(14(a) Gaussian Noise also known as normal noise model is used frequently in practice.

The PDF of a Gaussian random variable Z is given by $-\frac{(2-\mu)^2}{2}$

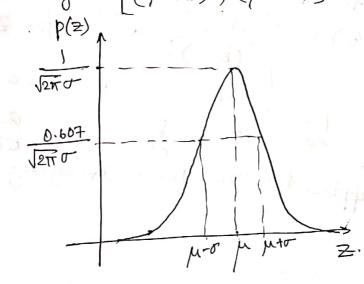
$$p(z) = \frac{1}{\sqrt{2\pi} \sigma} e^{-\frac{(z-\mu)^2}{2\sigma^2}}$$

2 = gray level.

 μ = mean of average value of 2.

T = standard deviation of Z.

In this distribution about 70% of its values will be in the range [(1-0), [1+05], and 95% in the roange [(1-20), (1+20)].



Anwian Chakraborty

· Rayleigh Noise

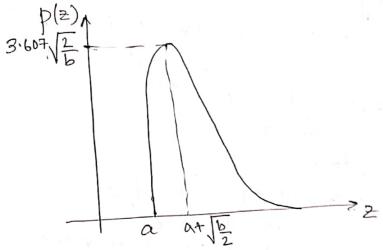
or which the writers

The PDF of the Rayleigh noise is given by
$$p(z) = \begin{cases} \frac{2}{b} (z-a) e^{-(z-a)/b}, & \text{for } z > a \\ 0 & \text{for } z < a \end{cases}.$$

$$\mu = a + \sqrt{\pi b | 4}$$

$$\sigma^2 = b (4 - \pi)$$
4.

The shape is skewed to the right



Gamma Noise This noise also known as Exlang

Gamma Noise This noise also knows for
$$2 \times a$$
.

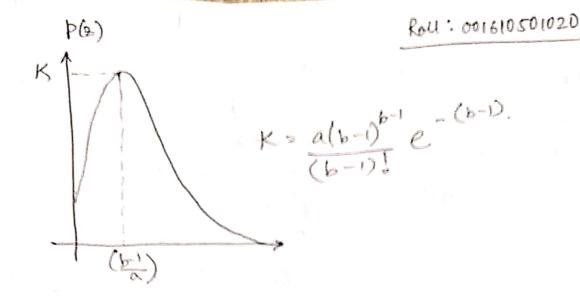
Noise has the PDF

 $(b-1)!$
 $(b-1)!$
 $(a \times b)$
 $(a \times b)$

a70, b= +ve integer

It has application in Laser Imaging.

Anucan Chakraborty



Q4(b). Harmonic Mean filter.

It is used for removing sattement promise It is given by the equation

 $f(a_1y) = \frac{mn}{\sum_{(s,t)\in S_{xy}} q(s_1t)}$

where Sny represents the set of coordinates in a rectangular subimage window of size mxn centered at the point (2,y). g(aiy) & = pixel value to corrupted image at coordinates (1,y) f(a,y) = value of restored image at coordinates (2, y).

The houmanic mean filter works well for salt noise byt fails for pepper noise. It does well also with other types of noise like Gaussian noise.

> Arman Chakraborty Scanned with CamScanner

Contraharmonic Mean filter

The contraharmonic mean filtering operation yields a restored image based on

f(x,y) =

(Sit) ESzy

where Q is called the order of the filter All other symbols have the same meaning

as in the prievious equation. This fitter is well suited for reducing or virtually eliminating the effects of saltand-pepper noise. For positive values of Q is eliminates som moise pepper noise. For regative values Q it elimenates salt noise. It cannot do both Simultaneously

if Q=0 it reduces to with metic mean filter

if Q=-4 it reduces to the harmonic mean filter.

Anwan Chakraborty

Algorithm for construction of a theffman Tree.

1. Begin

2. Create a table to store the frequency of each

character symbol.

3. Create a leaf node for each symbol and add it to the prescrity queue. Lower the probability higher the projonity. (higher forequency means lower priority)

4, while there is more than one node in the

queue:

4.1 Remove the two nodes of highest priority

from the queue

4.2 Create a new internal node with these two nodes and as children and with probability equal to the sum of two nodes probabilities.

4.3 Add the new mode to the queue

5. The remaining node is the root nocle and Deo the tree is complete-

B. End.

(1) (6) (6)

my name is anwan chakeraborty.

Frequency table				
	Frequence	4	requenc	<u>-</u>
	space n	5	·	
	Space	1	+	
	n		3 3 2 2	
	~ ~ \	1	3	
	1		2	
	e i		2	
	e		1	
	i		1	4 1
	S		1	
	S u c h		1	
			1	
			1 1 1 1	x /1=
			1	
	b		1	1
	0		1	
	t		1	

If the frequencies are same the characters are selected in the order

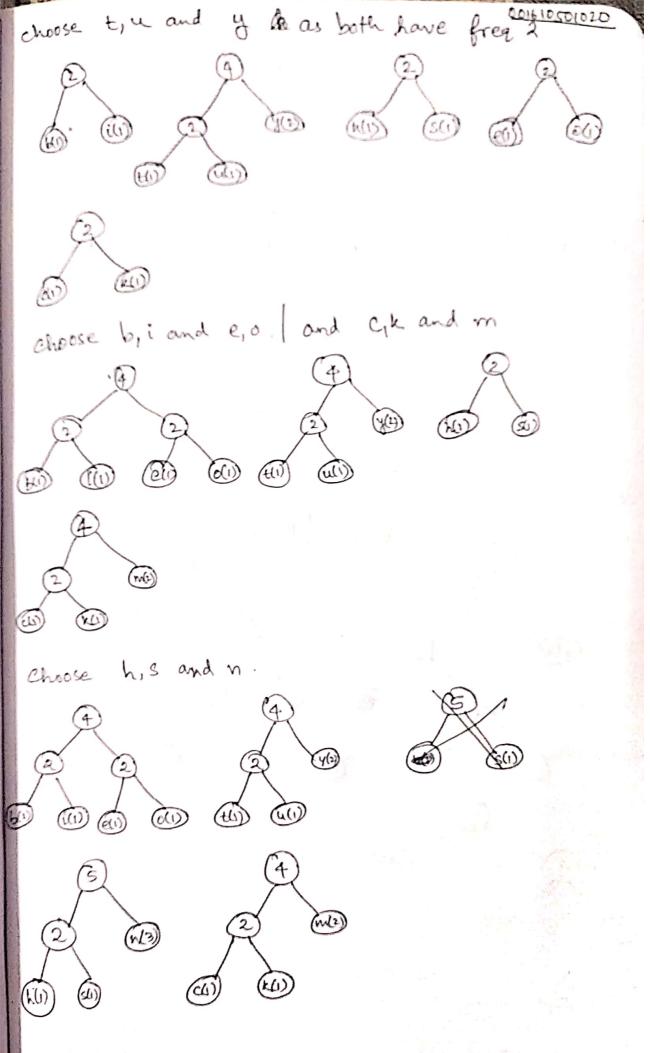


1. Select highest priority characters. choose b, i Remaining (P(1) (iii) choose to u. Par Cili (ta) enouse his. S(1) (10) choose e, o. (s(1)) (u(1) (ha) (00) (10) (10) choose c, k (30) (2(1) (un) (m)

(kn)

(ch)

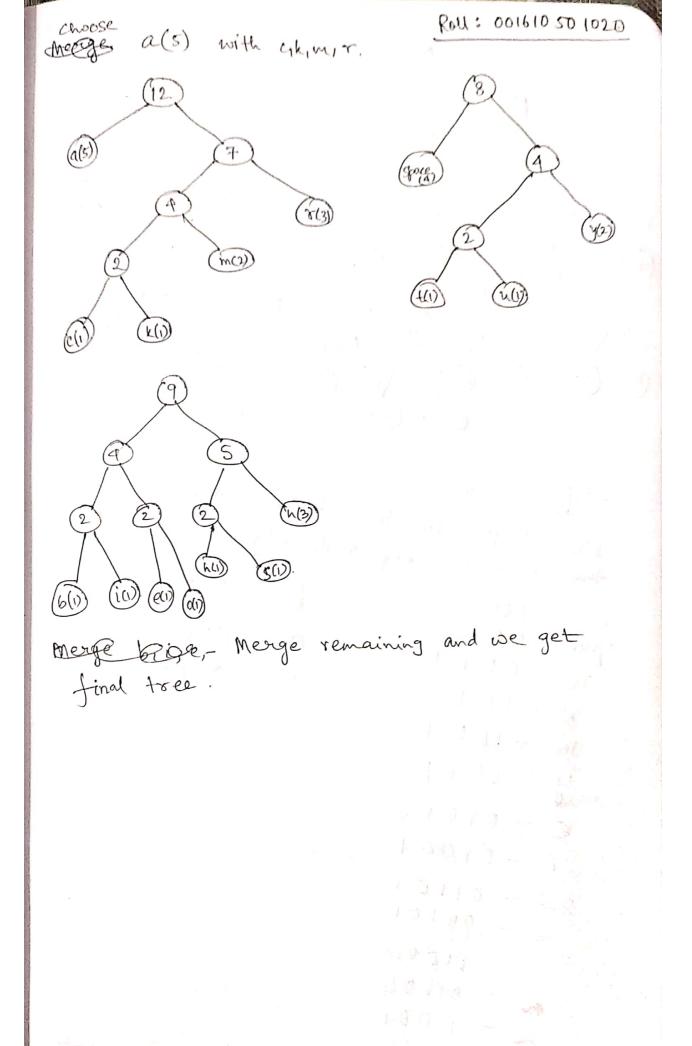


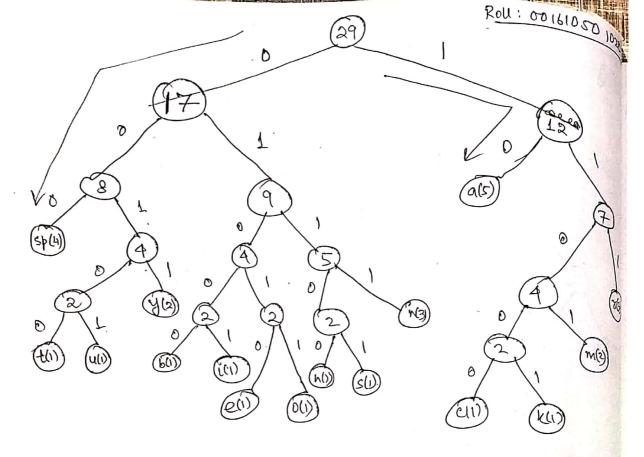


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Roll: 0016 (050 1020) Next choose cikim and r. **3**P (3) 0(1) (b(1) eui (ili) (I) (11) (WE (3(1) EU Choose space and t, y, u | Merge b, i, e, o pand h, s, n (5) (4) (y(2)) tu (i(i) Anusan Chak valoorty

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Huffman Codes.

Des To find the codes traverse the tree to the character. Left node edge is 0 else L.

a - 10 space - 000. n - 0111 ~ - 111

m -1101

y -0011

e - 01010.

1 -01001

5 - 01101

u - 00101

c - 11000. R - 01100.

- 11001

6 -01000

0 - 01011

t - 00100.

Anwan Crakraborty

my name is anwar chakuaberty

In Huffman coding two & codes of different length never have common & prefix. So in decoding the above code.

we see that the minimum length string having prefix 1101 @ at and also no other string has the same code.

1101 → m 0011 → y. 000 → space.

and so on. Substring whether it whenever we get a code use check whether it is a valid code then we is a valid code then we replace by that character as it will never be a prefix of a larger length code.

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