

10/27/22 RSA Signatures:-

$$\sigma = M^d \bmod N$$

$$\sigma^e = (M^d)^e \bmod N \quad (e \cdot d \bmod \phi(n) = 1)$$

$$= M \cdot \bmod N = \underline{\underline{M}}$$

11/01/22

$$\sigma = (\sigma_1 \cdot \sigma_2)^e \cdot \bmod N$$

$$= M \bmod N = \underline{\underline{M}}$$

(slide.no: 13)

1st change: $H(M) \rightarrow H(S, M); S \in \{0, 1\}^n$

2nd change: $H(\dots H(M))$

$$P.K = (P.K, \text{challenge set})$$

$$S.K = (S.K, \text{challenge set}).$$

Schnorr Identification Scheme:-

$$g^x \cdot y^{-c} = g^{cx+k} \cdot g^{-x \cdot c} = g^k = \underline{\underline{I}}$$

11/03/22 (slide.no: 24, 25)

DSA's ID scheme:-

Bob's verification:- (step 7 from diagram)

$$\begin{aligned} g^{xs^{-1}} \cdot y^{zs^{-1}} & \stackrel{?}{=} I = g^k \\ &= g^{xs^{-1}} \cdot g^{xzs^{-1}} \end{aligned}$$

$$= g^{s^{-1}} (x + x_2)$$

$$= g^k \cdot (x + x_2)^{-1} \cdot (x + x_2) = \underline{g^k}$$

$$\therefore S = k \cdot (x + x_2)^{-1} \pmod{q}$$

(Slide no: 26)

DSA:

Bob's verification:- (step 6 from diagram)

$$F(g^{H(m)} S^{-1} \cdot y^{x_2 S^{-1}}) \stackrel{?}{=} x_2$$

$$= F(g^{H(m)} S^{-1} \cdot g^{x_2 x_2 S^{-1}})$$

$$= F(g^{S^{-1} (H(m) + x_2 x_2)})$$

$$= F(g^{S^{-1} (H(m) + x_2 x_2)})$$

$$= F(g^{[k \cdot (H(m) + x_2 x_2)]^{-1} \cdot (H(m) + x_2 x_2)})$$

$$= F(g^k) = \underline{\underline{x_2}}$$

(Slide no: 33, 34)

Lamport Scheme:-

Ex: $m = 101$

$$s = (x_1, x_2, x_3)$$

$$A: SK_A[x_{10}, x_{11}, x_{21}, x_{20}, x_{31}]$$

