

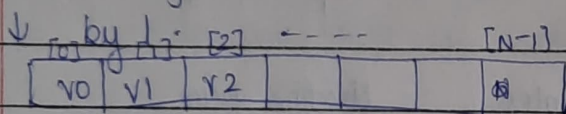
Review sheet

Q1) (a) → Reverse a string → Holding all data of a function / recursion
→ Memory management → Evaluating expression → Parsing
give example

(b) The use of 1-d array as a static stack means that the stack can become full. There are 2 ways of doing this

① To reserve the first element of the array for stack pointer, which would contain the index of the current top element in the stack. If the stack pointer is stored in array[0] then its initial value will be zero. If an element is pushed now, the stack pointer value increases to 1 and new element is stored in array[1]. $A[TOP] = \text{VALUE}$

② If the element is popped from stack, the top element in the array will be set to null and the pointer value will



Stack follows LIFO (last in first out)

isEmpty() → Reports if stack is empty or not and returns true if $TOP < 0$.

Q2) @ 4

② SET = False

is Full() → Returns true if stack is full. $TOP > N-1$ so true. return.

loop for I from 0 to N-1

loop for J from 0 to N-1

if $I = J$ then

if $MAT[I, J] \neq 0$ then

SET = True

I = 0

J = 0

end if

end if

if $MAT[I+1, J] \neq 0$ AND $MAT[I, J+1] \neq 0$ then

SET = True

$MAT[I+1, J] = 0$

$MAT[I, J+1] = 0$

end if

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if MAT[I,J] = 0 then
    SET = True
end if
end loop
end loop
output SET

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③ ③ mystery (MAT, 5)

↓
MAT[5][4] + mystery (MAT, 4)

↓
MAT[4][3] + mystery (MAT, 3)

↓
MAT[3][2] + mystery (MAT, 2)

↓
MAT[2][1] + mystery (MAT, 1)

↓
MAT[1][0] + mystery (MAT, 0)

5 + 7 + (-5) + 9 + 1
= 17

↓
0

① It finds the sum of elements in the lower subdiagonal.

Q3b) INVALID = False

R = 0

loop while R < N and not INVALID

C = 0

loop while C < N and not INVALID

if abs(R-C) ≥ 2 and A[R][C] != 0 or abs(R-C) < 2 ~~then~~
and A[R][C] == 0 then

INVALID = True

end if

C = C + 1

end while

$R = R + 1$

end while

return not INVALID