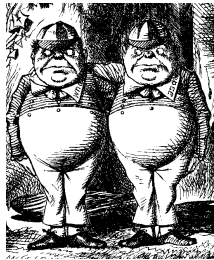


Nov 8

Simple Security Protocols continued

Mutual Authentication Attack



Trudy

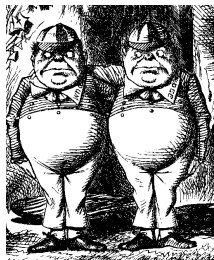
1. "I'm Alice", R_A

2. R_B , $E(R_A, K_{AB})$

5. $E(R_B, K_{AB})$



Bob



Trudy

3. "I'm Alice", R_B

4. R_C , $E(R_B, K_{AB})$

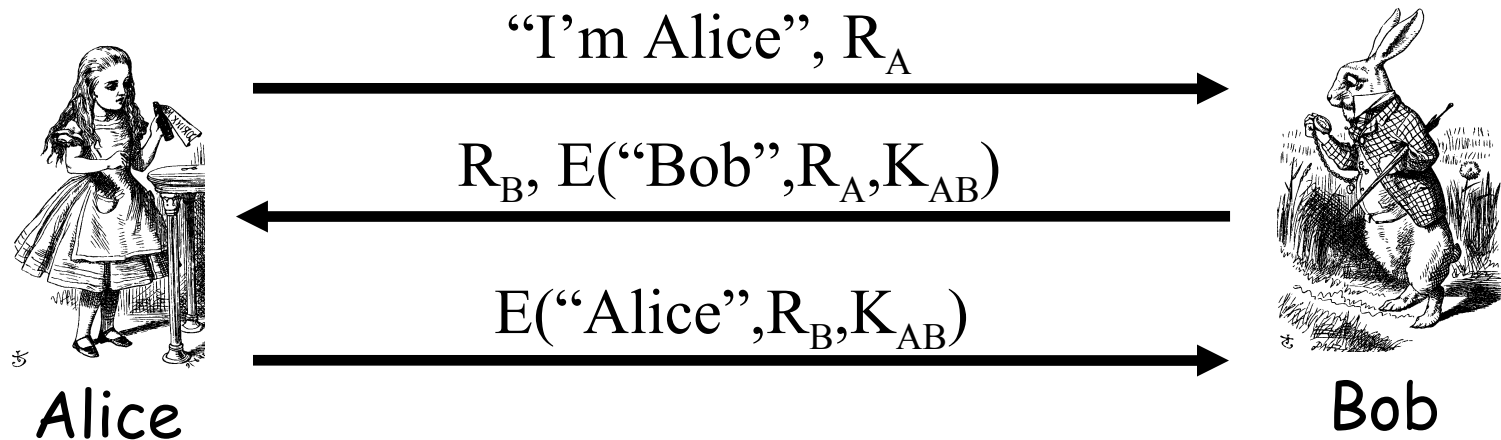


Bob

Mutual Authentication

- ❑ Our one-way authentication protocol **not** secure for mutual authentication
- ❑ Protocols are subtle!
- ❑ The “obvious” thing may not be secure
- ❑ Also, if assumptions or environment changes, protocol may not work
 - This is a common source of security failure
 - For example, Internet protocols

Symmetric Key Mutual Authentication

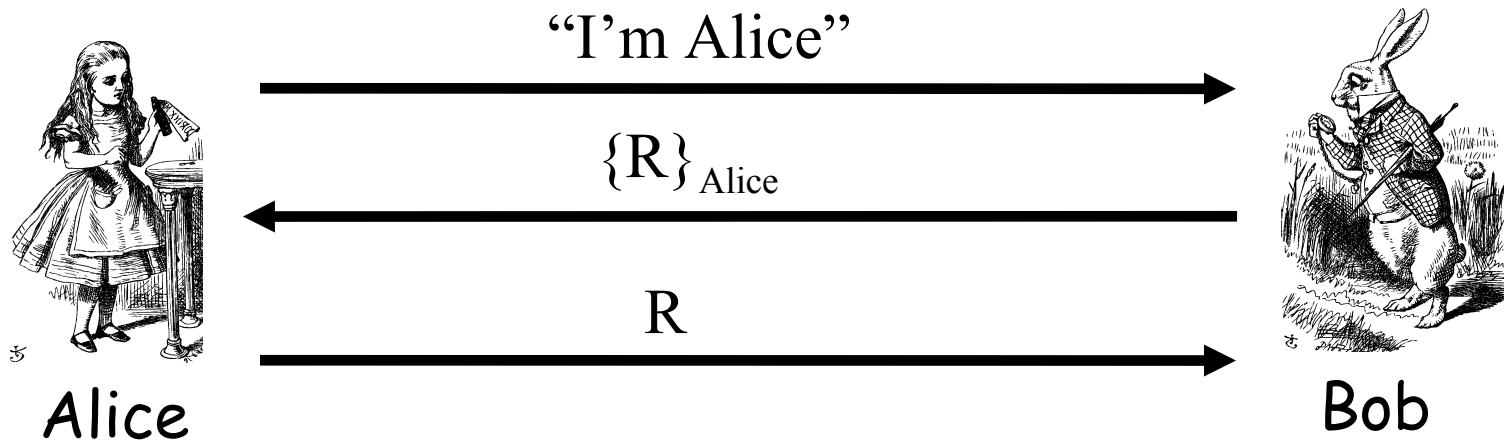


- ❑ Do these "insignificant" changes help?
- ❑ Yes!

Public Key Notation

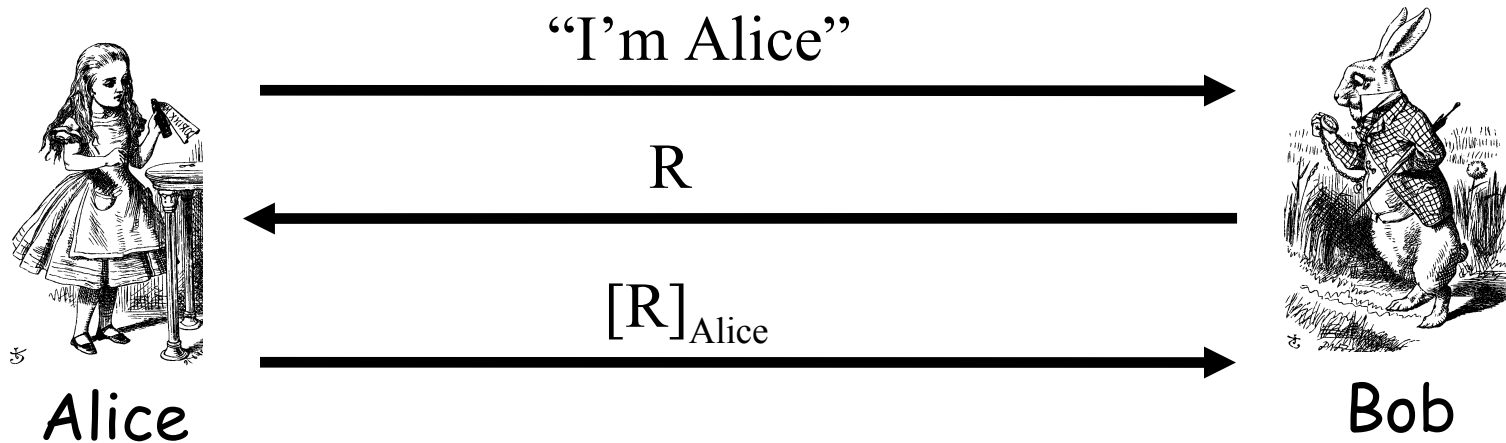
- ❑ Encrypt M with Alice's public key: $\{M\}_{\text{Alice}}$
- ❑ Sign M with Alice's private key: $[M]_{\text{Alice}}$
- ❑ Then
 - $[\{M\}_{\text{Alice}}]_{\text{Alice}} = M$
 - $\{[M]_{\text{Alice}}\}_{\text{Alice}} = M$
- ❑ **Anybody** can do **public key** operations
- ❑ Only **Alice** can use her **private key** (sign)

Public Key Authentication via Encryption



- ❑ Is Alice authenticated? Bob? Is this secure?
- ❑ Trudy can get Alice to decrypt anything! (how?)
 - Must have two key pairs (one for signature and one for encryption)

Public Key Authentication via Digital Signature



- Is this secure?
- Again ...Trudy can get Alice to sign (encrypt) anything!
 - Must have two key pairs

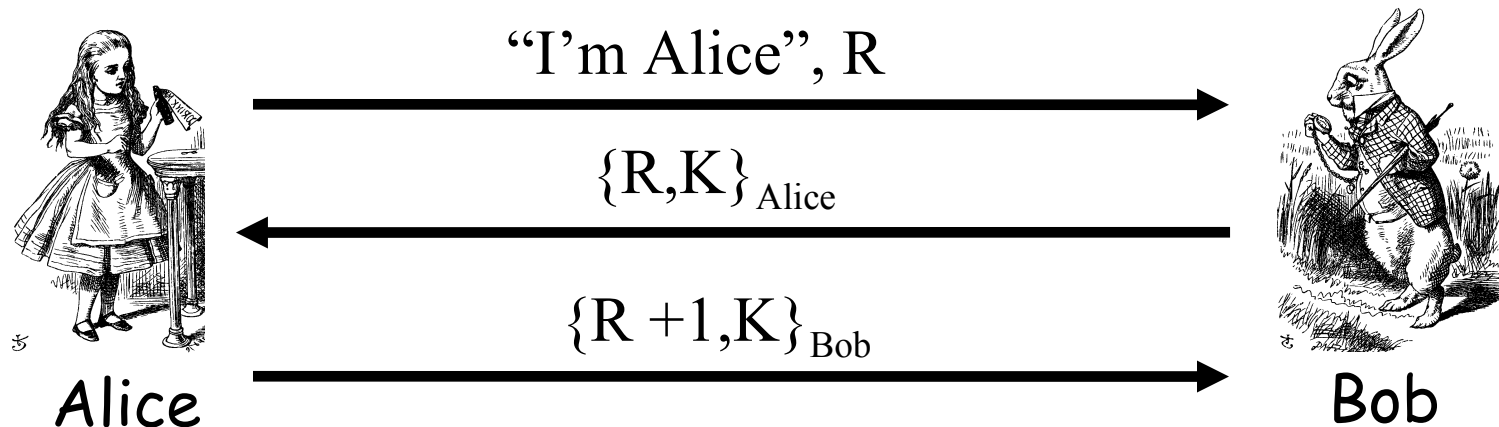
Public Keys

- ❑ Never use the same key pair for encryption and signing
- ❑ One key pair for encryption/decryption
- ❑ A different key pair for signing/verifying signatures

Session Key

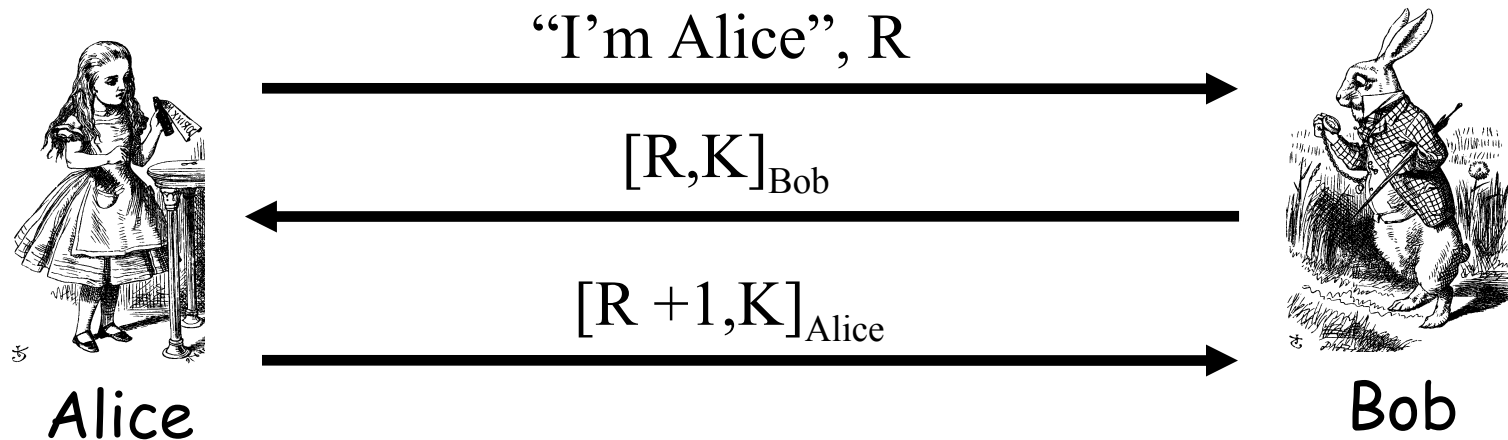
- ❑ Usually, a session key is required
 - Symmetric key for a particular session
- ❑ Can we authenticate and establish a shared symmetric key?
 - Key can be used for confidentiality
 - Key can be used for integrity
- ❑ In some cases, we may also require perfect forward secrecy (PFS)
 - Discussed later...

Authentication & Session Key



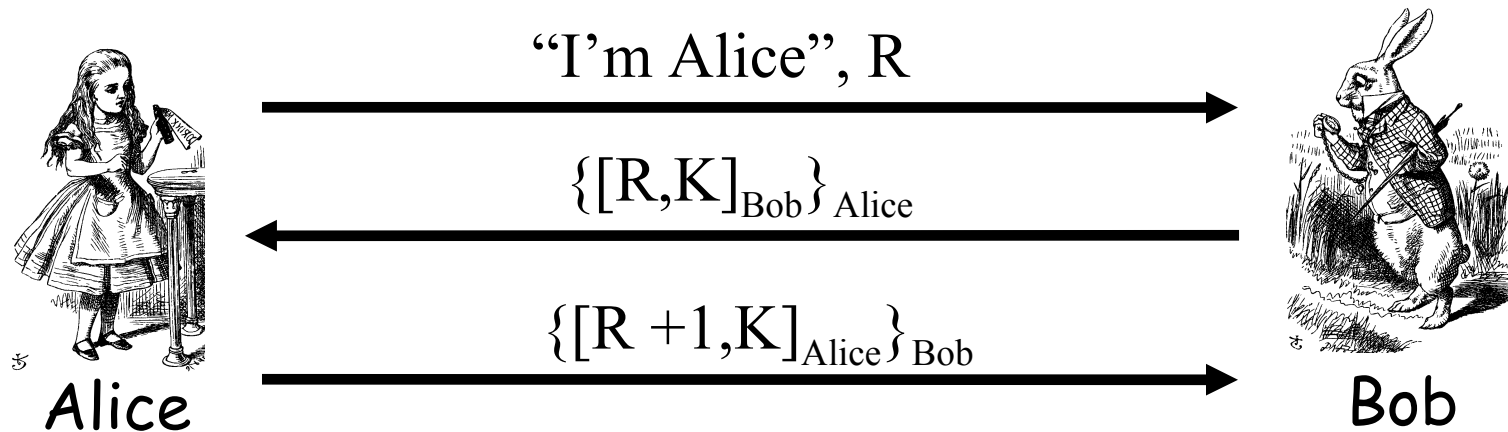
- Is this secure?
- OK for key, but no mutual authentication
- **Note** that K is acting as Bob's nonce

Public Key Authentication and Session Key



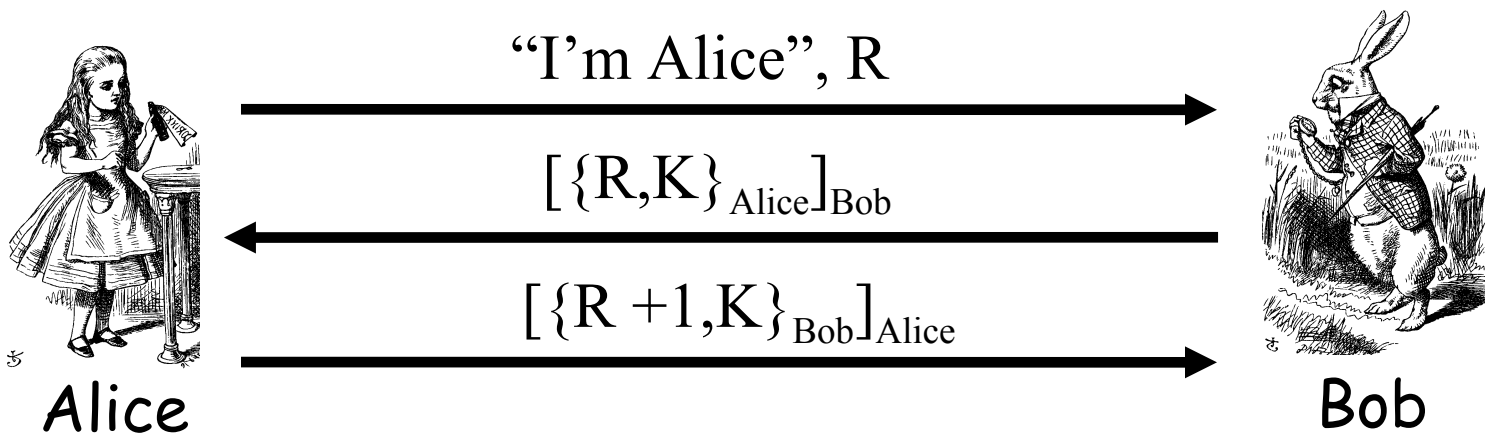
- ❑ Is this secure?
- ❑ Mutual authentication but key is not secret!

Public Key Authentication and Session Key



- ❑ Is this secure?
- ❑ Seems to be OK
- ❑ Mutual authentication and session key!

Public Key Authentication and Session Key



- ❑ Is this secure?
- ❑ Seems to be OK
 - Anyone can see $\{R, K\}_{\text{Alice}}$ and $\{R + 1, K\}_{\text{Bob}}$

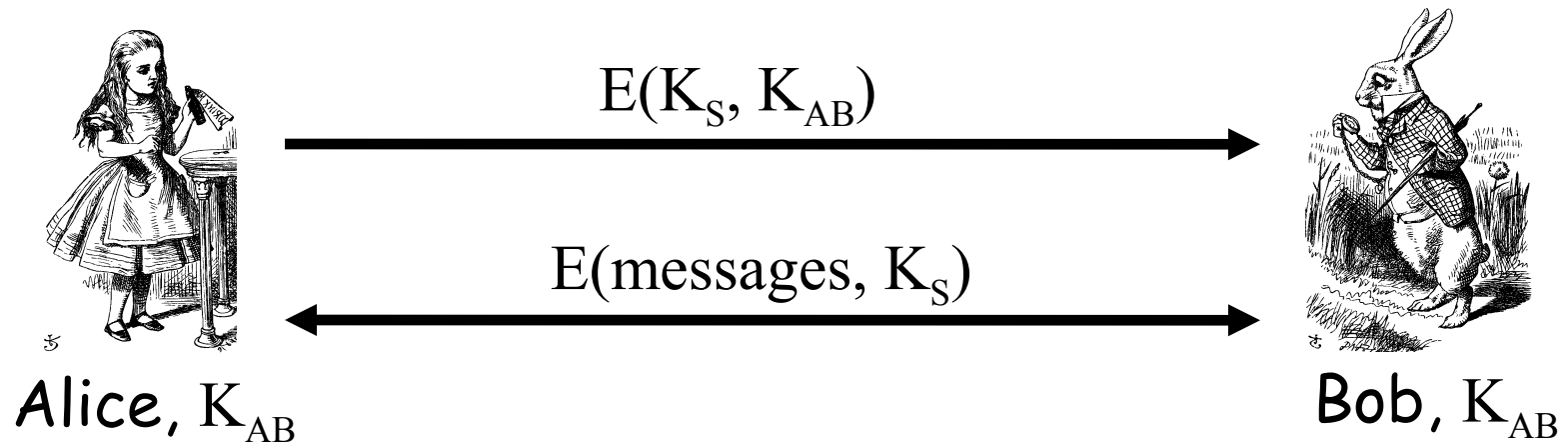
Perfect Forward Secrecy

- The concern...
 - Alice encrypts message with shared key K_{AB} and sends ciphertext to Bob
 - Trudy records ciphertext and later attacks Alice's (or Bob's) computer to find K_{AB}
 - Then Trudy decrypts recorded messages
- **Perfect forward secrecy (PFS):** Trudy cannot later decrypt recorded ciphertext
 - Even if Trudy gets key K_{AB} or other secret(s)
- Is PFS possible?

Perfect Forward Secrecy

- Suppose Alice and Bob share key K_{AB}
- For perfect forward secrecy, Alice and Bob cannot use K_{AB} to encrypt
- Instead they must use a **session key** K_s and forget it after it's used
- Problem: How can Alice and Bob agree on session key K_s and ensure PFS?

Naïve Session Key Protocol



- ❑ Trudy could also record $E(K_S, K_{AB})$
- ❑ If Trudy gets K_{AB} , she gets K_S