

15-01-2020

UCS 1403 - Design and analysis  
of algorithms.

Pg-1  
M. Rahul Ram,  
CSE - B,  
185001121.

1. a) Algorithm:

Let us consider that array is  
used to describe a stack of  $n$  pancakes.

Step 1: Iterate ~~through~~ from 0 to  $n-1$   
(using variable  $i$ ).

Step 2: Inside that loop, again  
iterate from 0 to  $i$ .  
(using variable  $j$ ).

Step 2.1: Inside first for loop ( $i+1$ )  
If  $\text{sizeof pancake } \text{index} < \text{sizeof pancake } (i)$ .  
(if inversion exists).

→ Flip pancakes from (0 to  $i+1$ ).  
→ (last will be first)

→ For  $k$  in range ( $i+1$ )

→ check for pancake with size  
less than the first one.

→ if found ~~intitalize~~ assign its  
index to flip-index.

→ if not assign ( $i+2$ ) to flip-index.

→ perform flip ( $0$ : flip-index - 1)

→ perform flip ( $0$ : flip-index - 2).

Step 3: Inside first for loop perform flip ( $0$ :  $i$ )

algorithm for flip: (starting always 0) pg-2

flip(start-index, end-index) :-

→ for  $i$  in range  $(end-index/2)$

→ swap( $i, end-index-i$ ).

b) consider - sequence ~~2 3 6~~

2	4	6	3	8
0	1	2	3	4

(i) inversion at 6, 3.

(ii) flipping (0, 3) → 

3	6	4	2	8
---	---	---	---	---

(iii) finding flip-index as 3.

(iv) flipping (0,  $\frac{3+1}{2}$ ) → 

4	6	3	2	8
---	---	---	---	---

(v) flipping (0,  $\frac{3-2}{1}$ ) → 

6	4	3	2	8
---	---	---	---	---

(vi) flipping (0, 3) → 

2	3	4	6	8
---	---	---	---	---

c) ~~A x 4~~  $(n-1) \times 4$

d) ~~A x 4~~  $(n-1) \times 4$ .

Here's the algorithm: / code

```
def isBalanced(s: str) :
    # w is the empty string
    if len(w) == 0 :
        return False True

    # else we create a empty stack
    stack = []

    # Now iterate traverse through the string.
    for character in w :
        if character == '(' :
            # append it to stack
            stack.append(character)
        # else -> character == ')'
        else :
            # no opening parenthesis
            if len(stack) == 0 :
                return False

            else :
                c = stack.pop()
                if c != character '(':
                    return False

    # if the control reaches here,
    # if stack is not empty
    if len(stack) != 0 :
        return False

    return True # every parenthesis is matched.
```

## Analysis

pg-4

- (i) if string is empty, return true
- (ii) return false when no matching '(' is found in the stack when ')' is found in the string.
- (iii) return false when stack has elements, after traversing through the string. Means no matching ')' for '(' found in stack.

This loop runs till the end of the string ( $n \rightarrow \text{len of string}$ ), if it is not returned false before the loop ends.

Worst case complexity:  $O(n)$ .

```
def longestSequence(w):
```

```
    if len(w) == 0:
```

```
        return 0
```

```
    stack = []
```

```
    max = 0
```

```
    length = 0
```

```
    for char in w:
```

```
        if char == '(':
```

```
            stack.append(char)
```

```
        else:
```

```
            # if no matching ')' is found,
```

```
            if len(stack) == 0:
```

```
                # make length as 0, subsequence ends
```

```
                length = 0
```

```
            # else add 2 to length ('(' and ')')
```

```
            else:
```

```
                stack.pop()
```

```
                length += 2
```

```
            # if length is greater than max
```

```
            if length > max:
```

```
                # assign max to length
```

```
                max = length
```

```
    # return max
```

```
    return max
```

## Analysis

pg-6

(i) if string is empty, length of longest subsequence = 0.

(ii) Here whenever a matching pair of '(' and ')' is found we increment ~~it~~ length and check if it is greater than max. if yes, assign length to max.

(iii) if no matching '(' found for ')', then the subsequence ends. we again start from (length) = 0.

(iv) finally, we return max.

Worst - time complexity =  $O(n)$ .



3)

Pg-7

Dimension of envelope is  $x \times y$ .

There are two cases of each envelope choose it or not select it.

Use dp array to hold maximum number of nests per envelope, then.

$$dp[x] = \max(dp[e_1] + 1, dp[e_2] + 1, dp[e_3] + 1, \dots)$$

where  $x \rightarrow$  envelopes  $e_1, e_2, e_3$ .

$dp[e_i] \rightarrow$  corresponds largest envelope representing  $n$ , the number of nests.

Algorithm:

- 1) sort the envelop,  
~~one~~ based on  $x_i$   
~~another~~ based on  $y_i$ .
- 2) dp array will hold the maximum value of each envelop, iterating over the array and update it.
- 3) when iterating, we need to check if all envelopes after this can be nested within it.
- 4) return max value.