 1 1	0 1 1 3	0 1 0 2
x	1 0 0 4	1 0 1 5	1 1 1 7	1 1 0 6

$\Sigma \rightarrow \text{SOP}$

$\pi \rightarrow \text{POS}$

↑ octat
quad
pair
single

$f(x, y, z, w)$

$x \backslash y$	$\bar{z}\bar{w}$	$\bar{z}w$	$z\bar{w}$	zw
$\bar{x}\bar{y}$	0000 0	0001 1	0011 3	0010 2
$\bar{x}y$	0100 4	0101 5	0111 7	0110 6
xy	1100 12	1101 13	1111 15	1110 14
$x\bar{y}$	1000 8	1001 9	1011 11	1010 10

Q. $f(x, y, z, w) = \Sigma 0, 1, 2, 3, 7, 8, 11, 12$

$x \backslash y$	zw	$\bar{z}\bar{w}$	$\bar{z}w$	$z\bar{w}$
$\bar{x}\bar{y}$	1 0	1 1	1 3	1 2
$\bar{x}y$			1 7	
xy	1 8			
$x\bar{y}$	1 8		1 11	

1 quad + 3 pair

$$\bar{x}\bar{y} + x\bar{z}\bar{w} + \bar{x}z\bar{w} + \bar{y}zw$$